

CLIVE EMBEREY & BOB TURNER

INVALUABLE UTILITIES for the COMMODORE

'The complete programmer's toolkit – essential programming aids for your micro'



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Clive Emberey and Bob Turner

Invaluable Utilities for the Commodore 64

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Introduction

This book as the title suggests, is a book of utilities for the Commodore 64. It has been written not only to provide a set of useful routines, but also to help you to begin to understand some of the more detailed workings of your 64.

We have tried to cover a reasonable spectrum and hope that through our examples you will attack areas other than those covered here with increased confidence. Towards this end we have covered in depth the development, background and implementation of each utility.

We have made no attempt to cover programming, in either BASIC or machine code, in this book because many other texts cover this in detail. We have also assumed that most serious 64 users will be in possession of a copy of the *Programmer's Reference Guide* and have adopted its nomenclature throughout, particularly with reference to memory locations and KERNAL routines.

Wherever possible a utility has been implemented in both BASIC and machine code. We felt that the BASIC versions, though sometimes crude, are easier to experiment with and should also help those readers unfamiliar with machine code to appreciate the workings of the equivalent routines. They also go to prove that it is not what you know, but how you use it. In some cases it might prove beneficial to use the BASIC rather than the machine code versions. Typical circumstances might be where only one or two features are required, or when you need the full 38K of RAM available to BASIC, or if you wish to switch to bank 2 when using graphics.

To facilitate entering the machine code it has been given in two forms: as BASIC loaders and assembler listings. The assembler listings are suitable for use with an extended monitor. For anyone not owning a monitor program, we have included Jim Butterfield's Supermon and instructions for its use in Appendices D and E. (Supermon for the 64 was first published in the January 1983 issue of *Compute*). However, before attempting the considerable task of typing it in and then getting it to work, ask around your friends and user groups as they may have a copy. If you do find one, you will save yourself a lot of time, effort and frustration. Jim Butterfield has also published a very complete Memory Map for the 64 in the October 1982 issue of *Torpet* which has since appeared in many other journals. This complements the one in the *Programmer's Reference Guide* as it gives the nominal entry points to most ROM routines. A copy of this map could save you a great deal of time when disassembling ROM routines to find out how they work.

To assist in entering the BASIC code all listings have been provided in an annotated form. This, we hope, will avoid the all-too-common problems associated with deciphering the symbols for cursor keys, function keys, colours, and so on when in quotes mode. A detailed list of all mnemonics used is given in Appendix H, but you should find that most are self-explanatory. The program we wrote to generate them is included in this book in the UTILITY as the CODER command. The listings as given in the text always have a maximum line length of 40 characters in their annotated form. Where a line exceeds forty characters it is continued on the next and subsequent lines always commencing in column 1. When looking at the listings you may find it helpful to compare the rightmost characters of continued lines. When the code is typed in, replacing the mnemonics with the correct key(s), no line will exceed 80 characters on the display.

Any of the BASIC utilities intended for use from within another program have been numbered in the 60000's to allow you to merge them with your own programs (using the simple technique described in Chapter 5 or the MERGE command of the UTILITY itself).

We have chosen to put our code at 32768 (\$8000), leaving the BASIC programmer with 30K free. There is no reason why the code could not be modified and relocated elsewhere in memory (the 4K block from \$C000 is not a bad idea) and the initialization routine adjusted to take advantage of the increased memory available. In fact, nearly all the routines were developed and tested in isolation, being enabled by a simple SYS call. They were then incorporated into the UTILITY by simply relocating and including a keyword and token to activate them. To conserve memory, common subroutines have not been duplicated. Often a pick 'n' mix approach was found useful to check out a range of extensions which relied heavily on common routines.

We have used 'hidden' RAM beneath BASIC to store data to conserve valuable user RAM and implemented a simple switching routine to access this data when necessary. Applications like setting up the function keys require access to this RAM, as does CODER. We have made extensive use of the ROM routines and RAM vectors available, but on some occasions found it more economic and faster to write our own code. The UTILITY, in the form given, occupies the same area as cartridge ROM and cannot therefore co-exist with cartridges. It was not written to run in conjunction with them and is intended as a standalone, extendable facility. As the owners of a disk unit will have received DOS 5.1. on the demo disk, the UTILITY has been written to co-exist with DOS 5.1. Some of the isolated routines will temporarily disable DOS as they make use of the same operating routine – CHRGET – but more about that later. However, a simple SYS call will restore DOS 5.1 commands. There are many commercial utilities and BASIC extensions. These may be purchased at reasonable prices and for many applications there will be no better solution. However, if you are interested in the Commodore 64 and wish to get the most from it, you may appreciate having a range of routines which you can modify, extend and, indeed, improve upon. After all, you can pay upwards of the cost of this book for a fairly simple renumber routine.

Before we finish, we would like to leave you with two suggestions and an option:

1) Always save a program before running it

2) Always make backup copies

This is good advice for BASIC and essential where machine code is concerned.

3) It is very easy to wire a reset switch to your 64 and the necessary reset line is available at both the Serial I/O and User I/O (see the *Programmer's Reference Guide*, Appendix I). This is almost essential if you use machine code, but *don't attempt this if you are not sure what you are doing*.

Have fun (if that's the right word)!

1 BASIC on the 64

Introducing BASIC

On powering up your 64 you will find it ready and waiting to go in BASIC, as part of the power reset sequence is to initialize BASIC and leave the user in direct mode.

The implementation of BASIC that Commodore has chosen to use for the 64 is identical to that on the VIC20 and PET microcomputers prior to the 4000 series. BASIC 2, as it is often called, differs from the later version only in its disk operating commands, the latter having a greatly improved and simplified instruction set for disk control. In producing BASIC 4 Commodore did maintain 99.99% downward compatibility and in doing so allowed users to run any program on a higher series PET. It was, therefore, a little surprising to find the BASIC on the 64 to be only V2. This may have been done to avoid a conflict of interests in so much as the new CBM micro, though in many respects far more powerful, was not quite the same.

BASIC, or to give its full title Beginners' All-purpose Symbolic Instruction Code, runs on the 64 as a high-level interpreted language. It is a subset of Microsoft's BASIC (who wrote the first implementation for the early PETs and now produce MBASIC and the MSDOS operating system for all major microcomputers). The history of BASIC is nothing to do with this book, but it is interesting to note that regardless of environment, or cost of system BASIC will usually be in there somewhere. It may only run in compiled form, or it may be syntactically different, but it is reassuring to know that a knowledge of 64 BASIC should allow you to grasp quickly other BASICs on other machines.

BASIC has its critics, particularly of CBM BASIC, who would advocate the use of Pascal, or Pilot, or Forth, or . . . Each of these languages is particularly suited to a range of tasks, but perhaps none lends itself as well as BASIC to the task of rapid development of 'untidy' and 'unstructured' programs which, most importantly, work. Arguments for and against will no doubt long continue, but as we are supplied with BASIC, let us make the most of it.

As its name imples, BASIC was developed to allow beginners to acquire programming skills rapidly. It adopted a system of naming its commands and functions to indicate the action produced. For example, if we wish to halt the execution of a program we issue the all-too-clear command : STOP. For non English-speaking countries even the use of English is no problem on the 64 as not only is it a simple matter to redefine the character set, but it is also easy to redefine the keyword table itself.

On the 64 we have 75 (76 if you include GO) BASIC commands, functions and operators as standard. For many applications this is perfectly adequate. Life would be simpler if more commands were available. Increasing the number of commands has the benefit of providing a more versatile programming language, but the disadvantage of slowing down the execution of the existing commands. This is true of any interpreted programming language. However, the way in which the interpreter has been implemented does allow you to add to the keywords to your heart's content, providing you understand how it works and are capable of writing the necessary machine code. For the moment the interpreter will be considered simply as a means of translating our 'meaningless entries' into something which is executable by the 6510 microprocessor at the machine code level. This it does by taking an instruction, finding the appropriate machine code routine, carrying it out and then returning to implement the next. The process is slow but very flexible and even allows us to interrupt the execution and take control should we wish to do so.

One of the best features of the 64 must be the screen editor. It allows changes to be made directly to anything appearing on the screen and, more importantly, allows you to implement these changes. The disadvantage is that the maximum length of program line or direct statement is limited to 80 characters (or two screen lines). Use of the standard abbreviations of first character and second (or third) character shifted, instead of typing the full keyword, does allow program lines, on listing, to exceed this limit. They cannot, however, be easily edited. On pressing RETURN to acknowledge the end of the edit only two screen lines are accepted. Anything beyond this point is not included in the revised line. Still, this limitation is far outweighed by the speed at which it allows existing code to be edited and repetitive code to be entered by simply using the cursor keys, altering the line number, modifying the necessary part of the line and pressing RETURN to enter the new line. It even allows us to write programs which can generate their own program lines, as we will see in Chapter 5.

BASIC may be used in two modes. These are direct (when a command is typed in without a line number and executed immediately) and program (a command preceded by a line number which is not executed until the program is RUN).

Storage of BASIC code

If we wish to examine a program, we may do so with the LIST command. What we see has undergone many processes from the form in which it was stored in RAM. To view the code *in situ* we first of all need to know where to look. In the default mode on powering-up a BASIC program will be sorted from memory location 2049 (\$0801) upwards. We can examine a program by simply PEEKing out each of the locations used by

FOR I=START TO START+200:PRINT PEEK (I);:NEXT I

We would see a series of decimal numbers with only the fact in common that none was less than zero or greater than 255. We might also notice some sort of related pattern occurring, but not a great deal more. We could adopt another approach by moving an area of RAM used for program storage to the screen. This is easily accomplished, but in doing so we must also remember to set a colour at the screen location we are putting the data into for it to be visible. The resulting display is easier to decipher if the 64 is put in lower case mode by pressing the shift and logo keys together. The following line should be typed in direct mode:

S=0:FOR I=START TO START+800:POKE 1024+S,PEEK(I):POKE 55296+S,14:S=S+1:NEXT I

If you wish to start at the beginning of a program then 2049 must be used and it assumes, as does the first example, a program to be present which occupies memory at least to START+800 (or +200). This time we see a series of characters and where our program has text within quotes it appears almost unchanged as do variable names, punctuation and constants. If we combine the processes, and to produce a more consistent format express the numbers in hexadecimal format, we begin to see some sort of relationship. (You had better get used to using hexadecimal notation as we use it extensively, but to help you on your way there is a table of decimal to hex conversions in Appendix B). The following program does just this and may be used to examine itself. If you wish to experiment, simply enter new lines with numbers less than 60000. Those of you with extended monitors or who jumped straight in and typed in Supermon can use the 'memory display' option.

The program displays on each line the start address and the values held in this and the next seven locations. At the right of the line the characters with ASCII (CHR\$()) codes corresponding to the byte values are printed. To avoid confusion, only those characters which are easily discernible are printed; all others are expressed by a ".". Appendix C of the PRG gives the full range of ASCII and CHR\$() codes. If you want to display all the characters then some of the codes will have effects which will destroy the display, for example, cursor moves, clear screen, colours, and so on; so you will have to trap these. They do, however, occur in blocks and are therefore not too difficult to isolate.

As will be standard practice throughout this book, a description precedes most program listings.

LINE ACTION

- 130 Examine selected range in groups of eight.
- 140 Convert current start address to low/high byte format, that is, units (0–255) and lots of 256s.
 - Then convert to hex notation in two stages.
- 150 Get eight successive bytes from start and print two digit hex
- to values each time.
- 170
- **190** Convert eight bytes to ASCII characters if printable.
- to
- **220** Else replace with a '.' and build eight character string.
- 230 Print string. Recycle if not end else start again.
- 1000 Convert start address to hex in two steps.
- 2000 Convert byte to two digit hex.

```
100 PRINT"MEMORY DISPLAY"
110 INPUT"DISPLAY FROM";F
120 INPUT"[10SPC]TO";T:IF T<F THEN PRINT
 "LESS THAN FROM":GOTO 120
130 FOR I=F TO T STEP 8
140 X=I:GOSUB 1000
150 FOR J=I TO I+7
160 X=PEEK(J):GOSUB 2000:PRINT X$;" ";
170 NEXT J
180 PRINT " ";
190 A$="":FOR J=I TO I+7
200 X=PEEK(J): IF X(32 OR X)95 THEN A$=A$
+".":GOTO 220
210 A$=A$+CHR$(X)
220 NEXT J
230 PRINT A$:NEXT I:GOTO 100
1000 MSB=INT(X/256):LSB=I-MSB*256
1010 X=MSB:GOSUB 2000:PRINT X$;
1020 X=LSB:GOSUB 2000:PRINT X$;" ";
1030 RETURN
2000 X1=INT(X/16):X2=X-X1*16
2010 X$=CHR$(X1+48-7*(X1>9))+CHR$(X2+48-
7*(X2>9))
2020 RETURN
```

If the program is used to examine itself by entering a start of 2048 and an end of 2504 for the program as listed the following display is given:

0970	20	32	30	30	30	ЗA	99	20	2000:.
0978	58	24	ЗB	22	20	22	ЗB	00	X\$;" ";.
0980	86	09	06	04	8E	00	A0	09	
0988	D0	07	58	31	B2	B5	28	58	X1(X
0990	AD	31	36	29	ЗA	58	32	B2	.16):X2.
0998	58	AB	58	31	AC	31	36	00	X.X1.16.
09A,0	СВ	09	DA	07	58	24	B2	C7	X\$
09A8	28	58	31	AA	34	38	AB	37	(X1.48.7
09B0	AC	28	58	31	B1	39	29	29	.(X1.9))
09B8	AA	C7	28	58	32	AA	34	38	(X2.48
09C0	AB	37	AC	28	58	32	B1	39	.7.(X2.9
09C8	29	29	00	D1	09	E4	07	8E	››
09D0	00	00	00						• • •

Its exact format will vary depending on how you typed the program in. A number of things are immediately apparent. All text inside quotes, all variable names, all destinations, all constants and punctuation appear unchanged. From just this information we can work out the general area of each line. Taking line 100 as an example, "MEMORY..." is clearly visible from \$0806 to \$0815. Immediately preceding it is the value \$99 which, not surprisingly, is the tokenized value for PRINT. The two bytes before this are \$64 and \$00. \$64 is the hex for 100 which is the line number. Line numbers may range from 0 up to 63999 and, like many values on the 64, are stored in low/high byte format. The actual line number is \$64+\$00*\$0100 (100+0*256). If we look along the hex values for line 100, we see that the byte immediately following the closing quote is a zero. This is how BASIC marks the end of a program line. The two bytes preceding the line number are \$17/\$08 which is the address (low/high) of the byte immediately following this end of line zero. These two bytes are known as the link address and point to the link address (and start) of the next line. If we follow the link address through the program, the sequence runs \$0817/082D/0864....09D1 and finally 0000. A link address of zero marks the end of the program. which in this case is \$09D1. A pointer to this address+2 is held in zero page (locations \$00=\$FF) at VARTAB (\$2D/2E) and marks the start of the BASIC variables. A second pointer on zero page, TXTTAB (\$2B/2C), points to the start of the program. This is the location of the first link address and in the default setting this will always be \$0801. Location \$0800 holds a zero; the byte before the start of a program must always be zero for RUN to work. This becomes of more significance when the start location of BASIC is changed.

A program can therefore be thought of as a 'linked list' of individual program lines. It is of the form:

START	link	LINE		END		END PROG
00	low high	low high	BASIC line	00	low high	00 00 00

0800	00	17	0 8	64	00	99	22	4D	"M
8080	45	4D	4F	52	59	20	44	49	EMORY DI
0810	53	50	4C	41	59	22	00	2D	SPLAY"
0818	08	бE	00	85	22	44	49	53	"DIS
0820	50	4C	41	59	20	46	52	4F	PLAY FRO
0828	4D	22	3B	46	00	64	08	78	M";F
0830	00	85	22	20	20	20	20	20	••
0838	20	20	20	20	20	54	4F	22	TO"
0840	ЗB	54	ЗA	8B	20	54	B 3	46	;T:. T.F
0848	20	A7	20	99	20	22	4C	45	"LE
0850	53	53	20	54	48	41	4E	20	SS THAN
0858	46	52	4F	4D	22	ЗA	89	20	FROM":.
0860	31	32	30	00	76	0 8	82	00	120
8680	81	20	49	B2	46	20	A4	20	. I.F .
0870	54	20	A9	20	38	00	85	0 8	т. 8
0878	8C	00	58	B 2	49	ЗA	8D	20	×.I:.
0880	31	30	30	30	00	95	0 8	96	1000
0888	00	81	20	4A	B2	49	20	A4	J.I .
0890	20	49	AA	37	00	B1	88	A0	I.7
0898	00	58	B2	C2	28	4A	29	ЗA	.X(J):
08A0	8D	20	32	30	30	30	ЗA	99	. 2000:.
08A8	20	58	24	ЗВ	22	20	22	ЗВ	X\$;" ";
0880	00	B9	08	AA	00	82	20	4A	J
0888	00	C4	08	Β4	00	99	20	22	*
0800	20	22	3B	00	DA	0 8	BE	00	";....
08 C8	41	24	B2	22	22	ЗA	81	20	A\$."":.
08D0	4A	B2	49	20	A4	20	49	AA	J.I . I.
0 8D8	37	00	05	09	C8	00	58	B2	7X.
08E0	C2	28	4 A	29	ЗA	8B	20	58	.(J):.X
08E8	B 3	33	32	20	B0	20	58	B1	.32 . X.
08F0	39	35	20	A7	20	41	24	B2	95 . A\$.
08F8	41	24	AA	22	2E	22	ЗA	89	A\$.".":.
0900	20	32	32	30	00	14	09	D2	220
0908	00	41	24	B2	41	24	AA	C7	.A\$.A\$
0910	28	58	29	00	1 C	09	DC	00	(X)
0918	82	20	4A	00	2F	09	Eб	00	. J./
0920	99	20	41	24	ЗA	82	20	49	. A\$:. I
0928	ЗA	89	20	31	30	30	00	4E	:. 100.N
0930	09	E8	03	4D	53	42	B2	85	MSB
0938	28	58	AD	32	35	36	29	ЗA	(X.256):
0940	4C	53	42	B2	49	AB	4D	53	LSB.I.MS
0948	42	AC	32	35	36	00	65	09	B.256
0950	F2	03	58	B2	4D	53	42	ЗA	X.MSB:
0958	8D	20	32	30	30	30	ЗA	99	. 2000:.
0960	20	58	24	ЗB	00	80	09	FC	X\$;
0968	03	58	B2	4C	53	42	ЗA	8D	.X.ĹSB:.

The link addresses are not used when a program is run, but are important during listing and editing. We can alter their values without affecting the way a program runs, but on listing some strange effects are produced.

We can now look through the display and find the start and end of a line and the associated line number. Knowing these, we can start to deduce the tokenized values for the BASIC keywords used. By adding lines to the program we could find out all keyword values, but to save you the effort we have produced a complete list in Appendix A. This table has been extended to include the new token values used by the UTILITY. These values should be ignored for the moment. With a little practice, reading displays of this type becomes very easy.

From the table in Appendix A we see that all BASIC keywords have token values in excess of 127 (\$7F). The highest token used for standard BASIC is 203 (\$CB) for the GO command. (GO simply searches for a corresponding TO to ensure GOTO is equivalent to GO TO). When a line is entered from the keyboard it is transferred to the input buffer (BUF \$0200–0258) on pressing RETURN. The line is then tokenized in accordance with this table with keywords being processed first. If no line number is present, the BASIC interpreter immediately executes the statement(s). If one is present then the line is put in its numerically correct position and the link addresses for the whole program are recalculated and VARTAB updated. A similar process is carried out when a line is deleted. These operations are discussed in greater detail in Chapter 3.

Variables

General

BASIC allows three types of variable. These are real, integer and string. String and integer are distinguished from real by trailing '\$' and '%' characters respectively. The default is therefore to real. Variable names may be of any length, but only the first two characters and the last character are of significance. This means ABXXXX% will be considered equal to ABYYY% and to AB%, but different from ABXXXX or ABXXXX\$. The last character is used to distinguish the variable's type, and if it is not one of the special characters above then the variable is treated as real. The only limitation on naming variables is that the first character must be alphabetic and the subsequent character, alphanumeric, providing that they do not form reserved keywords. For example, PEND would be treated as 'P' plus the keyword END and on running would produce a syntax error. Reserved words are any of those occurring in Appendix A with a token value exceeding 127 (bearing in mind that with the UTILITY in place the number of reserved words will be increased).

To allow each of the three types of variable (four, really, when we

include function names) to be stored in two bytes, the high bit is set or unset on each of the name bytes to give the necessary four combinations. These are:

Name	
1st char	2nd char
ASCII	ASCII
ASCII+128	ASCII+128
ASCII	ASCII+128
ASCII+128	ASCII
	Name 1st char ASCII ASCII+128 ASCII ASCII+128

Where the name is only a single character, the second byte is zero or 128 as appropriate.

Each of the three types of variable may be used in multi-dimensional arrays. These subscripted variables follow the same rules as for simple variables with the addition of a '(' following the name and type. This tells the interpreter it is dealing with an array and is handled accordingly.

Storage of variables

Any variable created in either direct or program mode is stored after the program currently in memory. Variables are stored in the order in which they are created. Strings are stored slightly differently from numeric values mainly due to their dynamic nature and are held in two parts. The first is a pointer to the string's location and the second is the length of the string itself. Strings are stored at the top of BASIC memory (\$9FFF/40959) and grow downwards. The current lower limit of string storage is stored in FRETOP (\$33-\$34/51-52).

When searching for a variable the interpreter starts at the end of the program and searches upwards in memory for the named value according to the rules given above. If the variable does not exist, the next available space is allocated to it. Thus, if we define the more important values early on they can be accessed quicker and the time spent in 'garbage collection' reduced.

Variables are either simple or subscripted.

Simple variables

All non-subscripted variables use seven bytes of RAM. The first two hold the name in its adjusted form. For each real variable the remaining five bytes are used in the following way: one for its exponent and the remaining four for its sign and mantissa. Integers are stored in only two bytes with the remaining three unused. Strings use one byte to indicate the length and two bytes to point to the location of the characters, which is usually at the top of memory (though not always). A function also uses seven bytes, of which the third and fourth point to its definition (DEF FN), the next two point to the variable it uses and the last points to an initial value of the variable (zero).

The following table summarizes the storage of simple variables:

Byte

1 and 2: name	3	4	5	6	7
REAL INTEGER STRING FUNCTION	exp sign+high length pointer to [low	sign+M1 low ptr low DEF FN high	M2 unused ptr high pointer to low	M3 unused unused variable high	M4 unused unused initial value

If we add the following lines of code to our memory display program, all variable types are generated (including arrays):

1 DIM A(5),B%(5),C\$(5),D\$(1,5) 2 FOR I=0 TO 5:A(I)=I:B(I)=I:C\$(I)=CHR\$(64+I):D\$(2,I)=C\$(CI): NEXT 3 M\$=M\$+''STRING1'':Z\$=Z\$+''STRING2''

We can now dump the memory associated with simple variables. The area to be displayed can be worked out from VARTAB, ARYTAB and FRETOP. To do this the program is RUN twice, the first time from 2641(\$0A61) to 2732(\$0AAC) to display the variables and the second time from 40900(\$9FC4) to 40960(\$A000 – the start of BASIC ROM) to display the strings *in situ*:

Simple variables and string pointers

49	00	8C	25	10	00	00	4D	I%M
80	07	ED	9F	00	00	5A	80	Z.
08	E5	9F	00	00	46	00	8C	F
25	10	00	00	54	00	8C	2A	%T*
C0	00	00	58	00	00	00	00	×
00	00	4D	53	84	20	00	00	MS
00	4C	53	88	01	00	00	00	.LS
58	31	82	00	00	00	00	58	×1×
32	82	00	00	00	00	58	80	2X.
02	19	9D	00	00	4 A	00	8C	J
2A	20	00	00	41	80	88	D8	*A
9C	00	00	41	00	25	00	01	A.%
	49 80 25 C0 00 58 32 02 20 9C	 49 80 80 87 85 25 10 60 90 90 90 40 90 40 90 40 91 92 92 92 90 90 	49008C8007ED08E59F251000C00040004040004C5358318232820002199D2A20009C0000	49 00 8C 25 80 07 ED 9F 08 E5 9F 00 25 10 00 00 25 10 00 00 C0 00 00 58 00 00 4D 53 00 4C 53 88 58 31 82 00 32 82 00 00 02 19 9D 00 2A 20 00 41	49 00 8C 25 10 80 07 ED 9F 00 08 E5 9F 00 00 25 10 00 00 54 C0 00 00 58 00 00 00 4D 53 84 00 4C 53 88 01 58 31 82 00 00 32 82 00 00 00 02 19 9D 00 00 2A 20 00 00 41	49008C2510008007ED9F000008E59F00005400251000005400C0006058000000004D538420004C5388010058318200000032820000004A2A20000041809C0000410025	49008C251000008007ED9F00005A08E59F000054002510000054008CC0006058000000000053842000004C53880100005831820000005802199D00004A002A2000044180889C000041002500	49008C251000004D8007ED9F00005A8008E59F000046008C2510000054008C2AC00060580000000000004D5384200000004C538801000060583182000000588032820000004A008C2A200000418008D89C00004100250001

Strings in situ

9FC4	30	30	30	30	30	31	30	30	00000100
9FCC	31	32	35	35	32	38	43	43	12552800
9FD4	38	30	30	30	30	34	39	39	80000499
9FDC	34	35	31	31	35	30	41	41	451150AA
9FE4	30	53	54	52	49	4E	47	32	0STRING2
9FEC	29	53	54	52	49	4E	47	31)STRING1
9FF4	45	45	44	44	43	43	42	42	EEDDCCBB
9FFC	41	41	40	40	94	EЗ	7B	EЗ	ΑΑ

Real variables

Looking through the display above, it is quite easy to spot the real variables as their names are stored in unmodified ASCII. At \$0A51 we see 'I', the first non-subscripted variable to be used, with a zero second byte in its name. We can also spot F, T, X and all the others (noting that MSB and LSB are stored as MS and LS).

Real numbers are stored in binary floating point format, always to an accuracy of 31 bits. Due to the way in which they are stored in single precision form, rounding errors are introduced though these are usually not significant enough to affect the final results. Examples of this type of error are encountered all the time as in the 'X.000001'-type value. We can convert 'I' back to a decimal number quite easily.

The exponent is stored in byte 3 and is the power of two. A unit change in this doubles or halves the resulting value. Positive exponents are expressed as 129+EXP and negative, as 129-EXP. Therefore, the full range is from $2^{(-129)}$ to $2^{(127)}$ or in decimal, from about $10^{(-38)}$ to 10^{37} . The high bit of byte 4 indicates the sign and is set for negative numbers. To calculate the decimal value, we have to successively divide the mantissa starting at the right by 256, add the result to the next on the right and so on until we reach M1, when we only divide by 128 and finally add 1. The resulting number will lie between 1 and 1.999999. This must finally be adjusted for its exponent and sign. The values for 'I' used below are in decimal.

 $\begin{array}{cccc} M4 & 0/256 = 0 \\ M3 & (0+0)/256 = 0 \\ M2 & (0+16)/256 = .0625 \\ M1 & (.0625 + 37)/128 = 0.28955 \\ & +1.00000 = 1.28955 \end{array}$

If this is then multiplied by the exponent of $2^{(140-129)}=2048$, the value is $2048 \times 1.28955 = 2640.999$ (almost 2641). This is the upper limit for the first memory display. A general formula may be written to convert any real variable from its floating point to decimal form:

 $(-)^{(M1 \text{ AND } 128) \star 2^{(EXP-129) \star (1+((M1 \text{ AND } 127)+(M2+(M3+(M4/256)/256)),.../128)})$

Integer variables

These are stored in a signed high/low byte format and can range from -32768 to 32767. The high bit of byte 3 is again used to indicate the sign. The value is easily determined from the following:

(BYTE3 AND 127)*256+BYTE4+(BYTE3>127)*32768

String variables

These are the easiest of all to pick out. At \$0A59 in the display above is the variable M\$(its second byte is not used so is set to \$80). Byte 3 tells us it is seven characters long, and bytes 4 and 5 that it is located at \$9FED. The seven bytes from \$9FED are "STRING1" as would be expected. Strings therefore use seven plus the number of characters bytes of RAM. There is one important point to make before leaving strings. If line 3 had simply been M\$="STRING1", its pointer would have pointed to the byte at which it occurred within the program itself, that is, the byte immediately following the quote. Only computed strings are stored at the top of memory which is why the line was written M\$=M\$+"STRING1". This economizes on memory usage by only storing the string once. It does have the drawback that if another program is loaded in program mode all non-computed strings are lost.

Subscripted variables

Duto

Arrays may be of any type, but unlike their 'simple' counterparts, only the required number of bytes are used to store the associated values. Real are stored in five bytes, integer in two and strings in three plus their length. In addition to the savings in storing the values, the array name is only stored once. Arrays are also created in the order in which they are encountered.

The area of memory used for arrays immediately follows that for simple variables. As for the latter, it, too, is recorded at two zero page locations. The start, ARYTAB, has already been mentioned when dealing with simple variables. The end, STREND, is held in \$31–\$32/49–50. For each new simple variable this whole block must be moved up seven bytes in memory. There will, of course, come a time when array storage builds up to meet that of the descending strings with the resulting 'OUT OF MEMORY' error.

Each array is preceded by a detailed header of the form shown below:

byte			
1 and 2	3 4	5	6 7 N-1 N
NAME	OFFSET TO	NO.	LASTFIRST
	1ST VALUE	DIMS	DIM+1DIM+1
Adj. form	low high	<256	low highlow high

Bytes 1 and 2 hold the name in its adjusted form. Bytes 3 and 4 record

the overall memory requirement for the array (this does not include string data at the top of memory) and is the offset from its start to the next array. Byte 5 records the level of dimensioning and may not exceed 255 (a little difficult to visualize at anything more than two or three). If an undimensioned array is used, this value will default to the number of subscripts at the first occurrence. Successive pairs of bytes then hold the number of elements in each dimension (plus one for the zero subscript) in the reverse order of dimensioning. If no dimensioning has been used, these each default to 11 (10+1). The following bytes are then used to store the data.

If the program is again run and memory between ARYTAB and STREND displayed, the following results:

0AAC	41	00	25	00	01	00	86	00	A.%
0AB4	00	00	00	00	81	00	00	00	
ØABC	00	82	00	00	00	00	82	40	
0AC4	00	00	00	83	00	00	00	00	
9ACC	83	20	60	00	00	C2	80	13	
0AD4	00	01	00	66	00	00	00	01	
ØADC	00	02	00	03	00	04	00	05	
0AE4	43	80	19	00	01	00	06	01	С
0AEC	FF	9F	01	FD	9F	01	FB	9F	
0AF4	01	F9	9F	01	F7	9F	01	F5	
ØAFC	9F	44	80	ЗF	00	02	00	66	.D.?
0 B04	00	03	00	00	00	01	FE	9F	
080C	00	00	00	00	00	00	01	FC	
0B14	9F	00	00	00	00	00	00	0 1	
0B1C	FA	9F	00	00	00	00	00	00	
0B24	01	F8	9F	00	00	00	00	00	
0B2C	00	Ø 1	F6	9F	00	00	00	00	
0B34	00	00	01	F4	9F	00	00	00	
0 B3C	00	01	FF	FF	00	01	FF	FF	

The first array is 'A(' at \$0AAC. It occupies 37 bytes (\$0025), and has one dimension (\$01) of five elements (\$06-1). The six values are then held in 5 byte real format. The next array starts at \$0AAC+\$25=\$0AD1. This is 'B%(' which occupies only 19 (\$13) for its six values. The values are easily read out as 0, 1, 2, 3, 4 and 5. The next is 'C\$(', and looking at the previous display we can read out its values as @, A, B, C, D and E. A little care has to be exercised here as in the loop which generated them 'D\$(' was also defined each time. This last array is the most complex of all. In its dimension statement it was defined as D\$(2,5). However, in its header these are reversed (last....first). The values set were assigned to D\$(1,1) and from the display we can see these occur at the second and subsequently at every third byte. This shows us that multi-dimensional arrays are stored in the form X(0,0) X(1,0) ... X(N,0) X(0,1) X(1,1) X(N,N).

This just about concludes our section on variables, except to say that the default values are zero for numeric and null for strings.

Link addresses and line numbers

General

Knowing where and in what form these are stored, there is no reason why we cannot modify them from BASIC itself. This we can do using simple POKES to produce some interesting results.

Links

If we modify a link address, the program will continue to run. It will, however, list in an unusual fashion and be difficult to edit. We can use this fact to make our programs difficult to read and modify. This we can do by hiding lines (the whole program if we wish). Hiding line 110 of the display program, as its listing was originally given, can be done by

POKE 2049,45

This simply skips the link at \$0817. We could very easily write a short routine to eliminate whole blocks of line numbers. This we leave up to you.

Line numbers

We can change line numbers as we did link addresses. We could change the line number of 100 to any value we choose.

POKE 2051,110

will, on listing the program, give two line 110s. A little care should be taken here, because if the line number changed is the destination for a GOTO or a GOSUB, some confusion may result.

Saving modified code

BASIC'S SAVE command transfers to tape or disk a copy of the RAM between TXTTAB and VARTAB. This means any modifications are also saved. The modified code returns on loading.

Modifying BASIC

Changing the load address

We can change the point at which BASIC programs load simply by setting the value in TXTTAB. Wherever the new start is to be, a zero must be set in the byte immediately before it. Once the new start has been set, a NEW will tidy up all other pointers.

Changing the start of BASIC is useful if using sprites or programmable characters within bank Ø (the default). The *Programmer's Reference Guide* recommends lowering the top of memory to make room for the

necessary data. Instead, why not move up the start of BASIC and leave yourself with far more memory to use?

Chaining programs

Chaining in this context refers to loading one program from within another.

The question arises as to whether there is a bug in the chaining process. The answer is a qualified 'no' as there can be problems. The effect of a program LOAD is roughly equivalent to executing a GOTO the first statement of the chained program. Thus, the new program can use only real, integer and computed string variables from the first. The problem occurs when the incoming program is larger than the original. If this is the case, it will overwrite the start of the variables, causing utter confusion. Once this has happened you really need to issue a CLR at the start of the new program to tidy things up. You have apparently lost all the variables anyway so there is nothing to lose.

On some micros a CHAIN command exists in addition to the normal LOAD. The action of this command is to move all or only the specified variables out of the way during the loading process and then move them back and update the necessary pointers. On the 64 no such command is available. There are two solutions. The first is to ensure a larger program is never loaded from a smaller one. The second is to make the first program the largest. To do this we do not need to generate a 'large program'. All that is needed is a simple POKE and CLR sequence at the start of the first program to reserve the necessary memory. BASIC can be fooled by:

POKE 46, (SIZE OF BIGGEST PROG)/256+8+1:CLR

In this example we have not bothered to be exact and have simply reserved to the nearest page.

Speeding up program execution

There are many ways in which to increase the speed of a program. The speed of peripheral devices plays a major part when inputting or outputting, but the topics covered here are mainly concerned with the 64 itself.

There are commercial 'CRUNCH' or 'compactor' programs available. These traditionally remove all unnecessary spaces, REMs, combine lines not the destination of a GOTO or GOSUB and that is about all. Even these few changes can produce significant increases in execution speed. Some of this we can do from BASIC itself. A short routine at the start of a program can combine lines by eliminating link addresses and line numbers, and remove all REMs and all spaces not inside quotations. After each deletion, the remaining code is moved down in memory and VARTAB is updated. The end of line marker must be replaced by a ':' to separate the last statement from the leading statement on the crunched line. There are a number of problems here. If a line number is eliminated which is referenced by a THEN, GOTO or GOSUB, the run will fail. The second problem stems from all statements following an IF being ignored if the condition is false. The resulting compacted lines are so long that they cannot be edited. As such, a universal compactor program in BASIC is fraught with danger.

It is far more sensible to consider these points at the time of writing the program. A well-known technique is to ensure that all GOTOs have destinations as near the start of the program as possible as the interpreter starts its search there. If this cannot be done, then the destination should have its high byte greater than that of the GOTO line due to the search technique used which compares this byte first.

There are three other common methods used to optimize the code:

- (i) the use of variables rather than constants;
- (ii) the setting up of variables in the order of frequency of use;
- (iii) and, specific to the 64, turning off the video display when not in use (see *Programmer's Reference Guide*, Appendix N, 'Screen Blanking').

Using these techniques, the second program below runs almost 25% faster than the first:

```
10 PRINTTI
20 POKE 49152,0
30 FOR I=0 TO 5000
40 J=J+1
50 NEXT I
60 POKE 49152,16
70 PRINTTI
```

10 PRINTTI:POKE53265,0:P=1:FORI=0T05000: J=J+P:NEXT:POKE53265,16:PRINTTI

The second program does leave you in x and y scroll mode if you are wondering just what has happened.

Conclusion

This chapter should have given you one or two ideas to play around with. Before reading Chapter 5, you might like to think about how to write simple renumber, delete, dump, and recover NEWED program routines. You might also like to think about how to overcome the chaining problem by, as the last action on leaving a program, moving all variables as high as they can go in memory. The first action of the chained program should be to move them back down to the end of this program and reset the necessary zero page pointers.

2 Peripherals

Introduction

This chapter deals with some of the more common peripherals for the 64. Also included here are the keyboard and screen even though they are not quite peripherals in the same sense as a disk drive, cassette or printer.

It is not our intention to go into any of these in great detail as the subject could fill a book of its own. We have tried to look at features of more immediate use.

Keyboard

Use of the keyboard, its ROM drive routines and RAM vectors is covered in Chapter 4. Programming of the keyboard is used extensively in Chapter 5. The following are a few useful points not directly covered elsewhere,

Keyboard as a device

The keyboard is viewed as device 0 by the 64's operating system and is the default for input. As such it may be used like any other device and a file OPENed to it. This file may only be for input and any attempt to output to it will result in an error. Once opened for input the 'annoying' question mark prompt is removed from the INPUT command display. When information is being obtained from it using INPUT# all warning messages, such as the double question mark for insufficient data, also disappear. The open format is the same as for any other device:

OPEN 1,0 or OPEN 1,0,129,"QWERTY"

In the second example everything following the Ø will be ignored. Use of the keyboard in this way is highlighted in the DISK utility facility in this chapter (see below).

Auto-repeat See Chapter 10.

Key detection See Chapter 4 and Appendix I.

Keyboard Buffer - KEYD (\$0277-\$0280/631-640)

The 64 provides type-ahead of up to ten characters. The buffer operates on the principle of first in/first out. However, once full no new characters will be accepted until it has been partially emptied. Characters are taken singly by a GET, up to the first RETURN on an INPUT and the buffer is emptied on an END.

The length of the buffer is determined by XMAX (\$0289/649) and as this is in RAM it may be changed. Theoretically, the buffer could be lengthened, but in practice this cannot be done as the RAM immediately following it is used for other purposes. The size, however, may be decreased and is perhaps most useful when the length is set to 1 where a program requires careful, restricted input or type-ahead is to be discouraged. Setting XMAX to 0 is quite a good way to prevent unwanted user input (the STOP key is still active as it is scanned by a different routine – see Chapter 10 to disable).

As the buffer is in RAM we can put data directly into it by simply POKEING the ASCII codes of the characters required. To complete the process NDX (\$C6/198) must be set to tell the system how many characters are in the buffer. This type of approach is used extensively in the BASIC utilities in Chapter 5 and we refer you there for examples of using the keyboard in this way.

One final point before leaving the keyboard which many of you may have already discovered. Pressing the 'Control' key with any other simply sets bit 6 of the ASCII code of the character low. For example, CNRL/T is the same as the DEL key (\$14/20).

Cassette

The 64 does not have to use a CBM Datassette. There are, to our knowledge, two manufacturers of interfaces which allow standard cassette recorders to be used. These interfaces duplicate the part of the interface normally resident within the Datassette. It is even possible to use a standard cassette through a suitable edge connector, but do not expect a high success rate in loading back saved programs or to get anything from recordings produced on a CBM recorder (see *Programmer's Reference Guide*, Appendix I for connection details).

Many consider the cost of the dedicated cassette high. However, it avoids the need to adjust the volume and tone controls to ensure an accurate save (a problem on many other micros where even saving twice is not guaranteed to work). It also seems slow, but perhaps is not as slow as it at first appears. Data is transferred between the 64 and the cassette at about 300 baud (some micros offer an optional fast 1200 baud rate). When a program is saved two copies are made. On reloading the first copy is put into memory and this is then compared with the second to check for and possibly recover load errors. In our experience it has proved worth the additional expense to buy the Datassette for peace of mind and to avoid the loss of many hours of hard work.

The speed of operation of the cassette has been chosen for reliability, but like most things on the 64 we can even change that. For many years superfast, jet, turbo, fast or whatever you care to call them, operating systems have been available for the PET, more recently for the VIC20 (ARROW) and now for the 64. The machine code listing of the original PET version has even been published. Many games now come with a high-speed load (some without the option for a normal load which has proved annoying when your cassette cannot cope with fast loads). These fast operating systems can be made to run the cassette at a higher speed than the standard operating speed of the disk drive (even this can be increased). There is no secret as to how it is done, but as many software houses pay a royalty for its use, or even sell their own versions of a high-speed loader, we have decided not to include a version of our own.

LOAD and SAVE with cassette (see Programmer's Reference Guide, Chapter 2)

These are dealt with in detail in the *Programmer's Reference Guide*, so we will deal with them briefly here. The general syntax for SAVE is (where square brackets denote optional parameters):

SAVE["program" or string variable] [,device] [, secondary address]

If no parameters are specified, the BASIC program currently in memory will be copied to the default of cassette without a name.

The secondary address is the more interesting. A secondary address of 2 will write an end of tape marker and one of 3 appears at first sight to do exactly the same. Using either of these will prevent the tape being read beyond this point without being physically wound on. There is, however, a world of difference on loading (see below). With an address of 3 not only is the end of tape set, but an end of tape header is written which is a duplicate of the program header with a type of 5.

The area of RAM saved is that between the values held in TXTTAB and VARTAB. These pointers are automatically kept up to date by the operating system whilst a program is being edited. Should we wish to save an area of memory other than the BASIC program, we can set these up by POKEing in the appropriate values (remember low/high format). This allows us to save machine code from BASIC or even the screen itself. Data stored in memory is more economically saved this way as only single bytes are saved and not the ASCII characters which make up each number (saves at least two bytes per number between Ø and 255). The problem is that on returning from the save, the current BASIC program and variables are lost until these pointers are restored. If you

are going to play with TXTTAB and VARTAB from BASIC, put the original values out of harm's way, say below \$0800 or above \$C000, to allow them to be recovered.

The syntax for LOAD is identical to that for SAVE:

LOAD ["program" or string variable] [,device] [,secondary address]

LOAD reads the next program from tape. If a program name is specified, then the named program will be searched for and if found loaded or if an end of tape marker is found first the cassette will stop. Again it is the secondary address which is of major importance. A 1 requests the operating system to put the program at the same location from which it was saved. If no secondary address is specified, then providing the program came from an address above the current start of BASIC it will return to its original location, but after the load TXTTAB will still hold the start of BASIC whereas VARTAB will hold the end address. The same is true when 1 is used, but in this case a load may be carried out below the start of BASIC. Typically, when loading machine code from BASIC an 'OUT OF MEMORY' error results if the code locates above \$9FFF due to the setting of VARTAB. A save with a secondary address of 3 ensures the code is reloaded to its original address, regardless of the syntax of the LOAD command (extended monitors use 3).

Tape Buffer

The tape uses a 192 byte I/O buffer, TBUFFR, which in its default setting extends from \$033C-03FB/828-1019. TBUFFR need not reside here and may be relocated, as a pointer to its start is held in RAM at SAL (\$AC-AD/172-173). To move it, simply POKE in the new location in the usual low/high byte format (STOP/RESTORE will reset it). We have found this of use when storing sprite data blocks in bank 0 when memory is tight (\$C000 is yet again a good place to put it). Usage of the buffer is very different between program and data files. Programs only use the buffer to store their header information (see below) and the transfer of memory is direct from the I/O port without passing through the buffer. Data files, on the other hand, use the buffer initially for the header then subsequently to hold 191 byte blocks (the first byte is used as a marker). This avoids continual starting and stopping of the tape motor and by using this block system the tape is more reliable as it is allowed to pick up speed between each read/write operation. Another zero page location, BUFPNT (\$A6/166), holds the current position within the buffer.

Tape Headers

All files are stored on tape with an initial header which is the length of the buffer. The exact format depends on the syntax of the SAVE or OPEN command (secondary address of 2 on an OPEN also writes an end of tape

marker). Each is made up of an identifier, two addresses and a file name, the format of which is given below:

Program headers

ID	START	END	FILE NAME (spaces to pad)
1	18	25 16	65 66 67 32 32 32

Data headers

ID	START	END	FILE NAME
4	60 3	252 3	68 65 66 32 32

The ID identifies the file type and for a program may also take a value of 3. The two bytes immediately following it are the start load address in low/high format and the next two the end address. The file name is not limited to 16 characters and in fact can be up to 187 characters. This allows machine code to be embedded in a header to add additional security to a program. When the name is printed out by LOADING only the first 16 characters are displayed. The header to a data file also contains the start/end bytes but these hold the start and end of TBUFFR itself.

The last operation on completion of a save or write is to store a duplicate header. If the command had a secondary address indicating an end of tape marker, then the ID would be changed to a 5 before writing. On loading or reading to the end of a file the last operation is to get back this trailing header (which remains until the next tape operation).

Tape directories

Tape directories as such do not exist unless you are using an improved cassette operating system such as ACOS+. There are times when it is necessary to catalogue a tape. The process is time-consuming as it is, not surprisingly, directly proportional to the length of the tape. The following program may be used to do the job. It is best left running whilst you go away to do something else.

Any header will be read with an OPEN statement. CLOSEing it immediately ceases tape operation and program execution continues. The parameters are then pulled from the buffer and stored for later use. The process is repeated for the next header. When the end of tape is reached or you stop the program, a simple GOTO 260 will display the file information. This is the file type, up to 16 characters of its name with non-alphanumeric characters replaced by a ".", and if a program its start and end addresses (in hex).

```
100 DIM F$(50),FT$(50),SA$(50),EA$(50):C
B=828
110 PRINT"[CLS]PRESS PLAY ON TAPE"
120 IF PEEK(1) <>7 GOTO 120
130 I=I+1:0PEN 1:CLOSE 1:PRINTF$(I-1)
140 FT$(I)=RIGHT$("[5SPC]"+STR$(PEEK(CB)
),4)
150 IF PEEK(CB)=4 THEN SA$(I)="[2SPC]***
*":EA$(I)="[2SPC]****":GOTO 200
160 X=PEEK(830):GOSUB 360:SA$(I)=X$
170 X=PEEK(829):GOSUB 360:SA$(I)=" $"+SA
$(I)+X$
180 X=PEEK(832):GOSUB 360:EA$(I)=X$
190 X=PEEK(831):GOSUB 360:EA$(I)=" $"+EA
(1)+X_{5}
200 A$="":FOR J=833 TO 848
210 X=PEEK(J): IF X(32 OR X)95 THEN A$=A$
+".":GOTO 230
220 A = A + CHR (X)
230 NEXT J
240 F$(I)=LEFT$(" "+A$+"[18SPC]",17)
250 GOTO 130
260 H$="[CLS]TYPE FILENAME[9SPC]START[3S
PC]END":PRINTH$
270 FOR J=1 TO I:PRINTFT$(J);F$(J);SA$(J
);EA$(J)
280 IF INT(J/20)<>J/20 GOTO 320
290 PRINT"PRESS RETURN FOR NEXT PAGE"
300 GET A$: IF A$<>CHR$(13) GOTO 300
310 PRINTH$
320 NEXT J
330 INPUT "REVIEW AGAIN";Y$:IF Y$="Y" GO
TO 260
340 IF Y$<>"N" GOTO 330
350 CLOSE 1:END
360 X1=INT(X/16):X2=X-X1*16
370 X$=CHR$(X1+48-7*(X1)9))+CHR$(X2+48-7
*(X2>9))
380 RETURN
```

Unfortunately, during tape I/O the internal clock variable (TI\$) is not updated as the interrupt is used exclusively for tape timing. Had this not been the case, a read of this variable could have been used to calculate the value of the tape counter. The best suggestion we can come up with is if the file is a program then the difference in its start and end addresses could be used to determine the loading time. For a

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data file bytes could be taken until the status is set to the end of file, the number of bytes read being an indication of the time. We might as well do this as the tape is running anyway. The time taken may be used to work out an approximate counter reading.

Auto-running

Generating programs which auto-run is also discussed in Chapter 10. There are many ways to accomplish this, most of which involve fairly detailed knowledge of the operating system. The following are suggestions only for you to pursue. All but one are suitable for disk or tape.

