## MACHINE LANGUAGE AND THE 8563

rickmk.com/rmk/Com/8563.html
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This shows how to take whatever is in the memory area 8192-16191 (the normal HIRES screen on the 128), such as a DOODLE picture (which can be loaded with LOAD "DDfilename", B0, p7168), or a picture made by 128-BASIC, and put it up on the RGB screen.

The first part of this file shows to do it with a C-128 with only 16K video memory (unenhanced). The picture will be in black-and-white. The second part of this file shows how to do the same thing with 64 K video memory (as in a C-128D), to produce a color picture.

To run the program, simply BLOAD"HIRES80", B15, P4864, make sure something is in the HIRES bitmap at 8192, switch to 80-columns, and type SYS4864, or whatever start address you load the program in at (it is completely relocatable: you can BLOAD"HIRES80", B15,P5000, the SYS5000, or use any available starting address, but you Must remain in Bank 15!). Then you can switch back to the Text screen simply by pressing the RUN/STOP key.

This was written using only the 128's built-in ML monitor. I will explain each step.

The first thing we will want to do is switch to FAST mode. This is a simple process. (This is actually simply BASIC's FAST routine). First we set Bit 4 in the Control Register at 53265 (HEX=\$D011) to 0, by reading the value there with LDA and then ANDing the value there with Binary 11101111 (HEX=\$6F), which turns off Bit 4 but does not affect anything else. This is the Bit which blanks the 40 -column screen.

| 1300 | AD | 11 | D0 | LDA | \$D011 |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 1303 | 29 | $6 F$ |  | AND | $\# \$ 6 F$ |
| 1305 | 8D | 11 | D0 | STA | $\$ D 011$ |

Then we go to the Clock Rate Register at 53296 (HEX=\$D030) and turn on Bit 0. This changes the Clock Rate from 1 MHz (SLOW) to 2 MHz (FAST). The other bits are unimportant and need not be retained, so we can simply put a 1 in that register.

| 1308 | A9 | 01 |  | LDA | $\# \$ 01$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $130 A$ | $8 D$ | 30 | D0 | STA | $\$ D 030$ |

The next step is to set up the 8563 to Bitmap display. This is controlled by Register \#25 (HEX=\$19). The Register will also turn off Attribute memory, which we must also do. Attribute memory controls individual character colors, as well as Reverse and Blink. Since we will not have any room available for any attributes, they must be turned off.

To read or write to the 8563, we use the doorway at memory locations 54784 and 54785 (HEX=\$D600 and \$D601). It is only at these two memory locations that we can communicate with the 8563 Video Chip. This is done by storing the Register we want to read or write to $\$ D 600$, then waiting
for Bit 7 at $\$ \mathrm{D} 600$ to be set to 1 (BPL will branch while Bit 7 is off), indicating that $\$ \mathrm{D} 601$ is ready for action. As soon as that happens, we go to work on $\$ D 601$, reading or writing our value.

Register $\$ 19$ will contain a different value depending on the particular 8563 chip in the particular 128 being used (there are two versions). Therefore, we cannot simply put a value there with Attributes-Off and Bitmap-On. Instead, we must first read the value that is there.

| $130 D$ | A2 | 19 |  | LDX | $\# \$ 19$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $130 F$ | $8 E$ | 00 | D6 | STX | $\$ D 600$ |
| 1312 | 2C | 00 | D6 | BIT | $\$ D 600$ |
| 1315 | 10 | FB |  | BPL | $\$ 1312$ |
| 1317 | AD | 01 | D6 | LDA | $\$ D 601$ |

Then we will put this normal value for Register \$19 away. Let's use location $\$ 16$, a Zero-page location normally used by BASIC, but available to us now since we're not using BASIC.

131A $8516 \quad$ STA $\$ 16$

Now we must turn off Attributes and turn on Bitmap. Bits 6 and 7 of the register control these two 8563 things. To turn on the Bitmap, we must set Bit 7 to 1. We'll do this by ORing it with Binary 10000000 (HEX=\$80) to turn on Bit 7 without affecting any other bit. To turn off the Attributes, we've got to get Bit 6 to 0 . Let's do that by ANDing it with Binary 10111111 ( $\mathrm{HEX}=\$ \mathrm{BF}$ ), which will turn Bit 6 off without affecting any other bit.

| 131 C | 09 | 80 | ORA | \#\$80 |
| :--- | :--- | :--- | :--- | :--- |
| 131 E | 29 | BF | AND | \#\$BF |

Now we will use BASIC's built-in routine to write a vaue to a register, which is at 52684 (HEX=\$CDCC). To do this, the register must be in the X-register, and the value to write must be in the Accumulator. The routine writes the register in $X$ to $\$ D 600$, waits for Bit 7 of $\$ D 600$ to be on, then writes the value in the Accumulator to \$D601. (The X-register still contains \$19.)

132020 CC CD JSR \$CDCC

The next thing we must do is set up the 8563 to start reading data, and storing it at location $\$ 000$ in the 8563 's RAM memory. We have to start at $\$ 0000$ because we need every byte from $\$ 0000$ to the highest location, $\$ 3 F F F$. The Screen Memory's location, which is controlled by Registers 12 and 13 (HEX=\$0C and \$0D) are already set to 0 to start the screen at $\$ 0000$, so we don't have to worry about them. But we do need to set the Current Memory Address registers 18 and 19 (HEX=\$12 and \$13) to \$0000, so when we start writing to the 8563 's RAM, the data will start at $\$ 0000$. Once these have been set for the address at which we want to start writing, it is updated automatically for each byte we write (nice of it to
do that for us, isn't it?). (Unlike all other addresses with the 6502 or 6510 or 8510 microprocessors, addresses in the 8563 are written high-byte First, then low-byte, the opposite order. In this case, though, since the high- and low-bytes of the address are the same, this peculiarity is not visible here.) To do this, we will again use our built-in write-to-the-8563-chip routine at $\$ C D C C$. We will set the Accumulator to 0 and the X-register to $\$ 12$. Then we will simply use INX to increase the X-register to write the low-byte to \$13.

| 1323 | A9 | 00 |  | LDA | \# \$00 |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 1325 | A2 | 12 |  | LDX | $\# \$ 12$ |
| 1327 | 20 | CC | CD | JSR | \$CDCC |
| $132 A$ | E8 |  |  | INX |  |
| $132 B$ | 20 | CC | CD | JSR | \$CDCC |

The next byte in the program is just an NOP, separating the preliminaries from the actual main routine to write the one memory into the other.

## 132E EA NOP

Now the real work begins! The hardest part of this project is reading the VIC HIRES screen, because the VIC stores Bitmap memory in vertical blocks of 8 bytes in horizontal rows of 40 blocks across, whereas the 8563 uses 80 sequential bytes across each row. What we must do is read bytes from the VIC Bitmap in horizontal rows, and not sequentially. Here is a BASIC way to do this: FOR $A=8192$ TO 16191 STEP 320: FOR B $=0$ TO 7: FOR $\mathrm{C}=0$ TO 312 STEP 40: X = PEEK (A+B+C) : NEXT C,B,A

That's what we have to do in Machine Language! Not an easy task. It is worth it, though, Because this BASIC routine is very slow. I decided to convert this exact routine into machine language. Here's how I did it:

I picked a couple of BASIC's Zero-page locations for a counter for 8192 to 16191. I chose locations $\$ 10$ and $\$ 11$, and stored 8192 ( $\mathrm{HEX}=\$ 2000$ ) there to begin.

| 132 F | A9 | 20 | LDA \#\$20 |
| :--- | :--- | :--- | :--- | :--- |
| 1331 | 85 | 11 | STA $\$ 11$ |
| 1333 | A9 | 00 | LDA $\# \$ 00$ |
| 1335 | 85 | 10 | STA $\$ 10$ |

Here is where the biggest loop begins! We're going to copy the address stored in $\$ 10$ \& $\$ 11$ to a couple more of BASIC's Zero-page locations (I'm sure BASIC won't mind), \$12 and \$13. This way, we can update the address in groups of 8 until we get to the end of a 40 -byte row, and then update the address by 320 to go on to the next row.

| 1337 | A5 | 10 |  | LDA | $\$ 10$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 1339 | 85 | 12 |  | STA | $\$ 12$ |
| 133B | A5 | 11 |  | LDA | $\$ 11$ |
| 133D | 85 | 13 |  | STA | \$13 |