The stack

During LOAD the return address is placed on the stack. As this is an area of RAM, there is no reason why we cannot load through this area and put our own address on instead. This could then go to our own machine code routine. The file type should be 3 to ensure a load to its original position. The same would apply to disk or tape if loaded with a secondary address of 1.

BASIC warm start - \$0302

After a load in direct mode BASIC is warm-started. Again as this vector is held in RAM we can load through it. The new value it then contained could jump to our machine code or straight to RUN (for BASIC programs).

IRQTMP – \$029F

This stores the current IRQ vector during tape I/O which is restored after the tape operation. Again we can do the same to this as in the above. On the first normal interrupt the action will be taken. This, of course, can only be used with tape.

CHAIN command of the UTILITY See Chapter 8.

Screen

The utilities in this section are confined to the text screen.

The screen on the 64 is a 40 column by 25 line memory-mapped display. Chapter 3 and Appendices B to D of the *Programmer's Reference Guide* cover in great detail all aspects of the screen and it is to there that we refer you. All the following utilities assume that you are familiar with or know the following.

- i) The screen may be moved from its default position.
- ii) There are two character sets.
- iii) The screen has an associated colour map at \$D800 on.
- iv) The display codes differ from the ASCII codes.
- v) Commodore 'ASCII' is not true ASCII which only ranges from 0 to 127. (Consult your printer manual.)

Printer dump

There are two routines, both of which output the current display in standard ASCII to a printer. One is a BASIC subroutine and the other is machine code. The second is noticeably faster than the first, as would be expected.

Both routines take account of whether the 64 is in upper or lower case mode as well as checking for the location of the screen.

64 owners with Commodore printers need not concern themselves with the conversions to standard ASCII.

BASIC printer dump subroutine

The version given here is for an RS232 printer running at 300 baud without auto-line feed. For this reason the output logical file is assigned at the start of the program. The out put file is designated 'P' to avoid specific reference to allow for easier change to other printers. The display is centred on an 80 column display by printing 20 spaces at the start of each line.

The program first examines the lower/upper case register at 53272 by calling the subroutine at 60090. If in lower case, LC is assigned a value of 32 (note lower case 'a' in character set 2 has a PEEK value of 1 which is standard ASCII is 97 - that is, bit 5 set). This adjustment will be applied to all letters between 'a' and 'z'. The whole dump is enclosed within two loops: I for the rows and J for the columns. All screen codes are ANDed with 127 to reduce them to values in the range 0 to 127 to eliminate reversed characters. If the screen code is <32, we have to add 64 and the LC adjustment. If it lies between 32 and 65, we can print it unchanged. Only if in lower case mode do we need to check for upper case letters. If we were in upper case, these would be nonprinted graphic characters. If in LC then the ANDed code is already in standard ASCII. If all the tests have failed, we have a graphic character so we replace it by a space to maintain the layout of printable text. Once a screen line has been processed we print it preceded by 20 spaces and recycle for all remaining 24 lines.

```
10 OPEN 129,2,0,CHR$(6)

60000 GOSUB 60090

60010 FOR I=0 TO 24:A$="":FOR J=0 TO 39:

CH=PEEK(S+I*40+J)

60020 CH=CH AND 127

60030 IF CH(32 THEN CH=CH+64+LC:GOTO 600

70

60040 IF CH(65 GOTO 60070

60050 IF CH(91 AND LC GOTO 60070

60060 CH=G

60070 A$=A$+CHR$(CH):NEXT J

60080 PRINT#P,SPC(SP);A$:NEXT I:CLOSE P:

RETURN
```

60090 P=129:SP=20:G=32 60100 LC=0:IF PEEK(53272)=23 THEN LC=32 60110 S=PEEK(648)*256:RETURN

Whenever a dump is required, simply GOSUB 60000. This could be actioned by, say, a GET statement, but should not add to the display, or if it does then only 24 lines should be printed. To improve the presentation, blank lines or a form feed should be issued at the end of the dump.

Machine code printer dump

The logic of this routine is identical to that above and is therefore not described in detail. The differences are that it is much faster and it does not pad a line with 20 leading spaces.

The routine as written assumes logical file 2 is open to the printer at the time of calling. To change this, simply alter the byte at \$C001 with a POKE. It works by changing the output device through the CHKOUT KERNAL call to that associated with file #2 (the equivalent of a CMD from BASIC). This then allows us to use the KERNAL routine CHROUT to output the data. There is a routine in ROM which could be used to do most of the conversion, but for this exercise the technique used here is adequate and easier to follow. The device need not be the printer and could be the disk or tape depending on the OPEN statement. We do not recommend you use this routine with anything other than a printer as far better screen saves follow. Once the dump is complete, the default device for output is restored to the screen before returning to BASIC.

The routine is used by at some point including an OPEN 2,4 or OPEN 2,2,CHR\$() if using RS232. A simple sys 49152 will perform the dump. If your printer requires a forced line feed, make the necessary adjustment to \$C001 for a value greater than 127.

BASIC loader for the machine code

The following must be loaded and run. Once this has been done the code remains present until overwritten by something else. Once run the machine code may be saved using an extended monitor for ease of loading later.

```
1 DATA 162, 2, 32, 201, 255, 173, 136, 2
, 133, 88, 169
2 DATA 0, 133, 87, 173, 24, 208, 201, 21
, 208, 6
3 DATA 169, 0, 133, 89, 240, 4, 169, 32,
133, 89
4 DATA 169, 32, 133, 90, 24, 165, 88, 10
5, 3, 133
5 DATA 91, 162, 4, 160, 0, 177, 87, 41,
127, 24
```
6 DATA 201, 31, 176, 7, 24, 105, 64, 101 , 89, 144 7 DATA 24, 24, 201, 64, 176, 2, 144, 17, 24, 165 8 DATA 89, 240, 10, 177, 87, 24, 201, 91 , 176, 3 9 DATA 24, 144, 2, 169, 32, 32, 210, 255 , 224, 1 10 DATA 208, 4, 192, 232, 240, 23, 230, 96, 201, 40 11 DATA 208, 9, 169, 13, 32, 210, 255, 1 69, 0, 133 12 DATA 96, 200, 208, 187, 230, 88, 202, 208, 182, 169 13 DATA 13, 32, 210, 255, 162, 0, 32, 20 1, 255, 96 14 DATA 0, 255, 255, 0, 0, 255, 255, 0, 0, 255 15 FOR I=49152 TO 49292:READ A:POKE I.A: NEXT I

Here is the assembly listing which is fully annotated to allow you to follow it:

C000	A202	LDX	#\$02	log file to printer
C000	20C9FF	JSR	\$FFC9	perform CMD2 via CHKOUT
C005	AD8802	LDA	\$0288	screen start from HIBASE
C008	8558	STA	\$58	set start registers
C00A	A900	LDA	#\$00	
C00C	8557	STA	\$57	
C00E	AD18D0	LDA	\$D018	check upper/lower case
C011	C915	CMP	#\$15	is it upper
C013	D003	BNE	\$C01B	no
C015	A900	LDA	#\$00	set adjustment value
C017	8559	STA	\$59	for ASCII
C019	F004	BEQ	\$C01F	skip lower case
C01B	A920	LDA	#\$20	lower case set adj flag
C01D	8559	STA	\$59	as ASCII a=97 etc.
C01F	A920	LDA	#\$20	set non-printable flag
C021	855A	STA	\$5A	to a space
C023	18	CLC		
C024	A558	LDA	\$58	set MSB end of screen
C026	6903	ADC	#\$03	
C028	855B	STA	\$5B	
C02A	A204	LDX	#\$04	almost 4 pages/screen
C02C	A000	LDY	#\$00	counter within page

C02E	B157	LDA	(\$57),Y	get byte
C030	297F	AND	#\$7F	eliminate high bit 7
Remem	ber dif	ference	between s	creen and ASCII codes
C032	18	CLC		start checks
C033	C91F	CMP	#\$1F	less than a space
C035	B007	BCS	\$C03E	no go to next check
C037	18	CLC		
C038	6940	ADC	#\$40	make ASCII by adding 64
C03A	6559	ADC	\$59	add lower case adj.
C03C	9018	BCC	\$C056	always taken
C03E	18	CLC		
C03F	C940	CMP	#\$40	check for upper case in
C041	B002	BCS	\$C045	1/c mode & branch > 65
C043	9011	BCC	\$C056	1-? same in both sets
C045	18	CLC		check upper case
C046	A559	LDA	\$59	if zero
C048	F00A	BEQ	\$C054	branch to avoid graphic
C04A	B157	LDA	(\$57),Y	get 1/c byte again
C04C	18	CLC		check not gt Z
C04D	C95B	CMP	#\$5B	-
C04F	B003	BCS	\$C054	if so avoid graphic
C051	18	CLC		<u> </u>
C052	9002	BCC	\$C056	valid A-Z so skip space
C054	A920	LDA	#\$20	
C056	20D2FF	JSR	\$FFD2	print char
C059	E001	CPX	#\$01	on last page
C05B	D004	BNE	\$C061	no – so branch
C05D	C0E8	CPY	#\$E8	yes so check end \$**E8
C05F	F017	BEQ	\$C078	branch all done
C031	E660	INC	\$60	end of screen line reg
C063	C928	CMP	#\$28	is it 40 dec
C065	D009	BNE	\$C070	no so skip next bit
C037	A90D	LDA	#\$0D	output next bit
C069	20D2FF	JSR	\$FFD2	print it
C09C	A900	LDA	#\$00	rezero end of line reg
C09E	8560	STA	\$60	
C070	C8	INY		continue current page
C071	DØBB	BNE	\$C02E	branch if not finished
C073	E658	INC	\$58	inc next page register
C075	CA	DEX		
C076	D0B6	BNE	\$C02E	always taken
C078	A90D	LDA	#\$0D	RETURN for last line
C07A	20D2FF	JSR	\$FFD2	
C07D	A200	LDX	#\$00	restore screen output
C07F	20C9FF	JSR	\$FFC9	
C082	60	RTS		

To improve this, why not patch into the interrupt routine to, for example, dump the screen whenever a designated key is pressed rather than using the sys command? Chapter 4 explains the interrupt in detail and Chapter 10 gives an example of its use. If you decide to do this, remember to include a routine to disable the patch. The necessary enable and disable routines can be added at the end of the code as given. The logical file will still have to be OPENED unless the appropriate KERNAL routines are called.

Screen dumps

Three ways are given to save the screen and its associated colour map in this section. Two are in BASIC and the third is in machine code. Both BASIC programs use a sequential file to store the data, but differ in the length of file produced. The machine code saves the screen as a program file and is the most economical and by far the quickest.

A few points should be made before discussing the routines in detail. Any area of memory may be saved from BASIC by setting TXTTAB and VARTAB to its start and end addresses. The problem is that once we have changed these pointers we have temporarily lost our program. Another problem is that a LOAD will cause BASIC to warm-start, which is this case will be at the newly set TXTTAB address. The screen is an area of memory and may be loaded and saved in this way. Unfortunately, its default position is below the normal start of BASIC so a 'crash' or 'hang-up' is usually the result. Try it and see. So from a practical viewpoint we must resort to other means.

All the following routines check HIBASE for the current screen location. The resulting screens will always reload to the current screen position regardless of its location at the time of saving. The reloaded screen will be identical to that saved in both characters and colours.

Screen save using numbers

This routine firstly PEEKS out the border and background colours and writes them, as numbers, to a disk file (change the OPEN command for tape). It then proceeds, writing alternate screen and colour values until finished.

To save a screen: GOSUB 60000 To load a screen: GOSUB 60050

```
60000 OPEN 2,8,2,"@0:TEST,S,W"
60010 S=PEEK(648)*256:C=55296
60020 PRINT#2,PEEK(53280);",";PEEK(53281
);CHR$(13);
60030 FOR I=S TO S+999:PRINT#2,PEEK(I);"
,";PEEK(C+I-S);CHR$(13);:NEXT I:CLOSE 2
```

```
60040 RETURN
60050 OPEN 3,8,3,"TEST,S,R"
60060 INPUT#3,A,B:POKE 53280,A:POKE 5328
1,B
60070 S=PEEK(648)*256:C=55296
60080 FOR I=S TO S+999:INPUT#3,A,B:POKE
I,A:POKE C+I-S,B:NEXT I:CLOSE 3:RETURN
```

Because numbers are written as their ASCII codes three to five bytes are used for each value (spaceXXXreturn). Therefore, using this method we will generate a sequential file of between 6 and 10K, which seems rather excessive. The second method reduces the size of this file.

Screen save using characters

This time a single byte is used to store each value in the screen and colour maps. This is done by simply PEEKing the value and generating the corresponding CHR\$() character with the ASC() function. Zero values must be trapped as ASC(0) will give a syntax error. The resulting file uses only one byte for most values and the file size is therefore about 2K. This is obviously far faster to generate and restore.

To save a screen: GOSUB 60000 To load a screen: GOSUB 60050

```
60000 OPEN 2,8,2,"20:TEST,S,W"
60010 S=PEEK(648)*256:C=55296
60020 PRINT#2, CHR$(PEEK(53280)); CHR$(PEE
K(53281));
60030 FOR I=S TO S+999:PRINT#2,CHR$(PEEK
(I));CHR$(PEEK(C+I-S));:NEXT I:CLOSE 2
60040 RETURN
60050 OPEN 3,8,3,"TEST,S,R"
60060 GET#3,A$:IF A$="" THEN A$=CHR$(0)
60070 POKE 53280,ASC(A$)
60080 GET#3,A$:IF A$="" THEN A$=CHR$(0)
60090 POKE 53281,ASC(A$)
60100 S=PEEK(648)*256:C=55296
60110 FOR I=S TO S+999:GET#3,A$:IF A$=""
 THEN A = CHR (0)
60120 POKE I, ASC(A$):GET#3, A$:IF A$="" T
HEN A = CHR (0)
60130 POKE C+I-S,ASC(A$):NEXT I:CLOSE 3:
RETURN
```

Machine code screen save

This is by far the best method. It is very simple to use the KERNAL LOAD and SAVE for both the screen and colour maps. Using these as they stand, two files would be generated – one for the colour map and one for the screen. This is no hardship, but a relocated load would be required if a screen is being restored to a different location from whence it came. This is not difficult, but perhaps is not the best way.

We have approached the problem slightly differently. Before performing the save, the screen and colour maps are combined into a 2K block at a convenient address. This has to be out of the way of BASIC to avoid corrupting program or data areas. This could be a reserved area at the end of BASIC or even under BASIC ROM if a switch like that used in the UTILITY is implemented to throw out and restore ROM. This is possible as no BASIC ROM calls are made. For this example we have chosen to move the screen from its current position to \$C400 and the colour map to \$C800. The routine also saves the sprite pointers and if you do not wish it to do so then you will have to modify the code to move 1000 rather than its current 1024 bytes from or to each area.

All the routine does to save is to move both screen and colour maps then use the KERNAL SAVE from \$C400 to \$CC00. To restore the screen it is reloaded to \$C400 and moved back to the colour map and the current screen position.

BASIC loader for screen save

The following must be loaded and run before it can be called. Once run it may be saved using an extended monitor for ease of loading later.

1 DATA 32, 253, 174, 201, 76, 208, 6, 16 9, 0, 133, 87 2 DATA 240, 11, 201, 83, 240, 3, 32, 8, 175, 169 3 DATA 255, 133, 87, 32, 115, 0, 32, 253 174, 201 4 DATA 34, 240, 3, 32, 8, 175, 32, 115, 0, 165 5 DATA 122, 133, 187, 165, 123, 133, 188 , 160, 0, 177 6 DATA 122, 201, 34, 240, 8, 200, 192, 1 9, 208, 245 7 DATA 32, 8, 175, 132, 183, 152, 24, 10 1, 122, 133 8 DATA 122, 144, 2, 230, 123, 32, 115, 0 32, 253 9 DATA 174, 144, 3, 32, 8, 175, 56, 233, 48, 133

10 DATA 186, 169, 1, 133, 184, 133, 185, 169, 0, 133 11 DÁTA 88, 133, 90, 133, 92, 133, 94, 1 73, 136, 2 12 DATA 133, 89, 169, 196, 133, 91, 169, 216, 133, 93 13 DATA 169, 200, 133, 95, 165, 87, 240, 43, 160, 0 14 DATA 162, 4, 177, 88, 145, 90, 177, 9 2, 145, 94 15 DATA 200, 208, 245, 230, 89, 230, 91. 230, 93, 230 16 DATA 95, 202, 208, 234, 166, 94, 164, 95. 169. 196 17 DATA 133, 91, 169, 90, 32, 216, 255, 32, 115, 0 18 DATA 96, 165, 87, 32, 213, 255, 160, 0, 162, 4 19 DATA 177, 90, 145, 88, 177, 94, 145, 92, 200, 208 20 DATA 245, 230, 89, 230, 91, 230, 93, 230, 95, 202 21 DATA 208, 234, 32, 115, 0, 96, 0, 255 255.0 22 FOR I=49152 TO 49362:READ A:POKE I.A: NEXT I nb"basic2"

To save a screen: SYS 49152,S, "filename",DEVICE SYS 49152,S, "@0:filename",DEVICE to replace on disk To load a screen: SYS 49152,L,"filename", DEVICE

All parameters are required and an illegal or missing parameter will produce a SYNTAX ERROR. A file name or a minimum of "" is required, even with cassette. Remember all bytes following a SYS command are ignored.

The following is the assembly listing, which should be selfexplanatory. The first part of the routine is our own version of 'GET PARAMETERS' and is a useful technique when passing parameters to a routine enabled with a SYS. CHRGET (see Chapter 3) is used to gather the necessary bytes.

ASSEMBLY LISTING

Set up	o the para	meters	s – commo	n to both LOAD and SAVE
C000	20FDAE	JSR	\$AEFD	check for comma
C003	C94C	CMP	#\$4C	is next char L for LOAD
C005	D006	BNE	\$C00D	no then test for SAVE
C007	A900	LDA	#\$00	set flag at 57
C009	8557	STA	\$57	to zero for later use
C00B	F00B	BEQ	\$C018	always taken if LOAD
C00D	C953	CMP	#\$ 53	is it S for SAVE ?
C00F	F003	BEQ	\$C014	if so continue
C011	2008AF	JSR	\$AF08	not L or S=SYNTAX ERROR
C014	A9FF	LDA	#\$FF	set flag to FF for save
C016	8557	STA	\$57	-
C018	207300	JSR	\$0073	inc CHRGET (see ch.4)
C01B	20FDAE	JSR	\$AEFD	next comma
C01E	C922	CMP	#\$22	opening quote of name
C020	F003	BEQ	\$C025	OK so continue
C022	2008AF	JSR	\$AF08	no quote so SYNTAX ERROR
C025	207300	JSR	\$0073	inc CHRGET to name
C028	A57A	LDA	\$7A	set FNADR low byte
C02A	85BB	STA	\$BB	
C02C	A57B	LDA	\$7B	do same for high
C02E	85BC	STA	\$BC	
C030	A000	LDY	#\$00	find length of name
C032	B17A	LDA	(\$7A),Y	by searching for closing
C034	C922	CMP	#\$22	quote
C036	F008	BEQ	\$C040	found it so exit
C038	C8	INY		
C039	C013	CPY	#\$13	limit of 16 chars +3 for
C03B	D0F5	BNE	\$C032	*30:" toreplace on disk
C03D	2008AF	JSR	\$AF08	>limit so SYNTAX ERROR
C040	8487	STY	\$B7	store length in FNLEN
C042	98	TYA		
C043	18	CLC		
C044	657A	ADC	\$7A	set CHRGET to end quote
C046	857A	STA	\$7A	low byte
C048	9002	BCC	\$C04C	
C04A	E67B	INC	\$7B	inc high if page crossed
C04C	207300	JSR	\$0073	get next byte
C04F	20FDAE	JSR	\$AEFD	comma ?
C052	9003	BCC	\$C057	OK
C054	2008AF	JSR	\$AF08	no then SYNTAX ERROR
C057	38	SEC		
C058	E930	SBC	#\$30	make byte a number 0-9
C05A	85BA	STA	\$BA	store current device-FA
C05C	A901	LDA	#\$01	set secondary add to 1

C05E	8588	STA	\$B8	and store at SA
C030	8589	STA	\$B9	same for logical file-LA
Set u	p pointers	to be	e used in	the move
CØ62	A900	LDA	#\$00	
C034	8558	STA	\$58	
C036	855A	STA	\$5A	
C038	8550	STA	\$5C	
C03A	855E	STA	\$5E	
C09C	AD8802	LDA	\$0288	find the current screen
C03F	8559	STA	\$59	start from HIBASE
C071	A9C4	LDA	#\$C4	
C073	855B	STA	\$5B	
C075	A9D8	LDA	#\$ D8	
C077	855D	STA	\$5D	
C079	A9C8	LDA	#\$ C8	
C07B	855F	STA	\$5F	
C07D	A557	LDA	\$57	0 for LOAD
C07F	F02B	BEQ	\$C0AC	do LOAD
Move	screen and	colou	ir to one	block and perform SAVE
C081	A000	LDY	#\$00	SAVE use Y within page
C083	A204	LDX	#\$04	and X for page counter
C085	B158	LDA	(\$58),Y	read byte from screen
C087	915A	STA	(\$5A) Y	store at combined area
C089	B15C	LDA	(\$5C),Y	get char colour
C08B	915E	STA	(\$5E),Y	store in comb area+\$0400
C08D	C8	INY		
C08E	D0F5	BNE	\$C085	cycle for one page
C090	E659	INC	\$59	inc all high bytes
C092	E65B	INC	\$5B	
C094	E65D	INC	\$5D	
C096	E65F	INC	\$5F	
C098	CA	DEX		dec X and
C099	DØEA	BNE	\$C085	repeat till 4 pages done
C09B	A65E	LDX	\$5E	X holds low end of save
C09D	A45F	LDY	\$5F	Y the high byte
C09F	A9C4	LDA	#\$C4	use 5A/5B on zero page
C0A1	855B	STA	\$5B	for start of SAVE
C0A3	A95A	LDA	#\$5A	A must hold offset 5A
CØA5	20D8FF	JSR	\$FFD8	do SAVE
C0A8	207300	JSR	\$0073	must inc CHRGET before
CØAB	60	RTS		returning to BASIC
Perfor	rm LOAD and	lispli	t block i	nto char and colour maps
CØAC	A557	LDA	\$57	read flag - A=0 for LOAD
CØAE	20D5FF	JSR	\$FFD5	do LOAD
C0B1	A000	LDY	#\$00	
C0B3	A204	LDX	#\$04	

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C0B5	B15A	LDA	(\$5A),Y	reverse of SAVE
C087	9158	STA	(\$58),Y	even if the screen
C0B9	B15E	LDA	(\$5E),Y	was at a different
COBB	915C	STA	(\$5C),Y	location at the time of
COBD	C8	INY		SAVE it will go to the
CØBE	D0F5	BNE	\$C0B5	current position
C0C0	E659	INC	\$59	
C0C2	E65B	INC	\$5B	
C0C4	E65D	INC	\$5D	
6000	E65F	INC	\$5F	
C0C8	CA	DEX		
C0C9	DØEA	BNE	\$C0B5	
COCB	207300	JSR	\$0073	must inc CHRGET before
C0CE	60	RTS		returning to BASIC
Disk				

This section deals with the 1541 disk drive though much is directly applicable to the 3040 and 4040 units.

The manual supplied with the 1541 contains all the information that most users will require. Perhaps the most difficult to master are the direct access programming commands such as BLOCK-READ, and so on. There is only one way to become proficient in their use and that is to experiment. When experimenting we suggest you use a disk containing unwanted information as disasters can happen.

We supply only one utility in this section which we like to think of as an expandable disk utility. Once direct access programming is mastered, there are all sorts of fun things you can do. To use it to its best advantage you have to know something of how the disk operates and how information is stored. To this end we give below a very short introduction and would refer you to the 1541 manual itself.

Introduction

The 1541 is a self-contained, intelligent device. It has two processors and its disk operating system (DOS) in ROM along with an area of RAM used for input/output (buffering) operations. This differs from, say, the BBC Micro where the interface is within the micro itself and, depending on the type of interface, removes RAM from the user area. The disadvantage to this self-contained arrangement is that you cannot use non-Commodore units (there are one or two now available) as most disk manufacturers do not supply a suitable controller and DOS.

Almost the whole of the disk's capacity can be used to store data except for one or two reserved areas. The disk is divided into tracks which are further subdivided into a varying number of sectors. Tracks are numbered from 1 through to 35 whereas sectors start at zero. The following is the arrangement for both the 1541 and 4040 units (see Table 6.1, 1541 manual):

TRACK	SECTOR	TRACK	SECTOR
1–17	0-20	25–30	0 –17
18–24	0 –18	31-35	Ø –16

In order to know where to find information, the disk uses an index or directory track. This is track 18 and it has two special areas. The first is track 18 sector 0 and is the Block Availability Map (BAM). This keeps a record of all sectors in use unless direct access programming operations have been used without a BLOCK-ALLOCATE command. If this is the case then the information is there, but can be overwritten as it is empty as far as DOS is concerned. The entries in BAM are made up as follows (see Table 5.1 of manual):

BYTES **CONTENTS**

START BYTES

000001	Pointer to start of directory 18/01
002	Holds an 'A' for 4040 format
004–143	Four bytes for each of the 35 tracks
	indicating whether in use 1=free Ø=allocated
144–161	Disk name plus shifted spaces to make 16 in total
162–163	Disk ID
165–166	Disk version of 2A

Each track uses four bytes in BAM. The first stores the number of free sectors on a track and is used in computing BLOCKS FREE. The remaining three are used to indicate whether a particular sector is allocated (bit set low and one bit per sector). As the maximum number of sectors is 21, not all bits are used. The following is a dump of BAM from which you can pick out the information given above. All values are given in hex with the byte position within the sector given first followed by this byte's value and the next seven and at the end of the line the equivalent ASCII characters (if printable):

START	BY	TES							ASCII
00	12	01	41	00	15	ff	ff	1f	a
Ø8	15	ff	ff	1f	15	ff	ff	1f	
10	15	ff	ff	1f	15	ff	ff	1f	
18	15	ff	ff	1f	15	ff	ff	1f	
20	15	ff	ff	1f	15	ff	ff	1f	
28	15	ff	ff	1f	00	00	00	00	
30	00	00	00	00	00	00	00	00	•••••
90	42	4f	4f	4b	20	50	52	4f	book pro
98	47	52	41	4d	53	aØ	aØ	aØ	grams
aØ	aØ	aØ	31	31	aØ	32	41	aØ	Ĩ1 2a

The file information starts on track 18 sector 1 and can continue throughout the remainder of the track. Each file uses 32 bytes. Therefore, one sector can hold information for eight entries. With a possible 20 sectors available, information could be held for 160 files. This is unlikely to happen as each file would have to be less then 1K. The directory format is such that bytes 0–31 hold file 1, 32–63 hold file 2, and so on. Each entry is divided up as follows (see Table 5.3 of the 1541 manual):

BYTE CONTENTS

- 000–001 Next directory track and sector. A track of 0 indicates last sector in use. These bytes only used for the first entry.
- 002 The type of file \$00=scratched or not in use.
 \$80=DELeted (scratched unclosed) \$81=sEQuential
 \$82=PROgram \$83=USER \$84=RELative \$1-4=unclosed
- 003–004 Starting track and sector of file
- 005–020 Name padded with shifted spaces
- 023 Record size of relative file
- 028–029 New track and sector for disk ops with replacement @
- 030-031 Number of blocks file uses in low/high byte format

Below is a typical dump of the first directory sector, track 18 sector 1, for two file entries:

START	CC	NT		ASCII					
00	12	Ø 4	82	11	00	44	55	4d	dum
Ø 8	50	2e	4d	49	4b	aØ	aØ	aØ	p.mik
10	aØ	aØ	aØ	aØ	aØ	00	00	00	
18	00	00	00	00	00	00	Ød	00	
20	00	00	82	11	03	44	55	4d	dum
28	50	2e	4d	aØ	aØ	aØ	aØ	aØ	p.m
30	aØ	aØ	aØ	aØ	aØ	00	00	00	
38	00	00	00	00	00	00	01	00	

It is worth noting that directory sectors do not follow sequentially. The same is true for file storage, as can be seen when using the disk utility.

Just to round things off, here is a dump of a BASIC program which occupies less than one block. It is in fact the loader for the UTILITY at the end of Chapter 9.

START CONTENTS

00	00	4B	Ø1	Ø8	24	Ø8	ØA	00	.K\$
Ø8	41	B2	41	AA	31	3A	8B	41	A.A.1:.A
10	B2	31	A7	93	22	55	54	49	.1"UT1
18	4C	49	54	59	20	44	41	54	LITY DAT
20	41	22	2C	38	2C	31	00	3C	A'',8,1.<
28	Ø8	14	00	8B	41	B2	32	A7	A.2.
30	93	22	55	54	49	4C	49	54	."UTILIT
38	59	22	2C	38	2C	31	00	47	Y'',8,1.G
40	Ø8	1E	00	9E	33	32	37	36	3276
48	38	00	00	00	A2	00	00	00	8

As it is less than one block, the linking track is zero, denoting the end. It is a straight copy of the RAM and like the memory dump in Chapter 1, we can pick out the link addresses, end of program, and so on.

Disk utility

This utility offers many of the housekeeping commands and provides a number of more interesting options. It is rather long as most of the subroutines are complete in themselves (to allow you to extract only those you want for your own programs). The listing has been left in lower case and when you are typing it in, it is easiest to put the 64 into that mode with the shift and logo keys. It makes extensive use of direct access programming so we suggest you use the information given above and the relevant sections of the 1541 manual to follow it. It has been run through the UTILITY'S CODER command to produce the mnemonics. Most annotated characters are cursor moves, colours or simply capital letters.

The usual options of NEW, VALIDATE, SCRATCH, INITIALIZE, RENAME, COPY (within a drive) and READ DISK ERROR are all present. The directory option is unusual in that everything is input or displayed in hex notation. A much shorter way to get a directory is given in the BACKUP utility at the end of this chapter. The option also displays the first track and sector of a file, and if it is a program, also its load address. The listing is further split up into directory sectors and will display even SCRATCHEd or DELETED files if the disk has not been VALIDATED. Two values are given for the BLOCKS FREE – the usual value exclusive of erased files and another inclusive of erased files. An erased file simply has its associated BAM set to 1 (not allocated).

The TRACE option follows a file through displaying its associated tracks and sectors. It will also check to see if the file it is following is scratched. If this is the case, it will ask whether you wish to recover it. If your answer is 'yes', then as it traces it will also allocate each block. Providing that all the blocks found were free it has recovered the file. If an allocated block is found then the original area of the file has been

overwritten and recovery is not possible and you will be told. If the scratched file has been successfully traced, all that remains to be done is to use the MOD/DISP BLOCKS option to change its 'file type' byte (third byte in its entry) from 0 back to \$81 or \$82. The revised directory block must be rewritten to complete the process.

The MOD/DISP BLOCKS option is similar to the demonstration disk's program 'DISPLAY TRACK AND SECTOR'. The main difference is that it also allows the block to be modified and rewritten to disk. When the block is written it is also allocated. The usual options to review again and get the next track and sector are available. The subroutine called at 1680 is a little unusual and merits some explanation. Earlier in the chapter under the heading 'Keyboard' it was mentioned that a file could be opened to it. This eliminates the '?' prompt and also releases the cursor. The cursor may then be moved to the appropriate line, hex values changed and on pressing RETURN the revised values are processed and written to the disk buffer. The same 128 bytes are then redisplayed by reading from the buffer. At the end of a block the option to write the changes must be taken to change it on disk. It is also possible to recover files using this option by following a file through taking the next track and sector option (first two bytes) and always writing the changes. Unlike TRACE, it does not check to see if a file can be recovered. Files in which a READ ERROR occurs may also be reconstructed. This we discovered when the EASY SCRIPT appendices file of this book was corrupted. All we did was modify the next track and sector bytes of the preceding block to skip the corrupt block. The resulting file could be read with the loss of only 256 bytes (and was immediately saved on another disk).

The APPEND for program files is the same as that in Chapter 5 (where it is fully explained). The APPEND for sequential files (and SCRATCH) builds the command string (separated by commas) before actioning on RETURN with no input.

The final option is to modify the disk's header. This is done by simply reading BAM, moving the buffer pointer and writing in the new values. It is worth noting that whenever a byte is read, the buffer pointer is moved forward one position. So in order to write to the same position at which the read started, the pointer must be set using a 'B-P' command.

The utility is not foolproof, but with a little attention to detail, may be used to advantage. Our last comment before the listing is to point out that when you try to allocate an already allocated block error 65 NO BLOCK occurs. This must be checked for and trapped as in MOD/DISP BLOCKS. The locations of all the subroutines may be read from the IF statements in lines 210–340.

```
100 poke 53272,23:poke 53280,6:dim a$(10
0),t$(5)
110 for i=0 to 5:read t$(i):next i
120 data del,seq,prg,usr,rel,???
130 print"[c]s][q>d][q>i][q>s][q>k] [q>u
][q>t][q>i][q>i][q>i][q>i][q>t][q>y]":print"[
cd][rev][q>n][off] new disk"tab(20);"[re
v][q>h][off] change header"
140 print"[cd][rev][g>v][off] validate d
isk";tab(20);"[rev][g/d][off] directory"
150 print "[cd][rev][g>t][off] trace fil
e";tab(20);"[rev][g>s][off] scratch file
(s)"
160 print"[cd][rev][g>r][off] rename fil
e";tab(20);"[rev][g>e][off] read disk er
ror"
170 print"[cd][rev][g>c][off] copy file"
;tab(20);"[rev][g>a][off] append files"
180 print"[cd][rev][g>b][off] backup fil
e";tab(20);"[rev][g>m][off] mod/disp blo
cKs"
190 print"[cd][rev][g>i][off] initialize
 disk";tab(20);"[rev][g)x][off] exit"
200 gosub 2360
210 if y$="n" then gosub 530:goto 130
220 if y$="v" then gosub 580:goto 130
230 if y$="r" then gosub 620:goto 130
240 if y$="s" then gosub 670:goto 130
250 if y$="e" then gosub 740;goto 130
260 if y$="a" then gosub 360:goto 130
270 if y$="c" then gosub 840:goto 130
280 if y$="d" then gosub 890:goto 130
290 if ys="h" then gosub 1240:goto 130
300 if y$="x" then end
310 if y$="m" then gosub 1400:goto 130
320 if y$="i" then gosub 1740:goto 130
330 if y$="t" then gosub 1760:goto 130
340 if y$="b" then gosub 2130:goto 130
350 goto 130
360 rem append
370 print"[cls][q>a][2q>p][q>e][q>n][q>d
][cd]":print"[rev][q>p][off] prg files[c
d]":print"[cd][rev][g>s][off] seq files"
380 gosub 2360:if y$<>"p" and y$<>"s" go
to 520
390 if y$="s" then gosub 770:return
```

```
400 rem prg files
410 print"[cd]append prg files - sure":g
osub 2360:if y$<>"y" goto 520
420 input"[cd]combined prg";f$:input"[3s
pc]first prg";x1$:input"[2spc]second prg
";×2$
430 open 3,8,3,"0:"+f$+",p,w":open 2,8,2
,"0:"+x1$+",p,r"
440 get#2,y$
450 x$=y$:qet#2,y$:if st<>0 goto 470
460 gosub 2330:print#3,x$;:goto 450
470 close 2:open 2,8,2,"0:"+x2$+",p,r"
480 get#2,y$:get#2,y$
490 get#2,x$:if st<>0 goto 510
500 gosub 2330:print#3,x$;:goto 490
510 close 2:print#3,chr$(0);:close 3
520 return
530 rem new disk
540 print"[cd]new disk - sure":gosub 236
0:if y$<>"y" goto 570
550 input"[cd]name";f$:input"i.d.";y$:f$
=left$(f$,16)+","+left$(y$,2)
560 open 15,8,15,"n0:"+f$:close 15
570 return
580 rem validate
590 print"[cd]validate - sure":gosub 236
0:if y$<>"y" goto 610
600 open 15,8,15,"v":close 15
610 return
620 rem rename
630 print"[cd]rename - sure":gosub 2360:
if y$<>"y" goto 660
640 input"[cd]old file";f$:input"new fil
e";y$:f$="r0:"+y$+"="+f$
650 open 15,8,15,f$:close 15
660 return
670 rem scratch file
680 print"[cd]scratch - sure":gosub 2360
:if y$<>"y" goto 730
690 fs="":print"[cd]use * or ? for patte
rn matching"
700 print"hit return to delete[cd]"
710 input "delete";y$:if y$<>"" then f$=
f$+","+y$:y$="":goto 710
720 f$="s0:"+mid$(f$,2):open 15,8,15,f$:
close 15
```

```
730 return
740 rem error
750 open 15,8,15:input#15,x$,f$,x1$,x2$:
close 15:print*[cd][g>e][2g>r][g>o][g>r]
:";x$,f$,x1$,x2$
760 gosub 2360:return
770 rem append seq files
780 print"[cd]append seq files - sure":g
osub 2360:if y$<>"y" goto 830
790 print"[cd]hit return to append[cd]"
800 f$="":input"[3spc]new";x$
810 input"append";y$:if y$<>"" then f$=f
$+","+"0:"+y$:y$="":goto 810
820 f$="c0:"+x$+"="+mid$(f$,2):open 15,8
,15:print#15,f$
830 close 15:return
840 rem copy same disk
850 print"[cd]copy - sure":gosub 2360:if
 y$<>"y" goto 880
860 input"[cd]new";f$:input"old";y$:f$="
c0:"+f$+"="+"0:"+y$
870 open 15.8.15.f$
880 close 15:return
890 rem directory
900 print"[cd]directory - sure":gosub 23
60:if y$<>"y" goto 1230
910 open 15,8,15:open 1,8,2,"#":t=18:s=0
:f$="":bf=0:bu=0
920 print#15,"u1";2;0;t;s;:print#15,"b-p
";2;144:print"[c]s]";tab(10);"[rev]";
930 for i=1 to 16:get#1,x$:gosub 2290:pr
intx$;:next i:print",";
940 print#15,"b-p";2;162
950 for i=1 to 2:get#1,x$:gosub 2290:pri
ntx$;:next i:print"[off][b]k]":t=18:s=1
960 print"$blk file[13spc]type $tk $st $
add[l blu]"
970 print#15,"u1";2;0;t;s:i=0:x=s:gosub
2200:print"[b]k]trk 12";" sct ";x$;"[] b
lu]"
980 get#1,x$:gosub 2330:t=asc(x$):get#1,
x$:gosub 2330:s=asc(x$)
990 i=i+1:print#15,"b-p";2;(i-1)*32+2:ge
t#1,x$:gosub 2330:x=asc(x$):y$=""
1000 for j=0 to 5:if x=j then y$=t$(j)
1010 if j=0 then x=x-128
```

```
1020 next j:if y$="" then y$=t$(5)
1030 get#1,x$:gosub 2330:x=asc(x$):gosub
 2200:t$=x$
1040 get#1,x$:gosub 2330:x=asc(x$):gosub
 2200:s$=x$
1050 for j=1 to 16:get#1,x$:gosub 2290;f
$=f$+x$:next j
1060 if y$<>"prg" goto 1090
1070 open 3,8,3,f$+",s,p":get#3,x$:gosub
 2330:x=asc(x$):gosub 2200:1$=x$
1080 get#3,x$:gosub 2330:x=asc(x$):gosub
 2200:1$=x$+1$:close3
1090 print#15,"b-p";2;(i-1)*32+30
1100 get#1.x$:gosub 2330:j=asc(x$)
1110 get#1,x$:gosub 2330:k=asc(x$)
1120 x=k:gosub 2200:bf$=x$:x=j:gosub 220
0:bf$=bf$+x$
1130 bu=bu+256*k+j:if y$<>"del" then bf=
bf+256*k+j
1140 printbf$;" ";f$;"[2spc]";y$;"[2spc]
";t$;"[2spc]";s$;" ";]$:bf$="";f$="";]$=
11 11
1150 get y$:if y$<>"" then gosub 2360
1160 if i<8 goto 990
1170 if t<>0 goto 970
1180 bf=664-bf:x=bf/256:gosub 2200:y$=x$
:x=bf-256:gosub 2200:y$=y$+x$
1190 print"[yel]";y$;" blocks free[] blu
יי נ
1200 bu=664-bu:x=bu/256:gosub 2200:y$=x$
:x=bu-256:gosub 2200:y$=y$+x$
1210 print"[yel]";y$;" blocks free inc d
el files[] blu]"
1220 gosub 2360:close 15:close 1
1230 return
1240 rem change header
1250 print"[cd]change header - sure":gos
ub 2360:if y$<>"y" goto 1390
1260 open 15,8,15:open 1,8,2,"#":t=18:s=
0:bf=0:f$=""
1270 print#15,"u1";2;0;t;s;:print#15,"b-
p";2;144:print"[cd]current:[2spc]":
1280 for i=1 to 16:get#1,x$:gosub 2290:p
rintx$;:next i:print",";
1290 print#15,"b-p";2;162
```

```
1300 for i=1 to 2:get#1,x$:gosub 2290:pr
intx$::next i:print
1310 input"[4spc]name";f$:f$=left$(f$+"[
16g/spc]",16)
1320 input"[4spc]i.d.";y$:y$=left$(y$+"x
x",2)
1330 print#15,"b-p";2;144
1340 for i=1 to 16:print#1,mid$(f$,i,1);
:next i
1350 print#15,"b-p";2:162
1360 print#15,"b-p";2;162
1370 for i=1 to 2:print#1,mid$(y$,i,1);:
next i
1380 print#15,"u2";2;0;t;s:close 15:clos
e 1
1390 return
1400 rem modify and display blocks
1410 print"[cd]modify and display blocks
- sure":gosub 2360:if y$<>"y" goto 1670
1420 open 15,8,15:open 1,8,2,"#":f$=""
1430 input"[cd] track[cr]$[2c1]";t$:x$=t
$:gosub 2250:t=x:if x<0 or x>40 goto 143
Й
1440 input"sector[cr]$[2c1]";s$:x$=s$:go
sub 2250:s=x:if x<0 or x>20 goto 1430
1450 print#15,"u1";2;0;t;s;:print#15,"b-
p";2;0
1460 get#1,x$:gosub 2330:x=asc(x$):tn=x:
gosub 2200:tn$=x$
1470 get#1,x$:gosub 2330:x=asc(x$):sn=x:
oosub 2200:sn$=x$
1480 nt$="[3spc][g>n][g>e][g>x][g>t]:[g>
t][g>r][g>a][g>c][q>k] "+tn$+" [q>s][q>e
][q>c][g>t][q>o][q>r] "+sn$
1490 ct$="[g>c][g>u][2g>r][g>e][g>n][g>t
]:[g>t][g>r][g>a][g>c][g>k] "+t$+" [g>s]
[g>e][g>c][g>t][g>o][g>r] "+s$
1500 print"[cls]";ct$:printnt$
1510 print#15,"b-p";2;0
1520 for i=0 to 15:f$="":x=i*8:qosub 220
0:printx$;" ";:for j=0 to 7
1530 get#1,x$:y$=x$:gosub 2330:x=asc(x$)
:gosub 2200:printx$;" ";:
1540 x$=y$:gosub 2290:f$=f$+x$:next j:pr
intf$:next i
1550 gosub 1680:if y$="y" goto 1500
```

```
1560 print"[c]s]";ct$:printnt$
1570 print#15,"b-p";2;128
1580 for i=16 to 31:f$="":x=i*8:gosub 22
00:printx$;" ";:for j=0 to 7
1590 get#1,x$:y$=x$:gosub 2330:x=asc(x$)
:gosub 2200:printx$;" ";:
1600 x$=y$:gosub 2290:f$=f$+x$:next j:pr
intf$:next i
1610 gosub 1680:if y≸="y" goto 1560
1620 print"review again":gosub 2360:if y
$="y" goto 1500
1630 print"write changes":gosub 2360:if
y$="y" then print#15,"u2";2;0;t,s
1640 if y$="y" then print#15,"b-a";0;t;s
:gosub 2390
1650 print"next t/s":gosub2360:if y$="y"
thent=tn:s=sn:t$=tn$:s$=sn$:aoto 1450
1660 close 1:close 15
1670 return
1680 rem any changes
1690 print"any changes":gosub 2360:if y$
<>"y" goto 1730
1700 open 9,0:input#9,a$:close 9
1710 x$=left$(a$,2):gosub 2250:print#15,
"b-p";2,x
1720 for i=0 to 7:x$=mid$(a$,4+i*3,2):go
sub 2250:print#1.chr$(x)::next i
1730 return
1740 rem intialize disk
1750 open 15,8,15,"i0":close 15:return
1760 rem trace file
1770 print"[cd]trace file - sure":gosub
2360:if y$<>"y" goto 2120
1780 input"[cd]file";bf$
1790 open 15,8,15:open 1,8,2,"#":t=18:s=
1:f$=""
1800 print#15,"u1";2;0;t;s:i=0:f$=""
1810 get#1,x$:gosub 2330:t=asc(x$):get#1
,x$:gosub 2330:s=asc(x$)
1820 i=i+1:print#15,"b-p";2;(i-1)*32+2:g
et#1,x$:gosub 2330:x=asc(x$):y$=""
1830 for j=0 to 5:if x=j then y$=t$(j)
1840 if j=0 then x=x-128
1850 next j:if y$="" then y$=t$(5)
1860 get#1,x$:gosub 2330:x=asc(x$):gosub
 2200:t$=x$
```

```
1870 get#1,x$:gosub 2330:x=asc(x$):gosub
2200:s$=x$
1880 f$="":for j=1 to 16:get#1,x$:gosub
2290:f$=f$+x$:next j
1890 if left$(f$,len(bf$))=bf$ goto 1930
1900 if i=8 and t>0 goto 1800
1910 goto 1820
1920 if t=0 ooto 2120
1930 print"[cls][rev]trace of[off] ";bf$
;" file type ";y$:print:ft$=y$:no=0
1940 if ft$="del" then print"recover fil
e":gosub 2360
1950 print#15,"b-p";2;(i-1)*32+3
1960 get#1,x$:gosub 2330:t=asc(x$):get#1
,x$:gosub 2330:s=asc(x$)
1970 x=t:gosub 2200:t$=x$:x=s:gosub 2200
:s$=x$
1980 print"track ";t$;" sector ";s$
1990 if y$<>"y" goto 2030
2000 print#15,"b-a";0;t;s:input#15,e:if
e<>65 goto 2030
2010 print"recovering not on as a suppos
edly"
2020 print"free block is allocated.":no=
1
2030 print#15,"u1";2;0;t;s
2040 get#1,x$:qosub 2330:t=asc(x$):get#1
,x$:gosub 2330:s=asc(x$)
2050 if t=0 goto 2070
2060 goto 1970
2070 if ft$="del" and y$="y" and no=0 go
to 2090
2080 goto 2110
2090 print"recovery ok remember to chang
e"
2100 print"directory track and file type
2110 gosub 2360
2120 close 1:close 15:return
2130 rem backup
2140 print:print"backup file"
2150 print:print"[g>t]o allow larger fil
es to be backed up"
2160 print"on both disk and cassette a s
eparate"
```

```
Peripherals 55
```

```
2170 print"utility has been provided.[3s
pc][q>f]or "
2180 print"smaller files code could be i
ncluded here."
2190 gosub 2360:return
2200 rem dec-hex
2210 x=x and 255:x1=int(x/16):x2=x and 1
5
2220 x1$=chr$(48+x1):if x1>9 then x1$=ch
r$(55+x1)
2230 x2$=chr$(48+x2):if x2>9 then x2$=ch
r$(55+x2)
2240 x$=x1$+x2$:return
2250 rem hex-dec
2260 x$=right$("00"+x$,2)
2270 x1 = asc(x$) - 48:x2 = asc(mid$(x$,2)) - 48
2280 x=16*(x1+7*(x1>9))+x2+7*(x2>9);retu
rn.
2290 rem convert to ascii
2300 if x$=chr$(160) goto 2320
2310 if x${" or x$><sup>"</sup>z" then x$="."
2320 return
2330 rem eliminate null string
2340 if x$="" then x$=chr$(0)
2350 return
2360 rem wait
2370 get y$:if y$="" goto 2370
2380 return
2390 rem error on b-a check
2400 input#15,en$,em$,et$,es$:if en$<"20
  or en$="65" then return
2410 close 15:print:printen$,em$,et$,es$
:gosub 2360:run
```

Many more options could be provided and some of the existing options could be made fully automatic. These are exercises for you to carry out.