I'm going to use another Zero-page location, $\$ 14$, to count from 0 to 7. This will be used as a Y-index. We have to have a memory location to store it, because we'll be needing the Y-register later on.

| $133 F$ | A9 | 00 | LDA | \#\$00 |
| :--- | :--- | :--- | :--- | :--- |
| 1341 | 85 | 14 | STA | $\$ 14$ |

One more of BASIC's Zero-page locations, \$15, we'll use to count from 1 to 40 , to tell us when we've reached the end of a HIRES row. But let's use it as a countdown, starting by putting 40 (HEX=\$28) and wait 'till we reach zero.

| 1343 | A9 | 28 | LDA | \# $\$ 28$ |
| :--- | :--- | :--- | :--- | :--- |
| 1345 | 85 | 15 | STA | $\$ 15$ |

Now let's get started, at last! We'll put the value in $\$ 14$ (something between 0-7) into the Y-register, and read a byte from the VIC HIRES screen at 8192-16191.

| 1347 | A4 | 14 | LDY | $\$ 14$ |
| :--- | :--- | :--- | :--- | :--- |
| 1349 | B1 | 12 | LDA | $(\$ 12), Y$ |

Now comes the really interesting part. Since the 8563 Bitmap display is $640 X 199$ bits and the VIC Bitmap is only $320 X 199$ bits, every bit of the VIC display must be doubled to fill the entire 8563 Bitmap. This is one of those rare cases of numerical manipulations which are actually much easier to do in Machine Language than in BASIC! What will happen is each byte will be doubled into two byte, one with each of the first 4 bits doubled, and the other with the other 4 bits of the original number doubled (8 bits doubled $=16$ bits, or 2 bytes).

The first thing we'll do is set up the X-register as a counter for the high 4 bits and then the low 4 bits of each byte (that is, we must run through this routine twice: once for each set of 4 bits). This will also serve as a Zero-page Index!

134B A2 01 LDX \#\$01

Now, we'll copy the byte into two Zero-page locations. Let's give BASIC a break and use \$FD and \$FE.

| $134 D$ | 85 | $F D$ | STA | $\$ F D$ |
| :--- | :--- | :--- | :--- | :--- |
| $134 F$ | 85 | FE | STA | $\$ \mathrm{FE}$ |

Now we'll put a 4 in the Y-register to count down the 4 bits to double at a time.

1351 AO 04 LDY \#\$04

Now is the time to start doubling! We'll use the ML instruction ROL. The

Carry will hold whether the bit is 0 or 1 . We don't have to worry about CLC or ASL, because we'll be using all 8 bits, and it won't matter what ends up in $\$ F D$ and $\$ F E$ when we're done! As each bit is rolled out into the carry, we'll roll it into a Zero-page location. The 2 bytes which will hold the final two bytes will be \$FB and \$FC, pointed to by the X-register! We have the byte in two memory locations, so it will be a simple matter to double the bit simply by rolling each bit out twice.

| 1353 | 26 | FD | ROL \$FD |
| :--- | :--- | :--- | :--- |
| 1355 | 36 | FB | ROL \$FB, X |
| 1357 | 26 | FE | ROL $\$ F E$ |
| 1359 | 36 | FB | ROL \$FB, X |
| $135 B$ | 88 |  | DEY |
| $135 C$ | D0 | F5 | BNE \$1353 |

Now the top four bits of $\$ F D$ and $\$ F E$ contain what used to be the low 4 bits. It makes no difference what's now in the 4 low bits of \$FD and \$FE, we'll never see them.

| $135 E$ | CA | DEX |  |
| :--- | :--- | :--- | :--- |
| 135 F | FO | FO | BEQ $\$ 1351$ |

Now \$FC and \$FB contain the value from the VIC screen with each bit doubled. So now we're ready to write two bytes to the 8563 Bitmap. The Memory read/write gateway to the 8563 is Register 31 (HEX=\$1F). Again, we'll use the routine at $\$ C D C C$ to write the two bytes to the 8563 RAM memory, and thereby put them up on the Bitmap display.

| 1361 | A2 | 1F |  | LDX | \# $\$ 1 F$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 1363 | A5 | FC |  | LDA | \$FC |
| 1365 | 20 | CC | CD | JSR | \$CDCC |
| 1368 | A5 | FB |  | LDA | \$FB |
| $136 A$ | 20 | CC | CD | JSR | \$CDCC |

Time now to go to the next horizontal byte of the VIC Bitmap. This byte is 8 away from the previous byte. We'll do this with a simple addition routine, adding 8 to the base address in \$12 \& \$13.

| $136 D$ | 18 |  | CLC |  |
| :--- | :--- | :--- | :--- | :--- |
| $136 E$ | A9 | 08 |  | LDA $\# \$ 08$ |
| 1370 | 65 | 12 |  | ADC $\$ 12$ |
| 1372 | 85 | 12 |  | STA $\$ 12$ |
| 1374 | A9 | 00 |  | LDA $\# \$ 00$ |
| 1376 | 65 | 13 |  | ADC $\$ 13$ |
| 1378 | 85 | 13 |  | STA $\$ 13$ |

Now we'll use our countdown counter we've set up in $\$ 15$ to see if we've reached the end of a row yet, and look back if we haven't.

137A C6 15 DEC \$15

Now we've got one line of bytes. We have 8 more lines of 40 bytes each to get. Remember, we use location $\$ 14$ to store our Y-index offset. Each byte of each row will be 1 byte higher than the previous row. First, we're getting the first byte of each 8-byte block, so we have one horizontal row. The next horizontal row will consist of the second byte of each 8 -byte block. The third row will be the third byte of each block, and so on until we've reached all 8 bytes. As we go to each byte of the block, we must maintain the base address as the first byte of the first block, so our offset will point to the right byte.

| $137 E$ | E6 | 14 | INC |  |
| :--- | :--- | :--- | :--- | :--- |
| 1380 | A5 | 10 | LDA |  |
| \$10 |  |  |  |  |
| 1382 | 85 | 12 |  | STA $\$ 12$ |
| 1384 | A5 | 11 |  | LDA |
| 1386 | 85 | 13 |  | STA $\$ 13$ |

Now we check to see if we've finished the 8-byte block.

| 1388 | A5 | 14 | LDA | 14 |
| :--- | :--- | :--- | :--- | :--- |
| 138 A | C9 | 08 | CMP | \# \$08 |
| 138 C | D0 | B5 | BNE | \$1343 |

Now the time has come to jump to the next row of 8 -byte blocks. We do this by adding 320 ( $\mathrm{HEX}=\$ 0140$ ) to the base address. This is a simple addition routine.

| 138 E | 18 |  | CLC |  |
| :--- | :--- | :--- | :--- | :--- |
| 138 F | A9 | 40 | LDA $\# \$ 40$ |  |
| 1391 | 65 | 10 | ADC $\$ 10$ |  |
| 1393 | 85 | 10 | STA $\$ 10$ |  |
| 1395 | A9 | 01 |  | LDA $\# \$ 01$ |
| 1397 | 65 | 11 | ADC $\$ 11$ |  |
| 1399 | 85 | 11 | STA \$11 |  |

Now we check to see if we're finished. We'll be finished when we have worked our way up to location 16191 ( $\mathrm{HEX}=\$ 3 \mathrm{~F} 3 \mathrm{~F}$ ). We already have the high byte in the accumulator, so we'll check and see if that's $\$ 3 F$ yet. If it is, then we'll check to see if the low byte is higher than $\$ 3 F$.

| 139B | C9 | 3 F | CMP \#\$3F |
| :--- | :--- | :--- | :--- | :--- |
| 139D | D0 | 98 | BNE $\$ 1337$ |
| 139F | A5 | 10 | LDA $\$ 10$ |
| 13A1 | C9 | 3 F | CMP $\# \$ 3 \mathrm{~F}$ |
| 13A3 | 90 | 92 | BCC $\$ 1337$ |

Hurrah! The VIC Bitmap is now on display on the 8563 RGB screen! Now we'll keep this on the screen, until the STOP key has been pressed. Location 145 ( $\mathrm{HEX}=\$ 91$ ) is constantly updated by the Kernal to contain the
value of the column of the Keyboard scan which has the STOP key. Bit 7 is cleared to 0 whenever the STOP key is pressed. BMI loops as long as Bit 7 is 1, so it will loop until the STOP key is being pressed.