BACKUP

We have produced this utility to allow selective backing-up of files between disk and tape, from disk to disk and from tape to tape. Commodore provide an excellent 1541 BACKUP program on the demonstration disk, but it only backs up whole disks. The following allows selective backing-up of single files, whether they be program or sequential. It could be modified to do more than one file when going between disk and tape, providing that the details of each file were known in advance. We wrote the program to avoid the need to produce a special program for sequential files and the use of an extended monitor to copy machine code. BASIC programs can, of course, be duplicated by a simple load and save sequence.

The program is in two parts: the machine code and the BASIC driver which uses the machine code for program files. The following describes the driver:

LINE ACTION

- 100 Set top of memory to \$1800/6144
- 110 Set source device and check whether valid
- **120** Do same for destination device
- 130 Contents of disk or tape
- 150 Set prog or seq file, if not known use 130
- 170- Go to appropriate subroutine for source, destination and file type

Subroutines

- **250** Seq file from disk to tape so read byte/write byte until status says end of file. Requires a file name.
- **290** Seq tape to disk
- **290** Read header and display info with sub 700 and if non-ASCII in name then offer chance to change name.
- 310 Final check on name.
- 330- Read and write bytes until status says end of file
- 370– Seq tape to tape. Has a limited capacity.
- 370 Check for non-ASCII and option to change name.
- **410** Read in bytes until end of file or until RAM filled eliminating ASC(0) on the way to avoid errors.
- 450 Warning only part of file read.
- 480 Write to destination tape
- **490** Pause for destination tape
- 510– Seq disk to disk. Same principle as for tape above
- 610 Pause for destination disk
- 630 Simple wait for any key
- 650 Print TAPE OF DISK in set up
- 690 Contents of next file on tape and prompt to rewind
- 700 Display tape buffer in full highlighting any non-ASCII
- 750 Get file name

- 760 Eliminate trailing spaces in file name
- 800 Display directory of disk
- 940- Prog file disk to disk.
- 940 Get file name and set it up in cassette buffer
- 950 Fill up rest of buffer with spaces
- 960 Set length of name register FNLEN and enter m/c to do a relocated load
- 980 Delete file from destination disk
- 990 Do relocated save
- **1010** Prog disk to tape. Initial part as for disk to disk
- 1040 Write header created
- 1045 Write RAM
- **1060** Prog tape to tape
- **1100** Prog tape to disk checking as before

BASIC PROGRAM

```
100 POKE 52,24:POKE 56,24:CLR
110 PRINT"[CLS]BACKUP UTILITY":PRINT"[CD
JFROM T/D"::GOSUB 650:F$=Y$
120 PRINT"[CD][2SPC]TO T/D";:GOSUB 650:T
$=Y$
130 PRINT"[CD]CONTENTS T/D":GOSUB 630:IF
 Y$="D" THEN GOSUB 800:GOTO 130
140 IF Y$="T" THEN GOSUB 690:GOTO 130
150 PRINT"[CD]TYPE P/S";:GOSUB 630:IF Y$
<>"P" AND Y$<>"S" GOTO 150
160 FT$=Y$:PRINT " ";FT$
170 IF F$="D" AND T$="T" AND FT$="S" THE
N GOSUB 250:RUN
180 IF F$="T" AND T$="D" AND FT$="S"
                                     THE
N GOSUB 290:RUN
190 IF F$="D" AND T$="D" AND FT$="S" THE
N GOSUB 510:RUN
195 IF F$="T" AND T$="T" AND FT$="S" THE
N GOSUB 370:RUN
200 IF F$="D" AND T$="D" AND FT$="P"
                                      THE
N GOSUB 940:RUN
210 IF F$="D" AND T$="T" AND FT$="P" THE
N GOSUB 1010:RUN
220 IF F$="T" AND T$="T" AND FT$="P" THE
N GOSUB 1060:RUN
230 IF F$="T" AND T$="D" AND FT$="P" THE
N GOSUB 1100:RUN
```

```
240 RUN
250 GOSUB 750:0PEN 2,8,2,N$+",S,R":0PEN
1,1,1,N$
260 GET#2,Y$:IF ST GOTO 280
270 PRINT#1,Y$;:GOTO 260
280 CLOSE 2:CLOSE 1:RETURN
290 OPEN 1,1,0:GOSUB 700:IF E=1 THEN GOS
UB 750:GOTO 310
300 GOSUB 760
310 PRINT"[CD]FILE NAME ";CHR$(34);N$;CH
R$(34):" OK Y/N":GOSUB 630
320 IF Y$="N" THEN GOSUB 750
330 OPEN 3,8,3,"20:"+N$+",S,W"
340 GET#1,Y$:IF ST GOTO 360
350 PRINT#3,Y$;:GOTO 340
360 CLOSE 1:CLOSE 3:RETURN
370 Y=6144:0PEN 1,1,0:GOSUB 700:IF E=1 T
HEN GOSUB 750::GOTO 390
380 GOSUB 760
390 PRINT"[CD]FILE NAME ";CHR$(34);N$;CH
R$(34):" OK Y/N":GOSUB 630
400 IF Y$="N" THEN GOSUB 750
410 GET#1,Y$:IF ST GOTO 460
420 IF Y$="" THEN Y$=CHR$(0)
430 POKE Y,ASC(Y$)
440 IF Y<40959 THEN Y=Y+1:GOTO 410
450 PRINT"[CD]FILE TOO BIG ONLY 34K COPI
ED"
460 CLOSE1
470 GOSUB 490
480 OPEN 1,1,1,N$:FOR I=6144 TO Y:PRINT#
1,CHR$(PEEK(I));:NEXT I:CLOSE 1:RETURN
490 PRINT"[CD]DESTINATION TAPE Y":GOSUB
630:IF Y$<>"Y" GOTO 490
500 RETURN
510 GOSUB 750:0PEN 2,8,2,N$+",S,R":Y=614
4
520 GET#2,Y$:IF ST GOTO 570
530 IF Y$="" THEN Y$=CHR$(0)
540 POKE Y,ASC(Y$)
550 IF Y<40959 THEN Y=Y+1:GOTO 520
560 PRINT"[CD]FILE TOO BIG ONLY 34K COPI
ED"
570 CLOSE2
580 GOSUB 610
```

.

```
590 OPEN 3,8,3,"20:"+N$+",S,W":FOR I=614
4 TO Y:PRINT#3,CHR$(PEEK(I))::NEXT I
600 CLOSE 3:RETURN
610 PRINT"[CD]DESTINATION DISK Y":GOSUB
630:IF Y$<>"Y" GOTO 610
620 RETURN
630 GET Y$:IF Y$="" GOTO 630
640 RETURN
650 GOSUB 630:IF Y$="T" THEN PRINT" TAPE
":GOTO 680
660 IF Y$="D" THEN PRINT" DISK":GOTO 680
670 GOTO 650
680 RETURN
690 PRINT"[CD]":OPEN 1:CLOSE 1:GOSUB 700
:PRINT"[CD][REV]REWIND TAPE[OFF]":RETURN
700 PRINT TYPE FILENAME ": SPC(10); "BUFFER
 START : I=PEEK(828) : E=0
710 Y$=" PRG ": IF Y=4 THEN Y$=" SEQ "
720 PRINTY$;"[REV]";:FOR I=833 TO 1019:X
=PEEK(I):IF X<32 OR X>95 THEN X=63:E=1
730 PRINT CHR$(X);:IF I=849 THEN PRINT"[
OFF1<":
740 NEXT I:PRINT">[REV]END[OFF]":RETURN
750 INPUT "[CD]FILENAME";N$:N$=LEFT$(N$,
16):RETURN
760 N$="":FOR I=848 TO 833 STEP -1:X=PEE
K(I)
770 IF X=32 AND N$="" GOTO 790
780 N$=CHR$(X)+N$
790 NEXT I:RETURN
800 PRINT"[CLS]";:OPEN 1,8,0,"$0":GET#1,
Y$,Y$
810 I=0:GET#1,Y$,Y$,Y$,X$:IF Y$<>"" THEN
 I=ASC(Y$)
820 IF X$<>"" THEN I=I+ASC(X$)*256
830 PRINTRIGHT$("[2SPC]"+STR$(I),3);" ";
:I=0
840 GET#1,Y$:IF ST GOTO 930
850 IF Y$=CHR$(34) THEN I=I+1:PRINT CHR$
(34);:GOTO 840
860 IF I=0 GOTO 840
870 IF I=1 THEN PRINT Y$;:GOTO 840
880 IF I=2 THEN PRINT TAB(22);:I=I+1
890 IF I=3 AND Y$=" " GOTO 840
```

900 IF Y\$<>"" THEN PRINT Y\$;:GOTO 840

```
910 PRINT:GET Y$:IF Y$<>"" THEN GOSUB 63
0
920 IF ST=0 GOTO 810
930 PRINT "BLOCKS FREE":CLOSE 1:GOSUB 63
0:RETURN
940 GOSUB 750:FOR I=1 TO LEN(N$):POKE 83
2+I,ASC(MID$(N$,I,1)):NEXT I
950 FOR I=833+LEN(N$) TO 1019:POKE I,32:
NEXT I
960 POKE 183, LEN(N$): SYS 49244
970 PRINT"[CD][REV]FILE WILL BE DELETED[
OFF]":GOSUB 610
980 OPEN 15,8,15,"S0:"+N$:CLOSE 15
990 POKE 183,LEN(N$):SYS 49343
1000 RETURN
1010 GOSUB 750:FOR I=1 TO LEN(N$):POKE 8
32+I,ASC(MID$(N$,I,1)):NEXT I
1020 FOR I=833+LEN(N$) TO 1019:POKE I,32
:NEXT I
1030 POKE 183, LEN(N$): SYS 49244
1040 SYS 49203
1045 SYS 49206
1050 RETURN
1060 SYS 49152
1070 GOSUB 490
1080 SYS 49203
1085 SYS 49206
1090 RETURN
1100 SYS 49152
1110 GOSUB 700:IF E=1 THEN GOSUB 750:GOT
0 1130
1120 GOSUB 760
1130 PRINT"[CD]FILE NAME ";CHR$(34);N$;C
HR$(34);" OK Y/N":GOSUB 630
1140 IF Y$="N" THEN GOSUB 750
1150 PRINT"[CD][REV]FILE WILL BE DELETED
[OFF]":GOSUB 610
1160 OPEN 15,8,15,"S0:"+N$
1170 POKE 183, LEN(N$): SYS 49343
1180 RETURN
```

The following is the BASIC loader for the machine code and must be loaded and run before using the above program.

1 DATA 32, 44, 247, 173, 60, 3, 133, 255 , 169, 0, 133 2 DATA 193, 169, 24, 133, 194, 56, 173, 63, 3, 237 3 DATA 61, 3, 170, 173, 64, 3, 237, 62. 3, 168 4 DATA 24, 138, 101, 193, 133, 174, 152. 101, 194, 133 5 DATA 175, 32, 162, 245, 165, 255, 141, 60, 3, 96 6 DATA 32, 183, 247, 169, 0, 133, 193, 1 69, 24, 133 7 DATA 194, 56, 173, 63, 3, 237, 61, 3, 170, 173 8 DATA 64, 3, 237, 62, 3, 168, 24, 138, 101, 193 9 DATA 133, 174, 152, 101, 194, 133, 175 , 32, 124, 246 10 DATA 96, 169, 96, 133, 185, 169, 1, 1 41, 60, 3 11 DATA 133, 184, 169, 8, 133, 186, 169, 0, 133, 195 12 DATA 133, 147, 169, 65, 133, 187, 169 , 3, 133, 188 13 DATA 169, 24, 133, 196, 164, 183, 32, 175, 245, 32 14 DATA 213, 243, 165, 186, 32, 9, 237, 165, 185, 32 15 DATA 199, 237, 32, 19, 238, 141, 61, 3, 32, 19 16 DATA 238, 141, 62, 3, 32, 232, 244, 1 65, 174, 141 17 DATA 63, 3, 56, 165, 175, 233, 24, 14 1, 64, 3 18 DATA 24, 173, 61, 3, 109, 63, 3, 141, 63, 3 19 DATA 173, 62, 3, 109, 64, 3, 141, 64, 3, 96 20 DATA 169, 97, 133, 185, 169, 1, 133, 184, 169, 8 21 DATA 133, 186, 169, 65, 133, 187, 169 . 3. 133. 188 22 DATA 165, 185, 164, 183, 32, 213, 243 , 32, 143, 246 23 DATA 165, 186, 32, 12, 237, 165, 185, 32, 185, 237

```
24 DATA 169, 0, 133, 172, 169, 24, 133,
173, 56, 173
25 DATA 63, 3, 237, 61, 3, 133, 174, 173
, 64, 3
26 DATA 237, 62, 3, 133, 175, 24, 169, 2
4, 101, 175
27 DATA 133, 175, 173, 61, 3, 32, 221, 2
37, 173, 62
28 DATA 3, 160, 0, 32, 33, 246, 96, 255,
255, 0
29 FOR I=49152 TO 49432:READ A:POKE I,A:
NEXT I
```

Once this has been run it could be saved as its machine code for later ease of loading. A detailed description of the machine code follows.

MACHINE CODE

The machine code is called by the driver as required. It consists of four parts:

- i) Read any header and do relocated load
- ii) Write to tape current header and write relocated code
- iii) Load from disk retaining original details but relocate
- iv) Save to disk relocated code with original details

Chapter 5 of the *Programmer's Reference Guide*, 'The KERNAL', discusses the use of LOAD and SAVE in detail. The entry points given are for complete loads and saves (it is possible to do a relocated load, but not a relocated save using these). Unfortunately, as we are using an all-purpose BASIC driver, these entry points may overwrite it. To overcome this problem, every operation is carried out in two stages. The first is to read or write the file's details which are always stored in or taken from the cassette buffer. This avoids having to do too much moving of information. The second is to perform a relocated load or save with the correct amount of code going to or being taken from \$1800 on.

To do this we must enter the load and save routines at much later points with the parameters already set. It would consume too much space to describe these in detail, so we leave it up to you to follow them through. The only tricks are to prevent a forced load to its original address with a tape marker of 3 and to prevent a header being written with a marker of 5 (when an end of tape has been written – see Tape Headers).

TAPE				
Read	any tape	header	withou	it loading
C000	202CF7	JSR	\$F72C	read any header
C003	AD3C03	LDA	\$033C	get sec add
C003	85FF	STA	\$FF	and store for later
C008	A900	LDA	#\$00	set start of load
C00A	8501	STA	\$C1	to \$1800 by setting
C00C	A918	LDA	#\$18	STAL
C00E	85C2	STA	\$C2	
C010	38	SEC		subtract MSBs
C011	AD3F03	LDA	\$033F	of start and end
C014	ED3D03	SBC	\$033D	of original load
C017	AA	TAX		put result in X
C018	AD4003	LDA	\$0340	do same for LSBs
C01B	ED3E03	SBC	\$033E	putting answer in Y
C01E	A8	TAY		
C01F	18	CLC		find overall length
C020	8A	TXA		_
C021	65C1	ADC	\$C1	and add to STAL
C023	85AE	STA	\$AE	to give the new
C025	98	TYA		end i.e. \$1800
C026	65C2	ADC	\$C2	plus result
C028	85AF	STA	\$AF	
Load	from \$180	0 on		
C02A	20A2F5	JSR	\$F5A2	do the relocated load
C02D	A5FF	LDA	\$FF	restore the sec add
C02F	8D3C03	STA	\$0330	in case end of tape 5
C032	60	RTS		load complete
Write	to tape	in two	parts	the correct header and then
the c	ode from	\$1800 d	n	
C033	20B7F7	JSR	\$F7B7	write header in orig form
C036	A900	LDA	#\$00	reset STAL as it
C038	85C1	STA	\$C1	has been changed by
C03A	A918	LDA	#\$18	writing the header
C03C	85C2	STA	\$C2	
C03E	38	SEC		recalc the relocated end
C03F	AD3F03	LDA	\$033F	MSB
C042	ED3D03	SBC	\$033D	
C045	AA	TAX		
C046	AD4003	LDA	\$0340	LSB
C049	ED3E03	SBC	\$033E	
C04C	A8	TAY		
C04D	18	CLC		
C04E	8A	TXA		
C04F	65C1	ADC	\$C1	set up EAL

CØ51 85AE STA \$AE C053 98 TYA C054 65C2 ADC \$C2 C056 85AF STA \$AF Save RAM for reloading CØ58 207CF6 JSR \$F67C save RAM C05B 60 RTS complete DISK Load from disk relocated to \$1800 C05C A960 LDA #\$60 C05E 85B9 STA \$B9 set sec add 0603 A901 LDA #\$01 put type 1 in tape buffer C062 8D3C03 STA \$0330 C065 85B8 STA **\$**B8 make log file 1 C067 A908 LDA #\$08 make device 8 C069 \$BA and put in FA 85BA STA C04B A900 LDA #\$80 C03D 85C3 STA \$C3 A=0 for load C06F 8593 STA \$93 C071 A941 LDA #\$41 set pointer to file name to file name C073 85BB STA \$BB in FNADR to TBUFFR + 5 C075 A903 LDA #\$03 C077 \$BC 85BC STA C079 MEMUSS set to \$1800 A918 #\$18 LDA C07B 85C4 STA \$C4 for relocated load C07D A4B7 LDY \$B7 read len name set in BASIC C07F 20AFF5 JSR \$F5AF print SEARCHING C082 20D5F3 JSR \$F3D5 print LOADING C085 A5BA \$BA LDA get current device C087 2009ED JSR \$ED09 send talk C08A A5B9 LDA \$B9 get sec add C08C 20C7ED **JSR** \$EDC7 send talk sec add C08F 2013EE JSR \$EE13 receive from serial C092 8D3D03 STA \$033D and store LSB in TBUFFR C095 2013EE JSR \$EE13 do same for MSB of start C098 8D3E03 STA \$033E C09B \$F4E8 do relocated load 20E8F4 JSR | C09E A5AE LDA \$AE get end LSB COAO 8D3F03 \$033F STA put in TBUFFR C0A3 38 SEC subtract relocated start C0A4 A5AF LDA \$AF and end C0A6 E918 SBC #\$18 and put in appropriate CØA8 8D4003 \$0340 locations of TBUFFR STA

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Peripherals

COAB	18	CLC		
CØAC	AD3D03	LDA	\$033D	leaving a header
COAF	6D3F03	ADC	\$033F	suitable for
C0B2	8D3F03	STA	\$033F	a tape write
C085	AD3E03	LDA	\$033E	•
C088	6D4003	ADC	\$0340	
C0BB	8D4003	STA	\$0340	
CØBE	60	RTS		back to BASIC
Save	relocated	code	to relo	ad at correct address
CØBF	A961	LDA	#\$61	set parameters
C0C1	85B9	STA	\$B9	
C0C3	A901	LDA	#\$01	
C0C5	85B8	STA	\$B8	
C0C7	A908	LDA	#\$88	
C0C9	85BA	STA	\$BA	
COCB	A941	LDA	#\$41	
COCD	85BB	STA	\$BB	
CØCF	A903	LDA	#\$03	
CØD1	85BC	STA	\$BC	
C0D3	A5B9	LDA	\$B9	
C0D5	A4B7	LDY	\$B7	
C0D7	20D5F3	JSR	\$F3D5	send sec add
CODA	208FF6	JSR	\$F68F	print SAVING
CØDD	A5BA	LDA	\$BA	send listen
CØDF	200CED	JSR	\$ED0C	device 8
C0E2	A5B9	LDA	\$B9	send listen
C0E4	20B9ED	JSR	\$EDB9	sec add
C0E7	A900	LDA	#\$00	set up SAL
C0E9	85AC	STA	\$AC	with \$1800
CØEB	A918	LDA	#\$18	
CØED	85AD	STA	\$AD	
CØEF	38	SEC		
CØFØ	AD3F03	LDA	\$033F	calculate prog length
CØF3	ED3D03	SBC	\$033D	and put
CUFG	85AE	STA	\$AE	in EAL
C0F8	AD4003	LDA	\$0340	
CNER	ED3E03	SBC	\$033E	
LUFE	8585	SIA	\$AF	
0100	18	ULU		
0101	A918	LDA	#\$18	add \$1800 to
0103	60AF	ADC	\$AF	give relocated
0105	80AF	514	\$AF	ena
C107	AD3D03	LDA	\$033D	
C10A	20DDED	JSR	\$EDDD	send serial deferred
C10D	AD3E03	LDA	\$033E	send actual start in A and Y

C110	A000	LDY	#\$00				
C112	2021F6	JSR	\$F621	save	RAM	to	reload
C115	60	RTS		back	to I	BAS	IC

The utility is only intended for your own files. It will not as it stands backup relative files.

As a point of interest, Supersoft's ZOOM monitor offers not only the option to perform relocated loads and saves, but to save in a form suitable for reloading on a PET, which eliminates an ID of 3 not used on the PET.

3 A token approach to BASIC

Introduction

In this chapter we deal with the five main routines BASIC uses in interpreting your programs or commands. One of these, CHRGET, picks up single bytes from the program and is called by the majority of routines in your 64. The other four routines covered are concerned with keywords – converting from ASCII to tokens, the reverse process (LISTING), and directing them to their respective routines.

Other than using sys commands a knowledge of these routines is essential if you wish to extend existing commands or add further ones. Those of you owning a disk drive will be familiar with the program DOS 5.1 and may know that this modifies CHRGET to trigger its commands.

CHARGET

BASIC gets its information from the input or program lines through a routine called CHRGET (CHaracterR GET). A copy of this routine is held in the KERNAL operating system and is copied into zero page on power-up. Each time BASIC wants a character it calls this routine.

The routine is held at locations \$0073-\$008A (KERNAL is \$E3A2-\$E3B9) and is as follows:

\$0073	Еó	7A		INC	\$7A
\$0075	D0	02		BNE	\$0079
\$0077	Еó	7B		INC	\$7B
\$0079	AD	00	02	LDA	\$2000
\$007C	C9	ЗA		CMP	#\$3A
\$007E	B0	0 A		BCS	\$008A
\$0080	69	20		CMP	#\$20
\$0082	F0	EF		BEQ	\$0073
\$8884	38			SEC	
\$0085	E9	30		SBC	#\$30
\$0087	38			SEC	
\$0088	E9	D0		SBC	#\$D0
\$008A	60			RTS	

Bytes 0073-0077

Every time the CHRGET routine is called it increases, by one, the location

68 A token approach to BASIC

from where it gets its information. After it increments the LSB, in location 007A, it checks whether the page has been crossed, that is, from \$FF to \$00. Only if it has will the MSB be increased.

Bytes 0079-007B

Here it takes the information from store and puts it into the accumulator. The store location is present before the initial entry to the routine. It is always set one byte less due to initial increment. If you were going to use the routine yourself, it would be these bytes you would change, as we shall see later.

Bytes 007C-007F

Here it checks to see if the character is a numeral. It is testing to see if it is greater than ASCII numeral 9 (\$39). If so then the routine is left via \$008A with the carry set.

Bytes 0080-0083

Here is a straightforward test to see if a space was picked up; if so the routine is carried out again. CHRGET cannot be left on encountering a space.

Bytes 0084-0089

These successively subtract two numbers from the original byte and end up with the same number. You may say that is senseless but it will set two flags in the status register that help us later. These are the carry flag and the zero flag.

The carry flag. If this is clear on exit the byte will be a numeral in ASCII form. If it is set we have something else.

When subtracting two numbers in machine code the carry flag must always be set first. If the number we are subtracting from is the larger then the carry will remain set. On the other hand, if the number we are subtracting is the larger the carry will clear.

In this case we have already eliminated any byte that has a higher ASCII value than the numeral 9, in bytes 007E and 007F. It now subtracts \$30 (ASCII for digit zero) from the accumulator in preparation for setting the final flags which can be used for testing for numbers. The carry flag at this point does not matter as it is set again anyway. Bytes with ASCII values lower than numerals now range between \$D0 and \$FF.

With the next subtraction the original value is restored. The carry flag will now be set or unset as BASIC requires. As numerals are the only ones less than \$D0 (the last figure subtracted) they will be the only ones to clear the carry flag.

The zero flag. This could be set in two instances. First, in bytes 007E and 007F where we tested our byte against \$3A, the ASCII value for a colon. If it was a colon then the zero flag would have been set as it was equal (the carry would also have been set). Secondly, the flag would be set after the second subtraction if, and only if, the original byte was

zero (not ASCII digit zero). In that instance after the first subtraction the accumulator would hold \$D0 and subtracting the same value would set our zero flag.

A colon in a BASIC line signifies an end of statement and a zero, the end of a line, and hence an end of statement. Therefore, by testing the zero flag, we can quickly tell if we have reached the end of that particular instruction.

CHRGOT

Keyword routines are entered immediately after a call to CHRGET. Before using that byte it may require the accumulator for something else. To recover the CHRGET byte we can use the CHRGOT routine. This is a shortened version of CHRGET. If, instead of entering the routine at \$0073, we enter at \$0079, we miss out the instructions which update the pointer and get the original byte again.

Wedges

If we want to patch in our own machine code routines to work alongside BASIC, one way to do this is to insert a wedge. Simply, this is a routine which diverts CHRGET to check whether it is one of our additions. From this a decision can be made whether to revert to the normal CHRGET flow or to a routine of our own.

Let us say that we put in some routines all to be actioned on the character '@'. For instance, we could have a renumber routine and a delete routine. The command for renumber might be '@R' and delete, '@D' and could place a wedge at \$C000.

The first thing we have to do is alter the CHRGET routine. We want to change it after it has collected a byte but before it starts checking and manipulating it. The alteration would therefore be at S007C with a JMP to our coding. The routine starts with six bytes of data with our changes and the original bytes. (The latter is for restoring CHRGET if you require so to do later). We can load them by using the load with the x register. Our first instructions will look like this:

C000	4C	11 C0	JMP	\$C011
C003	C9	3A	CMP	#\$3A
C005	BØ		BCS	
C006	A2	02	LDX	#\$02
C008	BD	00 C0	LDA	\$C000,X
C00B	95	7C	STA	\$7C,X
C00D	CA		DEX	
C00E	10	F8	BPL	\$C008
C010	60		RTS	
C011	Our	coding	will sta	art here

This is the routine to initialize our wedge and is called immediately after loading the program by using SYS 49158 ($combol{s}$). This loads a byte x

places from \$C000 and stores it x places from \$7C. We decrease the counter x and the branch to collect the next byte will work until we decrement it below zero, that is, no longer a positive number. If the branch fails, we go back to BASIC through the RTS.

The beginning of the CHRGET now looks like this:

0073	INC	\$7A			
0075	BNE	\$0079			
0077	INC	\$7B			
0079	LDA	\$0200	(this	number	varies)
007C	JMP	\$C011			

Now each time CHRGET is used it will go to our routine. As all our commands would be triggered with '@', the first thing we would do is to check to see if it is present:

C011	C9 40	CMP	#\$40
C013	FØ Ø3	BEQ	\$C018

If it is, we can branch to do further checks, and if not, we continue with the next code. Here we will have to revert to the normal course of events. We have two options which we can take. First, we include the bytes of CHRGET we changed into our program – we do not want to change CHRGET itself as we want it to use again – and jump back to CHRGET at \$0082. Secondly, we can use the CHRGET routine in the KERNAL ROM, jumping in at \$E3AB, and BASIC will continue as if nothing has happened.

The first method would be like this:

CØ15 4C AB E3 JMP \$E3AB

And the second like this (of course, the routine address will change from \$C018 to \$C01D):

C013	F0 07	BEQ	\$C01D
C015	C9 3A	CMP	#\$3A
C017	B0 03	BCS	\$C01C
C019	4C 80 0	a JMP	\$0080
C01C	60	RTS	

We now want to find out if it is one of our routines. This we do by checking the next character without updating the CHRGET pointer (in case it isn't). The code would look like this:

C018	08	PHP
C019	48	PHA
C01A	98	TYA
C01B	48	PHA
C01C	8A	TXA
------	----------	---------------------------
C01D	48	PHA
C01E	A6 7A	LDX \$7A
C020	E8	INX
C021	BD 00 02	LDA \$0200,X
C024	C9 52	CMP #\$52
C026	FØ ??	BEQ - TO RENUMBER ROUTINE
C028	C9 44	CMP #\$44
C02A	F0 ??	BEQ - TO DELETE ROUTINE
C02C	68	PLA
C02D	AA	TAX
C02E	68	PLA
C02F	A8	TAY
C030	68	PLA
CØ31	28	PLP
CØ32	4C AB E3	JMP \$E3AB

Bytes C018–C01D

Here we are preserving our registers, including the status register, on the stack in case it is not destined for our own routines.

Bytes C01E-C023

Location \$7A has the LSB of the pointer used by CHRGET and this is one less than the next character we want. So if we load it into the x register and increase it by one, we will have the position of the next byte. We can now load the byte using x as a pointer.

Bytes \$C024-\$C2B

A check is made to see if it is the letter R, signifying the renumber routine. If not we then check for the letter D and if so go to delete.

Bytes \$C2C-\$C034

It not, we restore our registers from the stack (in the reverse order we put them on). Then we continue the KERNAL routine as before.

One last point is that in the routines, such as renumber or whatever, it would be advisable to call a subroutine to remove the bytes from the stack (placed there in \$C018 to \$C01D). We do not require them, but as your routines are called the stack will become fuller and fuller, resulting in an 'OUT OF MEMORY' error.

Keywords

A more professional approach to adding routines than altering CHRGET is the use of *Keywords*. This approach holds with the idea behind the BASIC language that actions can be performed by using words which are indicative of the desired action.

Keywords can be divided into two types: commands and functions.

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Functions get information, for example, PEEK returns the contents of a location, and will always supplement a command keyword. For example, we use PRINT PEEK(xx) but never PEEK(xx)PRINT.

On the 64 it is possible to incorporate new routines actioned by keywords as they go to the relative routines through vectors held in RAM. As the vectors are in RAM we can change them to go to routines of our own.

The other items we will have to add are three tables of data. One will have the new keywords in ASCII code. This will be used when LISTING and tokenizing. The end of a word is indicated by adding \$80 (128) to its last letter. To signify the end of the table, a zero is used. If we had a table of two keywords, say END and NOT, it would look like this:

45 4E C4 4E 4F D4 00

The other two tables will have the addresses of our routines, one for command keywords and one for functions. These will hold the address of each routine less 1. The reason for this is that we will put them on the stack for an RTS instruction. The program counter will add 1 when it takes them from the stack, thus getting the correct address for the routine. The table will have two bytes for each routine, the LSB and then the MSB. For example, if we had a routine at \$C4DF, on the table it will look like this: DE C4.

There are four vectors we will have to change (three if not adding functions).

ADD OF VECTOR ADD OF ROM DESCRIPTION

\$0304/5	\$A57C	Tokenize basic
\$0306/7	\$A71A	Print Tokens(LIST)
\$0308/9	\$A7E4	Token Dispatch – Command words
\$030A/B	\$AE86	Token Dispatch – Function words
	T (·

Tokenize BASIC Text The object of this subroutine is to take an input line, check it for keywords, tokenize them and condense the line. It does this by taking every byte from the input buffer, not using CHRGET, and then checking through the keyword table for a match. If the letters do not make a keyword, it stores them as variables, meaning that variables cannot have keywords in them.

There are two ways we could approach the problem of incorporating our own keywords and tokens. First, we could copy the BASIC from ROM into RAM, and alter the tokenize routine within BASIC so that if it cannot find a match it jumps to a routine to check through our table of keywords. This would mean we would not have to change the vector but would lose the RAM area under the BASIC ROM which is useful for storing hires screens, data tables, and so on. The second way would be to change the vector to a routine of our own. Here we have a copy of the ROM routine altered slightly to be able to search our table as well. We would only use the ROM routine when we finished tokenizing a whole line.

We shall describe the second method, which we think is the better in the long run. A description follows the code.

10		LDX	\$7A
20		LDY	#\$04
30		STY	\$0F
40	ANOTHER	LDA	\$0200,X
50		BPL	SPACE
60		CMP	#\$FF
70		BEQ	STORE
80		INX	
90		BNE	ANOTHER
100	SPACE	CMP	#\$20
110		BNE	STORE
120		STA	\$08
130		CMP	#\$22
140		BEQ	QUOTE
150		BIT	\$0F
160		BVS	STORE
170		CMP	#\$3F
180		BNE	NUMBER
190		LDA	#\$99
200		BNE	STORE
210	NUMBER	CMP	#\$30
220		BCC	CONT
230		CMP	#\$3C
240		BCC	STORE
250	CONT	STY	\$71
260		LDY	#\$00
270		STY	\$0B
280		DEY	
290		STX	\$7A
300		DEX	
310	NEXT LETTER	INY	
320		INX	
330	CONT 1	LDA	\$0200,X
340		SEC	
350		SBC	\$A09E,Y
360		BEQ	NEXT LETTER
370		CMP	#\$80
390		BNE	NEXT WORD
400	STORE A	ORA	\$0B
410	FOUND	LDY	\$71

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420	STORE	INX					
430		INY					
440		STA	\$01FB,Y				
450		LDA	\$01FB,Y				
460		BEQ	END				
470		SEC					
480		SBC	#\$3A				
490		BEQ	COLON				
500		CMP	#\$49				
510		BNE	DATA				
520	COLON	STA	\$AF				
530	DATA	SEC	+ • •				
540	01111	SBC	#455				
550		BNE					
540		CTA	400				
570			4000 40700 V				
590			STOPE				
500			400				
200			700 CTODE				
410	OUNTE		STURE				
420	QUUIE	CTA	401ED V				
420		TNIX	⊅01rD,1				
210			1 111				
250	NEVE LIODS	DINE					
440	NEXT WURD		⊅/H #በቦ				
470			₽0D				
490			theon v				
200			PH070,1				
700							
710		DNE	PONT 1				
720			HALE				
720			₩₽ГГ				
740	NEYT						
750							
740	NEVT D		40000 V				
700	NEXT D	CDH CEC	¥0200,×				
700		320		000			
700		350	≫SIAKI UF	UUK	WURD	IABLE	,X
000		BEU	NEXI				
000			#>80				
910		BNE	NEXI NEW				
020		BEU	STURE A				
030	NEXT NEW	LDX	¥/A ≠2₽				
040		INU	⊅ 0R				
820	NEXIA	INY		-			
866		LDA	START OF	OUR	WORD	TABLE	-1,X
870		BPL	NEXT A				

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880		LDA	\$START	0F	OUR	WORD	TABLE,X
890		BNE	NEXT B				
900		LDA	\$0200,>	<			
910		BPL	FOUND				
920	END	JMP	\$A609				

LINES 10–30: Initialization. Location $$^{7}A$ will have a value the same as the position within the input buffer. In immediate mode this will be 0, the start of the buffer. If inputting a program line, it would be after the line number, which has already been taken care of. Now x will be our pointer to the original contents of the buffer and y will be the pointer to our new buffer set up. Y is stored in \$^{9}F just to initialize that location. The value does not matter as long as it was not over \$^{3}F, as we shall see.

LINES 40–50: We load in a byte and check to see if it is under \$80 (128). If it is, we branch off to line 100.

LINES 60–90: Values of \$80 and over arrive here and we check to see if it is 'PI' (\$FF). If it is, we branch further into the routine to store it. If not, we branch back to get another byte. This means we cannot use the first character of a keyword as a shifted letter because it would be greater than \$80.

LINES 100–140: If a space is found we branch off to store it, otherwise we store the accumulator in a register, in case the following check succeeds, for comparison later. Now we check to see if the byte is a quote. Items between quotes do not require tokenizing and therefore we branch and continually store them until we come across another quote or the end of the line.

LINES 150–160: Here we are checking to see if location \$0F has bit 6 set or not. Amongst other things, the BIT instruction takes bit 6 of \$0F and places it in bit 6 of the status register, the overflow register. This bit in \$0F will only be set in this routine if we tokenize DATA later in the routine. It means that after DATA all characters will be stored and do not go through the keyword table. Bit 6 can be unset by a colon if outside quotes and for this reason colons have to be in quotes within DATA statements. A colon outside quotes will mean that information after is tokenized. BASIC instructions can be placed at the end of a data line and will be actioned as the line is encountered. BASIC differentiates between DATA and REM. On REM it will go to the next line whereas with data it will search through it for another BASIC instruction.

LINES 170–200: The *c* is the shortened version of the keyword PRINT. It is the only keyword which can be shortened to a single character. These lines check for the question mark and if found place the token for PRINT, \$99, in the accumulator and go to store it.

LINES 210-240: Here we find out if the byte is a numeral, colon or

semicolon. If it is off we go to store it, if not then we continue the routine.

LINES 250–300: We now set up for the search through the table of keywords. We store our Y register, which, if you remember, is the pointer to the 'new look' buffer. Location some will be our counter to the number of keywords we encounter; it will not hold the token value but helps determine it. We store x in s7A. This is part of CHRGET, but we are only using it as a store. We decrease both x and y as the first part of the next section will increment them.

LINES 310–390: This section of the routine explains why on the Commodore 64 we can use shortened keywords using shifted letters. There are two things to remember here. First, the last letter of the keyword in the table is the value of the letter plus s80 (128) which when loaded will set the negative flag. Secondly, the value of a shifted letter (not logo shift) is also the value of the letter plus s80.

Back in the routine we increase the registers and load in our byte again; it will later load the next byte. We set the carry for subtraction and subtract the value of a letter, from the keyword table, from our byte. If we are left with zero then we have a match and we go back to get the next letter from the buffer. If it fails, we check to see if we have the value ^{\$80} left. This would indicate we have a match by either the input letter being the shifted letter or we have reached the end of the keyword in the table and have also matched. Failing this second check, it branches off to find the end of that word so we can check the next for a match.

As the second letter can be shifted to give our match this explains not only shortened keywords, but also why some require at least three. As an example, take the keywords CLR and CLOSE. CLR comes before CLOSE in the table so that C and shift L will match with the former before it gets to CLOSE, which will therefore need two standard letters before a shifted letter. We can also explain why there is no shortened version for INPUT. INPUT# comes before INPUT and so any shortened version will always match with INPUT#.

LINES 400–460: Back to the routine now. We have in the accumulator the value \$80 and have found a match in the table. Here we perform the logical OR of the accumulator and location \$08. Later on we will find out that every time we pass through a keyword that does not match, we increase the value of \$08. As we started at 0, \$08 will have the number of the match word, the first word in the table being 0. The instruction OR forces bits into the accumulator if they are not set. In this case it has the same effect as adding but saves bytes doing it this way. The accumulator always has \$80 and if you OR it with a value of one you get \$81. This is how we arrive at the token value. The keyword GO is the last word in the table and that is the 76th, giving a value in \$08 of \$48 (75), and ORing it

with \$80 gives a value of \$CB which is the token value of GO.

Having got our token value, we load back into the Y register from \$71 – the pointer to the new buffer layout. We have now reached the point where we store a lot of the characters into the new buffer layout. This is the point where many of the earlier branches arrive. First, we increase both buffer pointers. The base location to store is \$01FB indexed with Y. As we started with Y equal to \$04 and increase it immediately, our first location will be \$01FB + 5, giving \$0200. We will not overwrite anything in the original line we have not checked for two reasons. First, if there was a line number at the beginning it has been dealt with and is no longer needed. Secondly, the routine only shortens and never lengthens.

Once the byte has been stored, it is loaded back in and checked for zero which signifies the end of the line.

LINES 470–550: This section is going to test if the byte we have stored is either a colon or one of the keywords DATA or REM. If it is a colon, it unsets bit 6 of location \$0F which we discussed earlier, then goes to get more bytes. DATA will set the 6th bit of that location before getting more bytes.

REM is slightly different in that nothing after it requires tokenizing. We set the location of \$0B to zero, as that is the result of the subtractions, which we will not actually need for checking but which stops us from branching out of the next section for any reason other than the end of line.

LINES 560–640: These lines are used only in two instances: on encountering a quote or encountering REM. All this does is to move bytes from their original to their new position in the input buffer, until we reach the end of the line or, in the case of a quote, we find a closing quote. Finding either of these, we branch back to the normal store lines of 420–460. The branch in line 640 is enforced as x can never be zero as before we get here we have looked at a minimum of two bytes so the least x can be is 1 and the maximum \$58 (the maximum input line length).

LINES 650–710: This is the section we come to if we did not find a keyword match. All this does is to search for the next character in the table that has a value greater than \$80, and return with its position. The pointer will be increased before we start another match. It also checks for a 0, signifying the end of the table.

So far the routine is the same as in ROM but now we change the course of events. In the ROM routine, when it finds the end of the table it assumes the characters are variables and stores them as such. We, on the other hand, want to see if it is one of our keywords, so on getting zero we have to search our table. The y register is loaded at the beginning of the section because either way we want to get back the

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first character of this particular check from the original input buffer line up.

LINES 720–910: This is a repeat of the checking of the standard keywords except it will have the address of our keyword table. We will now be checking for our keywords.

LINE 920: When we have found and stored the end of line zero, we get here. The routine now jumps off to end or to the original ROM routine. There it will reset the CHRGET pointers to their initial setting of \$01FF and continue the normal flow of BASIC to either store the line or carry out its instructions if in direct mode.

Print tokens

This routine is part of the LIST routine in BASIC. It takes the token value, finds the keyword and prints it to the screen, or other device. The ROM print token routine is not a subroutine by itself but an integral part of LIST, but thankfully it is vector-started. The vector points to the next instruction in the ROM routine. What we would have to do is to change the vector to a routine of our own, PRINT keywords of either the standard ones or our own, and then jump back to the LIST routine at an appropriate point. The coding for such a routine is as follows:

10		BPL	ROM 1					
20		CMP	#\$FF					
30		BEQ	ROM 1					
40		BIT	\$0F					
50		BMI	ROM 1					
60		CMP	#\$CC					
70		BCC	CBM T	DKEN				
80		SEC						
90		SBC	\$CB					
100		TAX						
110		LDA	# LSB	START	0F	OUR	KEYWORD	TABLE
120		STA	\$22					
130		LDA	#\$MSB	START	0F	OUR	KEYWORD	TABLE
140		STA	\$23					
150		BNE	START					
160	CBM TOKENS	SEC						
170		SBC	#\$7F					
180		TAX						
190		LDA	\$9E					
200		STA	\$22					
210		LDA	#\$A0					
220		STA	\$23					
230	START	STY	\$49					
240		LDY	#\$FF					
250	NEXT WORD	DEX						

260			BEQ	WORD	FOUND
270	NEXT	Char	INY		
280			LDA	(\$22)),Y
290			BPL	NEXT	CHAR
300			BMI	NEXT	WORD
310	WORD	FOUND	INY		
320			LDA	(\$22)	Y,Y
330			BM1	ROM 2	2
340			JSR	\$AB47	,
350			BNE	WORD	FOUND
360	Rom	1	JSR	\$A6F3	}
370	Rom	2	JSR	\$A6EF	

LINE 10: This tests the negative flag. A value of \$80 (128) or over is signalled as negative. As all tokens are \$80 or over, this branch will succeed; values under \$80 go back to LIST unchanged.

LINES 20: \$FF is the value of 'PI'. If it is that value we again return to LIST.

LINES 40–50: What we are doing here is putting bit 7 of location \$0F into the negative flag, although it does other things which are of no concern to us. Location \$0F is the flag used by the LIST routine to signal if it is listing in quotes or not. If bit 7 of \$0F is 1, then it is listing in quotes and we do not want to print tokens but the ASCII of the bytes. Therefore, if the negative flag is set, we branch to go back to ROM.

LINES 60–70: Here we find out whether it is one of the standard tokens or one of ours. It will branch off if it is standard.

LINES 80–150 OUR TOKENS

80–100: Here we subtract a number that is one less than our first token value. The result is then transferred to the x register to act as a counter. The value of x is one greater than the position in the table (starting at 0) but x will be decreased before we start the search.

110–140: We store the start address of our table in what will be our search registers.

150: This is enforced as the last figure in the accumulator, the MSB of our table, will not be zero. We are hardly likely to have a keyword table in the zero page which has many important BASIC locations.

LINES 160–220 CBM TOKENS: This is a duplicate of lines 80–140 except it is for the standard tokens and keyword table.

LINES 230–240: So far we have not used or altered the y register but we store it here in location \$49 for the LIST routine as that is where it will expect to find it later. We initialise y with \$FF but will increase it before our search so it will start a zero.

It may be worth a note here that we will not alter the values of the search registers, \$22 and \$23, as the 64's keyword table is not longer than

256 bytes and it is unlikely that ours would be. Therefore, incrementing γ through its 256 range (\emptyset -FF) will serve our purpose. It also saves bytes and time.

LINES 250–260: Every time we read a word from the table we will come here and decrease the x register. If x is zero, then we have found the position one byte before the keyword we want. In that case we branch off to PRINT the keyword.

LINES 270–300: Here we increase our table pointer, the Y register, and then load in the next character from the table. Remember that the last character of a keyword is its ASCII form plus \$80 (128) and this is what we look for. This will set the negative flag in the status register.

The first check is to see if the negative flag is unset signifying a branch back to get the next character. If the negative is set, then the end of the word is found and we branch back to test x to see if we have come far enough. One of these two branches must work as a byte is determined as either negative or positive.

LINES 310–350: We have found our word and now have to print it out. First we increase our pointer to pick up the first character. We load it and test to see if it is the last character. If it is we go to the BASIC ROM to have it printed through the LIST routine. Failing this, we go to another ROM routine to have the character printed. We will return with the same character in the accumulator. As a keyword will not have a byte of zero value, the branch in line 350 is enforced, to get another character to print.

LINE 360—ROM 1: The character was not a token at the beginning so we go back to the LIST routine to have it printed and continue with the listing.

LINE 370—**ROM 2:** We have here the last character of the keyword in the accumulator. Now we go back to the LIST routine where it will be turned into the proper ASCII value, printed and the listing continued.

BASIC token dispatch

This is the routine that BASIC uses on finding a token to get the address for the routine. It deals only with command keywords, such as PRINT. It is a subroutine in itself. What we need to do is put in a routine that it goes to first, through the vector.

SR ≇i	0073			
MP #	\$CC			
CC R	OM			
MP #	HIGHEST	COMMAND	TOKEN	VALUE
CS R	OM			
SR D	ISPATCH			
MP \$4	A7EA			
	SR ≢ MP # CC R MP # CS R SR D MP \$	SR \$0073 MP #\$CC CC ROM MP #HIGHEST CS ROM SR DISPATCH MP \$A7EA	SR \$0073 MP #\$CC CC ROM MP #HIGHEST COMMAND CS ROM SR DISPATCH MP \$A7EA	SR \$0073 MP #\$CC CC ROM MP #HIGHEST COMMAND TOKEN CS ROM SR DISPATCH MP \$A7EA

80	DISPATCH	SEC					
90		SBC	#\$CC				
100		ASL					
110		TAY					
120		LDA	START	OF	OUR	VECTOR	TABLE+1,Y
130		PHA					
140		LDA	START	OF	OUR	VECTOR	TABLE,Y
150		Pha					
160		JMP	\$0073				
170	ROM	JSR	\$0079				
180		JMP	\$A7E7				

LINE 10: Get the token from the input buffer or program line through the CHRGET routine.

LINES 20–30: We check to see if it is one of our tokens. If it is not, we branch off to the normal routine in ROM.

LINES 40–50: Now we find out if it is a command or a function token of ours. If it is a function vector, then it is a 'SYNTAX ERROR' so we branch to ROM to print it.

LINES 60–70: Here we go to our subroutine for dispatch. When the keyword routine has been completed, the program flow will come back here where we shall jump back to BASIC for continuation.

LINES 80–150: We subtract our lowest token value from the value we have. This will give us values of \emptyset upwards. Now as each routine has a two byte address, we must double our 'new' token value to get its proper place in the vector table. The instruction ASL does just this by shifting all bits one place left and putting a zero in bit \emptyset . This new value is transferred to the γ register as a pointer in the table. What we are going to do is to put a new return address on top of the stack (the program counter expects the LSB on top with the MSB underneath). Therefore, we take the second byte of the table first, put it on the stack, and then the first. Remember our vector table is made up of LSB then the MSB.

We now JMP to the CHRGET routine to pick up the next byte. The RTS at the end of CHRGET will now be to our keyword routine as we have just put its address on the stack. We came to these lines (80–150) by a JSR command so its return address was originally on the top of the stack. We then put another address on top of that which was pulled off in CHRGET leaving our original return address once more at the top. At the end of the keyword routine this address will be pulled off and we will return to line 70 of this routine.

LINES 170–180: Here we go to the normal dispatch routine. This is not the address normally found in the Token Dispatch Vector because we

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will miss out the first instruction which is to get the next byte. We go to CHRGOT first not to get the byte we have already got but to set the flags that the ROM routine wants to test.

BASIC function dispatch

This is the routine that will find the routine addresses of function keywords.

10		LDA	#00
20		STA	\$0D
30		JSR	\$0073
40		CMP	#\$LOWEST FUNCTION TOKEN VALUE
50		BCC	ROM
60		CMP	#\$HIGHEST TOKEN VALUE
70		BCS	ROM
80		JSR	DISPATCH
90		RTS	
100	DISPATCH	SEC	
110		SBC	#\$LOWEST FUNCTION TOKEN VALUE
120		ASL	
130		TAY	
140		LDA	#\$START OUR FUNCT VECT TABLE+1,Y
150		PHA	
160		LDA	#\$START OUR FUNCT VECT TABLE,Y
170		PHA	
180		JMP	\$0073
190	ROM	JSR	\$0079
200		JMP	\$AE8D

This is basically the same as the previous routine. The return addresses to ROM are different, as will be the table address. The first two lines load a location which BASIC uses to decide whether to accept numeric or string data, the latter value would be \$80.

The other difference is that on return from the function routine we will arrive back at line 90. The previous routine went back to BASIC for another command, but here we RTS as functions will be performed as part of a command routine and therefore we go back to it.

4 Keyboard revisited – making use of the wasted keys

On the far right of your keyboard there are four keys that do not really do much, at least at the moment. They are, of course, the function keys. In this chapter we are going to show you how to make use of them. First we thought it a good idea to describe the ROM routine in the 64 which services the keyboard. In doing so we will also come across the locations that appertain to the keys.

The hardware interrupt vector

Every 1/60th second the computer hands control to an interrupt system. When the microprocessor receives an interrupt signal it will not do anything until the present instruction has been completed. The processor will then save the program counter and the status register. The program counter is then loaded with the contents of locations \$FFFE and \$FFFF. This will start a routine at \$FF48 which saves the register contents on the stack before doing an indirect jump to the vector at \$0314 and \$0315.

The interrupt routine found at this vector points to address \$EA31. This KERNAL routine performs several housekeeping operations such as the update of the system clock, but it also scans the keyboard. The key that you press is picked up by the Complex Interface Adapter #1, and in particular the Data Port B within that chip. From this the value of the key pressed, and shift keys if used, is calculated and stored.

There seems to be some doubt from what we have read about which location the current key value is stored in. The current value is stored in \$CB (203) and the last in \$C5 (197). This to the BASIC programmer does not make a lot of difference unless the key buffer is full when the key value is not logged except in \$CB. The shift, logo and CTRL keys have the same system, with the current location being \$028D (653) and the last press in \$028E (654).

Having stored your current input it will check to see if it is the same as the last key press. Its next action will depend on whether it is the same, or not. If the same, it will see if it is a repeat function such as the cursor keys or if location \$02BA (650) has been set for all keys to repeat. Failing this, the value will not be placed into the keyboard buffer. Where the key values are processed it does so by looking up a table to obtain the ASCII code for your key press. This value is placed in the keyboard buffer and its counter updated. The keyboard buffer is situated at \$0277-\$0280, a size of ten characters. It operates on the system that the first character in will also be the first out. The pointer for the number of characters in the buffer at a particular time is \$C6 (198). The size of the buffer can be reduced from its initial value of ten by setting register \$0289 (649).

Earlier we said that every key has a value. These are from 0 to 64, the latter being no key press. A table of these, and the shift key values are given in the appendices.

Here is a summary of keyboard locations:

\$CB	203	Current key press.
\$C5	197	Last key pressed.
\$028D	653	Current shift etc.
\$028E	654	Last shift etc.
\$028A	650	Repeat flag: \$80 all, \$00 normal.
\$0277-\$0280	631–640	Keyboard buffer.
\$0289	649	Size of keyboard buffer.
\$C6	198	No of chars in buffer.

The Function Keys

These keys have values and ASCII codes like any other key. They are:

	Value	
Function key	(\$CB and \$C5)	ASCII
F1	4	133
F2	4	137
F3	5	134
F4	5	138
F5	6	135
F6	6	139
F7	3	136
F8	3	140

Knowing these values and the locations mentioned earlier, we can make use of the function keys.

Function keys within a BASIC program

One of the most used BASIC statements for evaluating a key press is the GET function. This function returns the ASCII code for the first key in the keyboard buffer, or the latest key if the buffer is empty. It will not wait for a key press. A BASIC routine could look like this:

100 GET A\$ 110 IF A\$=''[F1]'' THEN 1000:REM ACTION ON F1 PRESS 120 REM ACTION IF ANY OTHER KEY PRESSED This routine will not stop and wait for a key press. It will only branch off to line 1000 if key F1 is pressed at the same time as the GET statement is actioned or the next character in the keyboard buffer is the ASCII for F1.

We could adapt this so that it will wait until a key is pressed – any key.

100 GET A\$: IF A\$ = ""THEN 100

Here line 100 will be repeated until one key is pressed or there is a value in the key buffer that has not been read.

The next thing we could add is a line to clear the input buffer before we GET a character. The easiest way is to set the register for the number of buffer characters to zero.

90 POKE 198,0

At the moment the routine actions on any key. If we wanted it to action on only two keys, say F1 and F7, we would have to alter line 120 to:

120 IF A\$ <> ''[F7]'' THEN 90 130 REM ACTION ON F7 PRESSED

Now the routine will wait until a key is pressed. Once a key is pressed it goes to 110 to see if it was F1 and branches if so. Failing that it goes to line 120 where we look to see if it was not F7. On F7 the program will continue its flow. Now lines 90–120 will keep repeating until either F1 or F7 is pressed.

The only other alteration we could do is to rid ourselves of the graphic characters in the quotes that represent the function keys. This would make it easier for someone else to read and on a non-Commodore printer the graphic character would not print. This we can do by using the CHRS(function when checking AS. Line 110 would now look like:

110 IF A\$ = CHR\$(133) THEN 1000: REM ACTION ON F1 PRESSED

In the GET statement all eight function keys can be tested in the same way, either by changing the character in the quotes or changing the CHR\$ value.

Another way of testing for the keys is by examining the key press registers set in the interrupt routines. From a BASIC programmer's viewpoint it does not really matter whether you test the current or the last key register. The snag with this method is that without checking the shift register only four of the function keys can be detected. On the other hand, by checking the shift register with all its combinations you can have up to 32 function key combinations. Here is a routine that tests for function key F1: 90 POKE 198,0: REM CLEAR KEY BUFFER

100 IF PEEK(203)=4 THEN 1000 : REM F1 VALUE.

110 REM PROG CONTINUES IF NOT F1.

This will not wait for F1. To wait, line 110 will have to be changed to:

110 GOTO 90

We have now set up a loop and the only exit is F1 being pressed. Now if we wanted to test for F2, the shift flag would have to be introduced. Line 100 could look like this:

100 IF PEEK(203)=4 AND PEEK(648)=1 THEN 1000 : REM ACTION ON F2 PRESSED

If you wanted to go to line 1000 on any key, or no key, apart from F2, then the equals sign should be replaced by greater than and less than signs.

Programming the keys in immediate mode

Our interrupt routine

The routines that follow will allow you to program the function keys with commands or phrases to be actioned as if you typed them in full, but using only one keystroke.

Most of the routines we have seen to do this operation change the vector address of the Hardware Interrupt Routine in \$0314 and \$0315. They alter it to point to their routine, which when finished will return direct to the normal interrupt routine. This course of action has drawbacks. First, it adds to the length of the interrupt, especially if the user's routine has to be completely followed through. Secondly, it means that you have to set up your own registers for checking to see if it was the same action as the last time or not, to avoid auto-repeat. A further drawback is that if you want to use the data assigned to function keys within quotes, it is more difficult to suppress the graphic character that is generated in the quotes mode along with your phrase.

So how are we going to achieve this desirable routine of making the function keys really useful? Earlier we described the interrupt routine and how your key presses are interpreted. What we did not say was that there is a vectored jump within it. This occurs after the value from the Data Port is put into the current key registers but before it is actioned. The vector is held in addresses \$028F and \$0290 (655 and 656) and is known as the 'Keyboard Table Setup Vector'. If we change the address in this vector to point to a routine of our own we can process the data first. If the data concerns us we can process it jumping back to the normal interrupt routine at a point which misses out the normal key press routine. When the data does not concern us, control will be handed back to the normal flow of things.

The use of the vectors by Commodore has allowed us an easy way to program the keys. This cannot be said of the values that have been assigned to the function keys. It would have been easier if F1 had a value of 1 and F3, of 3, but this is not the case.

We are going to have 16 programmed function keys. To get this number, you have to use the keys in conjunction with the shift and logo keys as follows:

KEYS F1, F3, F5, F7	– THE KEY ONLY
KEYS F2, F4, F6, F8	– THE KEY + SHIFT
KEYS F9, F11, F13, F15	– THE KEY + LOGO
KEYS F10, F12, F14, F16	– THE KEY + SHIFT + LOGO

This gives us keys in the range of 1 to 16, but for the routine it is easier to use 0 to 15. We shall load the data into the keyboard buffer so we are limited to ten characters. We also require a marker for the end of data for each key, which will be a zero, meaning a maximum 11 bytes storage for each. It is easier, and quicker, to use 16 bytes per key. This wastes five bytes but as we are going to store the data in the RAM under the BASIC ROM this is unimportant. This will mean the value of the key needs to be multiplied by 16 to get the start of its data. Multiplying by 16 for the low numbers we are using, 0 to 15, simply involves moving the four lower bits to the four higher bits and filling the lower ones with zeros, four ASL instructions will achieve this. Sixteen bytes of data for the 16 keys will take one page, 256 bytes, exactly.

To summarise:

- i) Find out if the key is a function key, yes continue, no go to interrupt.
- ii) Calculate key number less 1.
- iii) Multiply key number by 16 for table position.
- iv) Get data off the data table and store in the key buffer.

ASSEMBLY LISTING

9	*=\$ 8722		
10		LDY \$CB	! CURRENT KEY PRESS
20		CPY #\$03	! IS IT A FUNCTION KEY
30		BCC NORMAL	! NO
40		CPY #\$97	! IS IT A FUNCTION KEY
50		BCC CONT	! YES
60	NORMAL	JMP \$EB48	! NORMAL INTERRUPT
			KEY ROUTINE
70	CONT	LDA \$028D	! CURRENT SHIFT PRESS
80		CPY \$C5	! IS CURRENT KEY=LAST
90		BNE CONT2	! NO
100		CMP \$028E	! IS CURRENT SHIFT=LAST
110		BEQ NORMAL	! KEY AND SHIFT AS LAST

120	CONT2	STY	\$05	ļ	STORE CURRENT KEY
					IN LAST REGISTER
130		STA	\$028E	!	STORE CURRENT
					SHIFT IN LAST REG
140		CPY	#\$04	!	IS IT F1
150		BEQ	F1+1	!	YES
160		CPY	#\$05	ļ	IS IT F3
170		BEQ	F3+1	!	YES
180		CPY	#\$93	I	IS IT F5
190		BED	E5+1		YES
200		INY	#\$97	i	IT IS F7
210	E1	RIT	\$01 <u>0</u> 0	1	
210	F1		#03A0		
220	F3	BII	\$03A0	;	VALUE FUR F3
230	F0	BII	90-004	!	VALUE FUR F5
240		UMP	#\$02	!	WHAT SHIFT
250		BCC	NOCHANGE		NONE OR SHIFT -
					VALUES CORRECT
260		BEQ	CBM+1	!	LOGO KEY
270		LDA	#\$09	ļ	VALUE FOR SHIFT+LOGO
280	CBM	BIT	\$08A9	ļ	VALUE FOR LOGO
290	NOCHANGE	STY	\$BB		
300		DEC	\$BB	I	ONLY WANT NO'S 0-15
310		CLC			
320		ADC	\$BB	ı	GET FINAL VALUES
338		ASI	A	1	MULTIPLY VALUE BY 16
340		ASI	A	•	
350		451	Δ		
340		AGI	Δ		
270				1	HIGH ADDR VEY TARLE
200		CTV	##F11	·	HIGH HOOK KET THEEE
200		011	#400 #400		
370		CTY	H	,	
400		511	¥14	:	LUW ADDRESS
410		IAT		:	IRANSFER IU TAS
					PUINTER
420		LDX	#\$00	!	COUNTER KEY BUFFER
430	NEXT	JSR	\$81FB	ļ	SWITCH OFF BASIC
440		LDA	(\$14),Y	!	GET BYTE OF DATA
450		PHA		ļ	STORE TEMP
460		JSR	\$8202	ł	SWITCH ON BASIC
470		PLA		1	GET BACK DATA
480		RED	\$FXIT	ł	
400		CMP		·	ARRAL FOR RETURN
500			4CTOPE	ı	
510			₽01UKE ##015	:	I DAD CODE EOD DETUDAL
210			#₽UU #0077 \/	:	CTODE IN VEYDOADD
320		51A	₽0∠//,X		BUFFER

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530 540 550 560	EXIT	INX INY BNE NEXT STX \$C6	! ! ! !	INCREASE COUNTER INCREASE POINTER FORCED-GET NEXT DATA NO OF CHARS IN KEY
				BUFFER
570		LDA #\$7F	!	RESET CIA DATA PURI
580		STA \$DC00		
590		RTS		
875F	CBM	872F CON	Г	
873B	CONT2	8791 EXIT	Γ	
874E	Fi	8751 F3		
8754	F5	8778 NEX1	Г	
8762 878A	NOCHANGE STORE	872C NOR	141	L

LINES 10–60: What we do here is to get into the Y register the value of the current key press and see if it is a function key or not. Function keys have values from 3 to 6 inclusive. Line 60 has the normal address of the Keyboard Table Setup Vector and if we do not find a function key this is where we direct the flow.

LINES 70–110: This part of the routine checks to see if the last KEY is the same as the current KEY. If it is, then off to the standard routine to avoid auto-repeat. At this point we have the current key value in the Y register and the current shift value in the accumulator.

LINES 120–130: This is part of the housekeeping. We copy the current values we have obtained into the last key registers. This is not only for our routine but also for the normal key interpreting routine.

LINES 140–230: We now take our key value, find which key it is and give a number corresponding to the number on the key itself. The BIT commands will not alter any data at all except for the status flags (which we are not testing here). They allow us to 'hide' an instruction within the address, in these cases loading the Y register, saving bytes and branch instructions. For instance, the BIT address in line 210 is \$01A0 which is stored in memory as A001, which is the code for LDY #\$01.

LINES 240–280: We now do the same for the shift value. If there is no shift or just the standard shift, there will be no need for any alterations so they would branch off in line 250. The logo key requires the value of 8 (1+8 giving key 9 and so on) and both shifts 9. We again do this using the BIT function.

LINES 290–360: Here we subtract one from the key value and then add the result to the shift value, ending up with a value between 0 and 15.

This total will be in the accumulator which is then increased 16 times by the four ASL instructions. We now end up with a value between 0 and 15 which will be the pointer to the data for that particular key.

LINES 370–420: The start position of the data table is put in registers $14 and $15 . We also transfer the pointer in the accumulator to the Y register. Lastly, we initialize the x register to zero to use as a counter to the number of characters we put in the keyboard buffer.

LINES 430–550: At last we can get our data and use it. Earlier we said that we were going to put our data in the RAM under the BASIC ROM. To read it back, we have to 'remove' the ROM to access it. This we do by a call to an earlier routine in the UTILITY which you will come to later. Now we pick up a byte of data and put it on the stack for temporary safe keeping, as we require the accumulator for re-enabling the BASIC ROM. With the ROM back, and having recovered the byte, we have two checks before storing it. The first in line 480 is to see if the byte is zero, signifying that all the relevant data has been collected and we can finish up. The second is a check for the 'left arrow', which signifies the user wants a return to be included (more of this in programming the keys). If this succeeds, we will change the byte to the ASCII code for return.

The data is stored in the keyboard buffer starting at the beginning and working upwards – it will be removed in the same order. We do not need to check for overflow as we are only allowed ten characters to be programmed (see next section). Therefore, the zero, which is not stored, cannot be later than the eleventh byte.

Having stored our byte, the two registers are increased by one and we branch back to get a further byte. The branch is enforced as we will not increase Y enough to return it to a zero. The highest value Y will achieve is \$FB (251 dec).

LINES 550–590: The end is near. Having stored all our data, the x register will hold a number equal to the total number of characters we put into the buffer. This is put into the register denoting how many characters are in the buffer and the operating system will only take that many off. The following two instructions are again housekeeping in that we reset the data port for collection of the next press. A return follows, but didn't we come by a JMP? This is true, but the whole key routine is entered by a JSR where the vectored jump is found. We do not now need the use of the normal key interpreting routine so we can go straight back to the main interrupt.

Key

COMMAND SNYTAX KEY

Displays the current data assigned to the keys in a form which can be amended.

```
KEY[number between 1 and 15], "[data]"
```

Assign data to a particular key. If a return is required, type a " \leftarrow " to signify this. Quotes cannot be used as data. A typical command could look like this:

KEY 7, "LIST←"

Here is a full list of the key numbers and how to achieve them:

KEY	1	– F1 ONLY
KEY	2	– F1 + SHIFT
KEY	3	– F3 ONLY
KEY	4	– F3 + SHIFT
KEY	5	– F5 ONLY
KEY	6	– F5 + SHIFT
KEY	7	– F7 ONLY
KEY	8	– F8 + SHIFT
KEY	9	– F1 + LOGO KEY
KEY	10	– F1 + SHIFT + LOGO KEY
KEY	11	– F3 + LOGO KEY
KEY	12	– F3 + SHIFT + LOGO KEY
KEY	13	– F5 + LOGO KEY
KEY	14	– F5 + SHIFT + LOGO KEY
KEY	15	– F7 + LOGO KEY
KEY	16	– F7 + SHIFT + LOGO KEY

KEY 0... will generate a SYNTAX ERROR. We had thought about using this as a way of turning off the key routines, but decided on a separate command. This makes it more of a conscious decision rather than a typing error. The command will be OFF, which is discussed later.

We have seen that we can make use of the four 'mystery' keys by getting data output on their use and in fact having 16 keys when used with the shift and logo keys. Now we have a routine to program the data, in which the user can decide what data to apply. This operation is acted upon through the keyword κ_{EY} .

KEY will perform three functions. It will 'switch' on the keys if they are off. This is performed in both of the following options. The choices are to program a key or to display the data applied to all the keys, which can then be amended on the display.

As we have said, there are two routines included in this. There is one routine to program individual keys and one routine to display the data

assigned to all keys. The latter is very similar to the interrupt routine discussed earlier except that the data goes to the screen rather than a buffer. The former in many ways is the reverse: we take data from a buffer – the input buffer – and put it in a table.

ASSEMBLY LISTING

9	*= \$864D				
10		LDA	\$805B	ļ	CHECK IF INTERRUPT SET FOR KEYS
20		CMP	#\$87		
30		BEQ	START	ļ	YES
40		LDA	#\$87		
50		STA	\$805B		
60		LDA	#\$22		
70		STA	\$8056		
80		JSR	\$8054	ł	SET INTERRUPT
90	START	JSR	\$0079	!	GET LAST BYTE AGAIN
100		BEQ	DISPLAY		
110		JSR	\$81F5	!	GET PARAMETER
120		JSR	\$AEFD	!	CHECK FOR COMMA
130		LDA	\$14		
140		BEQ	SYNTAX	ļ	NO KEYØ
150		CMP	#\$11	!	HIGHEST KEY IS 16
160		BCS	SYNTAX		
170		DEC	\$14		
180		LDA	\$14	ļ	SET TO CALCULATE
190		ASL	A	ļ	CALCULATING POINTER
200		ASL	A		
210		ASL	A		
220		ASL	A		
230		TAY			
240		LDA	#\$A1	ļ	HI ADD FOR KEY TABLE
250		STA	\$15		
260		LDA	#\$00	ļ	SET LO ADD FOR KEYS
270		STA	\$14		
280		LDX	\$0A	ļ	COUNTER MAX NO OF CHARS
298		JSR	\$0079	ļ	GET LAST BYTE AGAIN
300		CMP	#\$22	ļ	IS IT A QUOTE
310		BEQ	CONT2	!	YES
320	SYNTAX	JMP	\$AF08	ļ	PRINT SYNTAX ERROR
330	CONT2	JSR	\$0073	ł	GET NEXT BYTE
340		BED	ZERO	ļ	END OF DATA INPUT
350		CMP	#\$22	ļ	IS IT A QUOTE
					50 T

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360	1	BEQ ZERO	!	END OF DATA INPUT
370	l	STA (\$14),Y	!	STORE DATA IN
				TABLE
380	I	INY	!	INC TABLE POINTER
390	T	DEX	!	DEX CHAR COUNT
400		BNE CONT2	ļ	IF ZERO MAX NO
				CHARS REMAINDER
				IGNORED
410	ZERO	LDA #\$00	ļ	END OF WORD MARKER
420		STA (\$14),Y		
430		JSR \$0073	!	GET NEXT BYTE
440		RTS	ļ	FINISHED
450	DISPLAY	LDX #\$00	ļ	SET COUNTER
460		STX \$5F		
470		INX		
480		LDA #\$20	ļ	SPACE AS NO TEN'S
				DIGIT
490		STA \$22		
500		LDA #\$31	ļ	ASCII FOR ONE
510		STA \$23		
520		LDA #\$00	ļ	LO BYTE OF DATA TABLE
530		STA \$14		
540		LDA #\$A1	!	HI BYTE OF DATA TABLE
550		STA \$15		
560	PD1	JSR PRINT		
570		INC \$23	ļ	INCREASE NUMERAL
580		INC \$5F	ļ	INCREASE KEY COUNT
590		INX		
600		CPX #\$0A	ł	HAVE WE DONE KEYS1-9
610		BCC PD1	!	NO
620		LDA #\$31	!	NOW HAVE A TEN DIGIT
630		STA \$22		
640		LDA #\$30		
650	55.0	STA \$23		
660	PD2	JSR PRINT		
670		INC \$23		
880		INC \$5F		
690		INX		
700		CPX #\$11	!	HAVE WE DONE 13
710		BCC PD2	!	NO
720		RTS	!	YES
730	PRINT	LDY #\$05	!	COUNTER
740	NEXTA	LDA PDATA,Y	!	PRINT " KEY "
750		JSR \$FFD2		
760		DEY		
770		BNE NEXTA		

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780		LDA	\$22		
790		JSR	\$FFD2	!	PRINT TEN'S
					NUMERAL OR SPACE
800		LDA	\$23	ļ	PRINT LOW NUMERAL
810		JSR	\$FFD2		
820		LDA	#\$2C		
830		JSR	\$FFD2	ļ	PRINT COMMA
840		LDA	#\$22		
850		JSR	\$FFD2	!	PRINT QUOTE
860		LDA	\$5F	ļ	CALC TABLE POINTER
870		ASL	A		
880		ASL	A		
890		ASL	A		
900		ASL	A		
910	CONT	TAY		!	PUT POINTER IN Y
920	NEXT	JSR	\$81FB	!	SWITCH OFF BASIC
930		LDA	: (\$14), Y	ļ	GET CHAR OFF TABLE
948		Pha	I	!	TEMP STORE
950		JSR	\$8202	!	SWITCH ON BASIC
960		PLA	I	!	RETRIEVE CHAR
970		BEG	EXIT	!	FOUND END OF WORD
980		JSR	\$FFD2	!	PRINT CHAR
998		INY			
1000		BNE	NEXT	!	ENFORCED
1010	EXIT	LDA	#\$22		
1020		JSR	\$FFD2	ļ	PRINT A QUOTE
1030		RTS	5		
1040	PDATA	BYT	`\$20,′Y,	Έ	.,′K,\$20,\$0D
				_ .	
8691	CONT2		86A8 DIS	PL T	AY
8716	EXIT		8/04 NEX	1	
86E0	NEXTA		86BD PD1	.	
86D1	PD2		871C PDA	TA	
86DE	PRINT		8661 STA	RT	-

LINES 10–80: Earlier in the UTILITY, a routine will exist that is used when the extension is initialized or when STOP/RESTORE is used. This sets the Keyboard Table Setup Vector to where we want it to point to. These addresses can be changed by the OFF command. Here we look to see if the high byte of the address is pointing to our interrupt routine. If not, we change the address in the setting routine to point to our interrupt routine and then call the setting routine to initialize.

86A0 ZERO

868E SYNTAX

LINES 90-270: A call first to the CHRGOT routine to get the byte after the

KEY token. This is necessary as we have used the accumulator and so overwritten the byte. If the byte has set the zero flag, then there are no further parameters and a display of the key data is required. The program in that case branches to the display which starts at line 450.

Knowing we have got some parameters, off we go to our 'GET PARA-METER' routine (Chapter 6) and to a ROM coding to see if the byte after the key number is a comma. This coding will not only update the CHRGET address but will generate a SYNTAX ERROR if a comma is not found.

The parameter we want is now held in location \$14 - the key number. This value is put in the accumulator and checked for two things. If it is zero or greater than sixteen, it is out of bounds, so an error message is required, and therefore we branch off to get this printed. As in the interrupt routine it is easier to work in numbers 0 to 15 rather than 1 to 16 so we decrease the value in \$14 by one and then reload back into the accumulator.

To get the pointer to the required position in the data table, the number is multiplied by 16 with the ASL instructions. The y register will be the pointer so the value is transferred to it. Next we load two registers with the address of the data table start. Now we are in a position to get, and store, the data.

LINES 280–440: The data generated by using the function keys will be placed in the keyboard buffer. This buffer is only ten characters in length so we have to limit the input to that number. This is achieved by setting the x register to ten (\$0A). We said earlier that the comma check updates the CHRGET address so a call to the CHRGOT routine will get the next byte we want. This should be a quote; if not a SYNTAX ERROR is generated (remember that CHRGET skips spaces).

Now to get the data and store it. To get the data we make use of CHRGET. If the zero flag is set, the end of the command has been reached with either a colon or a zero placed by the BASIC input routine. The second quote is checked which also signifies the end of data input. If any of these are found, we branch off to end the routine at line 410. We can now store our data in the table under the BASIC ROM. We do not have to disable the ROM as you cannot store data in ROM so it is automatically stored in the RAM underneath. We increase the Y register which points to the table position. We decrease X which checks for overflow of data. If x reaches zero, the maximum number of characters has been stored. The flow only branches back to get the next byte if x is greater than zero.

To finish off, we store a zero after the last byte of data. This will help when retrieving the data to signify all data has been gathered.

We do another visit to CHRGET to get the next byte as BASIC expects this. This will cause a SYNTAX ERROR if you have input more than ten characters of data though the first ten bytes will have been logged.

The RTS returns us to BASIC for further operations.

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DISPLAYING THE KEY DATA

LINES 450–550: These instructions set up the registers used in the display itself. The x register is again used as a counter. Location \$5F will have the value of the key number less one and will be used to calculate the pointer for data collection. Locations \$22 and \$23 hold the ASCII values of the key number. As keys up to and including 9 have only one digit location \$22 is loaded with the ASCII for a space character. \$23 starts with the ASCII for 1 and will be incremented. Finally, we load up the address of the start of the data table into registers \$14 and \$15.

LINES 560–610: Call the coding to print key data for keys 1 to 9. After calling, the ASCII value in \$23 is increased along with the key number register \$5F. Register x is also increased and checked to see if it has reached \$0A (ten). If so, we would have to reset the ASCII numbers before printing further data. If x has not reached this value, we branch back to call the print coding for the next key.

LINES 620–720: First we reset locations \$22 and \$23. Key numbers from 10 to 16 have to be displayed so we have two digit numbers, the first always being one. Therefore, \$22 is loaded with \$31, the ASCII for one. The other is initialized to zero in ASCII format. We now continue to print out the key data, incrementing \$5F, and x each time, until x reaches the value of 17 just after being incremented. This value of x signifies we have finished the display so we exit from the routine and hand control back to BASIC.

The Print Routine to Display the Key Data

This part of the command is entered 16 times in total to print the data to the screen. The value for calculating the pointer, held in \$5F, is set before these lines are implemented, as are the ASCII values of the key number. We use the KERNAL routine at \$FFD2 to print a character to the screen. The data is printed out in the same format as it was entered. This is done so that it can be changed, just like normal screen editing, if required.

LINES 730–770: The start of every key display line will be the same. These lines will print this from the area of data at the end of the routine (Line 140). We start with a return so it starts on a new line, then a space to give better clarity if the border and screen are different colours, especially if the border and text colours are the same. KEY is printed next, followed by a space for presentation.

LINES 780–850: The key number is printed, followed by a comma and the first set of quotes.

LINES 860-910: The key number, less one, is taken from \$5F and

increased 16 times with the now familiar four ASL instructions. The result is transferred to the Y register for the data pointer.

LINES 920–1030: Get the key data. First we switch off BASIC to get the data underneath. After returning the BASIC, we print the data as long as it is not the 'end of data' zero. Printing finished, we update the pointer and go back to get the next byte. When the zero is found we exit and print the closing quote. Then it's back to the main key display routine.

OFF – Turn off the keys

COMMAND SYNTAX OFF There are no parameters with this command.

If you want to use the function keys within a program simply as keys, you will want to be able to disable the programming they have been given. The command that enables you to do this is OFF. All we do is to alter the addresses in the routine that sets the Keyboard Table Setup Vector back to its normal address. Once changed, we call the routine to change them in the BASIC work area. Do not forget that they can be re-enabled with any KEY command.

ASSEMBLY LISTING

9	*=\$ 8799				
10		LDA	# \$48		
20		STA	\$8056	ļ	CHANGE LOW ADD IN
					SETTING ROUTINE
30		LDA	#\$EB		
40		STA	\$805B	!	CHANGE HIGH ADD IN
					SETTING ROUTINE
50		JSR	\$8054	ļ	CALL SETTING ROUTINE
60		RTS			

Stand alone programmable function keys

Perhaps this chapter would have been better located between Chapters 6 and 7. It was difficult to decide on its position as it also uses information from both Chapters 2 and 3, but will not work as it stands.

To provide programmable function keys without using the keyword enable routine, the 'get parameter' and 'switch off BASIC' routines have to be copied from Chapter 6. The whole routine may then be relocated and the actions of KEY and OFF performed using SYS commands.

5 Utilities in BASIC

General

This chapter includes many of the utilities in the form of BASIC subroutines and programs. You do not really need any of that which follows if you load-up the UTILITY each time. In time we suspect that the simple routines contained here will not only prove useful, but will also give you plenty of ideas of your own.

A number of the utilities require that you generate an ASCII file of a program on tape or disk. This produces a file in the same format as would be received at a printer or the screen itself. The resulting sequential file contains the program in 'un-tokenized' form. To do this, output must be directed to the desired device with an OPEN and CMD sequence. For a tape this is:

OPEN 1,1,1, "PROGRAM": CMD1:LIST[from – to] PRINT#1:CLOSE 1

and for a disk:

OPEN 2,8,2,"PROGRAM,S,W":CMD2:LIST[from – to] PRINT#2:CLOSE 2

Most of the utilities given here are in the form of subroutines and have been numbered in the 60000s to allow them to be easily added on to your own programs as and when appropriate. They may be included in whole, or in part, by a suitable merge or append technique. You may wish to combine a number of them together to form useful modules which in the future may save many hours of repetitive work. This you can easily do by using the mini-renumber and merge programs given. Many of the routines can be extended, but they have deliberately been kept as short as possible. Always try to adopt a 'house' format to simplify the creation of future programs. This may only be a simple line numbering sequence where: the working part of your program lies between lines 100 and 9999; the specific subroutines lie between 10000 and 50000; and your library routines are from 50000 on.

The information upon which much of the following is based is contained in Chapter 1 and we refer you to that chapter. The utilities that follow are arranged in alphabetical order.

Keyword – APPEND

Append 1

Function: To append two BASIC programs in memory (nose-to-tail)

In the past, whenever you have loaded a program, it has erased the one currently in memory. This need not be the case. BASIC can start at any address in memory and need not always be the default of 2049 (\$0801). The pointer (TXTTAB) to tell the 64 where BASIC begins is held in RAM and can therefore be changed. It is even possible to have two BASIC programs resident in memory concurrently by changing the necessary zero page pointers, though only one could be running at any time. We can manipulate these pointers to allow us to append one program to another.

With a program in memory change TXTTAB to point to its end (VAR-TAB-2) by:

A=PEEK (45) : POKE 43, A-2: POKE 44, PEEK(46) + A<2

The program to be appended will now be loaded at the end of the existing one. Resetting the start of BASIC will make the 64 see both programs as one by:

POKE 43,1:POKE 44,8 (assumes original start was 2049/\$0801)

The resulting program may then be edited or saved in the usual way. Many texts say the appended program should have line numbers higher than the original. This is not essential, but some confusion can result if this is not so. Try appending when the second program does not have higher line numbers and see.

The combined program will run correctly until a GOTO OF GOSUB references a line which occurs twice. By virtue of the way these commands work, the branch will always be taken to the first occurrence of a line.

Append 2

Function: To append two programs on disk (BASIC or machine code)

Program files on disk store an image of the memory which the program occupied. The first two bytes record the load address and the last byte is a zero to mark the end of file. They can, however, be read and written in a sequential manner. This allows us to append files in much the same way as we did above, but this time performing the operation solely on disk. The following program will append two programs which will load at the address of the first: 100 Utilities in BASIC

LINE ACTION

130 Open up 'Program' files for read and write.

140 Read first program and make a byte by byte copy

TO in the combined file. Skip the terminating zero byte

180 and jump to read the second program.

200 Read the load address and discard it.

210– Copy the remainder through to produce the combined file.

```
100 INPUT"FRONT PROGRAM":F$
110 INPUT"[2SPC]END PROGRAM";E$
120 INPUT"FINAL PROGRAM";R$
130 OPEN 2,8,2,F$+",P,R":OPEN 3,8,3,R$+"
, P, ₩"
140 GET#2,A$
150 B$=A$:GET#2,A$:IF ST AND 64 GOTO 180
160 IF A$="" THEN A$=CHR$(0)
170 PRINT#3,B$;:GOTO 150
180 CLOSE 2
190 OPEN 2,8,2,E$+",P,R"
200 GET#2,A$:GET#2,A$
210 GET#2,A$:IF ST AND 64 GOTO 240
220 IF A$="" THEN A$=CHR$(0)
230 PRINT#3,A$;:GOTO 210
240 PRINT#3,CHR$(0);
250 CLOSE 3:CLOSE 2
```

Append 3

Function: To reopen an existing closed sequential file on disk and continue writing data from the previous end of file.

This is a standard disk command which is not made clear in the disk manual. Its format is:

```
OPEN 2,8,2,''TEST,A''
```

Subroutine keyword – AUTO NUMBER

Function: To automatically generate line numbers as code is entered.

Initiation: RUN 60000

This allows the start line and increment to be set. The line number is printed, followed by any characters typed. When RETURN is pressed the program enters the line, resets the line number variables (as an edit destroys all variables) and reruns itself by forcing two RETURNS into the keyboard buffer. As written, the program will not accept any line not followed by BASIC code (equivalent of delete line).

LINE ACTION

60010 Position cursor to 3rd line down and print line number in black.

- 60020 Generate a flashing cursor not normally present on a GET
- 60040 Watch out for null lines
- 60060 Print line to GOTO 60010, reset variables, restart program and move to HOME.
- 60070 Set NDX for two characters in keyboard buffer. Put two returns in k/b buffer (KEYD). On END KEYD will be emptied and the returns will enter the line and execute line from 60060

```
60000 INPUT "START[4SPC]";LN: INPUT "INC
REMENT";I%
60010 B$="":PRINT CHR$(147);CHR$(17);CHR
$(17);CHR$(17);CHR$(144);LN;CHR$(154);
60020 POKE 204,0:POKE 207,0
60030 GET A$:IF A$="" GOTO 60020
60040 PRINT A$;:IF B$="" AND ASC(A$)=13
GOTO 60010
60050 B$=A$:IF ASC(A$)<>13 GOTO 60020
60060 PRINT "LN=";LN+I%;":I%=";I%;":GOTO
60010";CHR$(19)
60070 POKE 198,2:POKE 631,13:POKE 632,13
:END
```

The version below is a little more flexible. It will not only delete an existing line if RETURN is pressed after its number, but also allows you to change the printed line number to any value. Subsequent line numbers will increment from the new value until it is again changed. The main difference is the addition of the code to evaluate the current line number ((60060 and 60070). This is done by reading from the start of the fourth screen line until a non-numeric code is encountered and reassigning the line number variable 'LN'.

```
60000 INPUT "START[4SPC]";LN: INPUT "INC

REMENT";I%

60010 PRINT "[CLS][3CD][BLK]";MID$(STR$(

LN),2);"[L BLU]";

60020 POKE 204,0:POKE 207,0

60030 GET A$:IF A$="" GOTO 60020

60040 IF ASC(A$)<>13 THEN PRINT A$;:GOTO

60020

60050 PRINT:B$="":I=1143

60060 I=I+1:IF PEEK(I)>47 AND PEEK(I)<58

THEN B$=B$+CHR$(PEEK(I)):GOTO 60060
```

```
60070 LN=VAL(B$):PRINT "LN=";LN+I%;":I%=
";I%;":GOTO 60010[HOM]";
60080 POKE 198,3:POKE 631,13:POKE 632,13
:POKE633,13:END
```

Program keyword – BASES

Function: To convert hex to decimal, binary to decimal and vice versa

This program contains four useful inter-base conversion subroutines. The hex to decimal is most useful if you wish to use hex rather than decimal values in the DATA statements for a machine code BASIC loader. The *Programmer's Reference Guide*, Chapter 3 uses binary patterns for the sprite data in the 'BALLOON' program but pictorial data is also enlightening when setting up user-defined characters and makes for easier editing.

No explanation is given as the program is easy to follow.

```
100 PRINT"1 HEX/DEC":PRINT"2 DEC/HEX":PR
INT"3 BIM/DEC":PRINT"4 DEC/BIN"
110 PRINT: INPUT" SELECT ":N
120 ON N GOSUB 150,240,330,400
130 GOT0100
140 ON N GOSUB 150,240
150 PRINT: INPUT "HEX[4SPC] ":A$
160 IF LEN(A$) <4 THEN A$=LEFT$("0000"+A$
.4-LEN(A$))+A$
170 A=ASC(A$)-48
180 B=ASC(MID$(A$,2,1))-48
190 C=ASC(MID$(A$,3,1))-48
200 D=ASC(MID$(A$,4,1))-48
210 E=256*(16*(A+7*(A)9))+B+7*(B)9))+16*
(C+7*(C)9))+D+7*(D)9)
220 PRINT:PRINT"$ ";A$;" = D";E:PRINT
230 RETURN
240 PRINT: INPUT "DEC[4SPC] ";G:A=INT(G/256
):B=G-A*256:IF G<0 OR G>65535 GOTO 240
250 C=INT(A/16):D=A-16*C
260 C$=CHR$(48+C):IF C>9 THEN C$=CHR$(C+
55)
270 D$=CHR$(48+D):IF D)9 THEN D$=CHR$(D+
55)
280 E=INT(B/16):F=B-16*E
290 E$=CHR$(48+E): IF E>9 THEN E$=CHR$(E+
55)
```

```
300 F$=CHR$(48+F):IF F>9 THEN F$=CHR$(F+
55)
310 PRINT:A$=C$+D$+E$+F$:PRINT "D":G:" =
$ ":A$:PRINT
320 RETURN
330 PRINT: INPUT "BIN[4SPC] ";A$
340 A=0:A$=RIGHT$("000000000000000"+A$,
16)
350 FOR I=16 TO 1 STEP -1
360 B$=MID$(A$,I,1):IF B$="1" THEN A=A+2
^(16-I)
370 NEXT I
380 PRINT:PRINT"B ";LEFT$(A$,8);" ";RIGH
T$(A$,8);" = D";A:PRINT
390 RETURN
400 PRINT: INPUT "DEC[4SPC] ";A:IF A>65535
OR A<0 GOTO 400
410 B$="":D=A:FOR I=15 TO 0 STEP -1
420 B=INT(A/2^I):IF B=1 THEN B$=B$+"1":G
OTO 440
430 B$=B$+"0"
440 A=A-B*2^I:NEXT I
450 PRINT:PRINT"D";D;"= B ";LEFT$(B$,8);
" ";RIGHT$(B$,8):PRINT
460 RETURN
```

Program keyword – DATALINES

Function: To generate BASIC data statements for machine code programs.

Once again the keyboard buffer is used to generate program lines. This time there are more variables in use than would conveniently fit on a single assign line so they have been put 'out of the way' in the cassette buffer. Only variables in the normal BASIC variable storage area are lost by an edit. The resulting data values are generated to the nearest ten bytes.

LINE ACTION

60000-	Data	input.
--------	------	--------

- 60060- POKE values to TBUFFR.
- 60090 Recycle from here. Re-read next line number,
- 60100 step,
- 60110 start address,
- 60120 end address for current line,

```
60130 and end address of program. If finished STOP program.
60140- Print line number, DATA, the values and GOTO 60090.
60210- Increment line number, address, and set up k/b ready for END.
```

```
60000 INPUT"START ADDRESS";B
60010 INPUT"END ADDRESS[2SPC]":E
60020 F=B:L=F+10
60030 INPUT"START LINE[3SPC]":S
60040 INPUT"LINE INC[5SPC]";T
60050 PRINT"[4CD]"
60060 POKE831, INT(E/256)
60070 POKE832,E-INT(E/256)*256
60080 POKE828, T: GOT060160
60090 S=PEEK(826)*256+PEEK(827)
60100 T=PEEK(828)
60110 L=PEEK(829)*256+PEEK(830)
60120 E=PEEK(831)*256+PEEK(832)
60130 IFL>=EGOT060270
60140 F=L+1:L=L+10
60150 PRINT"[CU][14SPC]"
60160 PRINTS;
60170 PRINT"DATA":
60180 FORP=FTOL:PRINTPEEK(P);"[CL],";:NE
XTP
60190 PRINT"[CL][3SPC]"
30200 PRINT"GOT030090[4CU]":
60210 POKE198,2:POKE631,13:POKE632,13
60220 S=S+T
60230 POKE826, INT(S/256)
60240 POKE827, S-INT(S/256)*256
60250 POKE829, INT(L/256)
60260 POKE830,L-INT(L/256)*256:END
60270 STOP
```

Subroutine keyword – DELETE

Function: To remove unwanted program lines en masse

Two delete routines follow. Both use the link address and line number storage at the start of a BASIC line during execution to perform the deletion. The first deletes line numbers as they are encountered whereas the second only deletes one line as the final step in the process. The first line of each routine reads TXTTAB to find out the current start of BASIC.

Delete 1

This routine deletes lines using the all-too-familiar keyboard sequence and as such requires no explanation.

```
60000 TX=PEEK(43)+PEEK(44)*256
60010 INPUT"DELETE FROM";LL:M=256:INPUT"
[7SPC]TO[2SPC]";UL
60020 IF PEEK(TX+2)+PEEK(TX+3)*M(LLTHENT
X=PEEK(TX)+PEEK(TX+1)*M:GOTO 60020
60030 POKE 828,UL-INT(UL/M)*256:POKE 829
,UL/M:GOTO 60050
60040 M=256:TX=PEEK(830)+PEEK(831)*M:UL=
PEEK(828)+PEEK(829)*M
60050 IF PEEK(TX+2)+PEEK(TX+3)*M>UL OR P
EEK(TX)+PEEK(TX+1)*M=0 THEN END
60060 PRINT "[CLS][3CD]";PEEK(TX+2)+PEEK
(TX+3)*M:PRINT"GOTO 60040[HOM]"
60070 POKE830,TX-INT(TX/M)*M:POKE831,TX/
M:POKE198,2:POKE631,13:POKE632,13:END
```

Delete 2

This is, perhaps, a more refined way to carry out the task. It takes fullest advantage of the way programs are stored in RAM and in particular the function of link addresses. The routine scans the line numbers until the start of the block to be removed is found. It records the address of this link address and then continues to scan for the end line number for the delete. Once a line number equal or greater is found, this link address is substituted at the start of the block link. One very large line has thus been created in memory. A simple keyboard program is then used to remove the start line and all others go with it. This is without doubt a lot faster than the first method, but has the disadvantage that you cannot see the lines as they go.

```
60000 TX=PEEK(43)+PEEK(44)*256
60010 INPUT"DELETE FROM";LL:INPUT"[7SPC]
TO[2SPC]":UL
60020 L=PEEK(TX+2)+PEEK(TX+3)*256
60030 IF L(LL THEN TX=PEEK(TX)+PEEK(TX+1
)*256:GOTO 60020
60040 IF L=0 THEN PRINT"LOWER LIMIT";LL;
"NOT FOUND":END
60050 LL=L:D=TX
60060 L=PEEK(TX+2)+PEEK(TX+3)*256
60070 IF L=0 THEN PRINT"UPPER LIMIT";UL;
"NOT FOUND":END
60080 IF L(UL THEN TX=PEEK(TX)+PEEK(TX+1
)*256:GOTO 60060
60090 POKE D, PEEK(TX): POKE D+1, PEEK(TX+1
)
60100 PRINT"[CLS][3CD]";LL;"[HOM]":POKE
198,1:POKE 631,13:END
```

Subroutine keyword – DUMP

Function: To display the current values of all simple numeric, string and function variables

Initiation: Type GOTO 60000

This routine will display the values of all simple variables in use at the time of calling. The variables will be displayed in the order in which they were created by the program. The routine will not handle arrays nor will it work if editing has been carried out prior to its being called (simply because all variable pointers will be reset to the end of program). It also displays the values of the variables it uses – sv, v\$, and so on. As these are the last variables to be created they will be the final ones to be displayed. Output may be directed to a printer by a simple:

OPEN 4,4:CMD 4:GOTO 60000

The display may be stopped by holding down any key and will resume on the release of that key. Pressing the STOP key will 'break' into the program and allow you to use the cursor keys to move up and change values. If you resume program execution with a GOTO, then the amended values will be used. A simple CONT would re-enter the dump subroutine at the break and dump any remaining variables.

The routine makes extensive use of the information contained in Chapter 1 on the storage of BASIC variables. Remember the first two bytes are the variable name adjusted for its type. The following is a description of the routine:

LINE ACTION

- 60030 Read the current value of VARTAB.
- 60040 Do the same for ARYTAB.
- 60050 Default values.
- 60055 If equal then no simple variables, edit used, or finished. If not equal more variables exist so continue.
- 60060 Read the seven bytes used for variable.
- 60070 Determine the type from the two name bytes and GOTO the appropriate subroutine, these being real, integer, string or func-
- 60100 tion. The name bytes must be changed back to their unmodified ASCII values by the subtraction of 128, as necessary, and '%' or '\$' suffixes printed where required.
- 60105 Pause if key held down (64=no key at sFDX. Note this is the current key not LSTX as in the *Programmer's Reference Guide*).
- 60110 Increment 7 bytes to next variable and recycle.
- 61000 Subroutine to convert 5 floating point binary bytes to decimal.
- 61500 Subroutine to convert the 2 of the 5 bytes used to a signed integer.
- 62000 Subroutine to read string length and location then find and build string.
- 62005 Avoids the single pass through FOR/NEXT if null string.
- 62020 Surround a string with quotes required for changing its value on a break.
- 62500 Subroutine to detect a function and simply acknowledge the fact as its current value will be picked up by one of the other routines.

```
60000 :
```

```
60010 :REM DUMP VARIABLES
60020 :
60030 SV=PEEK(45)+PEEK(46)*256 :REM_STAR
T OF VARIABLES
60040 SA=PEEK(47)+PEEK(48)*256 :REM STAR
T OF ARRAYS
60050 V$="":VV$=":V=0:VV=0:REM DEFAULTS
60055 IF SA=SV THEN END: :REM NO SIMPLE
VARIABLES OR EDIT USED
60060 FOR V=0 TO 6:V(V)=PEEK(SV+V):NEXT
V:REM READ VARIABLE NAME AND VALUE
60070 IF V(0) (128 AND V(1) (128 THEN GOSU
B 61000:REM REAL
60080 IF V(0))128 AND V(1))127 THEN GOSU
B 61500:REM INTEGER
60090 IF V(0)<128 AND V(1)>127 THEN GOSU
B 62000:REM STRING
30100 IF V(0)>128 AND V(1)<128 THEN GOSU
B 62500:REM FUNCTION
60101 IF PEEK(203)<>64 GOTO 60101
60110 SV=SV+7:GOTO 60040:REM INCREMENT C
OUNTER AND DO NEXT
61000 V$=CHR$(V(0))+CHR$(V(1)):REM_REAL
NAME
61010 V=(-1)^(V(3)AND128)*2^(V(2)-129)
61020 VV=(1+((V(3)AND127)+(V(4)+(V(5)+V(
6)/256)/256)/256)/128)
61030 V=V*VV:PRINT V$;"=";V:RETURN
61500 V$=CHR$(V(0)-128)+CHR$(V(1)-128)+"
":REM INTEGER NAME
61510 V=(V(2)AND127)*256+V(3)+(V(2))127)
*32768
61520 PRINTV$;"=";V:RETURN
62000 V$=CHR$(V(0))+CHR$(V(1)-128)+"$":R
EM STRING NAME
62005 IF V(2)=0 GOTO 62020
```

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62010 FOR V=1 TO V(2):VV\$=VV\$+CHR\$(PEEK(V(3)+V(4)*256+V-1)):NEXT V 62020 PRINT V\$;"=";CHR\$(34);VV\$;CHR\$(34) :RETURN 62500 RETURN:REM FUNCTION PICKED UP BY O THER ROUTINES

An alternative approach might be to use the technique in RENUMBER (see below). Namely, print a line which reads: PRINT the variable name and GOTO the point at which program execution should be resumed. If we get the cursor movements right and POKE a RETURN into the keyboard buffer, a dump could be performed. To tidy up, we should really clear the line which says 'PRINT and GOTO' with more cursor movements and spaces, and so on.