| $13 A 5$ | A5 | 91 | LDA | $\$ 91$ |
| :--- | :--- | :--- | :--- | :--- |
| $13 A 7$ | 30 | FC | BMI \$13A5 |  |

The Bitmap is nice to look at, but eventually we'll want to be able to see characters on the screen again! Remember when we stored the original value of Register $\$ 19$ in location $\$ 16 ?$ Well, here's where we finally use it! This will turn the Bitmap off and turn the Attributes back on again.

| $13 A 9$ | A2 | 19 |  | LDX | \# $\$ 19$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $13 A B$ | A5 | 16 |  | LDA | $\$ 16$ |
| $13 A D$ | 20 | CC | CD | JSR | \$CDCC |

But we've still got us a problem here! Since we used the entire 16K to display our Bitmap, we have overwritten the entire 8563 character-set. We know have the text screen set up, but no character data to be able to print characters to print to it! The Kernal will come to our rescue here. When the 8563 is first initialized, VIC's character set is copied into the 8563 RAM, in the character memory storage area (which, by the way, is located at 8192 to 16383 ( $\mathrm{HEX}=\$ 2000$ to $\$ 3 \mathrm{FFF}$ ) in the 8563 's RAM memory). The 128 Kernal's Jump Table has a entry called INIT80, which carries out this copying procedure, at 65378 (HEX=\$FF62). It jumps to the actual routine which is in in the Screen Editor ROM at 49191 (HEX=\$C027).

13B0 2062 FF JSR \$FF62

Now the Text screen is being displayed by the 8563, and we have all our characters in memory so we can display them. But there's still one thing wrong. Attribute memory is turned on, but is filled with strange data, whatever was in the Bitmap display from 2048 to 4096 (HEX=\$0800 to \$1000) which is where the 8563 stores the Attributes in its memory. That's why you see that bizarre display for an instant after the Bitmap is switched out. Attribute memory is easiest to clear by simply clearing the screen, by printing the CLR/HOME character of CHR\$(147) (HEX=\$93) through the Kernal output routine at \$FFD2.

```
13B3 A9 93 LDA #$93
13B5 20 D2 FF JSR $FFD2
```

And now we're finished! Back to BASIC. Bye bye!

13B8 60 RTS

Switch to FAST mode

| 1300 | AD | 11 | D0 | LDA | $\$ D 011$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 1303 | 29 | $6 F$ |  | AND | $\# \$ 6 F$ |
| 1305 | 8D | 11 | D0 | STA | $\$ D 011$ |
| 1308 | A9 | 01 |  | LDA | $\# \$ 01$ |
| 130 A | 8D | 30 | D0 | STA | $\$ D 030$ |

Now we set up the VDC for BITMAP display, but do NOT turn off attributes!

130D A2 19 LDX \#\$19
130F 8E 00 D6 STX \$D600
1312 2C 00 D6 BIT \$D600
$1315 \quad 10$ FB BPL \$1312
1317 AD 01 D6 LDA \$D601
131A $8516 \quad$ STA $\$ 16$
131C 0980 ORA \#\$80
131E 20 CC CD JSR \$CDCC

Now, change the screen memory from $\$ 0000-\$ 4000$ up to $\$ 8000 € \$ C 000$. The screen memory is determined by the values in registers 12 and 13 (\$0C \& \$OD) in high-byte/low-byte format:

| 1321 | A9 | 80 |  | LDA | \# $\$ 80$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 1323 | A2 | 0C |  | LDX | \# $\$ 0 \mathrm{C}$ |
| 1325 | 20 | CC | CD | JSR | \$CDCC |
| 1328 | E8 |  |  | INX |  |
| 1329 | A9 | 00 |  | LDA | \# $\$ 00$ |
| $132 B$ | 20 | CC | CD | JSR | $\$ C D C C$ |

Now we set the current write-address to $\$ 8000$ by putting that value into registers 18 \& 19:

| 132 E | A9 | 80 |  | LDA | \# $\$ 80$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 1330 | A2 | 12 |  | LDX | $\# \$ 12$ |
| 1332 | 20 | CC | CD | JSR | $\$ C D C C$ |
| 1335 | A9 | 00 |  | LDA | $\# \$ 00$ |
| 1337 | E8 |  |  | INX |  |
| 1338 | 20 | CC | CD | JSR | $\$ C D C C$ |

No we convert the VIC bitmap to the VDC bitmap, just like in the 128 version

| $133 B$ | A9 | 20 | LDA $\# \$ 20$ |  |
| :--- | :--- | :--- | :--- | :--- |
| $133 D$ | 85 | 11 |  | STA |
| 133F | A9 | 00 |  | LDA | \# \$00


| 134B | A9 | 00 | LDA | \#\$00 |
| :---: | :---: | :---: | :---: | :---: |
| 134D | 85 | 14 | STA | \$14 |
| 134 F | A9 | 28 | LDA | \#\$28 |
| 1351 | 85 | 15 | STA | \$15 |
| 1353 | A4 | 14 | LDY | \$14 |
| 1355 | B1 | 12 | LDA | (\$12), Y |
| 1357 | A2 | 01 | LDX | \# \$01 |
| 1359 | 85 | FD | STA | \$FD |
| 135B | 85 | FE | STA | \$FE |
| 135D | A0 | 04 | LDY | \#\$04 |
| 135F | 26 | FD | ROL | \$FD |
| 1361 | 36 | FB | ROL | \$FB, X |
| 1363 | 26 | FE | ROL | \$FE |
| 1365 | 36 | FB | ROL | \$FB, X |
| 1367 | 88 |  | DEY |  |
| 1368 | D0 | F5 | BNE | \$135F |
| 136A | CA |  | DEX |  |
| 136B | F0 | F0 | BEQ | \$135D |
| 136D | A2 | 1F | LDX | \#\$1F |
| 136F | A5 | FC | LDA | \$FC |
| 1371 | 20 | CC CD | JSR | \$CDCC |
| 1374 | A5 | FB | LDA | \$FB |
| 1376 | 20 | CC CD | JSR | \$CDCC |
| 1379 | 18 |  | CLC |  |
| 137A | A9 | 08 | LDA | \# \$08 |
| 137C | 65 | 12 | ADC | \$12 |
| 137E | 85 | 12 | STA | \$12 |
| 1380 | A9 | 00 | LDA | \# \$00 |
| 1382 | 65 | 13 | ADC | \$13 |
| 1384 | 85 | 13 | STA | \$13 |
| 1386 | C6 | 15 | DEC | \$15 |
| 1388 | D0 | C9 | BNE | \$1353 |
| 138A | E6 | 14 | INC | \$14 |
| 138C | A5 | 10 | LDA | \$10 |
| 138E | 85 | 12 | STA | \$12 |
| 1390 | A5 | 11 | LDA | \$11 |
| 1392 | 85 | 13 | STA | \$13 |
| 1394 | A5 | 14 | LDA | \$14 |
| 1396 | C9 | 08 | CMP | \# \$08 |
| 1398 | D0 | B5 | BNE | \$134F |
| 139A | 18 |  | CLC |  |
| 139B | A9 | 40 | LDA | \# \$ 40 |
| 139D | 65 | 10 | ADC | \$10 |
| 139F | 85 | 10 | STA | \$10 |
| 13A1 | A9 | 01 | LDA | \#\$01 |
| 13A3 | 65 | 11 | ADC | \$11 |
| 13A5 | 85 | 11 | STA | \$11 |
| 13A7 | C9 | 3 F | CMP | \#\$3F |
| 13A9 | D0 | 98 | BNE | \$1343 |
| 13AB | A5 | 10 | LDA | \$10 |


| $13 A D$ | $C 9$ | $3 F$ | CMP | \# $\$ 3 \mathrm{~F}$ |
| :--- | :--- | :--- | :--- | :--- |
| $13 A F$ | 90 | 92 | BCC | $\$ 1343$ |

Now here is where the color is added. First we move Attribute memory from its normal location at $\$ 0800$ up to the unused area at $\$ 1000$. This move is done so that the old attributes will be preserved when we switch back to the text display. The location of Attribute memory is determined by the values in registers 20 \& 21 (\$14 \& \$15)

| 13B1 | A9 | 10 | LDA | \#\$ |
| :---: | :---: | :---: | :---: | :---: |
| 13B3 | A2 | 14 | LDX | \#\$14 |
| 13B5 | 20 | CC CD | JSR | \$CDC |
| 13B8 | A9 | 00 | LDA | \#\$00 |
| 13BA | E8 |  | INX |  |
| 3B |  |  |  |  |