An obvious extension would be to include arrays. The logic involved in determining and printing the values of subscripted variables is identical to the above and, with care, the same subroutines could be used. The tricky bit is deciphering the array header to determine the number of dimensions and the size of each dimension. If you do decide to try this, do remember integer array values are stored in only two bytes and string pointers in three bytes, unlike their simple variable counterparts. You must also check that arrays do exist by examining STREND and comparing it with ARYTAB+1. Array headers have also been covered in Chapter 1. Including arrays will greatly increase the size of DUMP and in applications where memory is tight, prove impracticable. It is also difficult, so do not worry if your efforts are not rewarded immediately as a simple error in the logic can cause some very unexpected results.

Program keyword – LISTER

Function: To produce dated, paged and neatly formatted listings

The version given below has been written for an RS232 printer operating at 300 baud, 1 stop bit and no parity (see *Programmer's Reference Guide*, Chapter 6: 'Input/Output Guide'). The printer used also required a carriage return/line feed sequence to be generated at the end of each line. Therefore, the logical file number used has to be greater than 127, in this case #129. When using any RS232 device, it is advisable to OPEN-up the file at the start of the program to allocate the input and output buffers. For other printers, the OPEN and PRINT# statements below will have to be amended to suit.

If your printer does not support the CBM special characters, the program to be listed should first be run through CODER before generating the ASCII file. With a cassette, the OPEN command to read sequential data on line 210 should read OPEN 2,1,0,A\$(I).

The listing produced is ideal for permanent record, though as the

process takes a little time it is not recommended for intermediate listings. The final listing will have all text inset to column 7 and any wrap-around lines will also be inset. Specifying a line width less than the maximum available has the benefit of allowing space for comments (can save a lot of time in the future). A brief description follows:

LINE ACTION

- 100 See above.
- 110– Set parameters.
- 160– Allocate files to be listed to array A\$().
- 210 See above.
- 220 This line is included to get any leading returns. The number of these will depend on exactly how the ASCII file was generated. Once a CMD has been issued all returns normally sent to the screen will go to the file. Typically this will be two for the LIST. If zeros appear on your output then you will have to adjust the program or the way you generate the file.
- 230 Create bottom margin.
- 250 Build one line into string As.
- 260 Same problem as 220 at end of file. Assume a line number of zero is the end.
- 270 Reset line for text to start in col 7.
- 290 If length < max then print it.
- 300– Else split it and print first part. Recycle each time adding 6 leading spaces to continuations.
- 340 Print blank lines to next top of form before next program.

```
100 OPEN 129,2,0,CHR$(6)
```

```
110 PRINT"LISTER UTILITY":PRINT
```

```
120 INPUT"DATE[10SPC]";D$
```

```
130 INPUT"LINES/PAGE[4SPC]";LP:IF LP=0 T
```

```
HEN LP=66
```

```
140 INPUT"MAX CHARS/LINE";CP
```

```
150 INPUT"NO.OF PROGS[3SPC]";N:DIM A$(N)
```

```
160 PRINT:FOR I=1 TO N
```

```
170 INPUT"PROGRAM[7SPC]";A$(I)
```

```
180 NEXT I
```

```
190 I=0
```

```
200 I=I+1:LC=0:IF I>N THEN END
```

```
210 Z=1:OPEN 2,8,2,A$(I)+",S,R"
```

```
220 Z=1:GET#2,A$:GET#2,A$:GOSUB 320
```

```
230 IF LC>=LP-8 THEN FOR J=1 TO LP-LC:PR
```

```
INT#129, "":NEXT J:GOSUB 320
```

```
240 J=0:B$=""
```

```
250 J=J+1:GET#2,A$:IF A$<>CHR$(13) THEN
```

```
B$=B$+A$:GOTO 250
```

```
260 IF VAL(B$)=0 THEN GOSUB 340:GOT0200
270 L$=STR$(VAL(B$)):B$=MID$(L$+"[6SPC]"
,2,6)+MID$(B$,LEN(L$))
280 L=LEN(B$)
290 IFL<=CPTHEN PRINT#129,B$:LC=LC+1:GOT
0 230
300 L$=LEFT$(B$,CP):PRINT#129,L$:LC=LC+1
310 B$="[7SPC]"+MID$(B$,CP+1):GOT0 280
320 PRINT#129,"PROGRAM ";A$(I);" LISTED
ON ";D$;" LISTING PAGE";Z:LC=2:Z=Z+1
330 PRINT#129,"":RETURN
340 FOR J=1 T0 LP-LC:PRINT#129,"":NEXTJ:
CLOSE2:RETURN
```

Subroutine keyword – MERGE

Function: To merge two BASIC programs

In all the following where line numbers are common to both the program in memory and the merging program those of the latter will take precedence.

Merge 1

Where a program is less than 22 screen lines when listed, it may be merged very easily indeed. Simply load the short program and list it. Type NEW and move the cursor to the line below the last line of the list. LOAD the main program and then move up and simply press RETURN on all lines to be included in the final program.

This is the reason for having short keyword routines, to allow the above technique to be used on many of them.

Merge 2

The following subroutine will merge programs of any length. The program (or part of) to be merged must be stored as an ASCII file on disk or tape. The program resident in memory must, of course, include the merge subroutine.

Initiation: RUN 60000

The resulting program will be an amalgamation of the two programs and unlike APPEND the lines will be in the correct numerical sequence. At the completion of the merge an 'OUT OF DATA' or 'SYNTAX ERROR' will be displayed depending on how the ASCII file was generated and which program had the highest line number, but who cares, as the result is exactly what we wanted. The program may then be saved in the normal way (after deleting lines 60000– if they are no longer needed). The version given is for disk and the necessary changes for cassette have been included in the description below, but should be only too familiar by now.

The program uses the keyboard programming technique for the most part. There is one problem and that is that whenever an edit is performed all OPEN files are CLOSEd. So in theory only one line may be read from the file. Any further attempts to obtain input will result in a 'FILE NOT OPEN' error. The solution is simple. BASIC is made to believe a file is open even though an edit has been carried out by POKEing the necessary values into the zero page file registers for current logical file (LA), secondary address (SA) and device number (FA).

LINE ACTION

- 60010 For tape OPEN 1,1,0,F\$
- 60020 Get bytes until numeric code. This overcomes the problem in LISTER and perhaps should also be used in that program.
- 60030 Set file parameters by poking into LA, SA and FA. For tape use 2, 0 and 1 (0=read 1=cassette).
- 60050 As the first numeric character has been found, mustn't forget to print it. B\$
- 60060- As all other programs using keyboard.
- 60080 Set up k/b buffer on END to enter printed line and GOTO 60030, the cycle repeating until all done.

```
60000 INPUT"PROGRAM ";F$
60010 OPEN2,8,2,F$+",S,R"
60020 GET#2,B$:IF VAL(B$)<1 GOTO 60020
60030 POKE 184,2:POKE 185,2:POKE 186,8:P
OKE152,1
60040 PRINT"[CLS][3CD]";
60050 PRINTB$;:B$="
60060 GET#2,A$:PRINTA$;:IF A$<>CHR$(13)
GOTO 60060
60070 PRINT"GOTO 60030[HOM]"
60080 POKE 198,2:POKE 631,13:POKE632,13:
END
```

Merge 3 (tape only)

This is the cleverest tape merge we have seen. It was originally worked out by J. Butterfield and B. Templeton for the PET and all we have done is to modify it for the 64.

Again, the program to be merged must be on tape in ASCII format. The statements may be typed in direct mode or, as in this case, be a subroutine. In direct mode the contents of the quotes should be typed after performing the necessary cursor moves and RETURN pressed at the end. Line 60030 is needed only in program mode.

Initiation: RUN 60000

The key to this is the POKE 153,1 (DFLTN) which changes the default input device after each line has been merged from the usual default of \emptyset (the keyboard) back to the cassette (1).

```
60000 INPUT "PROGRAM ";F$
60010 POKE 19,1:OPEN 1,1,0,F$
60020 PRINT"[CLS][3CD]POKE 153,1:POKE 19
8,1:POKE 631,13:PRINT CHR$(19)"
60030 POKE 198,1:POKE 631,13:PRINT "[HOM
]":END
```

The most common problem with merge is if a program line is in excess of 80 characters when listed (possible if abbreviations have been used). The merge will be unsuccessful as the cursor movements will be incorrect and also BASIC's input buffer will overflow.

Program keyword – OLD

Function: To recover NEWed programs

The command NEW does not actually erase the program in memory, it simply changes the first line's link address to 00 00 (2049 and 2050) and therefore fools BASIC into thinking that there isn't a program present. In addition, all variable pointers are reset to the end of the program, which in this case is the start of BASIC itself (action of CLR). The following uses these facts to recover the program by resetting the necessary pointers.

To use OLD, the start of BASIC (TXTTAB) must be set above the end of the NEWED program and TXTTAB-1 set to zero by:

POKE 43,01:POKE 44, no. of pages:POKE (no. of pages) *256,0:NEW

OLD may then be loaded and run. The erased program will be recovered and you are back in business. As a point of interest OLD is still present higher in memory and will remain so until overwritten by variable data or a larger program.

The program works by hunting from the input value of TXTTAB+4 (ignore first three zeros) for three consecutive zero bytes which mark the end of the erased program. *En route* the first link is changed to point to the second line. Once found, TXTTAB is changed to point to the specified start and VARTAB, to the end address. A CLR then tidies up and the original program is LISTEd.

60000 INPUT"TXTTAB(2049)";TX:MX=256 60010 POKE 828,TX-INT(TX/MX)*MX:POKE 829 ,TX/MX 60020 X=TX+4+J:IF PEEK(X)<>0 THEN J=J+1: GOTO 60020 60030 POKE TX,X+1-INT((X+1)/MX)*MX:POKE TX+1,(X+1)/256:TX=X+1 60040 X=PEEK(TX)+PEEK(TX+1)*MX:IF X<>0 T HEN TX=X:GOTO 60040 60050 TX=TX+1:POKE 830,TX-INT(TX/MX)*MX: POKE 831,TX/MX 60060 POKE 43,PEEK(828):POKE 44,PEEK(829):POKE 45,PEEK(830):POKE 46,PEEK(831) 60070 CLR:LIST

Subroutine keyword – PLOT

Function: To position the cursor to a specified screen location

There are many ways of positioning the cursor. The most common way is to include the necessary control characters inside quotation marks. This can be expensive on memory if a lot of cursor movement is used. The movement is also relative to the current location and not absolutely fixed to some reference point unless a clear screen or home cursor is first issued. Many micros have TAB(x,y), POS(x,y) or HTAB x and VTAB y functions within their BASICs to position the cursor. The following are just two ways of doing this on the 64 with its unmodified BASIC.

Plot 1

This uses a simple subroutine into which are passed the line and column position. First, two strings are defined – preferably at the start of the program as they remain unchanged throughout the run for speed of access. They are:

1 Y\$=''[24CD]'':X\$=''[40CR]''

and have been allocated line number 1. To position to any location, the x and y coordinates are passed to the subroutine which simply homes the cursor then prints the appropriate number of cursor downs and rights.

In the example below, lines 100 to 130 have been included for demonstration purposes.

```
1 Y$="[23CD]":X$="[40CR]"
100 INPUT"COLUMN";X
110 INPUT"[3SPC]ROW";Y
120 GOSUB 1000:PRINTX;",";Y
130 GOTO 100
1000 PRINT"[HOM]";RIGHT$(Y$,Y);RIGHT$(X$,X);:RETURN
```

The top left of the screen is considered as ' \emptyset , \emptyset '. The semicolon at the end of the print in 60000 is included to hold the cursor at the set location.

The idea of holding frequently used character patterns, control characters, and so on as string variables can reduce memory usage and also makes for easier-to-read code.

Plot 2

This second plot routine uses the same zero page locations as the KERNAL function PLOT see *Programmer's Reference Guide*. These are PNT (209/210), PNTR (211) and TBLX (214). If you look at the memory map in Chapter 5 of the PRG or Appendix K of this book, you'll notice locations from 200 to 245 all relate to the screen in some way or other. We are not going to run through them all, but try experimenting with them and see what happens. If in trouble, turn off the 64.

Let us look at the three locations we are going to use to accomplish PLOT in a little more detail.

PNT: contains the address of the start of the current line in low/high format. With the screen at its default start (1024–2023), this will hold a value $1024+40 \star row$ where row is in the range 0–24. Unusual results are produced if this does not correspond to the start of a physical screen line.

PNTR: holds the offset from the address held in PNT. It is the absolute screen column (0-39) when PNT holds the start of line address.

TBLX: holds the current physical screen row.

Using only PNT and PNTR, we can position the cursor to any X,Y location. The next print would occur at the specified point. However, when the cursor returns after the print, it reappears at or below the line it was on before PNT and PNTR were set. This is difficult to put into words and much easier to see. For example, if an input took place on line 23 and the cursor was then moved to line 10, column 10 and a PRINT took place without a semicolon, the cursor would reappear at the start of line 24 and not 11 as might be expected. To avoid this, we simply also set TBLX and all will be well. The routine below has the same effect as the first PLOT routine given. Again, lines 100 to 130 are included for demonstration only.

```
100 INPUT"COLUMN";X
110 INPUT"[3SPC]ROW";Y
120 GOSUB 1000:PRINTX;",";Y
130 GOTO 100
1000 POKE 214,Y:Y=1024+Y*40
1010 POKE 209,Y-INT(Y/256)*256:POKE 210,
INT(Y/256)
1020 POKE 211,X
1030 RETURN
```

Subroutine keyword – PRINT USING

PRINT USING is a very powerful output formatting command available in some BASIC languages. It allows numbers to be right or left aligned to a specified number of decimal places, or to be expressed in exponential format and much more. There are equally as many possibilities for strings. A routine to duplicate all the facilities would be very long, so here we have only considered the problem of formatting numbers.

Very quickly everybody picks up on the idea of:

 $X = INT(X \star 10^{\circ}W + .5)/10^{\circ}W$

to get numbers to a set number of decimal places, where W is the number of places. Unfortunately, due to the way numbers are stored, this is not guaranteed to produce the expected result. By way of a trivial example, try printing .01*649 and 649*.01 and see the difference. The result of any calculation is very much dependent on the order in which it was evaluated. To overcome the problem we have to resort to strings as these are the only type of variable we can fully format.

The following routine will format numbers not in scientific notation and will avoid the xx.x0001 type occurrence by not using any division. The value returned is right aligned to w decimal places and padded with leading spaces to a set width of L. The variable transferred is in x and the string x\$ is returned.

```
1 INPUT "X";X :INPUT"W";W:INPUT"L";L: G0
SUB 60000:PRINTX$
2 GOTO 1
60000 X$=STR$(INT(X*10^W+.5)):LE=LEN(X$)
60010 SZ$=".000000000000000":S2$="[31SPC]"
60020 IF LE(W+2 THEN X$=LEFT$(X$,1)+MID$
(S2$,1,W+2-LE)+RIGHT$(X$,LE-1)
60030 IF LE)=W+2 THEN X$=LEFT$(X$,LE-W)+
"."+RIGHT$(X$,W)
60040 X$=RIGHT$(S2$+X$,L):RETURN
```

To illustrate its use, the following display

COL1	COL2 COL3	COL4
99.000	100.00 .999	9.51456
100.091	98.22 .010	11.00000

would be produced by the program lines; where C1-4 represent the values to be PRINTED in cols 1-4 (set elsewhere within your own program).

```
1 C1=99.00001:C2=100.00123:C3=.99888:C4=
9.514569:REM EXAMPLES
10 L=10:W=3:X=C1:GOSUB 60000:PRINT X$;
20 W=2:X=C2:GOSUB 60000:PRINT X$;
30 L=5:W=3:X=C3:GOSUB 60000:PRINT X$;
40 L=14:W=5:X=C4:GOSUB 60000:PRINT X$
```

Where numbers are very large or very small, simply raise them to an appropriate power of ten prior to calling the routine and head the output ' $\star 10^{\circ}$ N'.

Subroutine keyword – RENUMBER

Function: To renumber a specified section of a program

It is not possible to write a full renumber program in BASIC which does not use ASCII disk files (somebody will no doubt wish to disprove this statement). There are many problems, the biggest of which is in renumbering GOTOS, GOSUBS and THENS line destinations. It is relatively easy, albeit slow, to hunt these out by their token values. The problem arises in correcting destinations which are held in ASCII form. For example, GOTO100 is held as 137 49 48 48 (\$89 \$31 \$30 \$30 in hex). If during the renumbering process the destination changes by a magnitude of ten or more (the overall length changes), we have to move all code from the byte following the reference up or down in memory, recalculating link addresses as we go. If all references are entered as five figures as standard, this problem is eliminated, for example, GOTO00100. Entering line numbers in this way is rather tedious and is considered impractical. Machine code renumber programs use the 'crunch tokens' routine and the necessary memory moves are performed as part of this routine when a line is added or removed. See RENUM in Chapter 7.

The program below only renumbers the lines. It will renumber all or only a set block. The new line numbers need not even be in sequence with the rest of the program, though problems will arise if they are referenced. The user will have to manually change all GOTOS, etc. This subroutine is really intended to allow you to put together a number of the shorter routines in this chapter.

```
60000 TX=PEEK(43)+PEEK(44)*256:MX=256
60010 INPUT"RENUMBER FROM";LL:INPUT"[9SP
C]TO[2SPC]";UL
60020 INPUT"[5SPC]NEW LINE";S:INPUT"[9SP
C]STEP";I
60030 IF PEEK(TX+2)+PEEK(TX+3)*MX(LL THE
N TX=PEEK(TX)+PEEK(TX+1)*MX:GOTO 60030
60040 S=S+J*I:IF TX=0 OR PEEK(TX+2)+PEEK
(TX+3)*MX)UL THEN END
60050 POKE TX+2,S-INT(S/MX)*MX:POKE TX+3
,S/MX:TX=PEEK(TX)+PEEK(TX+1)*MX
60060 J=J+1:GOTO 60030
```

Subroutine keyword – SQUASH

Function: To increase the speed of execution of BASIC programs

Many 'crunch' or 'compactor' programs are available, both commercially and in various journals. Their function is to increase the speed of execution of a BASIC program by the removal of redundant code.

There are many reasons why code is slower than it need be. Much of this code is useful at the time of developing the program, but is not required at run-time. Some examples have been given at the end of Chapter 1, but there are many more. Listing the more obvious:

Line numbers: When they are the reference for a GOTO, GOSUB OT THEN they are held in ASCII form. The shorter they are (that is, the lower the number), the quicker they are converted to numeric form. Therefore, a renumbering with an increment of 1 is advantageous.

REM: These are ignored at run-time and need only be retained if they are a destination. REMS also use valuable memory.

Spaces: Including spaces in a program makes for easier reading, but is unnecessary and wasteful (this is true only outside quotes).

Variable names: One-character names use less space and are found quicker.

Destinations: See Chapter 1 (page 23)

Screen: See Chapter 2 (page 32)

Print: Semicolons separating print lists are sometimes superfluous. They must be retained after a numeric variable and at the end of a PRINT list if a carriage return is to be inhibited.

Line length: Short lines use an extra five bytes each time (link=2 line=2 end=1) and also take time in working out the next line's details. Lines which are not destinations can be strung together, taking due care of the logic of any IF statements. Lines may be of any length, but are difficult to edit or generate once they exceed 80 characters (even if all the possible abbreviations are used, there is a limit to BASIC's input buffer).

FOR/NEXT loops: A surprising increase in speed is gained by omitting the variable on the NEXT statement. This eliminates the look-up operation for the variable name. Try timing:

FORI=1TO255:FOR J=1TO255:NEXTJ:NEXTI

and FOR.....NEXT:NEXT

Operating system: Once spaces have been eliminated, CHRGET itself may be modified to get rid of the test for spaces. (See Chapter 3).

The subroutine below will remove all unnecessary spaces, semicolons and REMS. Renumbering is left up to you. Once again, an ASCII file must first be generated of the program. The program is based on Merge 2 (see page 110) and only the differences from that program are described below:

LINE ACTION

- 60030 L is set to 1 to account for B\$ in first line.
- 60090 As Merge 2 line 60060, returning to k/b bit once a return found.
- 60100 Once a REM found, ignore all chars except return.
- 60110 Flag to indicate in or out of quote mode.
- 60120 Ignore spaces out of quotes.
- 60130 Keep spaces in quotes.
- 60140 Semicolons out of quotes require careful checking and this is carried out at 60210 on.
- 60150 Semicolons in quotes keep.
- 60160 If not an 'M', don't look for REM.
- 60170 Else see if preceding two chars were 'RE'.
- 60180 If they were, replace by a ':' and set RE to ignore everything following (see 60100).
- 60190 Build line to be printed.
- 60210 Handle the semicolon when out of quotes and eliminate if possible. Do this by getting next byte and if the list continues check for a preceding string or opening quote. Finally, re-enter the main body of the program where appropriate.

Initiation: RUN 60000

```
60000 INPUT"PROGRAM ";F$

60010 OPEN2,8,2,F$+",S,R"

60020 GET#2,B$:IF VAL(B$)<1 GOTO 60020

60030 L=1

60040 POKE 184,2:POKE 185,2:POKE 186,8:P

OKE152,1

60050 PRINT"[CLS][3CD]";

60060 GOSUB 60090

60070 PRINT"GOTO 60040[HOM]"
```

```
60080 POKE 198,2:POKE 631,13:POKE632,13:
END
60090 GET#2,A$:IF A$=CHR$(13) GOTO 60200
60100 IF RE=1 GOTO 60090
60105 IF A$="T" AND Q=0 AND RIGHT$(B$,4)
="PRIN" THEN P=-1
60106 IF A$=":" AND Q=0 THEN P=0
60110 IF A$=CHR$(34) THEN Q=NOT(Q)
60120 IF A$=" " AND Q=0 GOTO 60090
60130 IF A$=" " AND Q=-1 GOTO 60190
60140 IF A$=";" AND Q=0 AND P=-1 GOTO 50
210
60150 IF A$=";" AND Q=-1 GOTO 60190
60160 IF A$<> "M" GOTO 60190
60170 IF MID$(B$,L-1,2)<>"RE" GOTO 60190
60180 B$=LEFT$(B$,L-2)+":":RE=1:GOTO 600
90
60190 B$=B$+A$:L=L+1:GOTO 60090
60200 PRINTB$:RETURN
60210 GET#2,C$:IF C$=":" THEN A$=A$+C$:L
=L+1:P=0:GOTO 60190
60220 IF C$=CHR$(13) THEN B$=B$+A$:GOTO
60200
60230 IFRIGHT$(B$,1)="$"ORRIGHT$(B$,1)=C
HR$(34)ORC$=CHR$(34)THENA$=C$:GOTO60100
30240 A$=A$+C$:L=L+1:GOTO 30190
```

Conclusion

We hope that this chapter has given you food for thought. By way of a project, why not write a routine to recover as much of the data as possible after an edit or NEW has been performed?

6 Routines old and new

Introduction

In Chapter 4 we gave listings in machine code to make use of the function keys. These are actioned by keywords. At the present time, BASIC will not understand these. All the functions of the UTILITY, the remainder of which are in the following two chapters, require some sort of 'driving mechanism'. That is, routines which will not only recognize the keywords, but will action them. Those routines are the PRINT tokens, DISPATCH BASIC CHARS and BASIC EVALUATION. In Chapter 3 these were fully discussed, so we are only supplying in this chapter the coding that is particular to the UTILITY.

To initialize the UTILITY we need to change the addresses in certain locations. These fall into three categories. First, we have to change the vector addresses so that BASIC will go to our token routines; secondly, we need to protect the UTILITY from being overwritten by programs and strings; and lastly we need to retain its operation during a Non-Maskable-Interrupt, that is when RUN/STOP RESTORE is pressed.

There are certain subroutines which will be used by more than one command, so we include them in this chapter. These deal with getting parameters, the switching in and out of the BASIC ROM and memory moving.

That has dealt with the new, and now for the old. A few of the resident ROM routines are useful. Many of them will be covered when describing our new commands. The later part of this chapter describes some more.

Initialization

When you start up the UTILITY with SYS 32768 these instructions will be the first to be actioned. They will set up and protect the UTILITY. At the end of the four subroutines we return control back to you, with a screen message, and the UTILITY in operation.

ASSEMBLY LISTING

0 x-++0000

7	x			
10		JSR VECTOR	? !	CHANGE BASIC VECTORS
20		JSR KEYS	ļ	SET KEYBOARD VECTOR
30		JSR NMI	!	SET NMI AND BRK VECTORS

40		JSR	MEM	! SET TOP OF MEMORY
50		JMP	\$9200	! CLR AND MESSAGE
60	VECTO	r lda	#\$09	
70		STA	\$0304	! ICRNCH LOW
80		LDA	#\$BC	
90		STA	\$0306	! IQPLOP LOW
100		LDA	#\$02	
110		STA	\$0308	! IGONE LOW
120		LDA	#\$29	
130		STA	\$030A	! IEVAL LOW
140		LDA	#\$82	
150		STA	\$0305	! ICRNCH HIGH
160		STA	\$0307	! IQPLOP HIGH
170		LDA	#\$83	
180		STA	\$0309	! IGONE HIGH
190		STA	\$030B	! IEVAL LOW
200		RTS		
210	MEM	LDA	#\$FF	
220		STA	\$37	! MEMSIZ LOW
230		STA	\$33	! FRETOP LOW
240		LDA	#\$7F	
250		STA	\$38	! MEMSIZ HIGH
260		STA	\$34	! FRETOP HIGH
270		RTS		
280	NMI	LDA	#\$7E	
290		STA	\$0316	! BRK LOW
300		LDA	#\$61	
310		STA	\$0318	! NMI LOW
320		LDA	#\$80	
330		STA	\$0317	! BRK HIGH
340		STA	\$0319	!NMI HIGH
350		RTS		
360	KEYS	SEI		
370		LDA	#\$22	
380		STA	\$028F	! KEYLOG LOW
390		LDA	#\$87	
400		STA	\$0290	! KEYLOG HIGH
410		CLI		
420		RTS		
425	! NMI	ROUTINE		
430		Pha		
440		TXA		
450		Pha		
460		TYA		
470		Pha		
480		LDA	#\$7F	

490		STA	\$DD0	D !	CIA INTERRUPT
					CONTROL REG
500		LDY	\$DD0	D	
510		BPL	PLUS	3	
520		JMP	\$FE7	'2 !	RESET RS232
530	PLUS	JSR	\$F6E	3C !	SETS STOP AND RVS FLAG
540		JSR	\$FFE	1 !	CHECK STOP KEY
550		BEQ	BREA	чκ	
560		JMP	\$FE7	72 !	RESET RS232
565	! BRK	ROUTINE			
570	BREAK	JSR	\$FD1	5 !	KERNAL RESET
580		JSR	\$FD4	43 !	INITIALIZE I/O
000		00.0			CIA CHIPS
590		JSR	\$E51	18 !	INITIALIZE 1/0
400		JSR	KEYS	3	
A19		JSR	NMT	-	
470		.TMP	(\$40	192)	
020		0111	\ + / 1C		
807E	BREAK		8054	KEYS	
8034	MEM		8041	NMI	
8073	PLUS		800F	VECT	OR

We feel that this listing up to 430 is fairly self-explanatory, especially with a memory map. The remaining lines are dealt with in the next section.

BRK and NMI routines

These are included in the listing of the previous section, lines 430 to 620. When either of these are initiated, it will be to these lines they will come. The majority of these routines are copies of the equivalent ROM routines, plus a couple of directions to our set up routines to keep the UTILITY in service.

NMI

The NMI is initiated by the use of the RESTORE key (although there are means to initiate it through the cartridge slot). Not only does it tell the microprocessor it has been actioned, but it also sets a flag in the CIA #2. The processor will not action it immediately, but will wait until the present instruction is complete. The processor then saves the program counter and the status register on the stack. It will load the address stored at SFFFA and SFFFB into the program counter. This is normally SFF43. At this address it sets the interrupt flag, so that the other interrupt does not interfere, and then jumps to the vector address that we have changed. Note that the routine has so far not stored the A, X and Y registers.

The NMI in the UTILITY will end up at our routine, which is a series of subroutines. After saving the processor registers on the stack, it clears the NMI flag in the Interrupt Control Register of the CIA#2 chip, which deals with inputs and outputs of the computer. It then loads that location back into Y and if the NMI flag is still clear then it jumps ahead, missing out for the time being an RS232 reset. The following routine checks the STOP and RVS flags at location \$91. A call to the KERNAL routine to check for STOP follows. If on exit the accumulator is zero, then the STOP was initiated and we go to the BRK routine. Finally, we jump to the routine to reset the RS232 locations.

BRK

The first subroutine resets the KERNAL set up vector from \$0314 to \$0333 to their default values from a list held in the KERNAL ROM itself. This will reset two we have changed, the BRK and NMI vectors. The following routine will service the two CIA interface chips, by restoring them to their setup levels.

The routine at \$E518 performs the remaining functions of a BRK. It restores the output device to the screen and the input device to the keyboard. The video chip is next for the restoration treatment. The screen and character set are returned to their default positions, and sprite graphics turned off. After this it is the keyboard's turn, with the buffer, delays and set-up vector all returned to default values. The routine finishes off by resetting the input/output flags, clearing the screen, setting the colours and putting the cursor in the home position.

We now put in two calls ourselves so we can reset the NMI, BRK and keyboard vectors to those we require. Finally, there is the indirect jump of A002 which sets the stack pointer to its start, prints 'READY' and gives control back to the user.

Routine vectors and keywords

There is sufficient space, using the existing token system, for 51 further keywords. These will be split up into an area for command keywords and an area for function keywords. In the UTILITY we are supplying 34 commands and 1 function. Between the last command keyword vector and that of the function keyword there is space for a further nine commands (token values 238 to 246 (SEE-SF6)). Seven extra functions could be added within the space available. The vector table is positioned at \$8090 to \$80F5.

The keyword table is exactly 255 bytes long. Out of that our keywords use up 155 plus a zero byte to mark the end of the table. The amount of space available to you if wish to extend it is 58 bytes for command keywords and 41 bytes for functions. Remember that the last letter of each word has \$80 (128) added to it. In our table, the space

between our last command, TROFF, and the only function, DEEK, has been filled with bytes \$5A and \$EA to make up the nine unused token values.

MEMORY DUMP

.:8090	98	87	4C	86	B2	83	9F	84′
.:8098	EB	84	36	85	BE	85	14	84′
.:80A0	51	83	A6	83	AE	8F	Β4	8F′
.:80A8	80	83	AC	83	51	8E	C4	89′
.:80B0	43	8F	A6	87	92	8B	2D	84′
.:8088	D1	8F	ЗA	A9	D1	A8	30	86′
.:80C0	B6	91	39	8D	10	86	B5	92′
.:80C8	8C	91	80	91	4D	90	FB	85′
.:80D0	бE	88	60	8D	FF	FF	FF	FF′
.:80D8	00	FF	FF	FF	F6	FF	Bб	F7′
.:80E0	00	60	00	00	D6	83	D6	83′
.:80E8	00	00	00	68	00	00	00	40′
.:80F0	00	00	00	40	00	40	4F	46′OF
.:80F8	C6	4B	45	D9	44	4F	4B	C51fKEyDOKe
.:8100	54	45	CE	54	57	CF	48	45'TEnTWoHE
.:8108	D8	42	49	CE	4F	4C	C4	43′xBInOLdC
.:8110	4F	4C	4F	55	D2	57	52	49'OLOUrWRI
.:8118	54	C5	43	47	4F	54	CF	43'TeCGOToC
.:8120	47	4F	53	55	C2	50	4C	4F1GOSUBPLO
.:8128	D4	45	4E	54	45	D2	44	55'tENTErDU
.:8130	4D	D0	52	45	4E	55	CD	44'MpRENUmD
.:8138	45	4C	45	54	C5	4D	45	52'ELETeMER
.:8140	47	C5	43	4F	44	45	D2	41'GeCODErA
.:8148	55	54	CF	50	52	4F	СЗ	44'UToPROcD
.:8150	50	52	4F	СЗ	45	50	52	4F'PROcEPRO
.:8158	СЗ	50	4F	D0	51	55	49	D4′cPOpQUIt
.:8160	54	52	41	43	C5	52	45	53'TRACeRES
.:8168	45	D4	43	48	41	49	CE	4C′EtCHAInL
.:8170	4F	4D	45	CD	48	49	4D	45'OMEmHIME
.:8178	CD	49	4E	4B	45	59	A4	4D'mINKEY M
.:8180	45	CD	41	50	50	45	4E	C4'EmAPPENd
.:8188	54	52	4F	46	C6	5A	5A	5A'TROFfZZZ
.:8190	5A	5A	EA	5A	5A	5A	5A	5A'ZZ.ZZZZZ
.:8198	5A	EΑ	5A	5A	5A	5A	5A	5A'Z.ZZZZZZ
.:81A0	5A	EΑ	5A	5A	5A	5A	EΑ	5A'Z.ZZZZ.Z
.:81A8	5A	5A	5A	EΑ	5A	5A	5A	5A'ZZZ.ZZZZ
.:8180	EΑ	5A	5A	5A	5A	5A	5A	EA'.ZZZZZZ.
.: 81B8	5A	5A	5A	5A	5A	EΑ	5A	5A'ZZZZZ.ZZ
.:81C0	5A	5A	5A	5A	5A	5A	EΑ	44'ZZZZZZ.D
.:81C8	45	45	СВ	00	FF	FF	FF	FF'EEk

.:81D0	FF	FF	FF	FF	FF	FF	FD	FF'
.:81D8	FF	FF	FF	FF	FF	7F	FF	FF'
.:81E0	00	00	00	00	00	00	00	001
.:81E8	00	00	00	08	01	00	00	001
.:81F0	00	00	00	00	00	20	8A	AD'

This has been produced in upper case mode and as such the end shifted letter of each command is printed in lower case. If putting it into your computer in a way other than the dump, remember that they are shifted. The last letter in location \$817E is a shifted \$, giving the keyword INKEY\$.

Getting parameters and controlling BASIC

```
ASSEMBLY LISTING
```

9 *=\$ 81F5			
10	JSR \$AD8A	!	GET INPUT
20	JMP \$B7F7	!	CHECK AND TRANSFER
30	LDA \$01	ļ	6510 I/O PORT
40	AND #\$FE	!	TURN OFF BIT 0
50	STA \$01	ļ	BASIC OFF
60	RTS		
70	LDA \$01		
80	ORA #\$01	!	SET BIT 0
90	STA \$01	!	BASIC ON
100	RTS		

Parameters

Lines 10 and 20 hold the only two instructions that we need to incorporate, but they do a lot of work in getting our numeric parameters. Let us look at the instructions one at a time.

JSR \$AD8A

The first action of this is to call the evaluate expression routine at \$AD9E.

This is a complex routine which deals not only with numeric data, but also with strings. After setting the CHRGET pointer back one place, it proceeds to start picking up data after the command keyword. It will then go through checking to see whether a mathematical operator or a function keyword (such as PEEK), a variable or simply a number has been obtained. From the information obtained it will (after calculating if necessary) store the result or findings in the FAC#1. For numbers up to \$FFFFF, the relevant numbers will be in locations \$64 and \$65 of this accumulator.

We now return to our original subroutine at \$AD8A, where we check to see if the data received was numeric or not. The evaluate expression will set a flag in the zero page location sob. The value of \$FF indicates string data, whilst zero designates numeric data. If this subroutine finds the former, a 'TYPE MISMATCH' error is generated and the command, and program, is terminated.

JMP \$B7F7

We have our numeric parameter. This routine will do two checks and then transfer our data. The checks are to make sure that neither a negative number nor one over 65535 (\$FFFF) was given. In either case, failure will result in the 'ILLEGAL QUANTITY' error. The data is now transferred from \$64 and \$65 to locations \$14 and \$15. The reason for this is that the FAC#1 is used for many applications. The RTS at the end of this routine will return us to the place that called our complete GET PARA-METERS routine, that will most likely be a command routine.

The BASIC switch

As we said, when dealing with the function keys, the area of RAM under the BASIC ROM is a useful place for hiding data, or indeed routines which do not use the BASIC interpreter. To use this area, BASIC must be 'removed'. We have no trouble writing to the RAM as the computer, through its decoding logic, will select it when the processor sends a write signal. When reading, the ROM has priority unless we tell the electronics that it is not there. The main difference between the 6510 processor in the 64 and the normal 6502 is that the former has input/ output ports. The user can control these using locations \$0000 and \$0001. The first deals with the direction of the data, that is, whether the ports, of which there are six, are going to be input or output. The second location deals with the data itself, one bit for each port, either a one or a zero which gives a switching mode. Three of the ports are connected to the cassette port. The other three control three ROMS: BASIC, KERNAL and the Character ROM. A zero will switch all of these off. The one we are concerned with, BASIC, uses bit Ø of the data register and so by changing this bit, making sure not to disturb the others, we can remove or replace as required.

Lines 30 to 60 perform the switching-out of BASIC. We load the register and set bit 0 to zero. The AND instruction will do this without changing any other bit. After placing the result back, the ROM is no longer present as far as the computer is concerned.

Lines 70 to 100 reverse the process by using the ORA code which will only affect the bits according to the data with the instruction.

To switch off BASIC – JSR \$81FB To switch in BASIC – JSR \$8202

Dealing with the keywords

In Chapter 3 the routines that BASIC uses to deal with keywords and

tokens were fully described. Below are the listings to use with the UTILITY, which require no further explanation.

ASSEMBLY LISTING – CRUNCH TOKENS

_				1			
. 9	* =\$ 8209			410	STORE	INX	
10		LDX	(\$7A	420	l	INY	
20		LDY	′#\$04	430		STA	\$01FB,Y
30		STY	′\$0F	440		LDA	\$01FB,Y
40	ANOTHER	LDA	¥0200,X	450		BEQ	EXIT
50		BPL	. SPACE	460		SEC	
60		CMF	'#\$FF	470	1	SBC	#\$3A
70		BEG	STORE	480		BEQ	COLON
80		INX		490		CMP	#\$49
90		BNE	ANOTHER	500		BNE	DATA
100	SPACE	CMF	* #\$20	510	COLON	STA	\$0F
110		BNE	STORE	520	DATA	SEC	
120		STA	\$08	530		SBC	#\$55
130		CMP	#\$22	540		BNE	ANOTHER
140		BEQ	QUOTE	550		STA	\$08
150		BIT	\$0F	560		LDA	\$0200.X
160		BVS	STORE	570	LINE	BEQ	STORE
170		CMP	#\$3F	580		CMP	\$08
180		BNE	NUMBER	590		BEQ	STORE
190		LDA	#\$99	600	QUOTE	INY	
200		BNE	STORE	610		STA	\$01FB.Y
210	NUMBER	CMP	#\$30	620		INX	,
220		BCC	CONT	630		BNE	LINE
230		CMP	#\$3C	640	NEXTWORD	LDX	\$7A
240	00 T	BCC	STORE	350		INC	\$0B
200	LUNI	SIY	\$71	660	FIND	INY	
200		LDY	#\$00	670		LDA	\$A09D,Y
270		511	20B	680		BPL	FIND
200		DET	* 7 •	690		LDA	\$A09E,Y
270			⊅ 7A	700		BNE	CONT1
210				710		LDY	#\$FF
220	NEATLETTER			720		DEX	
220	CONTR		+0000 V	730	NEXT	INY	
240	CONTI	CDH	₩0200,X	740		INX	
250		SEC		750	NEXTB	LDA	\$0200,X
330 340		380	PHUYE, I	760		SEC	
370			NEATLETTEK	770		SBC	\$80F6,Y
370 200			# ₽ ႣØ NEVTLIODD	780		BEQ	NEXT
200	STOPEA	DINE		790		CMP	#\$80
370 100	STURCH		₽05 471	800		BNE	NEXTNEW
400	FUUND	LUT	⊅71 ·	810		BEQ	STOREA

870	LDA	\$80F6,Y	820	NEXTNEW	LDX	\$7A
880	BNE	NEXTB	830		INC	\$0B
890	LDA	\$0200,X	840	NEXTA	INY	
900	BPL	FOUND	850		LDA	\$80F5,Y
910 EXIT	JMP	\$A609	860		BPL	NEXTA
820F ANOTHER		8269 CO	LON			
8239 CONT		8245 CO	NT1			
826B DATA		82B9 EX	IT			
8286 FIND		8254 FO	UND			
8275 LINE		8294 NE	XT			
82A9 NEXTA		8296 NE	ХТВ			
8243 NEXTLETTER		82A5 NE	XTNEW			
8282 NEXTWORD		8231 NU	MBER			
827B QUOTE		8218 SP	ACE			
8256 STORE		8252 ST	OREA			

ASSEMBLY LISTING – PRINT TOKENS

9	*=\$82BC	
10		BPL ROM1
20		CMP #\$FF
30		BEQ ROM1
40		BIT \$0F
50		BMI ROM1
60		CMP #\$CC
70		BCC CBMTOKEN
80		SEC
98		SBC #\$CB
100		TAX
110		LDA #\$F6
120		STA \$22
130		LDA #\$80
140		STA \$23
150		BNE START
160	CBMTOKEN	SEC
170		SBC #\$7F
180		TAX
190		LDA #\$9E
200		STA \$22
210		LDA #\$A0
220		STA \$23
230	START	STY \$49
240		LDY #\$FF
250	NEXTWORD	DEX

80.

260		BEQ	WORDFOUND
270	NEXTCHAR	INY	-
280		LDA	(\$22),Y
290		BPL	NEXTCHAR
300		BMI	NEXTWORD
310	WORDFOUND	INY	
320		LDA	(\$2 2),Y
330		BMI	R0M2
340		JSR	\$AB47
350		BNE	WORDFOUND
360	ROM1	JMP	\$A6F3
370	ROM2	JMP	\$A6EF

82D8	CBMTOKEN	82EB	NEXTCHAR
82E8	NEXTWORD	82EC	
82FF 82F2	ROM2 WORDFOUND	82E4	START

ASSEMBLY LISTING – DISPATCH AND EVALUATION

- 9	*=\$ 8302		
10		JSR	\$0073
20		CMP	#\$CC
30		BCC	ROM3
40		CMP	#\$EE
50		BCS	ROM3
60		JSR	DISPATCH
70		JMP	\$A7EA
80	DISPATCH	SEC	
90		SBC	#\$CC
100		ASL	A
110		TAY	
120		LDA	8091,Y
130		Pha	
140		LDA	\$8090,Y
150		Pha	
160		JMP	\$0073
170	ROM3	JSR	\$0079
180		JMP	\$A7E7
190		LDA	#\$90
200		STA	\$0D
210		JSR	\$0073
220		CMP	#\$F7
230		BCC	ROM4
240		CMP	#\$F8

130 Routines old and new

	BCS ROM4
	JSR DISPATCH1
	RTS
DISPATCHI	SEC
	SBC #\$Fð
	ASL A
	TAY
	LDA \$80E5,Y
	Pha
	LDA \$80E4,Y
	PHA
	JMP \$0073
ROM4	JSR \$0079
	JMP \$AE8D
DISPATCH	833C DISPATCH1
ROM3	834C ROM4
	DISPATCH1 ROM4 DISPATCH ROM3

The start up message

This is the final subroutine called during the initialization of the UTILITY. It performs a CLR, to set all the variable addresses, changes the screen and text colours, and finally puts a message on the screen indicating that the UTILITY is in operation.

ASSEMBLY LISTING

- 9	*=\$9200			
10		JSR \$A663	!	CLR
20		LDA #\$93	ļ	CLEAR SCREEN
30		JSR \$FFD2		
40		LDA #\$00	!	SET COLOURS TO BLACK
50		STA \$D020	ļ	BORDER
60		STA \$D021	!	BACKGROUND
70		LDA #\$05		
80		STA 0286	!	GREEN TEXT
90		LDX #\$0A		
100		LDY #\$09		
110		JSR STARS		
120		LDX #\$0C		
130		LDY #\$0 9		
140		CLC		

Routines old and new 131

150		JSR	\$FFF0	ļ	SET CURSOR
160		LDX	#\$15	I	CHARS TO PRINT
170	CONT	LDA	DATA,X		
180		JSR	\$FFD2	I	PRINT
190		DEX			
200		BPL	CONT	1	NOT FINISHED
210		LDX	#\$0E		
220		LDY	#\$09		
230		JSR	STARS		
240		STA	\$05C1	ļ	FILL IN MISSING CHARS
250		STA	\$0611		
260		STA	\$05D6		
270		STA	\$0626		
280		LDA	#\$05	!	COLOUR MAP VALUE
290		STA	\$D9C1		
300		STA	\$DA11		
310		STA	\$D9D6		
320		STA	\$DA26		
330		LDA	#\$0D	!	PRINT RETURN
340		JSR	\$FFD2		
350		JMP	\$A474	!	READY FOR BASIC
360	STARS	CLC			
370		JSR	\$FFF0	ļ	SET CURSOR
380		LDA	#\$2A	ļ	ASCII FOR *
390		LDX	#\$16	!	NUMBER TO PRINT
400	NEXT	JSR	\$FFD2		
410		DEX			
420		BNE	NEXT		
430		RTS			
440	DATA	тхт	** YTILIT	U	CISAB NAP **
	COMIT				

9226	CONT	9267	DATA
9260	NEXT	9258	STARS

Memory moving

RENUMBER and CODER, described in Chapter 7, both require some manipulation of memory in the form of either gaining space or removing unnecessary bytes. This section deals with the two subroutines, CLOSE and OPEN, which perform these operations. CLOSE is self-contained whilst OPEN uses a ROM routine for the actual moving of memory. In the BASIC interpreter there are routines to both open and close up a BASIC program, used when you insert or delete lines, but we can only really use the opening routine. It is a subroutine on its own whereas the closing-up is integral with the inputting of a BASIC line. We have written coding that is virtually identical to the one in ROM as it is efficient enough.

Having moved the program about, all the link addresses, from the line the move started, will now be wrong by the amount of the move. There is a subroutine in the interpreter which changes the link addresses but we have not used it. The reason for this is one of speed as during the course of using CODER or RENUM, these subroutines may be called several times and would prove to be very slow.

The ROM routine for rechaining the lines goes through the whole program, byte by byte, to calculate the link addresses and store them. It has been done this way as it is a multi-purpose routine, catering for the lengthening and shortening of code. What we have done is to write separate routines for each direction of movement and place them immediately after the moving instructions. These will only rechain from the program line in which the alteration occurred. In addition, we only need to look at the link addressess as we know by how much they have changed so we can subtract or add as required.

To set the scene, as they say, here are the locations that need to be set before calling these subroutines:

\$FB and \$FC-	The address of the start of the current BASIC line.
\$49-	The number of the current position on that line. This will
	be where the replacement code will start.
\$FD and \$FE-	The address of the next BASIC line, that is, the link address of the line in \$FB and \$FC.
\$3E-	The number of bytes in the original code to changed.
register-	The number of bytes in the replacement code.

ASSEMBLY LISTING

9	*=\$ 888B				
10		STX	\$C2		
20		LDA	\$3E	!	FIND HOW MANY BYTES TO REMOVE
30		SEC			
40		SBC	\$02		
50		STA	\$BB		
60		CLC			
70		LDA	\$FB		
80		ADC	\$49	ļ	FIND START OF BLOCK TO MOVE
90		STA	\$5F		
100		LDA	\$FC		
110		ADC	#\$00		
120		STA	\$60		
130		LDA	\$5F		

140		ADC \$BB		
150		STA \$5A	! START + AM	OUNT OF
			REDUCTION	
160		LDA \$60		
170		ADC #\$00		
180		STA \$5B		
190		LDA \$2D	! END OF PRO	G
200		SEC		-
210		SBC \$5A	! CALCULATE	TOTAL
			AMOUNT TO	MOVE
220		STA \$58		
230		TAY	! NO OF BYTE	S OF
			INCOMPLETE	PAGE
240		LDA \$2E		
250		SBC \$5B		
260		TAX	! NO OF PAGE	ѕ то
			MOVE	
270		INX	! FOR EASIER	CHECKING
280		TYA		
290		BEQ PAGE	! NO SEPARAT	E BYTES
300		LDA \$5A	! MOVE SEPAR	ATE
			BYTES FIRS	Г
310		CLC		
320		ADC \$58		
330		STA \$5A		
340		BCC NOINC		
350		INC \$5B		
360		CLC		
370	NOINC	LDA \$5F		
380		ADC \$58		
390		STA \$5F		
400		BCC NOINCA		
410	NOTION	INC \$69		
420	NUINLA	11A 500 ##55		
430		EUR HAFF		
440				
430				
400		DEC \$35		
470	PAGE	UEL 700 IDA (45A) V		
100	PHUE	CTA (#CH),T		
770		JIH (DOF),T		
510		DNE DAGE		
570		TNC 450		
520		TING & ZA		
5/0		114L \$00		
540			. FUINIER - I	

550		BNE	PAGE		
560		SEC			
570		LDA	\$20	!	SET END OF PROG
580		SBC	\$BB		
590		STA	\$2D		
600		BCS	RECHAIN		
610		DEC	\$2E		
620		SEC			
630	RECHAIN	LDY	#\$00		
640		LDA	\$FD	ļ	GET LINK
650		SBC	\$BB	ł	CALC NEW ADDRESS
660		STA	\$FD		
670		STA	(\$FB).Y	1	STORE IN LINE
680		STA	\$57		
390		I DA	\$FF		
788		SRC	#\$00		
710		INY			
728		STA	 \$FF		
720		STA	458		
7/0		CTA	(4EP) V		
750			(#10),1		
730	MEATI		(457) V	ı	GET LINKS
700		CTA	(⊅J/)ji #D0	:	CTOPE THEM
770			PD7	:	STURE THEN
700			(#57) V		
770		CTA	(\$077,1 #DA		
000		DEO		ŧ	COMPLETED DECUAINING
810		BEW	EXII	:	COMPLETED RECHAINING
020					
040		JEL	* D0		
050		LUH	₽ 6 7 #DD		
830		580	⊅ ₽₿	:	LALL NEW LINK ADDS
800		IAX		:	IEMP STURE
870		SIA	(\$57),Y		
880		LDA	\$BA		
890		SBC	#\$00		
900		INY			
910		STA	(\$57),Y		
920		STA	\$58		
930		TXA			
940		STA	\$57		
950		JMP	NEXT1	ļ	GET NEXT LINE
960	EXIT	RTS			
970		TXA			
980		SEC			
990		SBC	\$3E	ļ	CALCULATE NO OF
					SPACES REQUIRED

1000		STA	\$BB		
1010		CLC			
1020		LDA	\$49		
1030		ADC	\$BB		
1040		BCS	ERROR	!	>255 CHARS IN LINE
1050		CMP	\$FE		
1030		BCC	CONT	ļ	ONLY 254 ALLOWED
					-NEED END MARKER
1070	ERROR	LDX	#\$17		
1080		JMP	\$A437	ļ	ERROR STRING TOO LONG
1090	CONT	LDA	\$2D		
1100		ADC	\$BB	!	ENOUGH MEMORY?
1110		TAX			
1120		LDA	\$2E		
1130		ADC	#\$00		
1140		CMP	\$38		
1150		BNE	CONT2	ļ	ENOUGH MEMORY
1160		CPX	\$37		
1170		BCC	CONT2		
1180		JMP	\$A435	ļ	ERROR OUT OF MEMORY
1190	CONT2	CLC			
1200		LDA	\$2D	ļ	SET ADDS FOR MOVE
1210		STA	\$5A		
1220		ADC	\$BB		
1230		STA	\$58		
1240		LDA	\$2E		
1250		STA	\$5B		
1260		ADC	#\$00		
1270		STA	\$59		
1280		LDA	\$FB		
1290		ADC	\$49		
1300		STA	\$5F		
1310		LDA	\$FC		
1320		ADC	#\$00		
1330		STA	\$60		
1340		JSR	\$A3BF	!	ROM ROUTINE TO OPEN UP MEMORY
1350		CLC			
1360		LDY	#\$00		
1370		LDA	\$2D	!	SET NEW END OF PROG
1380		ADC	\$BB		
1390		STA	\$2D		
1400		BCC	CONT3		
1410		INC	\$2E		
1420		CLC			
1430	CONT3	LDA	\$FD		

136 Routines old and new

1440		ADC \$BB	
1450		STA \$FD	
1460		STA \$57	
1470		STA (\$FB),Y	
1480		LDA \$FE	
1490		ADC #\$00	
1500		INY	
1510		STA \$FE	
1520		STA \$58	
1530		STA (\$FB),Y	
1540	NEXT3	DEY	
1550		LDA (\$57),Y	
1560		STA \$B9	
1570		INY	
1580		LDA (\$57),Y	
1590		STA \$BA	
1600		BEQ EXIT2	
1610		DEY	
1620		CLC	
1630		LDA \$B9	
1640		ADC \$BB	
1650		TAX	
1660		STA (\$57),Y	
1670		LDA \$BA	
1680		ADC #\$00	
1690		INY	
1700		STA (\$57),Y	
1710		STA \$58	
1720			
1730		SIA \$57	
1740		JMP NEX13	
1750	EXIIZ	RIS	
CONT	8949	CONT2	895D
CONTS	898B	ERROR	8944
EXIT	892F	EXIT2	8904
NEXT 1	890E	NEXT3	89A0
NOINC	: 88CA	NOINCA	88D4
PAGE	88DD	RECHAIN	88F7

CLOSE ROUTINE

LINES 10–270: Before we can move a block of memory, we have to determine three values: the start address of the block to move, the new start address and the amount of code to move. The first thing we

work out is the number of redundant bytes. This is done, obviously, by subtracting from the original amount of data to be changed the number of bytes of the replacement code. The resultant value is stored in location \$BB. We shall need this number later for rechaining the lines. The new start of the block will be obtained by adding the line pointer, \$49, to the address of the current BASIC line. To this value is added the contents of \$BB which will give us the location of the first byte in the block to be moved.

To get the amount of data to be moved, the result of the last calculation is taken away from the end of program address, held in \$2D and \$2E. The answer will be held in the processor registers, the high byte in the x and the low in the y. A page of memory is 256 bytes so the x register is therefore the number of pages to be moved, increased by one for easier checking on completion. We move a complete page and then decrease x. x will be zero when all done, checking immediately after decreasing. To summarize, we have found the amount to move, its current start and its destination.

LINES 280–470: This is the hardest part of the routine to follow, and we hope that we succeed in explaining it clearly.

We transfer the v register to the accumulator. To recap briefly, this will be the number of bytes, other than complete pages, of memory to move. If the value now in the accumulator is zero, only complete pages require moving, so we skip this section completely. In closing up memory we start from the low addresses, move them, and work to the higher end addresses. We do this by setting the address of the page and moving it up, using the v register as a pointer. If we have an odd number of bytes to start with, this causes a slight problem. For example, if we have \$10 bytes and the v is set thus we would move 246 bytes by increasing v. To compensate for this, what we do is to produce the 2's complement of the value. This is done in lines 430 to 450. The EOR #\$FF will change all the bits set to one to zero and vice versa. One is then added. So instead of \$10, we should now have \$F0. This means that if we now increase v from \$F0 until it becomes zero it will have been incremented \$10 times.

For the same reasons we have to alter the address of the start of the block and its new start address. We add to these the original number of odd bytes, held in \$58. Finally, we decrease the high byte of the address by one. The next effect of these changes is a stalemate as the locations along with the y pointer value are equivalent to the original values but now allow us to increase y the required amount.

LINES 480–550: Having set all the values we move the data, byte by byte, until both x and γ registers are zero. We simply load a byte from its position and store its new lower location.

LINES 560-620: The end of the BASIC program will now be shorter by

the value of location \$BB. The original end address is adjusted and reset.

LINES 630–960: All that remains is to change the values of the link addresses from the current BASIC line onwards. First, we change the links in the current line and as these are also held in \$FD and \$FE, used by the calling routine, we change these also.

We proceed through the lines gathering the addresses, subtracting the value in \$BB, and then we restore them. The end of the program is indicated when the MSB of a link address is zero. Finally, we return to the calling program, such as CODER.

OPEN ROUTINE

LINES 970–1080: We calculate the space required by subtracting the value in s_{3E}, the length of the old code, from the value in the x register, the length of the new code, and store the result in s_{BB}.

As a BASIC line may not exceed 255 bytes (to allow for a zero at the end making a maximum of 256), we check this by adding the line marker to the SBB value. A set carry flag will mean the maximum has been exceeded. We then check that there will be room for the end of line zero. Failure of either of these will generate the syntax error 'STRING TOO LONG'.

LINES 1090–1180: As we are creating space we must check that there is sufficient room available in the BASIC program area. These lines do just that by checking that we will not exceed the values in \$37 and \$38, which indicate its limit. If we do go over, we call a BASIC routine to generate the 'OUT OF MEMORY' error message.

LINES 1190–1340: Next on the agenda is to set the registers for the interpreter's OPEN-up memory routine at \$A3BF. On leaving this routine: \$5A and \$5B – This will hold the address of the end of the block to move. It will be the same as the end of program address before the move. \$58 and \$59 – These registers will hold the address of the end of the new block. It will also be the end of the BASIC program after the move. It is arrived at by adding the amount of move to the address in \$5A and \$5B. \$5F and \$60 – The start of the block to move. These hold the location of the first byte of the code to be changed. It is calculated by adding the line marker to the address of the current BASIC line to be processed.

LINES 1360–1420: Now that the data has been moved, we reset the end of program address to its new value.

LINES 1430–1750: A replica of the rechaining in lines 630 to 960, except that here we increase the addresses instead of reducing them.

This concludes the new routines that we planned to introduce in this

chapter. The remainder are descriptions of some of the ROM routines we use (and hope that you will come to use).

RECHAINING THE LINES

During our memory move routine, we did not use the ROM routine to rechain the link addresses because for our purposes it was inefficient due to the number of calls required. However, we do use the sub-routine, in DELETE for instance, where only one call is required. It serves another purpose in that from the addresses it exits with, one can calculate and set the end of program/start of variable registers.

ROM LISTING

A533	A5	2B	LDA	\$2B
A535	A4	2C	LDY	\$20
A537	85	22	STA	\$22
A539	84	23	STY	\$23
A53B	18		CLC	
A53C	A0	01	LDY	#\$01
A53E	B1	22	LDA	(\$22),Y
A540	F0	1 D	BEQ	\$A55F
A542	A0	94	LDY	#\$04
A544	С8		INY	
A545	B1	22	LDA	(\$22),Y
A547	D0	FB	BNE	\$A544
A549	C8		INY	
A54A	98		TYA	
A54B	65	22	ADC	\$22
A54D	AA		TAX	
A54E	A0	00	LDY	#\$00
A550	91	22	STA	(\$22),Y
A552	A5	23	LDA	\$23
A554	69	80	ADC	#\$00
A556	C8		INY	
A557	91	22	STA	(\$22),Y
A559	86	22	STX	\$22
A55B	85	23	STA	\$23
A55D	90	DD	BCC	\$A53C
A55F	60		RTS	

The routine commences by getting the program start address and placing it in registers for its own use. The carry flag is cleared for addition. The first byte of a line that it picks up is the high byte of the link address and it tests for the end of the program (a zero). The y register is loaded again so as to skip the addresses and line number. It

now proceeds through the line, searching for the end of line zero marker. When this is discovered, the y register will contain one less than the number of bytes in the complete line. This is immediately rectified by incrementing y by one. This value is added to the line start address and placed as the link address of the line. As this is also the address of the next line, it is loaded into the locations used by the routine. The flow now branches back, (the carry flag will be clear), to process the next line. Every BASIC line will be processed until the end of the program.

On exiting, the program locations \$22 and \$23 will hold the address of the two end zero bytes. If this address is increased by two then the end of program address can be derived, and hence the start of variables, as they are one and the same thing VARTAB.

Opening up memory

In our memory move routine we made use of a ROM routine when we required more space in a BASIC program. It will move a block up in memory even if its new start is within the original block. Six locations have to be set before entering the routine, which are, in low/high byte order:

\$5A and \$5B-	End address of present block
\$5F and \$6 0 –	Start address of present block
\$58 and \$59–	End address of the new block

ROM LISTING

A3BF	38		SEC
A3C0	A5	5A	LDA \$5A
A3C2	E5	5F	SBC \$5F
A3C4	85	22	STA \$22
A3C6	A8		TAY
A3C7	A5	5B	LDA \$5B
A3C9	E5	60	SBC \$60
A3CB	AA		TAX
A3CC	E8		INX
A3CD	98		TYA
A3CE	F0	23	BEQ \$A3F3
A3D0	A5	5A	LDA \$5A
A3D2	38		SEC
A3D3	E5	22	SBC \$22
A3D5	85	5A	STA \$5A
A3D7	BØ	03	BCS \$A3DC
A3D9	C6	5B	DEC \$58
A3DB	38		SEC

A5	58	LDA \$58
E5	22	SBC \$22
85	58	STA \$58
BØ	08	BCS \$A3EC
C6	59	DEC \$59
90	04	BCC \$A3EC
B1	5A	LDA (\$5A),Y
91	58	STA (\$58),Y
88		DEY
D0	F9	BNE \$A3E8
B1	5A	LDA (\$5A),Y
91	58	STA (\$58),Y
C6	5B	DEC \$5B
ርሪ	59	DEC \$59
CA		DEX
DØ	F2	BNE \$A3EC
60		RTS
	A5 E5 B0 C6 90 B1 91 88 D0 B1 91 C6 CA D0 60	A5 58 E5 22 85 58 B0 08 C6 59 90 04 B1 5A 91 58 88 0 D0 F9 B1 5A 91 58 C6 58 C6 58 C6 58 C6 59 C4 59 C50 59 C4 59 C4 59 C4 59 C50 50

The immediate action is to calculate the number of bytes to move. The number of low bytes is placed in the x register and location s22. The number of pages to move, the difference of the high bytes, is placed in the x register and immediately increased by one. This will be the counter where the zero state is checked to determine completion. As it is decreased before being checked, increasing by one will ensure that all pages will be done. If x was zero and was not incremented, then you would end up going around the circuit 256 times before a zero was discovered in the x register.

The low byte result is checked again; if there is no value, then a large chunk of instructions can be skipped. The bytes between addresses \$A3D0 and \$A3E7 deal with cases where there is an element of an incomplete page of data to move. These lines reduce the two end addresses by the number of low bytes to move. This will not effect the move as the data is loaded and stored with respect to y and this has the number that was the reduction. The incomplete page is moved first.

Except when y is zero, all the bytes are transferred within addresses \$A3E8 and \$A3EE. The y register will start at a high value and be decremented to zero. When that is reached, the next bytes are moved separately, before the high addresses are decreased. After this has been achieved, the x counter is reduced and checked, and if it is not zero, it's back to move the next page of data.

From this it can be seen that the transfer is done by taking the high addresses and moving them first. This means that the program will not overwrite itself.

Find a line

This routine finds the start address of a BASIC line, given the line number. We shall use it in our RENUMBER and DELETE. It uses all three processor registers and locations \$5F and \$60. On top of that the entry requirement is the line number in low/high byte form in locations \$14 and \$15.

ROM LISTING

A613	A5	2B	LDA \$2B
A615	Aó	2C	LDX \$20
A617	A0	01	LDY #\$01
A619	85	5F	STA \$5F
A61B	86	60	STX \$60
A61D	B1	5F	LDA (\$5F),Y
A61F	F0	1F	BEQ \$A640
A621	63		INY
A622	С8		INY
A623	A5	15	LDA \$15
A625	D1	5F	CMP (\$5F),Y
A627	90	18	BCC \$A641
A629	F0	03	BEQ \$A62E
A62B	88		DEY
A62C	D0	09	BNE \$A637
A62E	A5	14	LDA \$14
A630	88		DEY
A631	D1	5F	CMP (\$5F),Y
A633	90	0C	BCC \$A641
A635	F0	0A	BEQ \$A641
A637	88		DEY
A638	B1	5F	LDA (\$5F),Y
A63A	AA		TAX
A63B	88		DEY
A63C	B1	5F	LDA (\$5F),Y
A63E	B0	D7	BCS \$A617
A640	18		CLC
A641	60		RTS

Locations \$5F and \$60 are loaded with the start of BASIC. The high link address is again picked up first to see if the end of program has been reached. The high byte of the line number is checked first. If the value is greater than the required value, the carry will clear and the subroutine is left. If the two values are the same, we go forward to test the low byte values. Failure of either of these checks means that we have not reached the required line and have to go ahead and get the address
of the next line. When the low bytes are checked if they are equal, or the carry flag clear, the routine is terminated. On failing to find the desired line, or on finding a higher one, the link addresses are gathered in and we branch back to check the next line.

Due to the way the checks are made, the routine can be left in one of two states. In the first, the exact line number has been found, in which case the address in \$5F and \$60 will be what you require. The second state will be that there is no such line number and the routine returns the address of the next highest line. These conditions can be tested in the calling routine by examining the carry flag on return. If the carry is set then the actual line number was found, and if clear it was not.

7 Programming aid routines

Introduction

In Chapter 5 we gave routines to help in the preparation and editing of BASIC programs. These routines were themselves in BASIC, so were slow and had to be tagged onto the end of the resident program. This chapter not only puts these routines into 6502 machine code, but also extends their capabilities. In addition the following are included: OLD, RENUM, DELETE, MERGE, APPEND, DUMP, TRACE, CODER, HEX, BIN, TEN, TWO, AUTO, and MEM.

Our object has been to show you that with a little thought and perseverance, adding new BASIC commands is well within your grasp. Most of the routines start with an explanation of what we wish to achieve and how it is possible to do it. This is followed by the assembly listing and the label addresses used. These are provided for assemblers which do not allow the use of labels (Supermon) and with relocation in mind. Finally, a byte-by-byte explanation of the routine is given.

At the beginning of each routine, the command name and parameters are given for use in the UTILITY.

Renumber

COMMAND SYNTAX

RENUM start line number or Ø, increment, new start line value

Using \emptyset as the first parameter will indicate that the whole program requires renumbering. If a start line is set, it will renumber from that line to the end of the program:

for example, RENUM 0,10,100 RENUM 100,10,200

Later in the chapter we will discuss an AUTO routine. This is of use when typing in programs where the line numbers are sequential and of a fixed step. Renumbering a program makes it easier to read and opens up space to incorporate new lines.

The system we are going to use is known as a two pass system. The first pass will renumber commands that have line numbers associated with them. This is not as straight-forward as it might at first appear as the commands THEN and RUN have optional line numbers.

There are cases where we do not need to look for a 'renumbering command'. These will be after a DATA or REM token is encountered, or when inside quotes. In the latter case, we just loop until the next quote or the end of the line is found, whichever is soonest. The procedure on finding the tokens is simply to go to the next statement.

On finding a line number after a command, we convert it from its stored ASCII form to a two byte number. If it is less than the 'start line number', renumbering is not required. When it is not, we calculate its new value, convert it to ASCII, and overwrite the original.

Once all the directive line numbers have been dealt with, the simple task of actually changing the line numbers themselves is carried out.

We will be using many zero page locations in the routine and so to help you to follow the routine, a list of the main ones and what they control is given below:

\$FB and \$FC	Address of the current BASIC line being worked on
\$FD and \$FD-	Address of the next line – the links of the current line
\$49–	Stores the position in the current line, the line marker or y register
\$C9 and \$CA	Line number of first line to be renumbered
\$41 and \$42-	Address of the first line to be renumbered
\$BC-	Value of increment between new line numbers
\$BD and \$BE-	Value of the new start line number
\$B9 and \$BA–	Starts with the same values as \$BD and \$BE and is changed whilst calculating the new line number for directives after keywords.
\$58 and \$59–	Starts with the same values as \$41 and \$42 and is incre- mented to give the actual new number of a directive command

7	OPEN	=	\$8933		
8	CLOSE	=	\$888B		
9	*=\$ 89C5				
10		JSR	\$81F5	!	GET PARAMETER
20		JSR	\$AEFD	!	CHECK COMMA
30		LDA	\$14		
40		STA	\$C9		
50		LDA	\$15		
60		STA	\$CA		
70		JSR	\$81F5	!	GET PARAMETER - INC
80		JSR	\$AEFD	!	Check Comma
90		LDA	\$14		
100		STA	\$BC		
110		JSR	\$81F5	!	GET PARAMETER
					- NEW START LINE #

120		LDA	\$14		
130		STA	\$BD		
140		LDA	\$15		
150		STA	\$BE		
160		LDA	\$2B		
170		STA	\$FB		
180		I DA	\$20		
199		STA	\$FC		
200		I DA	\$CA		
210		RNF	FINDS		
228			\$69		
220		RNF	FINDS	ı	IS START I INE INPUT A
200		INY	#\$92	·	10 START EINE INFOT U
250			(4ED) V		
220		CTA		ı	GET EIDET PROG I INE #
200		TNIV	# C7	·	DET TIKST TROO EINE #
270			(#ED) V		
200		CTA	(PFD / ,)		
270	EINDO		₽UH d:CO		
300	F IND5	CTA	おし ア オートイ		
220		JIH I DA	#14 #CA		
220		CTA	⊅UH ¢15		
330		TCD	#A(17)	,	EIND START LINE ADD
250		DCC	PHOIS CTOPE	:	LINE FOUND
300			310RE	:	LINE FOOND
270			##1J #A427	ı	EPPOP - THECAL DIRECT
280	STORE		45F		STORE START I INF ADD
300	OTORE	CTA	\$41	•	STOKE STRACT EINE HDD
400			\$40		
410		STA	\$42		
420	START	I DY	#\$00		
438	011111	I DA	(\$FB).Y		
440		STA	4FD	ł	GET LINKS TO NEXT LINE
450		INY	4, 5	·	
460		LDA	(\$FB),Y		
470		STA	\$FE		
480		BNE	CONT	ļ	NOT END OF BASIC PROG
490		JMP	RENUM	!	CHANGE LINE NUMBERS
500	CONT	INY		ļ	SKIP LINE NUMBERS
510		INY			
520	NEXT	INY			
530		LDA	(\$FB).Y	ļ	GET CHAR OF LINE
540		BNE	CONT1	1	NOT END OF LINE
550	LINE	LDA	\$FD	ļ	PUT NEXT LINE IN
					LINE REGISTERS
560		STA	\$FB		

570		LDA \$FE		
580		STA \$FC		
590		BNE START	!	ENFORCED - NEXT LINE
600	CONT1	CMP #\$22	ļ	IS IT A QUOTE
610		BNE CONT2	ļ	NO
620	QUOTE	INY		
630		LDA (\$FB),Y	!	LOOK FOR NEXT
410			,	QUUTE OR LINE END
040 (50		BEW LINE	:	END OF PRUG LINE
470		UMP #\$22	:	QUUTE?
000		BNE QUUTE	!	
070	CONTO	BEQ NEXT	:	YES - NEXT CHAR
000	LUNIZ	UMP #\$8F	!	REM TUKEN?
070		BEU LINE	!	YES - NEXT LINE
700		CMP #\$83	!	DATA TOKEN?
/10		BEQ LINE	!	YES - NEXT LINE
720		CMP #\$A7	!	THEN TOKEN?
730		BEQ THEN	ļ	YES
740		CMP #\$8A	!	RUN TOKEN?
750		BEQ THEN	ļ	YES
760		CMP #\$89	!	GOTO TOKEN?
770		BEQ CONT3		
780		CMP #\$CB	!	GO TOKEN?
790		BNE NOGO	ļ	NO
800	SPACE	INY		
810		LDA (\$FB),Y		
820		CMP #\$20		
830		BEQ SPACE		
840		CMP #\$A4	ļ	TO TOKEN AFTER GO?
850		BEQ CONT3		
860	NOGO	CMP #\$8D	ļ	GOSUB TOKEN
870		BEQ CONT3	!	YES
880		CMP #\$E6	i.	RESET TOKEN
890		BEQ CONT3		
900		BNE NEXT	I	NO RELEVENT TOKEN
910	THEN	INY		
920		LDA (\$FB).Y	ī	GET NEXT BYTE
930		CMP #\$20	1	IS IT A SPACE
940		BED THEN	i	YES - SKIP IT
950		CMP #430	1	
,			•	IN ASCII
960		BCS NUMBER	ļ	FOUND A NUMBER?
970		DEY		
980		BNE NEXT	!	NOT LINE # AFTER THEN
990	NUMBER	DEY		
1000		CMP #\$3A	!	IS IT A NUMBER

1010		BCS	NEXT	!	NO LINE #
1020	CUNT3		(aco) V	,	CET NEVT DYTE
1030			(⊅FD/j) #∉20	:	A SPACES
1040				:	H SFHUE?
10/0		OEW	CUN13	:	TED CTODE I INE MADVED
1000		511	\$47	:	STURE LINE MARKER
1070		DET	***		COUNTED FOD NO OF
1080		LUX	# ⊅ 00	:	CHARS IN NUMBER
1898	DIGITS	INY			
1100		LDA	(\$FB).Y		
1110		CMP	#\$30	ļ	NUMBER IN ASCII ?
1120		BCC	CONT4	Ì	NO END OF LINE#
1130		CMP	#\$3A	Ì	NUMBER IN ASCII ?
1140		BCS	CONT4	1	NO END OF LINE#
1150		STA	\$0200.X	I	STORE IN INPUT BUFFER
1160		INX	, ,		
1170		BNE	DIGITS	ļ	ENFORCED GET NEXT BYTE
1180	CONT4	LDA	#\$3A	ļ	STORE COLON AS
					END MARKER
1190		STA	\$0200,X		
1200		STX	\$BF	ļ	NO OF CHARS IN LINE
1210		LDA	#\$02	!	SET CHRGET TO START OF BUFFER
1220		STA	\$7B		
1230		LDA	#\$00		
1240		STA	\$7A		
1250		JSR	\$81F5	ļ	CHANGE INTO REAL NOS
1260		LDA	\$14		
1270		STA	\$ C3		
1280		LDA	\$15		
1290		STA	\$C4		
1300		CMP	\$CA		
1310		BEQ	CHECK2		
1320		BCS	CONT5	ļ	> START LINE
1330	NORE	LDA	\$49	!	NO RENUM OF DIRECTIVE
1340		ADC	\$BF		
1350		TAY			
1360		JMP	Comma	!	CHECK FOR ON COMMAND
1370	CHECK2	LDA	\$C3		
1380		CMP	\$C9		
1390		BCC	NORE		
1400	CONT5	LDA	\$BD		
1410		STA	\$B9	!	TRANS START ADD TO WORKING REGISTER

1420		LDA \$41	TRANS START ADD TO
1430		STA \$58	WORKING REGISTERS
1440		LDA \$BE	
1450		STA \$BA	
1460		LDA \$42	
1470		STA \$59	
1480	FINDL	LDY #\$00	
1490		LDA (\$58).Y	SEARCH FOR LINE NO
1500		STA \$5A	SAVE LINKS
1510		INY	
1520		LDA (\$58).Y	
1530		STA \$5B	
1540		BNE CONT6	NOT END OF PROG
1550		LDY #\$02	
1560		LDA (\$FB) Y	GET LINE# FOR
		,	ERROR MESSAGE
1570		STA \$39	
1580		INY	
1590		LDA (\$FB),Y	
1300		STA \$3A	
1610		LDX #\$11	
1620		JMP \$A437 !	ERROR - UNDEF'D
			STATEMENT
1630	CONT6	INY	
1640		LDA (\$58),Y !	GET AND STORE LINE NO
1320		STA \$B7	
1660		INY	
1670		LDA (\$58),Y	
1680		CMP \$C4 !	COMPARE FOR SAME LINE
1690		BNE NEXTLINE!	NOT SAME
1700		LDA \$B7	
1710		CMP \$C3	
1720		BEQ FOUNDL	
1730	NEXTLINE	LDA \$B9	
1740		CLC !	INC REGS TO CALC
			NEW LINE NO
1750		ADC \$BC	
1760		STA \$B9	
1/70		BCC NOINC	
1780		INC \$BA	
1790	NOINC	LDA \$5A	
1800		STA \$58 !	PUT NEXT LINE ADD
			IN CURRENT REG
1810		LDA \$5B	
1820		STA \$59	

1830		BNE	FINDL	ļ	ENFORCED - CHECK NEXT LINE
1840	FOUNDL	LDX	\$B9		
1850		LDA	\$BA	ļ	MSB OF NEW LINE
1860		JSR	\$847F	ļ	CONVERT TO ASCII
1870		LDA	\$BF		1
1880		STA	\$3F		
1000		CPY	43E	ı	DOES MEM HAVE TO MOVE
1998		BED	NOMOUE	·	
1010		BCS	OPENUP		
1920		JSR	CLOSE	ł	REQUIRES LESS
1930		JMP	NOMOVE		011102
1940	OPENHP	JSR	OPEN	ł	REQUIRES MORE
1 / 10	01 21101	001	2.7 2.17	•	SPACE
1950	NOMOVE	LDY	\$ 49		
1960		LDX	#\$00		
1970	NEXTE	I DA	\$0200.X	ł	GET NEW NO IN ASCII
1980		BEQ	COMMA	1	END OF NUMBER
1990		STA	(\$FB).Y	i	STORE IN PROG
2000		INY	··· _ · , ·	•	
2010		INX			
2020		BNE	NEXTE	ļ	ENFORCED
2030	COMMA	LDA	(\$FB).Y	ļ	COMMA MEANS ON USED
2040		CMP	#\$2C		
2050		BEQ	ANOTHER		
2060		DEY			
2070		JMP	NEXT	ļ	GET NEXT TOKEN
2080	ANOTHER	JMP	CONT3	ļ	NEXT LINE - ON COMMAND
2090	RENUM	LDY	#\$00		
2100		LDA	(\$41),Y	!	GET AND, STORE LINKS
2110		STA	\$5A		
2120		INY			
2130		LDA	(\$41),Y		
2140		STA	\$5B		
2150		BNE	CONT8	!	NOT END OF PROGRAM
2160		PLA		!	REMOVE RETURN ADDRESS
2170		PLA			
2180		JMP	\$A474	ļ	GOTO READY FOR BASIC
2190	CONTS	INY			
2200		LDA	\$BD	ļ	NEW LSB LINE NO
2210		STA	(\$4 1),Y	!	CHANGE PROG
2220		INY			

M 111
•••
DV.
DI
LINE

LINES 10–150: The parameters are gathered here and put into their registers. Commas separating the inputs are also checked, giving 'SYN-TAX ERRORS' if not present.

LINES 160–190: The start of the BASIC program is now put into the current line registers, as it is also the address of the first line.

LINES 200–290: If the first parameter input was zero, indicating a full program RENUMBER, then the first line number is found and stored in its appropriate register.

LINES 300–410: Although we do not need this at the moment – here we find the address of the start line. We use the ROM routine FIND BASIC LINE

(see Chapter 6, page 142). If the carry flag is set, it will mean that the start line requested was not found and an 'ILLEGAL DIRECT' error will be printed. The address, if found, will be stored in \$41 and \$42.

LINES 420–1020: The byte-by-byte search for the appropriate keywords starts here. We start at the beginning of the BASIC program, no matter what the start line requested. As soon as the link addresses are collected and stored, the end of the program is checked for. Passing this means that only the actual line numbers require changing, so it is off to the final section of the whole routine, which is described later.

To continue finding the tokens, we skip the line numbers as they are not required here. Lines 540 to 590 set the values for the next line, after the end of line zero is discovered, and branches back to process the next line.

There is nothing of interest to us in quotes so on finding one we go into a loop to find a second quote or the end of the line. This is carried out in lines 620 to 670.

The next two tokens checked for are DATA and REM. Encountering these indicates that we can proceed to the next line as there will be nothing further to renumber in these lines.

There are two keywords – RUN and THEN – that may, or may not, have line numbers. These will therefore branch to check this possibility before proceeding. The standard Commodore directive commands are next in line: GOTO, GO TO and GOSUB. The centre keyword is checked in two stages, first for the GO. A loop is then set up to skip over spaces and then the TO token is looked for. All three keywords on being found will cause the routine flow to branch further ahead than RUN or THEN, as it assumes they will have line numbers. If not, then unless you have a line number of Ø, an error will be detected later.

The last keyword to be checked for is one of our new ones: RESET.

If line 900 is reached then we have not found a relevant keyword and therefore branch back to get the next program byte.

The last part in this section is to check if the next significant byte after RUN or THEN is a number, in ASCII. Spaces are skipped over and checks are made for values between \$30 and \$39 inclusive to continue.

LINES 1030–1390: The line numbers after the keywords will be in ASCII form and the line itself in two byte form. To do our calculations, we want both in the same two byte format. ASCII numerals therefore have to be changed.

After skipping spaces, we store our Y register (so we know where to write our new line number from), which is the line marker. Proceeding, we pick up bytes until a non-numeral is found, and store them in the input buffer. The x register is used to count the number of digits and is stored for later use. To convert the ASCII into the form we require, we use the GET PARAMETER routine. For this to work, we perform two operations. First, we make sure that after the last line number digit

there is a non-numeric character by storing a colon there. Secondly, we set CHRGET to point to the first numeral – \$0200.

The converted result is taken from registers \$14 and \$15 and stored in the two we have designated.

We only wish to renumber from the start line in the command. Lines 1290 to 1390 check this by comparing the two values. If no renumber of that particular line is required, we retrieve our line marker and increase it by the number of digits in the directive number, as we do not require to check them again. This is then transferred to the Y register. It will actually point to the byte following the last digit but this is taken into account in what follows. We jump further ahead to a position noted as COMMA (described later) starting at line 2030.

LINES 1400–1830: Having got this far, we have found a number which requires a different value. To find the new value, we have to go through the program from the designated start line and find the line that it points to (remember we have not changed the actual line numbers yet). At the same time, we calculate the new value.

To start with, we take the address of the start line and store it in \$58 and \$59. We then take the new value for the start line, the third parameter in the command, and place it in \$89 and \$BA. The line number we are checking for is held in \$C3 and \$C4.

As before, we get and store the link addresses, but here if we discover the end of the program has been reached (high link address of zero), an error is present as the line number of the directive was not found. The error produced is the same as when RUNING a program – 'UNDEF'D STATEMENT'. To make it easier for you, we also print out the line number with the error.

As long as this is not encountered, collecting a line number and comparing it for a match comes next. If it is not the one we want, the new start line number is increased by the increment value. This will calculate the value for the following line. The value is only increased on not finding a match, which conversely means that when the line searched for is located, the value is ready and waiting. After incrementing, we transfer the links to the line registers and branch back to check further lines.

LINES 1840–1940: We now have our new line number and need only to insert it into the line after the token, overwriting the original directive.

The new value is in two byte form and so requires converting to ASCII form. We do this by using a routine earlier in the UTILITY. This requires the accumulator to be loaded with the high byte and the x register with the low byte. We then call our 'convert to ASCII' at \$847F. This will do the conversion and store the answer in the input buffer, with a zero signifying the end. Also returned is the number of characters in the x register. If this value is the same as the original number,

the value stored in location SBF, we can just overwrite with no problems and proceed to line 1950.

The number of characters is transferred to location \$3E, via the accumulator. This location is for the memory move routine described in Chapter 6. By comparing the value in the x register (the new number of digits) with the accumulator figure (the old number of digits), we determine if a move is required. If the x value is less, the CLOSE routine is called; if it is greater, the OPEN routine is called; if it is the same, no move is required.

LINES 1950–2020: This leaves us one thing to do which is to write in the new number. First, we reload the y register with the line marker. This points to the position of the first digit in the number. By increasing x, starting at zero, and y, we take the digits from the input buffer and store them into the program line. This is repeated until we collect a zero from the buffer, when the branch in line 2020 will fail.

LINES 2030–2080: When we went through the lines checking for tokens, we did not look for the ON statement. The reason for this is that the only time a comma should be used after a line directive following a GOSUB or GOTO is when the ON keyword has been used.

On entering these lines the v register will point to the byte following the last one checked or stored. We load the accumulator with that byte to see if a comma is present. Finding it means that we branch back to the position just after the token search where it will commence with gathering in the line number directive for processing.

If the comma is not present, the γ register is decreased and we go back to start the search for the next appropriate token. The γ has to be decreased in case it was the end of the line, otherwise it would not have mattered.

LINES 2090–2350: RENUMBERING THE LINE NUMBERS

All the directive line numbers have now been checked and processed where required. The only thing which remains is to renumber the program lines themselves.

This is started from the line number requested in the first parameter of the command and whose address is held in \$41 and \$42. Its new starting value is held in \$BD and \$BE, and these will be incremented by the value in \$BC for each line. We progressively go through the program inserting the new numbers, in two byte form, until the end of the program is reached. After each line the number value is incremented ready for the next line.

When the renumbering is completed, we take the return address from the stack and jump to the start of BASIC to await further instructions. We do this to break out of the program in the unlikely event of RENUM being initiated within it.

A WORD OF WARNING

We cannot stress strongly enough that a copy of the original program should be saved to tape, or disk, prior to renumbering. We all make mistakes and if the RENUM finds a non-existent line number after, say, a GOTO, then an error is produced. This leaves the program only partially renumbered.

Auto

COMMAND SYNTAX AUTO first line number, increment

To escape from AUTO simply press RETURN immediately after the line number, as if deleting a single line.

This command removes the need to type in line numbers. The user just decides the start line number and the increment between consecutive lines. To achieve this we want a place to break into the normal flow of BASIC. This is made possible as every time an input line is processed it goes to a vectored jump before it is ready to receive the next. This is called the BASIC Warm Start Vector which is at \$0302 and \$0303. By changing the address in this vector to a routine of our own we can calculate the line number, put it into ASCII form and then insert it into the keyboard buffer just as if you typed it yourself. Having completed this, we then return to the input routine for you to make up the program line.

- 9	*=\$842E				
10		JSR	\$81F5	!	GET PARAMETER
20		JSR	\$AEFD	ļ	CHECK COMMA
30		LDA	\$14	ļ	1ST LINE# - LOW
40		STA	\$FB	!	SPARE ZERO PAGE
					LOCATION
50		LDA	\$15	!	1ST LINE# - HIGH
60		STA	\$FC	ļ	SPARE ZERO PAGE
					LOCATION
70		JSR	\$81F5	!	GET PARAMETER
80		LDA	\$14	i	INCREMENT
90		STA	\$FD	ļ	SPARE ZERO PAGE
					LOCATION
100		LDA	# <auto< td=""><td></td><td></td></auto<>		
110		STA	\$0302	!	START BASIC VECT LOW
120		LDA	#>AUTO		
130		STA	\$0303	ļ	START BASIC VECT HIGH
140	AUTO	LDA	\$0200	ļ	1ST CHAR IN BUFFER

150		BEQ	EXIT	ļ	TURN OFF AUTO
160		LDX	\$FB		
170		LDA	\$FC	ļ	NEXT LINE NUMBER
180		JSR	ASCII	ļ	PUT LINE # IN ASCII
190		STX	\$ C6	ļ	NO OF CHARS IN
					KEYBOARD BUFFER
200	NEXT	LDA	\$0200,X	ļ	PICK UP ASCII
210		STA	\$0277.X	ļ	PUT IN KEY BUFFER
220		DEX	,		
230		BPL	NEXT		
240		CLC			
250		LDA	\$FB		
260		ADC	\$FD	ł	INCREMENT LINE NUMBER
270		STA	\$FB		
280		BCC	NOINC		
290		INC	\$FC		
300	NOINC	JMP	\$A 483	!	READY FOR BASIC
310	EXIT	LDA	#\$83		
320		STA	\$0302		
330		LDA	#\$A4		
340		STA	\$0303		
350		JMP	(\$0302)		
360	ASCII	STX	\$63	!	LOW BYTE
370		STA	\$62	!	HIGH BYTE
380		LDX	#\$90		
390		SEC			
400		JSR	\$BC49		
410		JSR	\$BDDF		
420		JSR	\$B487		
430		JSR	\$B6A6		
440		LDX	#\$00		
450	AGAIN	LDA	\$0100,X		
460		STA	\$0200,X	!	PUT ASCII CHARS IN
. – .					INPUT BUFFER
470		BEQ	FINISH	ļ	ZERO CHAR WAS FOUND
480		INX			
490		BNE	AGAIN		
200	FINISH	RTS			
0404	A CA 1NI			,	
8474	AUAIN	{	847F ASCII		
8440	AUTU	{	34/2 EXIT		
8491	FINISH	{	8458 NEXT		
843F	NUINC				

LINES 10–130: These lines are only used when the AUTO command is active. We take the parameters of the command and place the start line number in \$FB and \$FC and the increment in \$FD. As we only take the low byte of the increment, the multiples over 256 are ignored. The only syntax check is made in line 20, where we check for a comma between the two parameters. Lastly in this section we change the vector address to point to the AUTO numbering which starts in line 140.

LINES 140–300: The first thing is to see if the first character in the input buffer is zero. This will signify no BASIC coding was inserted in the line and that you want to cancel the AUTO routine. When a BASIC line is typed, the line number is put into the input buffer. During the processing stage of inserting the line into the program BASIC takes out the line number and moves the rest of the line back up the buffer to overwrite it. This means that if the first input after a line number was a RETURN (BASIC inserts a zero for that as an end of line marker) then the first character in the buffer will be a zero after the line number has been removed. Therefore, on finding a zero we will branch off to exit the AUTO mode.

Assuming that we are still auto-numbering, we take the values in \$FB and \$FC and go off to convert them to ASCII form. This we will come across shortly. On returning from that subroutine the x register will have the number of ASCII characters in the line number and they will be in the input buffer. The x value is stored in the register which tells the operating system how many characters will be in the keyboard buffer. Having done that, we transfer the characters from the input buffer to the keyboard buffer.

We now set the line number for next time by adding the increment in \$FD to the values in \$FB and \$FC.

That is all there is to do, so we return you to the normal BASIC flow where the input routine will take the line number from the keyboard buffer, place it in the input buffer and print it on the screen.

LINES 310–350: These lines will be operated when you want to exit from AUTO. All we do is restore the BASIC Warm Start Vector to its initial value and then return you to BASIC to wait for your next instruction.

LINES 360–500: CONVERT TO ASCII: This subroutine will also be used by other commands when they require a one or two byte number converted into ASCII form. The subroutine is entered with the low byte of the number in the x register and the high byte in the accumulator.

The conversion is carried out by four ROM routines, but before we can call them we have four items to set. First, the number to be converted is transferred to locations \$63 (the low byte) and \$62. These are part of the floating point accumulator #1 (FAC#1) which is the main number manipulation area for BASIC. The other two prevent certain actions in the conversion process. Setting the carry flag will bypass a

routine that will perform the complement of the number and loading x with \$90 will set the Exponent byte of FAC#1 to avoid getting an answer in exponent form.

The first ROM routine visited clears all the bytes in the FAC#1 (or sets them to default values) which we have not dealt with. The next routine does the actual conversion. The remaining routine puts the result into a string and places it at the bottom of the stack area. The last byte placed there will be zero to mark the end.

We cannot leave it there as BASIC often uses this area. We therefore transfer it to the input buffer where it can be taken and used by the coding calling this routine. On exit the value in the x register will be the number of ASCII characters in the conversion.

Merge and append – combining BASIC programs together

COMMAND SYNTAX MERGE "program title", device APPEND "program title", device

The default device is tape and if there is no program title, the first program found on the tape will be used.

Merging programs together means that they are weaved together in program line order. The result is as if you typed in the lines of the merging program at the keyboard. This also means that if the programs both have lines with the same number, the ones in the merging program will overwrite the original.

Appending a program to another is simply a process of tagging it onto the end of the one in memory, irrespective of line numbers.

Merging is the more complicated of the two programs, but is not really complicated in itself. Both programs are initially loaded at the end of the current memory program, APPEND overwriting the last two bytes whilst MERGE comes just after them. The last two bytes of the program are the unique link address of zero, signifying the end of a BASIC program. By overwriting them on APPEND, we achieve our aim immediately and all that remains is to amend the link addresses to continue the program flow and to reset the end of program pointer.

In merging we take each new line in turn and insert it in the main program – if we had overwritten the original end links we would both merge and append which we do not want. To incorporate the new lines we make use of the normal BASIC input routine. After you input a program line and press RETURN, BASIC takes off the line number and then tokenizes all the BASIC keywords. At this point the line number is in two registers, in a two byte form rather in than the ASCII form typed in, and the line's content has been moved to the beginning of the input buffer. There is a counter of the number of bytes in the line which is four greater to incorporate the line number and link address. BASIC therefore knows the total space required for the line. We will enter the ROM routine at this point with the appropriate data set. Unfortunately, it is not a subroutine but finishes up waiting for an input. We therefore have to change the same vectors as in the AUTO routine (the BASIC Warm Start) to point us back to continued merging until we reach the end.

- 9	* ≕\$87A7		
10		JSR SETADD	! SET ADDRESSES
			FOR MERGE PROG
20		STX \$2B	SET MERGE PROG START
30		STY \$2C	SET HEROE FROD START
40		LDY #\$AA	PHT ZERO AT 1ST
			LOCATION
50		TYA	
60		STA (\$2B).Y	
70		JSR LOAD	
80		STX \$2D	' END OF MERGE PROG
90		STY \$2E	
100		JSR \$4533	PRECHAIN MERGE PROG
110		JSR RESET1	RESTORE POINTERS TO
120		LDA #KMERGE	
130		LDX #>MERGE	
140		STA \$0302	CHANGE WARM START
			VECTOR TO MERGE
150		STX \$0303	
160	JOIN	LDA #\$01	
170		STA \$7B	
180		LDA #≸FF	
190		STA \$7A	! SET UP CHRGET
200		LDY #\$00	
210		LDA (\$FB),Y	! GET AND STORE LINKS
220		STA \$FD	
230		INY	
240		LDA (≸FB),Y	
250		STA \$FE	
260		BEQ EXIT	! END OF MERGE PRG
270		INY	
280		LDA (\$FB),Y	! GET AND STORE LINE NO
290		STA \$14	! BASIC EXPECTS THEM
			IN THESE LOCATIONS
300		INY	
310		LDA (\$FB),Y	

320		STA	\$15		CET COUNTED TO
330		LDX	#≯04	:	TNCLIDE LINKS AND
					LINE#
340	NEXT	INX			
350		INY			
360		LDA	(\$FB),Y	!	GET LINE DATA
370		STA	\$01FB,X	!	STORE IN INPUT BUFFER
380		BNE	NEXT	ļ	NOT END OF LINE 0
390		TXA			
400		TAY		ļ	PUT COUNTER IN Y
410		JSR	\$A4A2	ł	BASIC AFTER CRUNCH
					ROUTINE
420	MERGE	LDA	\$FD	ļ	BASIC WARM START
					POINTS HERE
430		LDX	\$FE	!	PUT LINKS IN LINE
					REGISTERS
440		STA	\$FB		
450		STX	\$FC		
460		BNE	JOIN	ł	ENFORCED AS \$FE
					CHECKED FOR 0 EARLIER
470	EXIT	LDA	#\$83		
480		LDX	#\$A4		
490		STA	\$0302	ļ	RESTORE WARM START
500		STX	\$0303		
510		JSR	\$A474	ļ	READY FOR BASIC
520	LOAD	JSR	\$E1D4	!	GET LOAD PARAMETERS
530		LDA	#\$00	ļ	FOR RELOCATED LOAD
540		STA	\$B9	ļ	IN CASE A SEC ADD
550		I DX	\$28		
560		I DY	\$20	ı	SET LOAD START
570		JSR	\$EED5	i	KERNAL LOAD
580		BCS	ERROR	i	BAD LOAD
590		JSR	\$FFB7	I	READ I/O STATUS WORD
600		AND	#\$BF		
610		BEQ	EXIT1	I	GOOD LOAD
620		JSR	RESET	Ì	RESET POINTERS TO
					ORIGINAL PROG
630		LDX	#\$1D		
640		JMP	\$A437	ļ	LOAD ERROR
650	EXIT1	RTS			
660	ERROR	PHA		!	SAVE FOR ERROR
670		JSR	RESET		
680		PLA			
690		JMP	\$E0F9	ļ	ERROR DEPENDING ON A

700	RESET	LDA	\$FB		
710		SEC			
720		SBC	#\$02		
730		STA	\$14		
740		LDA	\$FC		
750		SBC	#\$00		
760		STA	\$FC		
770		LDA	#\$00		
780		TAY			
790		STA	(\$14),Y	!	RESTORE TWO ZEROS AT END OF PROG
800		INY			
810		STA	(\$14),Y		
820	RESET1	LDA	\$FD		
830		LDX	\$FE		
840		STA	\$2B		
850		STX	\$20		
860		LDA	\$FB		
870		LDX	\$FC		
880		STA	\$2D		
890		STX	\$2E		
900		RTS			
910	SETADD	LDA	\$2B		
920		STA	\$FD		
930		LDA	\$2C		
940		STA	\$FE		
950		LDX	\$2D		
960		LDY	\$2E		
970		STX	°\$FB		
980		STY	\$FC		
998		RTS			
1000	! APPEND	ROUTINE	Ē		
1010		JSR	SETADD		
1020		TXA			
1030		SEC			
1040		SBC	#\$02		
1050		STA	\$2B		
1030		TYA			
1070		SBC	#\$00		
1080		STA	\$2C		
1090		JSR	LOAD		
1100		STX	\$FB		
1110		STY	\$FC		
1120		JSR	RESET1		
1130		JSR	\$A5 33		
1140		RTS			

8830	ERROR	8803	EXIT
882F	EXIT1	87CA	JOIN
8810	LOAD	87F9	MERGE
87EB	NEXT	8838	RESET
884D	RESET1	885E	SETADD

This time we are not going to describe the program in line number order. There are three subroutines in the body of the program used both by MERGE and APPEND and we will deal with these first.