Reset the current write to memory location

| 13BE | A9 | 10 | LDA | \# \$10 |  |  |  | JSR | \$CDCC |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 13 C 0 | A2 | 12 | LDX | \#\$1213C2 | 20 | CC C |  |  |  |
| 13 C 5 | A9 | 00 | LDA | \#\$00 |  |  |  |  |  |
| $13 \mathrm{C7}$ | E8 |  | INX |  |  |  |  |  |  |
| 13 C 8 | 20 | CC | JSR | \$CDCC |  |  |  |  |  |

40-column color memory is stored starting at location $\$ 1 \mathrm{CO} 0$ (7168).
Let's use a dynamic routine to read the colors, by storing the address within the program

| $13 C B$ | A9 | 1C |  | LDA | \# \$1C |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 13CD | 8D | D7 | 13 | STA | \$13D7 |
| 13D0 | A9 | 00 |  | LDA | \# \$00 |
| 13D2 | 8D | D6 6 | 13 | STA | \$13D6 |

Now we get the value, then convert it to 24 -bit nibbles

| 13D5 | AD 00 1C | LDA \$1C00 |
| :---: | :---: | :---: |
| 13D8 | 48 | PHA |
| 13D9 | 4A | LSR |
| 13DA | 4A | LSR |
| 13DB | 4A | LSR |
| 13DC | 4A | LSR |
| 13DD | A8 | TAY |

There is a table of color translation
between 40-column and 80-column stored in ROM starting at \$CEC5. We'll put each nibble in the Y-register to use as an offset to get the equivalent 80-column color

```
13DE B9 5C CE LDA $CE5C,Y
13E1 85 FE STA $FE
```

| $13 E 3$ | 68 |  | PLA |  |
| :--- | :--- | :--- | :--- | :--- |
| $13 E 4$ | 29 | $0 F$ | AND $\# \$ 0 \mathrm{~F}$ |  |
| $13 E 6$ | A8 |  | TAY |  |
| $13 E 7$ | B9 9 E | CE | LDA | \$CE5C, Y |

The two halves of the byte (foreground and background colors) are combined

| 13EA | OA | ASL |
| :---: | :---: | :---: |
| 13EB | OA | ASL |
| 13EC | OA | ASL |
| 13ED | OA | ASL |
| 13EE | 05 FE | ORA \$FE |
| 13F0 | A2 1F | LDX \# 1 1F |



Check for STOP key

| 1407 | A5 | 91 | LDA | $\$ 91$ |
| :--- | :--- | :--- | :--- | :--- |
| 1409 | 30 | FC | BMI $\$ 1407$ |  |

Turn off bitmap and restore test screen

| $140 B$ | A5 | 16 | LDA | $\$ 16$ |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $140 D$ | A2 | 19 |  | LDX | $\# \$ 19$ |
| 140 F | 20 | CC | CD | JSR | $\$ C D C C$ |

Put Attribute memory back at \$0800

| 1412 | A9 | 08 |  | LDA | $\# \$ 08$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 1414 | A2 | 14 |  | LDX | $\# \$ 14$ |
| 1416 | 20 | CC | CD | JSR | \$CDCC |
| 1419 | E8 |  |  | INX |  |
| $141 A$ | A9 | 00 |  | LDA | $\# \$ 00$ |
| $141 C$ | 20 | CC | CD | JSR | $\$ C D C C$ |

Put screen memory back at $\$ 0000$

| $141 F$ | A9 | 00 | LDA | \# $\$ 00$ |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 1421 | A2 | $0 C$ | LDX | \# $\$ 0 \mathrm{C}$ |  |
| 1423 | 20 | CC | CD | JSR | \$CDCC |
| 1426 | E8 |  |  | INX |  |
| 1427 | 20 | CC | CD | JSR | \$CDCC |
| $142 A$ | 60 |  |  | RTS |  |

And return

End of file.

You can write to me at
rmkqGLphrc.grg