LINES 910–990: SETADD: This simply takes the start and end addresses of the original program and temporarily stores them. On coming out of the subroutine, the Y register will contain the high byte of the end address and the x register, the low value.

LINES 700–900: RESET: The first 12 lines will only be encountered when there is an error in loading the secondary program. These simply ensure that the end of program zeros are at the end of the original program. This will mean that when exiting from either command your original program is intact before starting.

The remaining lines are the reverse of SETADD, that is, they take the values in the temporary registers and place them in the program end and start registers. These last lines are called in the assembly listing as RESET1.

LINES 520–690: LOAD: The first thing this subroutine does is to call one resident in the BASIC ROM used by the standard LOAD and SAVE commands. It gathers up the parameters and sets various registers according to that information, and as it is there we also make use of it. We are going to do a relocated load and if a secondary address is present this will override our objective. To correct this we load location \$B9 with zero to bring back the state for a relocated load.

The KERNAL LOAD routine expects the start address of the load in the two processor registers, x and Y, with the former holding the low byte of that address. The accumulator is the flag for either a load or a verify operation. The value for load is zero, the other being one, which was set whilst confirming the secondary address. The KERNAL LOAD routine is situated at \$FFD5. Error checking comes now in the order of operations – you may have put in the wrong tape or disk. The first indicator to a bad load is the carry flag being set; if this is so, then we branch off to deal with it. We have to check the 1/O status word if the carry is clear. This is achieved by calling another KERNAL routine, at \$FFB7. The result coming out of this call is ANDEd with the value \$BF and everything is fine if the zero flag is set. The error given for any other outcome is 'LOAD ERROR'.

On the first check we go to line 660 if the carry was set. The value in the accumulator will be used for the error so we temporarily store this

on the stack. The reason for this is to reset all the pointers we altered to give you back your original program. This is done by a call to RESET, described above. Once done we retrieve the accumulator value and jump to part of the BASIC loading routine for an error to be generated, based on the accumulator value.

MERGE

LINES 10–150: First of all we call the SETADD routine. From this we can set the start address of the merging program to immediately after the memory program. BASIC expects the first byte to be zero and therefore we oblige it. Having done that, we load the program we want to merge. The address that is returned from loading is placed in the end of program address. At this point as far as BASIC knows the only program in memory is your merge program. The link addresses have to be set up so we know where the lines start. This is done by a call to the ROM routine at \$A533 which will do this for us. From now on we want BASIC to respond to the original program and therefore we reset all the registers back to their original values through RESET1.

To merge the lines into the master program we use the program line input routine in the BASIC ROM. This is not a stand alone routine but ends up at the BASIC Warm Start after each line has been processed. It follows that, as in the AUTO routine, we will have to alter the vector pointing to that position to divert to this routine until all lines are merged. This is done in the lines 120 to 150. The addresses will be in the location of line 420.

LINES 160–460: These are the instructions which actually combine your two programs. First, CHRGET is set to a position one place before the start of the input buffer. We now turn our attention to the merging program lines. The address of the first line will still be in locations \$FB and \$FC following the SETADD routine called at the beginning. Using this information, we get the link addresses and store them for later. We use this system a lot in our routines but in this case it is vital. We have left no room between the programs so that when the line is transferred the master program will be longer, overwriting the needed link addresses. The check is also made for the end of the merge program, the usual high byte zero link address.

Next we take the two byte line number and place it in registers so that BASIC will know where to find them, \$14 and \$15. The contents of the line, including the zero byte end marker, are now transferred to the input buffer starting at \$0200. The listing shows \$01FB but there is a reason for this. The x register will come out with a value of the number of bytes in the line, but to account for the four bytes holding the link addresses and the line number, it starts with a value of 4. This is one more than required but the ROM routine will compensate. This means that with the initial value of \$04 in x the first location written to will be \$0200. The ROM routine we are using wants the number of bytes in the y register rather than the x, so we oblige by transferring them via the accumulator.

We are now ready to use the ROM to merge our line into the main BASIC program. We join the ROM just after the coding that turns the keywords into tokens – ours are in that state already. As far as the ROM is concerned you have typed in the line and will put it in the program as such.

The BASIC Warm Start Vector will bring us back after inserting the line to LINE 420 in the above listing. We now put the address of the next line into the working registers and branch back to deal with it. The branch instruction will always succeed as we have checked previously for a zero value in SFE.

These lines will be repeated until all the merge program lines have been assigned to the master program.

LINES 470–510: The merge is complete so we restore the BASIC Warm Start Vector to its normal setting and return to BASIC where it will await further commands.

APPEND

LINES 1010–1140: The first thing, as in MERGE, is to call the SETADD routine. On coming out though, it is slightly different. The new BASIC program start has to be reduced by two so will overwrite the end of memory program links. The appending program will load directly after the final line of the master program. On completion of the load, we store the loading end addresses (not in the end of program registers), but overwrite the original stored values, set by SETADD. This means that on resetting the values, by RESET1, we end up with the original start and the end marker corresponding to end of the appended lines. The final thing to do is to set the link addresses to follow as one program. This is carried out by the ROM routine at \$A533. The two programs have been joined together with our own form of 'superglue'.

Delete

COMMAND SYNTAX DELETE first line number,[second line number]

The first line number to be deleted and the comma are essential and if missing will give errors. The second line number is optional in that if you want to delete to the end of the program you omit it; otherwise insert a number.

Deleting a line of BASIC program is easy, you just type in the line number followed by a return. There is no real hardship in deleting one or two lines but longer blocks become tedious and time-consuming. DELETE will rid you of a block of lines with one command. To do this we use the same ROM routine as if deleting one line. What the BASIC ROM does is to take the address of the line to be deleted and the link address within that line. It then takes the program starting at that link address to the end of the program and moves it to the address of the line to be deleted. For example, if we have a line whose address is \$0001, its link address is \$0025 (the start of the next line), and the end of the program is \$0A45, the block to move is \$0025 to \$0A45 with its new starting address of \$0001. Hence, the line at \$0001 is overwritten or deleted. By the way, the variables and the arrays are also moved along with the block.

It therefore goes to show that if we get the address of the next line after the 'second line number' and place it in the link address of the 'first line number', a whole block of lines will be overwritten at once. Where there is no 'second line number' we take the end of program address and deduct two from it. This will give us the address of the two zero bytes at the end of a program. The first line number is placed in the input buffer, with a zero at the end of it signifying no further data, and goes to ROM as if you typed it in.

Two listings follow for this command. The reason for this is explained later.

9	*=\$8F44				
10		BCC	DEL	ļ	PARAMETER A NUMBER
20	SYNTAX	JMP	\$AF08	ļ	GENERATE SYNTAX ERROR
30	DEL	JSR	\$81F5	!	GET 1ST LINE NUMBER
40		JSR	\$A613	!	FIND LINE LOCATION
50		BCS	FOUND	ļ	LINE NUMBER FOUND
60		LDX	#\$15	ļ	ILLEGAL DIRECT ERROR
70		JMP	\$A437	!	ERROR ROUTINE
80	FOUND	LDA	\$5F	!	PUT LINE ADDRESSES
90		STA	\$FB	ļ	IN STORAGE
100		LDA	\$60		
110		STA	\$FC		
120		JSR	\$0079	ļ	CHECK FOR COMMA
130		CMP	#\$2C		
140		BNE	SYNTAX	!	NOT FOUND
150		JSR	\$0073	ļ	GET NEXT BYTE
160		BNE	NUMERAL	!	A SECOND LINE NUMBER
170		SEC		i	PREPARE FOR SUBTRACT
180		LDA	\$2D	ļ	END OF PROGRAM
190		SBC	#\$02	ļ	DEDUCT BY TWO
200		STA	\$5F	!	READY FOR DELETION
210		LDA	\$2E	ļ	END OF PROGRAM

220		SBC	#\$00	ļ	IN CASE OF PAGE CROSSING
230		STA	\$60	ı	READY FOR DELETION
240		BNE	CONT	!	ENFORCED \$2E CAN'T BE ZERD OR 1
250	NUMERAL	BCS	SYNTAX	ļ	NO NUMBER
260		JSR	\$81F5	!	GET 2ND LINE NUMBER
270		INC	\$14	ļ	SO WE GET
					FOLLOWING LINE
280		JSR	\$A613	ļ	FIND LINE
290		LDA	\$FC	!	CHECK IF 1ST NO IS SMALLER THAN 2ND
300		CMP	\$60		
310		BCC	CONT		
320		BNE	SYNTAX		
330		LDA	\$FB		
340		CMP	\$5F		
350		BCS	SYNTAX		
360	CONT	LDY	#\$00		
370		LDA	\$5F		
380		STA	(\$FB),Y	!	STORE ADRESS
390		INY			
400		LDA	\$60		
410		STA	(\$FB),Y		
420		INY			
430		LDA	(\$FB),Y	!	GET 2 BYTE LINE NO
440		TAX			
450		INY			
460		LDA	(\$FB),Y		
470		JSR	\$847F	!	CONVERT TO ASCII
480				!	AND PUT INTO
					START OF INPUT BUFFER
490		PLA		!	REMOVE RETURN ADDRESS
500		PLA			
510		LDX	#\$FF	ļ	TO INITAILIZE INPUT BUFFER
520		LDA	#\$01		
530		JMP	\$927D	!	WILL DELETE LINE AND RETURN TO BASIC

8F91	CONT	8F49	DEL
8F56	FOUND	8F79	NUMERAL

8F46 SYNTAX

ASSEMBLY LISTING2

9	*=\$927D			
10		LDA \$0302	ł	SAVE WARM START
20		STA ADD+1		
30		LDA \$0303		
40		STA HADD+1		
50		LDA #(ADD	!	CHANGE WARM START
30		STA \$0302		
70		LDA #>ADD		
80		STA \$0303		
90		JMP \$A486	!	DELETE BLOCK
100	ADD	LDA #\$83	!	RESTORE WARM START
110		STA \$0302		
120	HADD	LDA #\$A4		
130		STA \$0303		
140		JSR \$A533	ļ	RECHAIN LINES
150		CLC	!	PREPARE FOR ADD
160		LDA \$22		
170		ADC #\$02	!	INCREASE FOR
				VARIABLE START
180		STA \$2D	ļ	START OF VARIABLES
190		LDA \$23		
200		ADC #\$00	ļ	IN CASE OF CARRY
210		STA \$2E		
220		JSR \$A660	ļ	CLR
230		JMP \$A474	!	READY FOR BASIC

9296 ADD

9298 HADD

LISTING1

LINES 10–110: These instructions deal with the 'first line number'. The routine first checks the carry flag, set or unset by CHRGET on entering, and if set a SYNTAX ERROR is generated as the first byte after the delete token was not a numeral. A call to our GET PARAMETERS routine is next, immediately followed by a visit to the ROM routine FIND BASIC LINE. The result from GET PARAMETERS is in the registers used to call FIND BASIC LINE. On returning from the latter, if the carry is not set, then the line was not found and we therefore generate a further error. The address of the first line is placed in locations \$FB and \$FC.

LINES 120–280: The remaining parameter is now dealt with. CHRGET is positioned to where the comma should be, so we call CHRGOT to see if it is there. A call to CHRGET now will get the first byte of the second line number. If no line number is present, the zero flag will be set. In that

case we gather in the address of the end of the program, deduct two from it, and store the result in registers \$5F and \$60.

If the zero flag was not set, we make a further check as earlier to see if the byte picked up by CHRGET was a numeral. GET PARAMETER is called to get the second line number and the low byte result in \$14 is increased by one. This is done as we do not require the address of that line but rather the one following it. After the visit to FIND BASIC LINE the address in \$5F and \$60 will be the next line, whether its line number is one or ten greater than the 'second line number'.

LINES 290–350: These instructions check to see that the address of the 'second line number' is higher than that of the first, otherwise a SYNTAX ERROR is given.

LINES 360–480: Here we insert the second address we found into the link address position of the first line. We then get the line number, in its two byte format, putting the low byte in the x register and the high into the accumulator. We now call another routine which we coded at \$847F in the UTILITY where the number will be converted to ASCII and placed in the input buffer, starting at \$0200, with a zero at the end.

LINES 490–530: From the stack we remove the return addresses which were placed there on entering DELETE. The x register and the accumulator are given the address of the input buffer less one which will be the CHRGET address. The final thing is to jump to the second listing.

LISTING2

LINES 10–230: The ROM routine that we will use is not a subroutine but ends up at the BASIC Warm Start Vector. We want to return here so we first store its present values and replace them to point back to these lines. We now go to ROM where it will treat the number in the input buffer as if you were deleting a single line from the keyboard, but as we have changed the link address it will delete more than one.

On returning, we restore the BASIC Warm Start Vector. We now subject the program to the rechain routine – not that it requires it, but from this routine we can calculate the end address. From the address the rechain routine ends with, we add two and set the end of program registers. A call to the CLR routine will set the remaining variable addresses. Finally, we jump to BASIC, printing 'READY' and give you back control.

The reason for two listings is due to the way in which the ROM memory moving routine sets the end of program address. We came across this when testing the UTILITY. The BASIC normally expects lines of around 80 characters and definitely no more than 255. Mainly for the latter reason the ROM routine only decreases the end address by the maximum of a page. It does not affect the deletion, but it did not make the required reduction in memory used. The second listing was added to overcome the times when the number of lines took more than 256 bytes. Thus in the second listing we were able to set the addresses ourselves.

Memory - Display number of bytes free

COMMAND SYNTAX MEM

There are no parameters in this command. The command is available only in direct mode. If found in a program the routine is not carried out.

BASIC has a command that prints out the amount of space available to it. It is FRE(x) where x is a dummy argument. Unfortunately it returns, when used with PRINT, an integer value which means any value over 32767 (\$7FFF) will be a negative number. For example, if the number of bytes free is 36500, the result printed would be -29035. If you add that, with the sign, to 65535 (\$FFF), you will arrive at the true figure of 36500. We produce here a short routine to print out the correct value straight away. Having said that, with the UTILITY in place, the maximum space available is less than 32768 and so FRE(x) will always print out the correct value.

The first thing to do is call a ROM routine to do a 'garbage collection'. It is at \$B526 and tidies up the variable and string area. It will reset the necessary registers after the compaction. The area of memory that is unused will be from the end of arrays to the beginning of the area used by strings. If we take the higher address from the lower, we will have the number of free bytes available.

The routine that we have used to print the result to the screen is a subroutine of the HEX command, which is described later. Suffice to say that on calling this subroutine with the low byte in the Y register and the high in the accumulator, it will convert it to ASCII and print the result to the screen.

To check whether you are in direct mode, we look at location \$9D(157). This will hold \$80 (128) for direct or \$00 for program mode.

9	*=\$85FC		
10		LDA \$9D	! DIRECT OR PROGRAM
20		BNE MEM	! DIRECT ONLY
30		RTS	! PROGRAM NOT EXECUTED
40	MEM	JSR \$B526	! ROM COLLECT GARBAGE
50		SEC	! PREPARE FOR SUBTRACT
60		LDA \$33	! POINTER START OF
			STRING STORAGE
70		SBC \$31	! POINTER END OF ARRAYS

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80	TAY	! TEMP STORE
90	LDA \$34	! POINTER START OF
		STRING STORAGE
100	SBC \$32	POINTER END OF ARRAYS
110	JMP \$85AD	! CONVERT TO ASCII
		AND PRINT TO SCREEN

8601 MEM

Coder

COMMAND SYNTAX CODER There are no parameters to this command.

How many times have you picked up a listing from a magazine and wondered what graphic symbol is in that PRINT statement? Is it a shifted N graphic or shifted L? How many have been used together, is it 2 or 3? You then come across a colour code and have to look it up in the manual to remember which colour to program. Owners of non-Commodore printers also have a problem as these symbols and graphics do not print.

We would like to introduce a routine that replaces these graphics with mnemonics. For example, the symbol for clear screen would be replaced by [CLS].

Except for one, all the codes we want to change appear within quotes. That means we have to look through the program for a quote and when found look for ASCII values that we want to change until the end of the line or the second quote appears. Having found one, we also have to look to see if it has been repeated. This done, we will either calculate the new code or find one in a data table. The codes produced will be of a different length from the original, but if it repeats, may be of shorter overall length. To accommodate this, we will use the memory move routines described in Chapter 6.

The one exception we mentioned earlier is the mathematical 'PI' (3.14159, etc). This we also found does not appear on some listings and is essential if in a mathematical equation. This is therefore coded whether in or out of quotes.

Most of the program operation is described after the assembly listing (see below), but before the listing we would like to say a word or two about the data make-up. This can be split into two sections. Graphics that are obtained by using the shift with most of the 'letter' keys can be calculated directly to the ASCII code of that particular letter. The remaining graphics and codes require the use of data tables.

We have employed two tables and stored them out of the way under the BASIC ROM. The first table, the data address table, has the three bytes for each character we are going to encode. The first byte is the ASCII value of the character and is followed by the address within the second table where the data is stored. The second table, the data table, holds all the data for those characters. The data will be the characters printed between the [] brackets, and may be of differing length. Because of this the first byte is the number of bytes of data.

What are we going to produce instead of all these graphics and codes? These are listed in Appendix I and are mainly self-explanatory. However, an explanation of two of them is required. If you look at the *Programmer's Reference Guide*, page 74, under 'Other Special Characters', you will see five functions available. Three of these can be achieved more easily than described in the PRG. These are SWITCH TO LOWER CASE, DISABLE CASE-SWITCHING KEYS and ENABLE CASE-SWITCHING KEYS. They can be obtained by simply holding down the CTRL key and appropriate letter. In quotes they will print the appropriate symbol. Out of quotes the action will be carried out. The remaining two 'special characters' are implemented in the way the PRG describes. We have given them codes of [CRG>M] and [CRG>N]. These stand for CTRL REVS GRAPHIC SHIFT RIGHT and the appropriate letter.

7	OPEN	= \$8933	
8	CLOSE	= \$888B	
9	∗= \$8B93		
10		LDX #\$00	! INITIALIZE QUOTE
			COUNTER
20		LDA \$2B	! GET AND STORE
			START OF BASIC PROG
30		STA \$FB	
40		LDA \$20	
50		STA \$FC	
60	LINKS	LDY #\$00	! SET Y TO BEGINNING
			OF LINE
70		LDA (\$FB),∖	/ ! GET ADD OF NEXT LINE
80		STA \$FD	! STORE FOR LATER
90		INY	
100		LDA (\$FB),	(
110		STA \$FE	
120		BNE CONT	! NOT END OF BASIC PROG
130		PLA	! REMOVE RETURN ADDRESS
140		PLA	
150		JMP \$A474	! GOTO "READY FOR
			BASIC"-END OF CODER
160	CONT	INY	! SKIP LINE NUMBER
170		INY	

180	NEXT	INY			
190		LDA	(\$FB),Y	!	GET BYTE OF PROG LINE
200		BNE	CONT1	ļ	ZERO SIGNIFIES END
					OF LINE
210		LDA	\$FD	!	GET NEXT LINE ADDRESS
220		STA	\$FB	ļ	PUT IN CURRENT
					LINE REGISTERS
230		LDA	\$FE		
240		STA	\$FC		
250		LDX	#\$00	Ţ	RESET QUOTE COUNTER
260		BEQ	LINKS	ļ	X SETS ZERO FLAG -
					BRANCH ENFORCED
270	CONT1	CMP	#\$FF	ļ	IS IT DI
280		BNE	NOPI	ļ	NO
290		STA	\$3E	ļ	STORE VALUE
300		BEQ	CONT2	I	ENFORCED
310	NOPI	CMP	#\$22	Ì	IS BYTE A QUOTE
320		BNE	CHECK	I	NO GO TO SEE LE IN
					QUOTES
330		INX		ļ	IT'S A QUOTE SO INC
					COUNTER
340		CPX	#\$02	ļ	IS IT SECOND QUOTE
350		BNE	INQUOTES	ļ	IN QUOTES CODER IN
					ACTION
396		LDX	#\$00	ļ	RESET COUNTER
370		BEQ	NEXT	ļ	ENFORCED
380	CHECK	СРХ	#\$01		
390		BNE	NEXT	ļ	NOT IN QUOTES
400	INQUOTES	STA	\$3E	ļ	STORE BYTE
410		CMP	#\$C0	!	IS IT LESS THAN 192
420		BCC	Compare	!	YES
430		SBC	#\$60	!	NO SUBTRACT 96
440	COMPARE	CMP	#\$60	ł	IS IT $>$ OR = TO 96
450		BCS	CONT2		
460		CMP	#\$21	ļ	IS IT LESS THAN 21
470		BCS	NEXT	ļ	CHARS 21 - 95
					DON'T REQUIRE CODING
480	CONT2	STA	\$3D	!	STORE REVISED CHAR
					VALUE
490		STY	\$49	ļ	STORE LINE MARKER
500		STX	\$30	ļ	STORE QUOTE COUNTER
510		LDX	#\$01		
520	NEXT1	INY			
530		LDA	(\$FB),Y	!	GET NEXT CHAR
540		CMP	\$3E	ļ	IS A REPEAT CHAR
550		BNE	NEXT2	!	NO

560		INX			
570		BNE	NEXT1	!	ENFORCED
580	NEXT2	STX	\$3E	ļ	REG - NO OF REPEATS
590		CPX	#\$02		
300		BCS	CONT3	ļ	MORE THAN ONE CHAR
610		DEX			
620		LDA	\$3D	i	GET CHAR BACK AGAIN
330		CMP	#\$20	ļ	IS IT A SPACE
640		BEQ	SPACE	!	DON'T CODE SINGLE SPACE
650		LDA	#\$00		
660		STA	\$40	ļ	RESET REG WITH ASCII FOR NO OF REPEATS
370		STA	\$3F		
680		BEQ	CONT4	ļ	ENFORCED
690	SPACE	JMP	RELOAD	I	RELOAD REGISTERS
					FOR GET NEXT BYTE
700	CONT3	LDA	#\$88	ļ	X HAS LOW VALUE -
				·	NO HIGH VALUE
710		JSR	\$847F	ļ	NO OF REPEATS INTO ASCIL FORM
720		LDX	#\$88		
730		LDA	\$0200.X	ł	GET ASCLL INTO REGS
748		STA	\$3F	·	
750		INX			
760		LDA	\$0200.X		
770		STA	\$40		
780	CONT4	LDA	\$3D		
790		CMP	#\$61		
800		BCC	CONVERTE		
810		CMP	#\$7B		
820		BCS	CONVERTE		
830	CONVERTA	SEC			
840		SBC	#\$20	!	REDUCE VALUE
850		STA	\$3D		
860		LDX	#\$07	ļ	TOTAL NO OF SPACES REQUIRED
870		LDA	\$40		
880		BNE	CONT5	ļ	MORE THAN 9 REPEATS
890		DEX			
900		LDA	\$3F		
910		BNE	CONT5	ļ	SOME REPEATS
920		DEX			
930	CONT5	СРХ	\$3E	!	FIND HOW MUCH ROOM
940		BEQ	NOMOVE	!	RIGHT AMOUNT OF SPACE

950		BCS	OPENUP	!	NEEDS SPACE IN LINE
960		JSR	CLUSE	!	GET RID UF
0.70		-			UNWANTED CHARS
970		JMP	NUMUVE	!	CUNI WITH PRUG
986	UPENUP	JSR	UPEN		NUT ENUUGH RUUM IN
000		1.5%	+ 10	,	
770	NUMUVE		\$47 #*ED	:	GET LINE PUINTER
1000		LUH	#FD) //	!	ASULI FUR L
1010		510	(\$F8),1	!	PUT IN LINE
1020		DED	2940 CONT /		NO TENE DIGIT
10.00			CUNIO	:	NU TENS DIGIT
1040		ATO			
10.20	CONITY	1 DA	(PFD/)) #05		
1000	CUNTO	DED	POF CONT7	,	
1070			CUN17	:	NU REFERIS
1000			(#ED) V		
1100	CONT7	1.64	(PFD/j)		ACCLL EDP C
1110	CONT		##47	:	HSCII FOR O
1120		CTA	(450) V		
1120					ACCTI EOP \
11/0			##JE	:	ASCII FOR /
1150			(4EB) Y		
11.60			\$3D	1	CHAR
1170			¥30	·	CHAR
1180		STA	(\$FB) Y		
1190			#\$50	ł	ASCII FOR 1
1200		INY	1700	•	
1210		STA	(\$ FB).Y		
1220		LDX	\$30	ł	RESET QUOTE COUNTER
1230		JMP	NEXT	i	NEXT BYTE
1240	CONVERTB	STA	\$3D	•	
1250		LDA	#\$50	Ţ	LSB OF DATA
					ADDRESS TABLE
1260		STA	\$62		
1270		LDA	#\$A3	ţ	MSB OF D.A.T.
1280		STA	\$63	-	
1290		LDX	#\$51	ļ	COUNTER MAX NO OF
					CHARS IN DATA
					TABLE
1300		LDY	#\$00		
1310	NEXT3	JSR	\$81FB	ļ	SWITCH OF BASIC
1320		LDA	(\$62),Y	i	GET ASCII CHAR NO
1330		PHA	,	ļ	TEMP STORE
1340		JSR	\$8202	ļ	SWITCH IN BASIC
1350		PLA		ļ	RETRIEVE

1360		CMP	\$3D	! IS IT THE SAME	
1370		BEQ	FOUND	! YES	
1380		INY		! SKIP UNWANTED ADD	RESS
1390		INY			
1400		INY			
1410		DEX		! DECREASE COUNTER	
1420		BPL	NEXT3	! GET NEXT ASCII NO	I
				UNTIL X<0	
1430		LDA	\$3D		
1440		CMP	#\$1B	! SEE IF IT USED	
				WITH CTRL KEY	
1450		BCC	CTRL	! YES	
1460		LDX	#\$0D	! NO	
1470		JMP	\$A437	! ERROR OUT OF DATA	
1480	CTRL	ADC	#\$40	ADD \$40 TO GIVE	
				ASCILLETTER	
1490		STA	\$ A448	STORE IN DATA TAR	IF
1599		ΙDΔ	#\$43	I ISB OF DATA FOR	
		2011	n+ /0		
1510		STA	4 .4.2	ICTRE: J	
1520			# 4 ∆4	I MER DE DATA EOR	
1020		LUA	7+177	COTION	
1520		CTA	# / 2	LURILYI	
1540		DUC	₽03		
1550	COLNIN	BINE	LUNI8		
1000	FUUND				
1296		JSR	\$81⊦B	SWITCH UUT BASIC	
12/0		LDA	(\$62),Y	GET LSB OF DATA	
				POSITION	
1580		PHA		! TEMP STORE	
1590		INY			
1300		LDA	(\$62),Y	GET MSB OF DATA	
				POSITION	
1310		STA	\$63		
1620		JSR	\$8202	SWITCH IN BASIC	
1630		PLA			
1640		STA	\$62		
1650	CONT8	LDY	#\$00		
1660		JSR	\$81FB		
1670		LDA	(\$62),Y	! GET NO OF DATA CH	ARS
1680		STA	\$C1		
1690		JSR	\$8202		
1700		LDA	\$01		
1710		CLC			
1720		ADC	#\$04	FOR BRACKETS AND REPEATS	
1730		TAX			

1740	LDA \$40	
1750	BNE CONT9	
1760	DEX	! NO TENS IN REPEATS
1770	LDA \$3F	
1780	BNE CONT9	
1790	DEX	! NO REPEATS AT ALL
1800 CONT9	CPX \$3E	! DO WE REQUIRE A
		MEMORY MOVE
1810	BEQ NOMOVE1	! NO
1820	BCS OPEN1	! MORE SPACE
1830	JSR CLOSE	! LESS SPACE
1840	JMP NOMOVE1	
1850 OPEN1	JSR OPEN	
1860 NOMOVE1	LDY \$49	! GET LINE MARKER
1870	LDA #\$5B	! ASCII FOR [
1880	STA (\$FB),Y	
1890	LDA \$3F	
1900	BEQ CONTA	! NO TENS DIGIT IN REPEATS
1910	INY	
1928	STA (\$FB),Y	! STORE IN PROG
1930 CONTA	LDA \$40	
1940	BEQ CONTB	! NO REPEATS AT ALL
1950	INY	
1960	STA (≸FB),Y	! STORE IN PROG
1970 CONTB	STY \$49	! STORE LINE MARKER
1980	LDY #\$00	
1990	JSR \$81FB	! SWITCH OUT BASIC
2000 DATA	INY	
2010	LDA (\$62),Y	! GET DATA FROM TABLE
2020	STY \$C2	! STORE DATA MARKER
2030	LDY \$49	! GET LINE MARKER
2040	INY	
2050	STA (\$FB),Y	! STORE DATA IN PROG LINE
2060	STY \$49	! STORE LINE MARKER
2070	LDY \$C2	! GET DATA MARKER
2080	CPY \$C1	! HAVE WE GOT ALL DATA
2090	BNE DATA	! NO
2100	JSR \$8202	! SWITCH IN BASIC
2110	LDY \$49	
2120	LDA #\$5D	! ASCII FOR]
2130	INY	
2140	STA (\$FB),Y	

2150 2160 2170 2180 2190	RELOAD	LDX \$3C JMP NEXT LDY \$49 LDX \$3C JMP NEXT		GET QUOTE COUNTER NEXT BYTE TO PROCESS GET LINE MARKER GET QUOTE COUNTER NEXT BYTE TO PROCESS	
8BD7	CHECK	8889	CLOSE		
8BE3	COMPARE	8BAF	CONT		
8BC2	CONT1	8BEB	CONT2		
8015	CONT3	8C27	CONT4		
8042	CONT5	8C5E	CONT6		
8065	CONT7	8000	CONT8		
8CDF	CONT9	8D04	CONTA		
8D0 B	CONTB	8C7E	CONVERT	ТВ	
8CAB	CTRL	8D12	DATA		
8CBA	FOUND	8BDB	INQUOTE	ES	
889D	LINKS	8BB1	NEXT		
88F3	NEXT1	8BFD	NEXT2		
8C8C	NEXT3	8051	NOMOVE		
8CF7	NOMOVE1	8BCA	NOPI		
8931	OPEN	8CF4	0PEN1		
8C4E	OPENUP	8D33	RELOAD		
8012	SPACE				

LINES 10–150: These set up the routine and if necessary return control back to you through BASIC. There are no parameters included in the command to pick up as it codes the whole program. The address of the first line is taken from the start of BASIC program variables at \$2B and \$2C. The x register is initialized to zero and is used as a quote counter. We get the link address to the next line and if it is the end of the program we remove the return address from the stack, placed there on entering, and go back to BASIC with the program in memory coded for listing or saving.

LINES 160–390: We are going to start to look for our trigger codes – quotes or pi. We skip the line number and start to scan. If the end of the line is encountered we transfer the links to the line register and start the next line.

We check for PI (ASCII value is \$FF). On finding it, we store it in a register for later use and branch further into the program. The check for the quote takes two forms. When one is found, we increase and check the x register. A two here will indicate that it is the second quote and therefore going out of the area we are interested in. It also means we go back to look for another quote.

If x is one, then the first quote has been found and we go forward to

check for codes. It will fail there the result is that we return to get the next byte.

On encountering a byte other than a quote or pi, we check to see if the x register is one, indicating that we are in quotes and it will require processing.

LINES 400–470: We first store the byte. This is done as we are going to manipulate this data and possibly alter it. The original value is needed later when checking for repeat characters.

Values over \$C0 (192) are reduced by \$60 (96) and we are in a position to weed out characters that do not require any action. These will be values of \$21 to \$5F inclusive. This is the position that the first quote will end up in. These characters cause the flow to go back and get the next byte.

LINES 480–690: Our character value is stored again, as it may be different, in another register. We also store the line pointer (the γ register) and the x quote counter. The latter is stored because if pi is being changed, the x register could be zero; at other times it will always be one.

The next procedure is to see if there is more than one character of the same type consecutively in the program. The x register will be used as a counter and as it is one already, it is already initialized. The following bytes are gathered in and checked against the original value. The x register is increased until a byte of a different value is found.

The routine now splits up. Where there are two or more repeat characters, we jump ahead to CONT3 to put that number into ASCII.

Continuing along, the x value will be one but we will not print out the number one as it implied. Registers \$3F and \$40 are set to zero, which as we shall see shortly will hold the ASCII value for repeats.

The action taken is to check to see if the character we are coding is a space or not. We do not want to code single spaces as it would clutter the listing unnecessarily. On finding a space the flow jumps further ahead to reload the registers and go back to get the next byte. For characters not spaces, and all single characters, the routine branches forward to skip the next section.

LINES 700–770: On finding more than one of the same character we want to convert the number into ASCII format. We already have the number in the x register. To use our own conversion routine at \$847F we need to set the accumulator to zero, as the high byte value. The result will be in the input buffer with a zero after the last digit. As a line of BASIC program when typed into the 64 cannot be more than 80 characters, it therefore means that the number of repeats cannot be any greater. This means that the number of ASCII digits will be two at a maximum.

We therefore pick up the first two digits from the buffer and store
them in \$3F and \$40. If there was a single digit, that is 2 to 9 repeats, \$40 will be zero.

LINES 780–820: We said at the beginning that some characters would require the use of the data tables whilst some can be coded by calculation. These few lines divide up the flow into these two areas.

We load back the value achieved in earlier calculations (lines 410–440) to the accumulator. Values of \$60 and under, or \$7B and over, will branch off to Conversion B, which uses data tables.

CONVERSION A

LINES 830–1230: The first task to undertake is to subtract \$20 (32) from our value and store it. This is now the same value as the ASCII code of the letter of the key it shares. They will all be achieved using the shift with the key rather than the logo key.

The maximum number of characters we could insert is seven, two for the brackets, two for the number of repeats and three for the code. This number is placed into x. We check the 'repeat' digits storage for the number of numerals. A zero will indicate that there is no digit in that column. The x register will be decreased accordingly. Location \$3E has the number of graphic characters to be coded and this is compared with x. From this we either open-up the program, close-up the program or leave it unchanged. The memory move routines are described in Chapter 6.

Now we are ready to insert the code in the order of:

- i) The [bracket.
- ii) The number of repeats if applicable.
- iii) The letter G.
- iv) The symbol >.
- v) The letter of the key, held in location \$3D.
- vi) The] bracket.

Once completed, we load the quote counter back into x and jump back to get the next character to code.

CONVERSION B

LINES 1240–1640: This is where we have to use data tables to find the relevant code. This part is entered with the character value in the accumulator and is put into \$3D for later use. The first table we look up is the data address table. The start address of the table is placed in locations \$62 and \$63. The x register has the total number of characters catered for and the y register is used as a general pointer.

As the table is in the RAM under the BASIC ROM we have to disable that ROM, get the byte we want and then switch back the ROM. The byte is placed temporarily on the stack during the enabling of the ROM. The byte we have collected is compared for equality with our character value. Succeeding forces a branch forward. Failure means we continue the search. The first thing is to increase the v register three times. This will skip the address of the rejected character in the data table. The v value will be in line with the next character value. If the x register has been decreased to a value below zero, that is, *SFF*, then all the data address table has been checked and a match not found. There is one further chance. It could be a character which uses the CTRL key along with a letter key. These will have values no greater than *STA* (26). This is checked and if it does not fall in, then an 'OUT OF DATA' error is generated and coder is exited. We think that this should never come about as we believe we have catered for all eventualities.

Supposing a CTRL value is the one found, then we add \$40 (64). This simply gives the value of the letter on the key. This is stored immediately in the data table. The start address of the start of CTRL data is placed into \$62 and \$63.

Now back to the other characters. A match has been found in the data address table and we have arrived at line 1550. The two bytes next in line in the address table are the data address in the second table. These are placed into registers \$62 and \$63.

LINES 1650–2160: We have now finished with the data address table and concentrate on the data table itself. This time we only require one byte, the first byte, which will give us the number of bytes of code. To this value we add four, the brackets and the 'repeat' digits, transfer it to x and decrease it if one or both repeat digits are redundant. This final value is compared with the number of characters to be replaced to determine whether more or less space is required. This and the moves, if needed, are achieved in lines 1800 to 1850.

The insertion of data is the only thing left to do. We reload the line marker and start. The left square bracket and, if required, the repeat digits are stored first. The data insertion is slightly complicated. The line marker is stored and the Y register is re-initialized. The BASIC ROM is switched off and a byte is taken from the table. The Y register is stored in \$c2 and the line marker is restored and incremented. The byte is now inserted in the program. Now the line marker is stored and the data marker placed back in Y. This is compared with the number of bytes of data in the code (\$c2). If it has not reached this number, we branch back to get further bytes to insert. Once all the data has been collected and stored, the BASIC ROM is switched back in.

Finally, the right hand square bracket is inserted and we jump to get the next character to be coded, after restoring the quote counter.

LINES 2170–2190: This simply restores the line marker and quote counter, after which the routine goes back to get the next character. Single space characters, which are not coded, are sent here.

The data table – a program

After much thought, we have decided to supply the data tables for CODER in the form of a BASIC loader. This is mainly due to the fact that it is stored under the BASIC ROM which makes it hard to check and correct using a monitor. With the loader program we can put in a checksum which helps to see if you typed in the correct values.

A further item that the loader program does is to clear the area used by the KEY command (see Chapter 4) for its data. So type in the program, check it and save it.

```
10 L=41472:T=0
20 READD: IFD=-1THEN40
30 T=T+D:POKEL,D:L=L+1:GOT020
40 IFT()51131THENPRINT"[REV] DATA INCORR
ECT": END
50 FORL=41216T041471:POKEL.0:NEXT
60 PRINT"[REV] DATA LOADED":END
70 DATA3,87,72,84,2,67,68
80 DATA3,82,69,86,3,72,79
90 DATA77,3,82,69,68,2,67
100 DATA82,3,71,82,78,3,66
110 DATA76,85,3,83,80,67,3
120 DATA71,62,42,3,71,62,43
130 DATA3,71,60,45,3,71,62
140 DATA45,1,126,3,71,60,42
150 DATA3,79,82,71,2,70,49
160 DATA2,70,51,2,70,53,2
170 DATA70,55,2,70,50,2,70
180 DATA52,2,70,54,2,70,56
190 DATA3,66,76,75,2,67,85
200 DATA3,79,70,70,3,67,76
210 DATA83,3,73,78,83,3,66
220 DATA82,78,5,76,32,82,69
230 DATA68,3,71,82,49,3,71
240 DATA82,50,5,76,32,71,82
250 DATA78,5,76,32,66,76,85
260 DATA3,71,82,51,3,80,85
270 DATA82,2,67,76,3,89,69
280 DATA76,3,67,89,78,5,71
290 DATA62,83,80,67,3,71,60
300 DATA75,3,71,60,73,3,71
310 DATA60,84,3,71,60,64,3
320 DATA71,60,71,3,71,60,43
330 DATA3,71,60,77,3,71,60
340 DATA92,3,71,62,92,3,71
350 DATA60,78,3,71,60,81,3
```

```
360 DATA71,60,68,3,71,60,90
370 DATA3,71,60,83,3,71,60
380 DATA80,3,71,60,65,3,71
390 DATA60,69,3,71,60,82,3
400 DATA71,60,87,3,71,60,72
410 DATA3,71,60,74,3,71,60
420 DATA76,3,71,60,89,3,71
430 DATA60,85,3,71,60,79,3
440 DATA71,62,64,3,71,60,70
450 DATA3,71,60,67,3,71,60
460 DATA88,3,71,60,86,3,71
470 DATA60,66,5,67,84,82,76
480 DATA65,5,67,84,82,76,66
490 DATA5, 67, 84, 82, 76, 72, 5
500 DATA67,84,82,76,73,5,67
510 DATA84,82,76,78,5,67,82
520 DATA71,62,78,5,67,82,71
530 DATA62,77,3,68,69,76,2
540 DATA80,73,255,0,0,0,0
550 DATA5,0,162,17,4,162,18
560 DATA7,162,19,11,162,28,15
570 DATA162,29,19,162,30,22,162
580 DATA31,26,162,32,30,162,96
590 DATA34,162,123,38,162,124,42
600 DATA162,125,46,162,126,50,162
610 DATA127,52,162,129,56,162,133
620 DATA60,162,134,63,162,135,66
630 DATA162,136,69,162,137,72,162
640 DATA138,75,162,139,78,162,140
650 DATA81,162,144,84,162,145,88
660 DATA162,146,91,162,147,95,162
670 DATA148,99,162,149,103,162,150
680 DATA107,162,151,113,162,152,117
690 DATA162,153,121,162,154,127,162
700 DATA155,133,162,156,137,162,157
710 DATA141,162,158,144,162,159,148
720 DATA162,160,152,162,161,158,162
730 DATA162,162,162,163,166,162,164
740 DATA170,162,165,174,162,166,178
750 DATA162,167,182,162,168,186,162
760 DATA169,190,162,170,194,162,171
770 DATA198,162,172,202,162,173,206
780 DATA162,174,210,162,175,214,162
790 DATA176,218,162,177,222,162,178
800 DATA226,162,179,230,162,180,234
810 DATA162,181,238,162,182,242,162
```

```
B20 DATA183,246,162,184,250,162,185
B30 DATA254,162,186,2,163,187,6
B40 DATA163,188,10,163,189,14,163
B50 DATA190,18,163,191,22,163,1
B60 DATA26,163,2,32,163,8,38
B70 DATA163,9,44,163,14,50,163
B80 DATA142,56,163,141,62,163,20
B90 DATA68,163,255,72,163,5,67
900 DATA84,82,76,67,-1
```

SAVING THE DATA AREA

The following listing will save the area we have used for both KEY and CODER routines. The saving of data through machine code is described in the *Programmer's Reference Guide*. The only extra coding is to switch the BASIC ROM out, so that we will save our data and not the BASIC interpreter. You could use this after setting up the function keys (see Chapter 4) so on reloading, the data is there and ready.

10		LDX	#\$08	ļ	DEVICE NO (TAPE=1)
20		LDA	#\$01	!	LOGICAL FILE NO
30		LDY	#\$FF	!	ND SEC ADDRESS
40		JSR	\$FFBA	ļ	SETLFS
50		LDA	#\$0C	!	CHARS IN FILENAME
60		LDX	#KNAME	ļ	LOW ADDRESS OF NAME
70		LDY	#>NAME	ļ	HIGH ADDRESS
80		JSR	\$FFBD	!	SETNAM
90		LDA	\$01	!	SWITCH OFF BASIC
100		AND	#\$FE		
110		STA	\$01		
120		LDA	#\$00	!	STORE START ADDRESS
130		STA	\$FB		
140		STA	#\$A1		
150		STA	\$FC		
160		LDX	#\$49	!	LOW END OF SAVE
170		LDY	#\$A4	!	HIGH END OF SAVE
180		LDA	#\$FB	!	LOCATION OF START ADD
190		JSR	\$FFD8	!	SAVE
200		LDA	\$01	!	SWITCH IN BASIC
210		ORA	#\$01		
220		STA	\$01		
230		RTS			
240	NAME	TXT	"UTILITY	ſ [DATA"

RELOCATING THE DATA TABLES

If you relocate CODER may also want to relocate the data. Here is one suggested way. Using the BASIC loader program for the data, change the value of L in line 10 to the new data start address. The data normally starts at \$A200 (41472) but the data address table starts at \$A350 (41808). From this calculate the data address table new address and put its value in lines 10 and 30 of the routine below. The end of the data address table is normally \$A442 (42050), so work out its new end, subtract one, and this is put in lines 210 and 240. The difference between the old address and the new address should be put in lines 80 and 120. The routine below is for a new table at a higher address; for one lower, change the addition to subtraction and set the carry instead of clearing it.

10		LDA	#\$50	ļ	START OF DATA
					ADDRESS TABLE
20		STA	\$14		
30		LDA	#\$A3		
40		STA	\$15		
50	NEXT	LDY	#\$01	1	POINTER
30		LDA	(\$14),Y	!	LOW ADD IN TABLE
70		CLC			
80		ADC	#\$60	ţ	ADD LOW DIFFERENCE
90		STA	(\$14),Y		
100		INY			
110		LDA	(\$14),Y	!	HIGH ADD IN TABLE
120		ADC	#\$25	!	ADD HIGH DIFFERENCE
130		STA	(\$1 4),Y		
140		CLC			
150		LDA	\$14	!	UPDATE TABLE ADDRESS
160		ADC	#\$03		
170		STA	\$14		
180		LDA	\$15		
190		ADC	#\$00		
200		STA	\$15		
210		CMP	#\$ A4	ļ	END OF DATA
					ADDRESS TABLE
220		BNE	NEXT	!	NO
230		LDA	\$14		
240		CMP	#\$43		
250		BNE	NEXT	ļ	NO
260		RTS			

```
Old
```

COMMAND SYNTAX OLD

There are no parameters with this command.

There are four ways to 'lose' a BASIC program. The first way is by switching off and then there is absolutely no way of recovering it. Another two ways are by doing a system cold start or a reset. This is as if you have just switched on but retaining data held by the RAMS. This can be achieved by typing SYS64738 or by a reset button, if you have fitted one. The final way to lose a program is by issuing the BASIC command NEW.

To lose a program the operating system of the 64 sets the first two bytes of the BASIC program area to zero. This would normally be \$0801 and \$0802 (2049 and 2050) and would be the link address in the first line of a BASIC program. This means that as far as BASIC is concerned no program is present as it would encounter the zeros straight away.

Now as long as no further lines of BASIC are typed in, we can reverse the process, but will lose all the variables. The way it is done is made clear in the description of the coding.

9	*= \$8415				
10		LDA	#\$FF		
20		LDY	#\$01		
30		STA	(\$2B),Y	!	PUT ANY LINK IN 1ST LINE
40		JSR	\$A533	ļ	RECHAIN LINES
50		LDA	\$22		
60		CLC			
70		ADC	#\$02		
80		STA	\$2D	!	SET END OF PROG ADDRESS -LOW
90		LDA	\$23		
100		ADC	#\$00	1	IN CASE CARRY WAS SET IN 70
110		STA	\$2E	!	SET END OF PROG ADDRESS -HIGH
120		JMP	\$A660	!	PERFORM CLR

ASSEMBLY LISTING

LINES 10–40: If we change the first two bytes from zero, BASIC will no longer think it is at the end of the program. We put \$FF in those, and get the address from the start of BASIC variables in \$2B and \$2C. Now a call to the ROM routine RECHAIN LINES will achieve two things. First, it will

correctly set the link address in the first line, and secondly, we will be able to set the end of program variables.

LINES 50–120: Locations \$22 and \$23 are set to the beginning of the two zero bytes, which mark the end of the program, when the RECHAIN routine is finished. By adding two to those, we have the end of program and can set the respective registers, \$2D and \$2E.

The final thing is a jump to the CLR routine. This will set all the variable addresses to coincide with the recovered program. The BASIC program is now restored to its original state.

Dump

COMMAND SYNTAX

DUMP

There are no parameters with this command. It will also only operate in direct mode. If used within a program it will just skip out of the command. Hitting STOP will break out of DUMP and allows direct editing; typing CONT will resume at the break point in the BASIC program. Holding down any other key, apart from shift, will halt the routine until it is released.

The action of DUMP is identical to the BASIC subroutine in Chapter 5 except that as the routine is in machine code it does not add to the simple BASIC variables. The logic closely follows the BASIC routine. Output may again be directed to a printer by an OPEN and CMD sequence. The major departure is in the use of one or two ROM routines to carry out the mathematical conversions and convert the number to an ASCII string to be printed.

ASSEMBLY LISTING

9	*=\$8E52	
10	INOTE REAL ASC /	ASC DR 0
20	! STRING ASC /	ASC+128 OR 128
30	! INTEGER ASC+128 /	ASC+128 OR 128
40	! FUNCTION ASC+128 /	ASC OR 0
50	!	
30	!	
70	TEST FOR DIRECT MODE	
80	1	
90	LDA \$9D	!MSGFLG
100	CMP #\$80	IDIRECT ???
110	BEQ DIRECT	
120	RTS	PROGRAM MODE SO ABANDON
130	f	
140	INOTE CURRENT VARIABLE	IN FILE NAME POINTER

150 ! 160 DIRECT VARTAB LDA \$2D STA \$BB 170 LDA \$2E !VARTAB+1 180 190 STA ≸BC 200 ! 210 !RETURN TO HERE TO SEE IF ALL DONE 220 START LDA \$BC 230 CMP \$30 BNE CONT MORE VARIABLES 240 250 LDA \$BB CMP \$2F 260 BNE CONT 270 IDONE ALL OR WERE RTS 280 NONE TO START 290 ! 300 !FIND VARIABLE TYPE 310 ! LDY #\$00 320 CONT LDA (\$BB),Y !FIRST CHAR OF NAME 330 CMP #\$80 340 BCS INTER ! INTEGER OR FUNCTION 350 360 ! 370 !STRINGS AND REAL 375 ! JSR \$FFD2 !OUTPUT ASCII CHAR 380 INY 390 LDA (\$BB),Y !GET SECOND CHAR 400 CMP #\$7F 410 BCS STRING 420 430 ! 440 !REAL VARIABLE 450 ! JSR \$FFD2 !OUTPUT SECOND ASCII 460 CHAR JSR UPDATE !PRINT '=' AND 470 HPDATE POINTER POINT TO VARIABLE LDA \$BB 480 FOR MEMORY MOVE LDY \$BC !TO FPACC#1 490 GO DO IT JSR \$BBA2 500 JSR \$BDDD 'FPACC#1 TO STRING 510 AT \$0100 PRINT IT JSR \$AB1E 520 LDA #\$FF 530 BNE NEXT !SKIP TO NEXT VARIABLE 540

550	!		
560	STRING V	JARIABLE	
570	ļ		
580	STRING	AND #\$7F	MAKE ASCII
590		JSR \$FED2	
600		1 DA #\$24	1/4/
610		ICD CEENO	. ₽
629		JCD HDDATE	
430			
X40		100 40000	
650		1 DV #400	
550		LD1 #900 IDA (400) N	<i>,</i>
470		TAV	LIENCTH OF CTRING IN M
4070			LENGTH OF STRING IN X
400			INULL STRING
700		1001 1000 (#000) N	,
710		CDH (PDD), I	
720		51H #22 TNV	LSB LUCATION
720			,
740		CTA 400	
750		51F1 #25 1DV #4400	TISE LUCATION
7.50	CHAR	LDI #200 IDA (400) \	,
770	CIMAN	LUH (\$22);1	
780		JOK PFFUZ	
790			
800			
810	QUOTE		
820	40072	JSR \$FED2	
830		BED NEXT	
840		BNE NEXT	
850	HALFSTART	BCS START	
860	!		
870	INTEGER	AND FUNCTIONS	
880	!		
890	INTEN	AND #\$7F	CONVERT TO ASCIT
900		JSR &FED2	CONVERT TO HOUT
910		INY	
920		IDA (\$RR) Y	
930		CMP #\$7F	
940		BCS INT	INTEGER IE SECOND
			CHAR)128
945	1		
950	FUNCTION	DEFINITION	
960	!		
970		JSR \$FFD2	
980		JSR UPDATE	

```
!'F' PRINT FN
 990
                 LDA #$46
                 JSR $FFD2
1000
                 LDA #$4E
                             12'N2
1010
                 JSR $FFD2
1020
                 BNE NEXT
1030
1040 !
1050 !INTEGER VARIABLE
1060 !
                 AND #$7F !CONVERT TO ASCII
1070 INT
1080
                 JSR $FFD2
                          1111
1090
                 LDA #$25
1100
                 JSR $FFD2
1110
                 JSR UPDATE
                 LDY #$00
1120
                 LDA ($BB),Y !MSB
1130
                 STA $62 !FPACC#1
1140
1150
                 INY
                 LDA ($BB),Y !LSB
1160
                 STA $63
1170
                 LDX #$90
1180
                             CONVERT TWO BYTE
                 JSR $BC44
1190
                              INT TO REAL
                              !FPACC#1 TO STRING
1200
                 JSR $BDDD
                              AT $0100
1210
                 JSR $AB1E !PRINT IT
1220
1230 !SEE IF KEY HIT AND UPDATE POINTERS
1240 1
                            CARRIAGE RETURN
1250 NEXT
                 LDA #$0D
                 JSR $FFD2
1260
                 CLC
1270
                 LDA $BB
1280
                 ADC #$05
1290
                 STA $BB
1300
                 BCC WAIT
1310
                INC $BC
1320
                JSR $FFE4
                              !STOP
1330 WAIT
                 JSR $FFE1
1340
                 BNE NOT
1350
                 RTS
1360
                 LDA $CB
                             !CURRENT KEY
1370 NOT
                              !NO KEY=64
                 CMP #$40
1380
                 BNE WAIT
                              CYCLE WHILE KEY
1390
                              HELD DOWN
1400
                 SEC
1410
                 BCS HALFSTART
```

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190 Programming aid routines

1420	1				
1430	!PRINT '='	AND SET	POINT	ER IN	↓ \$BB/BC
	TO START	OF VARIA	BLE		
1440	!				
1450	UPDATE	LDA #\$:	3D	! '='	
1460		JSR \$F	FD2		
1470		CLC			
1480		LDA \$B	3		
1490		ADC #\$	82		
1500		STA \$B	3		
1510		BCC RE	TURN		
1520		INC \$B	2		
1530	RETURN	RTS			
SERC	CHAR	8E A P			

		OFOF	CONS
8E59	DIRECT	SECE	HALF
8EEE	INT	8ED0	INTFN
8F11	NEXT	8F2A	NOT
8EC5	QUOTE	8F43	RETURN
8E31	START	8E97	STRING
8F33	UPDATE	8F21	WAIT

The routine has been written to be easily relocatable. The only change necessary is to alter all JSR UPDATES to JSR (start address + \$E1).

The ROM routines used are as follows:

CHROUT (\$FFD2)

A full description of this function is given in the *Programmer's Reference Guide*, 'The KERNAL B-5'. It outputs the contents of A as an ASCII character to the screen.

STOP (\$FFE1)

See 'The KERNAL B-33'. Test for the stop key. UDTIM must be called before using this routine.

UDTIM (\$FFEA)

See the *Programmer's Reference Guide,* 'The KERNAL B–36'. This updates the system clock.

MEMORY TO FAC#1 (\$BBA2)

This routine takes a five byte real number and moves it to the floating point accumulator #1. *En route* the sign bit of the mantissa is stripped off and the sign register FACSGN (\$66) set, the exponent put at FACEXP (\$61) and the mantissa of FACHO (\$62–65). On entry A must hold the low and Y the high byte of the address of the bytes to be moved.

FAC#1 TO STRING (\$BDDD)

Converts FAC#1 to an ASCII string stored at the bottom of the stack (\$0100). On exit A holds #\$00 and Y holds #\$01.

PRINT STRING FROM MEMORY (\$AB1E)

This routine prints successive characters starting at the memory location whose address is held in A (low) and Y (high). The routine continues until a zero terminator is found (as will be the case at the bottom stack in this application). Note A and Y already hold the start address on exiting the previous routine and need not be changed.

EVALUATE TWO BYTE SIGNED INTEGER (\$BC44)

Evaluates a two byte signed integer held in FAC#1 and deposits the result in floating point form back in FAC#1. Before calling x must be set to #\$90, FACHO must hold the high and FACHO+1 the low byte of the integer (remember integers are held in high/low format unlike addresses). Once in this form the same routines as for real may be used to convert to ASCII and print.

LINES 1450–1530: JSR UPDATE

This will be used by all types of variables. It will be used directly after the variable name has been printed. All this does is to print out the equals sign and increase the address registers by two, so that they will point directly to the next byte to be collected – the first of the actual variable.

LINES 90–120: These check for direct mode. If program mode is discovered, then the routine is exited.

LINES 160–190: The locations we are going to use to step through the variable area are initialized with the start of variable address, which also happens to be the end of BASIC program address. We are now ready to start.

LINES 220–280: Locations \$2F and \$30 are the address of the end of the variable block that we are going to DUMP. By checking the values in those with our registers, we can find out if we have completed all.

LINES 320–350: The routine is divided up here and will be further divided later. When the first byte is picked up, we check if it is an integer or a function by seeing if the value is \$80 (128) or over. These are dealt with further into the routine.

LINES 380–420: The first byte we have already is printed – the first letter of the variable name. The next byte is collected and this will distinguish between real and string variables.

LINES 460–540: REAL VARIABLES: Again we print what we have in A, making the whole variable name output. A call is now made to JSR UPDATE. With the address of the present position placed in the A and Y

registers, we call three ROM routines, described at the start, to print out the variable to the screen or output device. The accumulator is loaded with \$FF just so the branch following will succeed.

LINES 580–840: STRING VARIABLES: Before we print out the accumulator, we remove the negative bit, bit 7, so that it is the pure ASCII code of the variable letter. The dollar sign is printed and UPDATE is called. As there are no separaters between stored strings, we cannot use a similar approach to the one used in real variables. The first thing printed is the start quotes. The length of the string is the byte after the name and this is placed in the x register. Now we gather in the address of where the string is stored. From this we can print out the characters directly, decreasing x each time, until the counter is zero. Finally, the closing quotes are output. One of the following two branches must succeed so we can continue.

LINES 890–940: Integers and functions start here. In these lines we distinguish between them and act accordingly. By stripping off the negative bit we can print out the first character of the name, and do so. The second byte is loaded and this will tell us what type to deal with. A value of \$80 (128) or over will signify integer.

LINES 970–1030: FUNCTION VARIABLE: We cannot print any value for the function, so after the second name character is printed, UPDATE is called, and then we just print the letters 'F' and 'N' and branch off.

LINES 1070–1210: INTEGER VARIABLES: After printing the second character of the name the integer sign of '%' is output. Once more UPDATE is visited. The next two bytes in the variable area are the integer value and these are transferred to FAC #1. With the x register set to normalize the result, we call a ROM routine at \$BC44. This will convert the integer value to a real number. We then convert to ASCII and print the result with the routine described at the beginning.

LINES 1250–1410: After dealing with any of the four types of variables, the flow is directed to this part of the routine. The return character is printed so that the next variable is printed on a new line. Each variable takes seven bytes of memory and as we added two bytes to our address registers in UPDATE, we only need to add five more to get to the start of the next variable in line.

The STOP key is now tested for and if the negative flag is set then DUMP is ended, as STOP was pressed. By examining \$CB we can see if any key is being pressed. As long as a key is held down we loop around here and then check for STOP.

To continue with DUMP we set the carry flag and we use a Branch with Carry Set to line 850 where the same happens, going further back to proceed dumping variables.

IMPROVEMENTS?

Obvious improvements are as for the BASIC subroutine. If you had trouble extending the BASIC subroutine to handle arrays, you haven't tried anything yet! Most dump routines (wisely) do not handle subscripted variables (probably because it is considered too difficult).

Trace and Troff

COMMAND SYNTAX TRACE and TROFF

Constant Constant
Speed Control
(n' recet single sten

'1' a relative delay of 2^0
'3' a relative delay of 2^2
$^{\prime}5'$ a relative delay of 2^4
$'7'$ a relative delay of 2^6
'9' a relative delay of 2^8

The space bar operates the single-stepping.

The delay is in addition to the normal time taken by BASIC to move to a new line and execute the common trace code. The delay may be changed at any time by hitting the appropriate key. It is, however, not possible to break into program execution in single step. If you wish to do this, hit a number other than 0 first.

TRACE is a diagnostic aid which provides useful information on the path taken through a BASIC program. In this particular version the previous and the current line numbers are displayed in reverse video at the top right of the screen.

We considered it far more important to allow the user to be able to vary the speed of the trace and have single-stepping capability. When called for the first time the default will be to single stepping and thereafter at each run it will continue at the last set speed. After being disabled with TROFF it will, on being enabled, revert back to single stepping. In single step mode program execution halts until the space bar or a speed change key is pressed. The keys for speed change are given above.

ASSEMBLY LISTING

9	*=8D3A						
100	!TRACE	ENABLE					
110	!						
120	!						
130	ENABLE	LDA	\$9D	!MSGFLG	CHECK	FOR	DIRECT

140		BEQ	PMODE	!\$00=PROG \$80=DIRECT
150		SEI		!OK-NOT IN PROG MODE
				SO DISABLE
130		LDA	#\$FF	INTERRUPT AND SET
				SINGLE STEP
170		STA	SSTEP	TO SEE FOR SINGLE STEP
180		I DA	#\$FF	DO SAME FOR TRACE FLAG
190		STA	TRELAG	I
200		LDX	\$0308	LIGONE GET LOW BYTE
210		STX	IGONE	IDISPATCH AND STORE AT
220			NEUL	INFLUTINE VECTOR LOW WHICH
238		STA	40308	POINTS TO TRACE
240			40300	ISAME END HIGH
250		CTV	TCONEL	t shale i ok hi oh
230				1
200		CTA	MLVH #0000	
270	DMODE	CIT	₽0307	IDECET INTERDURT
200	PHODE			IAND DETUDN TO DACIC
270		KI2		AND RETURN TO BASIC
300				
310	! I KUFF = I KI	ALE I	DISABLE	
320				
330	DISABLE	SET		REVERSE ENABLE PRUCESS
340		LDA	\$9D	CHECKING IN DIRECT MODE
300		REN	PMUDE	
360		LUA	#\$00	DISABLE TRACE FLAG
370		SIA	TRELAG	RESET TUKEN DISPATCH
380		LDA	IGUNE	TU VALUES AT THE TIME
390		STA	\$0308	!OF CALLING
400		LDA	IGUNE+	1
410		STA	\$0309	
420		CLI		
430		RTS		'BACK TO BASIC WITH TRACE OFF
440	!			
450	PERFORM TR	ACE	IGONE	POINTS HERE
460	1			
470	TRACE	STA	AREG	IF BASIC IS TO RESUME
480		PHP		THEN WE MUST SAVE A,X,Y
490		STX	XREG	AND STATUS FLAGS
500		STY	YREG	TO RESTORE THEM
510	!			
520 530	! ONLY PROCE	ED W	ITH TRA	ACE IF A PROGRAM RUNNING

540 LDA \$9D 550 BEQ RUNNING 560 ! 570 !RESTORE ENTRY VALUES BEFORE CONTINUING 580 ! 590 BASIC LDA AREG **!REVERSE ENTRY PROCESS** 600 LDY YREG **!TO ALLOW PROG** TO CONTINUE 610 LDX XREG !UNCORRUPTED 620 PLP !DON'T FORGET FLAGS!!! 630 JMP (IGONE) CONTINUE AT TOKEN DISPATCH 640 1 350 PROGRAM RUNNING SO CHECK IF TRACE ENABLED 630 !FROM TRACE FLAG = \$FF??? 670 ! 680 RUNNING LDA TRFLAG 690 BEQ BASIC !TRACE OFF SO RESTORE AND CONT 700 ! 710 !TRACE IS ON SO UPDATE DISPLAY 720 ! 730 SEC !READ CURSOR POSITION 740 JSR \$FFF0 !AND SAVE BY CALLING PLOT STX ROW 750 WITH CARRY SET 760 STY COL 770 CLC SET CURSOR POSITION 780 LDX #\$00 TO ROW 0 COLUMN 24 790 LDY #\$18 800 JSR ≸FFF0 810 LDX #\$0F 820 SPACE LDA #\$20 **!CLEAR PREVIOUS NUMBERS** 830 JSR \$FFD2 840 DEX 850 BINE SPACE ISET BACK TO ROW 0 COL 24 860 CLC 870 LDX #\$00 LDY #\$18 880 890 JSR \$FFF0 900 LDA #\$12 **!TURN ON REVERSE VIDEO** 910 JSR \$FFD2 920 LDA OLHIGH !LOW BYTE PREVIOUS LINE 930 LDX OLLOW !HIGH 940 JSR \$BDCD !PRINT LINE NUMBER 950 LDA #\$92 !REVERSE OFF JSR \$FFD2 930

970 LDA #\$20 BIT OF SPACE JSR \$FFD2 980 LDA #\$12 !REPEAT FOR CURRENT LINE 990 JSR \$FFD2 !GETTING ITS VALUES FROM 1000 LDA \$3A CURLIN LOW BYTE 1010 STA OLHIGH !NOW BECOMES OLD LINE 1020 1030 LDX \$39 !CURLIN+1 1040 STX OLLOW JSR ≸BDCD 1050 LDA #\$92 1030 1070 JSR ≸FFD2 1080 CLC PREPARE TO RESET CURSOR 1085 ! 1090 !IGNORE THIS BIT AS ONLY TO ALLOW BRANCH TO WORK 1095 ! 1100 BASIC1 BCS BASIC !HALFWAY BRANCH TO BASIC 1110 LDX ROW CONTINUE RESET CURSOR 1120 LDY COL 1130 JSR \$FFF0 !RESTORE CURSOR POSITION 1140 ! 1150 !CHECK FOR ANY KEYS PRESSED 1160 ! 1170 JSR \$FFE4 !GETIN 1180 CHCHAR BEQ SINGLE !NOTHING IN K/B BUFFER CMP #\$2F !KEY<0??? 1190 1200 BCC SINGLE !YES THEN OF NO INTEREST 1210 CMP #\$3A !KEY>9??? BCS SINGLE !YES - NO INTEREST 1220 1230 SBC #\$30 !BETWEEN 0 AND 9 SO -\$30 1240 BNE CHDELAY!1-9 1250 ! 1260 10 PRESSED SO RESET SINGLE STEP 1270 ! 1280 LDA #\$FF 1290 STA SSTEP 1300 BNE SINGLE INO NEED TO CALC DELAY 1310 ! 1320 !CALCULATE DELAY AS POWERS OF 2 1330 ! 1340 CHDELAY TAX !PUT 1-9 IN X 1350 SEC 11 IN CARRY 1360 LDA #\$00 1370 ROLL ROL A MOVE CARRY BIT X TIMES 1380 DEX !TO SET KEY-2 BIT BNE ROLL ! TO GIVE DOUBLING DELAY 1390 1400 STA COUNT ISTORE IT TO USE AS TIMER

1410 LDA #\$00 !DISABLE SINGLE STEP 1420 STA SSTEP BEQ DELAY !ALWAYS TAKEN 1430 1440 ! 1450 !SINGLE STEP PAUSE 1460 ! 1470 SINGLE LDX SSTEP !IS IT ON BEQ DELAY ! IF NOT GO TO DELAY 1480 CMP #\$20 !SPACE HIT ORIGINALLY??? 1490 1500 BEQ BASIC2 !YES THEN PERFORM LINE 1510 SSLOOP JSR \$FFE4 !GETIN WAIT FOR A CHAR 1520 BEQ SSLOOP !AND KEEP WAITING 1530 CMP #\$20 !SPACE??? BEQ BASIC2 !YES THEN SKIP DELAY 1540 1550 BNE CHCHAR !NO - WAS IT A SPEED CHANGE 1560 ! 1570 !TIMER COUNTDOWN FOR DELAY 1580 ! 1590 DELAY LDX COUNT !DO COUNT LOTS 1600 DL00P1 LDY #\$FF !OF 2561S 1610 DL00P2 DEY 1620 BNE DL00P2 1630 DEX 1640 BNE DLOOP1 1650 ! 1660 ! GUARANTEED BRANCH TO HALFWAY BACK TO DISPATCHING LINE 1100 1370 ! 1680 BASIC2 SEC !ENSURES BRANCH BASIC1 1690 BEQ BASICI 12 FLAG SET ALWAYS SET HERE 1700 ! 1710 !RESERVE TEMPORARY STORES AND FLAGS SET TO DEFAULTS 1720 ! 1730 TRFLAG BYT \$00 **!TRACE FLAG OFF** BYT ≸FF 1740 SSTEP SINGLE STEP ON 1750 COUNT BYT \$00 IND DELAY 1760 AREG BYT \$00 A ON ENTERING FROM BASIC 1770 XREG BYT \$00 !X " 1780 YREG BYT \$00 !Y * 1790 COL BYT \$00 CURSOR WHILE LINES PRINTED 1800 ROW BYT \$00 1810 OLLOW BYT \$00 PREVIOUS LINE LOW BYT \$00 1820 OLHIGH ı HIGH

1830 IGONE	BYT \$00,\$00	!STORE	FOR	ORIG	VECTOR
1840	END	!A JMP	TO I	HERE	

8E49	AREG	8D87	BASIC
8DEF	BASIC1	8E43	BASIC2
8DFD	CHCHAR	3E12	CHDELAY
8E4C	COL	8E48	COUNT
8E38	DELAY	8D61	DISABLE
8E3B	DL00P1	8E3D	DL00P2
8D3A	ENABLE	8E50	IGONE
8E4F	OLHIGH	8E4E	OLLOW
8D5F	PMODE	8E16	ROLL
8E4D	ROW	8D94	RUNNING
8E24	SINGLE	8D8F	SPACE
8E2D	SSLOOP	8E47	SSTEP
8D97	TRACE	8E46	TRFLAG
8E4A	XREG	8E4B	YREG

The ROM routines used are as follows:

CHROUT (\$FFD2) As DUMP (see page 190).

GETIN (\$FFE4)

See *Programmer's Reference Guide*, 'The KERNAL function B–11'. This removes one character from the current input device (usually the keyboard buffer) and returns its ASCII value in A. Zero is returned if none waiting.

PLOT (\$FFFØ)

See Programmer's Reference Guide, 'The KERNAL function B–19'. Reads the cursor position with the carry set and positions the cursor when the carry is clear. Misleadingly, x is used for the row and Y for the column.

PRINT LINE NUMBER (\$BDCD)

Useful little routine, this one, and well worth noting. Not only can it be used for line numbers, but also for a two byte unsigned integer (\$0000 to \$FFFF). Before calling it, x must hold the low and A the high byte. It also strips off the traditional leading and trailing spaces before printing.

HANDLE NEW LINE (\$A734)

This routine is vectored by the page 3 vector IGONE (\$0308) and \$A7E4 is the 64's default setting. This is BASIC's token DISPATCH routine and is covered in great detail in Chapter 3. When used with the UTILITY, IGONE has been modified and hence the reason why IGONE has been first read and

stored. Doing it this way means the routine will work with or without the UTILITY. IGONE is called to tokenize each new line and is thus the ideal point at which to patch our trace.

LINE 130–290: TRACE ENABLE: These set up TRACE ready for when you RUN a BASIC program. A scan is made for direct operation only, and only if it is direct do we continue. During this initialization the interrupt will be disabled. The default speed is single step and its value is stored in the appropriate location at the end of the routine. The original value of IGONE, BASIC Character Dispatch Vector (see Chapter 6), are stored for safe-keeping and the start of TRACE replaces them. After clearing the interrupt, we return you to BASIC until the RUN command is issued.

LINES 330–290: TROFF – Trace Disabled: The reverse of the TRACE set up procedure.

LINES 470–690: The BASIC dispatch is used each time a BASIC command is issued whether in direct or program mode. This means that the routine can be called when not required. To avoid this, we check for program mode, after preserving the processing registers, as in the set up. If still in direct mode, we restore the registers and jump to the normal DISPATCH routine. A final check is made before operating TRACE to ensure it is enabled by looking at the TRFLAG at the end of the routine.

LINES 730–1130: Display and updating are the purpose of these lines. We print at the top of the screen so as not to disrupt your display. We locate and save the current cursor position before setting it to the start of our print, top row and column 24, and clear the area we use by printing 15 spaces to the end of the line. After turning on reverse video, we gather in the values of the previous line number and visit the ROM routine to print it. To distinguish between the line numbers, we put a space between them, after turning off the reverse video. We now repeat the operation for the current line number. At the same time as getting the current line number, we store its value in our previous line store. As TRACE is called before every BASIC command is initiated, then when more than one command is on a line the previous and the present line numbers will be identical.

The instructions in line 1080 and 1100 are little tricks. Clearing the carry will ensure that the branch will fail. The branch is there for a later instruction when it will save a JUMP command.

Finally, we restore the original cursor position.

LINES 1170–1430: Is a speed change required? To find out, we use the KERNAL routine GETIN. If there is no character or it is not between the ASCII values of \$30 and \$3A, we branch out of this section. On remaining here, we subtract \$30 from the character value to convert from ASCII to a real value between 0 and 9.

The zero value signifies single-step mode is required again, so its flag is set to \$FF.

The remainder have to be acted upon in order to gain our power of two figures for delay purposes. The value we have is to be used as a counter and so is transferred to the x register. To start with, we set the carry and initialize the accumulator to zero. The accumulator is rotated right x times. The first time, the bit set in the carry is transferred to bit zero of the accumulator. The carry from this time on will be unset, except for the last time when nine is the value of x and that will not worry us. Every consecutive rotate will shift the set bit further to the right and so increase the power of two of the value. When nine is raised to a power the accumulator will end up with a value of zero, but when the delay is explained, lines 1590–1090, you will see that this is in effect 256, that is, 2[°]8. Finally, we disable the single-step mode, if set, by storing zero in its flag.

LINES 1470–1530: SINGLE-STEP MODE: To check if it is in operation, we test its flag. We not only look for the space character, which operates single-stepping, but also for others by branching back to lines 1170 to 1430. This means to continue in single-step the space bar or a numeral key will in fact cause the program to continue, the latter ones ending single-step at the same time.

LINES 1590–1620: DELAY: This consists of a loop within a loop. The inner loop is completed the number of times calculated earlier, thereby giving variable time delays. When speed nine has been selected COUNT is zero, the inner loop is carried out, and the COUNT is decreased before checking for zero. When zero is decreased it becomes \$FF (255) so the check will fail until it is decreased to zero once more. This therefore operates the inner loop 256 times.

Finally, we set the carry and branch back, as the accumulator will be zero, to line 1100 where the Branch with Carry Set will send it further back to jump in to the normal IGONE routine to carry out the BASIC command.

IMPROVEMENTS?

It is possible to modify the trace to list not only the previous and current line, but also to highlight the current statement being executed.

To list a line we can use the LIST routine in ROM which starts at \$A69C. There are two major problems if we try to use it. The first is that LIST uses a number of zero page locations also used during a run. The second is that on completion of LIST a warm boot of BASIC is carried out. (Try putting LIST in a program and running it.) We can overcome the first by copying zero page to elsewhere in RAM before calling LIST. The second requires that on return from LIST, we must re-enter our TRACE

routine at the next instruction after performing LIST which must restore zero page. To do this we must read and store the warm start vector IMAIN (\$0302) and set it to the next instruction after LIST is called. TROFF must, of course, reset this vector to its original value.

To highlight the statement within a listed line places even greater demands on our ingenuity and would require the TRACE routine to be rewritten to use CHRGET which has purposely been avoided (because of DOS 5.1). If it were to use CHRGET, the line could be re-listed each time with a marker character printed at the current byte held in TXTPTR (\$7A/\$7B) through the use of the PRINT tokens link.

Both additions seem of little point as we can use the STOP key followed by LIST line number(s). We have not even attempted to incorporate either possibility.

Numeric conversions

In the world of the Commodore 64 we come across three main numbering systems: that of decimal, to the base of 10, hexadecimal, to the base of 16, and binary which is to the base of 2 (octal, to the base eight is less common).

The binary number system is used because there are only two numerals used: 0 and 1. This matches the type of electronics used in the computer world, digital electronics, which has only two states, either on or off. These two positions are known as logical states. Logic 1 is on and obviously logic 0 is off. These, as you can see, fit well with the binary system.

The hexadecimal system was introduced because although binary matches the electronics, it is unwieldly and is not so easy to recognize in everyday form. Hexadecimal is easier to remember, using only two digits to the binary eight, and therefore faster to type in. Hexadecimal is nearly always entered in groups of two for example, \$FF.

Decimal is used in our everyday life and is therefore used in BASIC. One of its disadvantages is that numbers have varying amounts of digits. For instance, in numbers up to 255 there are one, two or three digits whereas with hex there is always two.

Some BASICS give you the option to enter numbers in forms other than decimal. The Commodore 64 BASIC does not. We are not going to rectify this but are giving you four conversion routines. These are converting decimal and binary, and decimal and hex.

TEN – Decimal to hexadecimal conversion COMMAND SYNTAX

TEN decimal number [,decimal number,....]

The maximum decimal number that can be converted is 65535 and then only positive numbers can be converted. Multiple conversions can be 202 Programming aid routines

done if they are separated by commas. The result will be a four digit hex number.

ASSEMBLY LISTING

*=\$84A0				
START	BEQ	SYNTAX		
	BCS	SYNTAX		
	JSR	\$81F5	!	GET PARAMETER
	LDA	#\$20		
	JSR	\$FFD2	ļ	PRINT A SPACE ON SCREEN
	LDY	#\$02		
NEXT	JSR	HEX	!	CONVERT A BYTE TO HEX
	DEY			
	BNE	NEXT	!	TWO CONVERTS FOR EACH BYTE
	LDY	\$14		
	STY	\$15	ļ	PLACE LOW BYTE FOR CONVERSION
	LDY	#\$02		
NEXT1	JSR	HEX		
	DEY			
	BNE	NEXT1		
	JSR	\$0079	ļ	GET LAST BYTE OF BUFFER AGAIN
	CMP	#\$ 2C	!	COMMA FOR MORE NUMBERS
	BNE	EXIT	ļ	NO MORE TO CONVERT
	JSR	\$FFD2	ļ	PRINT COMMA AS SPACER
	JSR	\$0073		
	JMP	START		
EXIT	RTS			
HEX	LDX	#\$04	ļ	SET COUNTER
	LDA	#\$00	ļ	INITIALISE ACC
AGAIN	ASL	\$15		
	ROL	A		
	DEX			
	BNE	AGAIN		
	CMP	#\$0A	!	IS IT 10 OR MORE
	BCC	DIGIT+1	!	ASCII ADDITION FOR NUMBER
	CLC			
	ADC	#\$37	ļ	FOR LETTER
DIGIT	BIT	\$3069		
	JSR	\$FFD2	!	PRINT RESULT
	RTS			
	*=\$84A0 START NEXT NEXT1 EXIT HEX AGAIN DIGIT	*=\$84A0 START BEQ BCS JSR LDA JSR NEXT LDY NEXT LDY NEXT1 LDY STY NEXT1 LDY STY NEXT1 LDY BNE JSR DEY BNE JSR JSR JSR JSR JSR JSR JSR JSR JSR JSR	 *=\$84A0 START BEQ SYNTAX BCS SYNTAX JSR \$81F5 LDA #\$20 JSR \$FFD2 NEXT LDY #\$02 JSR HEX DEY BNE NEXT LDY \$14 STY \$15 NEXT1 LDY #\$02 JSR HEX DEY BNE NEXT1 JSR HEX DEY BNE NEXT1 JSR \$6073 JSR \$6073 JSR \$6073 JSR \$6073 JSR \$6073 JSR \$6073 JSR \$6073 AGAIN AGAIN DIGIT CLC ADC #\$37 BIT \$3069 JSR \$FFD2 RTS 	*=\$84A0 START BEQ SYNTAX BCS SYNTAX JSR \$81F5 ! LDA #\$20 JSR \$FFD2 ! NEXT LDY #\$02 JSR HEX ! DEY BNE NEXT ! LDY \$14 STY \$15 ! NEXT1 LDY #\$02 JSR HEX DEY BNE NEXT ! LDY #\$02 JSR HEX DEY BNE NEXT1 ! JSR \$0079 ! CMP #\$2C ! BNE EXIT ! JSR \$0079 ! CMP #\$2C ! BNE EXIT ! JSR \$0073 JMP START EXIT RTS HEX LDA #\$00 ! AGAIN ASL \$15 ROL A DEX BNE AGAIN CMP #\$0A ! DIGIT BIT \$3069 JSR \$FFD2 !

360 SYNTAX JMP \$AF08

84D5	AGAIN	84E2	DIGIT
84DØ	EXIT	84D1	HEX
84AE	NEXT	84BA	NEXT1
84A0	START	84E9	SYNTAX

LINES 30–50: Here we pick up the first decimal number to convert. The high byte will be in \$15 and the low, in \$14. We then print a space on the screen so the first digit is a character away from the border or the last PRINTED statement.

LINES 60–90: Each part of the hexadecimal number will have two characters. As we will convert our decimal in two stages, the high byte first then the low, each will require two entries to the conversion routine. We therefore set a counter to two, in this case the y register. After going to the conversion routine we decrease the counter. If it is zero then we have done it twice, if not we go back again.

LINES 100–110: We have now converted the high byte. As the conversion subroutine uses the high byte register in the transposition, we transfer the low byte of the decimal to that register.

LINES 120–150: This is the same as lines 40–70 but for the low byte.

LINES 160–220: The GET PARAMETERS has already picked up the byte after the last decimal digit. Here we get that byte again by a call to CHRGOT. We want to find out if more than one conversion is required. The syntax of the command is for a comma as a separator, so we check for that. If the check succeeds, we print the comma to the screen and go back to convert the number. On failing the check, it is back to BASIC *via* the RTS.

THE CONVERSION ROUTINE

We use this routine four times for every decimal number in the command, twice for both the high and low bytes. We enter with the byte in location \$15. The hexadecimal number for a byte consists of two digits, one for the upper four bits and the other for the lower four. As we print on the screen from left to right, we print from the most significant hex digit and therefore want the high bits of the decimal number first.

Hex uses numerals 0 to 9 and letters A to F. Unfortunately, these do not follow in sequence in the ASCII table, as other characters lie between 9 and A. We therefore have to test for this when converting to ASCII for printing to the screen.

LINES 230–240: The x register is initialized to 4 as a counter for taking off the required bits for each hex digit. The accumulator is used to gather in the bits so is initialized to zero before we start.

LINES 250–280: This is the main part of the conversion. We use the instruction ASL to move all the bits of the decimal byte one place to the left. The most significant bit (bit 7) is moved to the carry flag. The least significant bit (bit \emptyset) is filled with zero (although that does not worry us). We need the bit we put into the carry back in the accumulator. This is achieved by the command ROL. This moves the accumulator bit one place to the left, filling bit \emptyset with the carry value, which we have just set (or unset). Bit 7 of the accumulator goes to the carry and again it is of no use to us here.

Now the counter is decreased and checked to see if we have done the bit shifting four times. We have now taken the four high bits of \$15 (the decimal byte) and put them in the same order in the accumulator but in the low bit positions.

LINES 290–350: The answer in the accumulator is now converted to ASCII form and printed to the screen. If it is less than \$0A, it is a number so we add \$30 to it. Greater than 10 means it is a letter, so we have to add \$37, giving letters from A to F.

```
HEX – Converting a hexadecimal number to decimal
COMMAND SYNTAX
HEX hex number [,hex number,....]
```

The hex number can be of two or four digits. More conversions can be added if separated by commas. The normal '\$' sign which preceeds hex numbers must not be used.

A four digit hex number can be split very conveniently into two parts. The two left digits are the high byte whilst the right are the low byte. Where a two digit conversion is required, we treat it as a low byte number. The two digits can be further split in that one represents the high nybble and the other the low nybble (a nibble is half a byte or four bits).

To do the conversion we collect two hex digits at a time and convert them to a one byte answer.

ASSEMBLY LISTING

9	*=\$85 37			
10	START	STA	\$63	
20		JSR	\$0073 !	GET NEXT BYTE
30		STA	\$62	
40		JSR	DECIMAL	
50		Pha	!	PUT HIGH ANS ON STACK
60		JSR	\$0073	
70		BEQ	LOWPRINT!	ONLY TWO BYTE HEX NO
80		CMP	#\$2C !	IS IT A COMMA
90		BEQ	Comma !	YES & ONLY 2 BYTE HEX
100		STA	\$63	

110		JSR	\$0073		
120		STA	\$62		
130		JSR	DECIMAL		
140		TAY		ļ	PUT LOW ANS IN Y
150		PLA		I	GET HIGH ANS IN ACC
160		JSR	PRINT	i	PRINT ANSWER
170		JSR	\$0073	i	GET NEXT BYTE
180		CMP	#\$20		
190		RED		•	YES
200		RTS	0011111	·	120
210		PIA		ı	GET HIGH ANS
228	Lowinin			1	PIT IT IN LOU AND DEG
220			****	ì	CET UICU AND TO ZEDO
230		TCD	##00 DD1NT	:	DETNIT ANGUED
250		DTC	FRINT	•	FRINT HNOWER
230	COMMA				
200		TAV		:	DUT IT IN LOU AND DEC
270			4400	:	PUT IT IN LOW ANS REG
200		LUH		:	SET HIGH ANS TO ZERU
270	0.00046.1	JSK	PRINI	:	PRNT ANSWER
300	LUMMAI	LDA	#\$20	!	ASULL FUR CUMMA
310		JSR	\$FFD2	:	PRINT IT AS SPACER
320		JSR	\$0073	!	GET NEXT BYTE
330		JMP	START		
340	DECIMAL	LDY	#\$01	!	COUNTER
350	AGAIN	LDA	\$0062,Y	!	GET LOW CHAR
360		CMP	#\$30	!	IS IT A NUMBER
370		BCC	SYNTAX	!	NOT NUMBER OR
200		-	***		
380			#947	!	15 II LEITER > F?
370		BLS	STNIAX		
400			#F⊅JHi DICIT	:	15 II A NUMBER?
410		BLL	DIGII	!	
420			#\$41 OV&FTAV	:	15 II A LEFTER?
430		BUU	STNTAX	:	NU AGAIL LETTER
440		SBC	#\$37	!	CUNVERI ASCII LEITER
					INTU REAL NUMBER
450		BNE	NEXT	!	ENFURCED
460	DIGII	SEC			
470		SBC	#\$30	!	CUNVERT ASCII NUMBER
					INTO REAL NUMBER
480	NEXT	STA	\$0014,Y		
490		DEY			
500		BPL	AGAIN	!	NEXT CHARACTER
510		LDY	#\$04	ļ	COUNTER
520	NEXT1	ASL	\$15	!	PUT HIGH ANS IN
					HIGH BITS

530		DEY
540		BNE NEXT1
550		LDA \$14
560		ORA \$15 ! JOIN BOTH TOGETHER
570		RTS
580	SYNTAX	JMP \$AF08
590	PRINT	JSR \$B391
600		JSR \$8401 ! CONVERT TO
		POSITIVE REAL NUMBER
610		JSR \$AD8D
620		JSR \$BDDD
630		JSR \$B487
640		JMP \$AB21
857E	AGAIN	856A COMMA
8571	COMMA1	857C DECIMAL
8595	DIGIT	8562 LOWPRINT
8598	NEXT	85A0 NEXT1
85AD	PRINT	8537 START
85AA	SYNTAX	

LINES 10–50: The routine is entered with the first digit and is stored. Calling CHRGET, we get the next byte and again store it. The decimal conversion routine is visited (this is described later), and the result comes back into the accumulator which we place on the stack.

LINES 60–90: The next byte of the command is now collected and two checks made, first to see if it is the end of the command and secondly for a comma. If the first succeeds we go off and print what we have already collected, but as a low byte answer. If it is a comma, we again print but will return to do further conversions.

LINES 100–130: This is a repeat of lines 10–50 except the result in the accumulator is put in the Y register instead of the stack.

LINES 140–200: The result of the four digit conversion is printed to the screen here. The v register has the low byte and the high byte is pulled off the stack into A. The print routine described at the end is now called. The conversion is complete and we now check to see if further conversions are required by getting the next byte. If a comma is not present, the routine is ended.

LINES 210–250: The low byte answer is printed here. The byte is pulled off the stack and placed in the y register and the accumulator set to zero. After printing, the HEX routine is left.

LINES 260–290: This is the same as lines 210–250 but instead of leaving,

we continue as we know there was a comma present when we arrived here.

LINES 300–330: A comma is printed on the screen to separate the answers. The first byte of the next conversion is gathered and we go back to the beginning to start converting again.

LINES 340–570: THE DECIMAL CONVERSION: The two bytes will be in locations \$62 and \$63. They will hold the ASCII values of either a numeral or a letter between A and F. SYNTAX ERRORS are given if they fall outside these limits.

Taking each location in turn, we determine what it is and deduct from it \$37 for a letter of \$30 for a numeral, the value ending up between \$60 and \$6F (0 to 15). These are placed in registers \$14 and \$15. We now have to combine these into one number. Address \$15 will have the high nybble but in the wrong bit positions. To get them into the upper four bits we shift the bits left four times. To join the two together, the byte in \$14 is copied to the accumulator and is OREd with location \$15. With the final result in the accumulator, we exit the subroutine.

LINES 590–640: PRINT RESULT TO SCREEN: Six subroutines are called here where the result of numeric calculations are converted to a string of ASCII characters and printed to the screen which except for one are all ROM routines. The one exception is a subroutine in the DEEK routine (see Chapter 8). For convenience, we reproduce it below. We enter this PRINT routine with numeric data, the y register holding the low and the accumulator the high byte.

ROUTINE \$B391 – FIX TO FLOAT

This sets the data flag in \$0D to zero signifying numeric data. The number we wish to convert is placed in FAC#1 registers \$62, meaning that numbers over 32768 (\$80) will be output as negative numbers.

ROUTINE \$8401 – CONVERT TO POSITIVE

10		LDA	\$66
20		BPL	EXIT
30		LDY	#>DATA
40		LDA	#≺DATA
50		JSR	\$BA8C
60		JSR	\$B86A
70	EXIT	RTS	
80	DATA	BYT	\$91,\$00,\$00,\$00,\$00

We check register \$66 of the FAC#1 to see if it is negative. If so we load FAC#2 with zero and set for no exponent. This is done through the ROM routine \$BA8C, entering with the data start address in A and Y.

Now by adding the two FACs together we will end up with a result in FAC#1 which is a real whole number; \$B86A will achieve this.

ROUTINE \$AD8A – CHECK

This just checks that the data is numeric, otherwise a 'TYPE MISMATCH' error is given.

ROUTINE \$BDDD - FAC#1 TO STRING

This converts the contents of FAC#1 to an ASCII string and places it at the bottom of the stack. The Y and A registers will hold this address when the routine is finished.

ROUTINE \$B487 – SET UP STRING

This sets various registers so that the **PRINT** routine knows where to print from and how long the string is.

```
ROUTINE $AB21 – PRINT
```

This takes the data from the bottom of the stack and prints it to the screen. We jumped to this routine, so when it is ended, the processor will be directed back to the position calling this whole subroutine.

This routine, being a separate routine, is therefore capable of being used by other commands as in the MEM command.

```
TWO – Decimal to binary conversion
COMMAND SYNTAX
TWO decimal number [,decimal number,....]
```

The maximum decimal number which can be converted is 65535 and must be positive. Multiple conversions can be done if they are separated by commas. The result will be two eight digit binary numbers separated by a space, unless the number is 255 or less, when only one binary result will be shown.

All we need to do is to test each bit and print a zero or a one according to its state.

ASSEMBLY LISTING

9	*=\$84EC			
10	START	BEQ SYN	VTAX	
20		BCS SYN	VTAX	
30		JSR \$83	IF5 !	GET PARAMETER
40		LDA #\$2	20	
50		JSR \$FF	FD2 !	PRINT SPACE
60		LDA \$15	5	
70		BEQ LSE	3	
80		LDX #\$0)8 !	COUNTER
90	NEXT	ASL \$15	5	

110 LDA #\$30 ! TU PRINT A 0 120 SET BIT \$31A9 ! SET+1 IS LDA #\$31 130 JSR \$FFD2 140 DEX 150 BNE NEXT 160 LDA #\$20 170 JSR \$FFD2 180 LDA #\$20 170 JSR \$FFD2 180 LDA #\$20 170 JSR \$FFD2 180 LSB 190 NEXT1 ASL \$14 200 BCS SETA+1 210 LDA #\$30 220 SETA BIT \$31A9 230 JSR \$FFD2 240 DEX 250 BNE NEXT1 260 JSR \$0077 270 CMP #\$2C 270 CMP #\$2C 280 BNE EXIT 280 BNE EXIT 300 JSR \$0073 300 JSR \$0073 310 JMP \$TART 320 EXIT 330 SYNTAX	100		BCS	SET+1	!	TO PRINT A 1
120 SET BIT \$31A9 ! SET+1 IS LDA #\$31 130 JSR \$FFD2 140 DEX 150 BNE NEXT 160 LDA #\$20 170 JSR \$FFD2 180 LSB LDX #\$08 190 NEXT1 ASL \$14 200 BCS SETA+1 210 LDA #\$30 220 SETA BIT \$31A9 ! SETA+1 IS LDA #\$31 230 JSR \$FFD2 240 DEX 250 BNE NEXT1 260 JSR \$0079 ! GET LAST BYTE OFF BUFFER AGAIN BUFFER AGAIN 270 CMP #\$2C ! IS IT A COMMA? 280 BNE EXIT ! NO COMMA 290 JSR \$FFD2 ! PRINT COMMA AS SPACI 300 JSR \$0073 ! GET NEXT BYTE 310 JMP START 320 EXIT 320 EXIT RTS 330 SYNTAX JMP \$AF08	110		LDA	#\$30	1	TO PRINT A U
130 JSR \$FFD2 140 DEX 150 BNE NEXT 160 LDA #\$20 170 JSR \$FFD2 ! PRINT A SPACE 180 LSB LDX #\$08 190 NEXT1 ASL \$14 200 BCS SETA+1 210 LDA #\$30 220 SETA BIT \$31A9 ! SETA+1 IS LDA #\$31 230 JSR \$FFD2 240 DEX 250 BNE NEXT1 260 JSR \$6079 278 CMP #\$2C 279 CMP #\$2C 280 BNE EXIT 280 BNE EXIT 300 JSR \$6073 300 JSR \$0073 310 JMP \$TART 320 EXIT 330 SYNTAX JMP \$AF08	120	SET	BIT	\$31A9	!	SET+1 IS LDA #\$31
140 DEX 150 BNE NEXT 160 LDA #\$20 170 JSR \$FFD2 ! PRINT A SPACE 180 LSB LDX #\$08 190 NEXT1 ASL \$14 200 BCS SETA+1 216 LDA #\$30 220 SETA BIT \$31A9 ! SETA+1 IS LDA #\$31 230 JSR \$FFD2 240 DEX 250 BNE NEXT1 260 JSR \$0079 ! GET LAST BYTE OFF BUFFER AGAIN 270 CMP #\$2C ! IS IT A COMMA? 280 BNE EXIT ! NO COMMA 290 JSR \$FFD2 ! PRINT COMMA AS SPACE 300 JSR \$0073 ! GET NEXT BYTE 310 JMP START 320 EXIT RTS 330 SYNTAX JMP \$AF08	130		JSR	\$FFD2		
150 BNE NEXT 160 LDA #\$20 170 JSR \$FFD2 ! PRINT A SPACE 180 LSB LDX #\$08 190 NEXT1 ASL \$14 200 BCS SETA+1 210 LDA #\$30 220 SETA BIT \$31A9 ! SETA+1 IS LDA #\$31 230 JSR \$FFD2 240 DEX 250 BNE NEXT1 260 JSR \$0079 ! GET LAST BYTE OFF BUFFER AGAIN 270 CMP #\$2C ! IS IT A COMMA? 280 BNE EXIT ! NO COMMA 290 JSR \$FFD2 ! PRINT COMMA AS SPACE 300 JSR \$0073 ! GET NEXT BYTE 310 JMP START 320 EXIT RTS 330 SYNTAX JMP \$AF08	140		DEX			
160 LDA #\$20 170 JSR \$FFD2 ! PRINT A SPACE 180 LSB LDX #\$08 190 NEXT1 ASL \$14 200 BCS SETA+1 210 LDA #\$30 220 SETA BIT \$31A9 ! SETA+1 IS LDA #\$31 230 JSR \$FFD2 240 DEX 250 BNE NEXT1 260 JSR \$0079 ! GET LAST BYTE OFF BUFFER AGAIN 270 CMP #\$2C ! IS IT A COMMA? 280 BNE EXIT ! NO COMMA 290 JSR \$FFD2 ! PRINT COMMA AS SPACE 300 JSR \$0073 ! GET NEXT BYTE 310 JMP START 320 EXIT RTS 330 SYNTAX JMP \$AF08	150		BNE	NEXT		
170 JSR \$FFD2 ! PRINT A SPACE 180 LSB LDX #\$08 190 NEXT1 ASL \$14 200 BCS SETA+1 210 LDA #\$30 220 SETA BIT \$31A9 ! SETA+1 IS LDA #\$31 230 JSR \$FFD2 240 DEX 250 BNE NEXT1 260 JSR \$0079 ! GET LAST BYTE OFF BUFFER AGAIN 270 CMP #\$2C ! IS IT A COMMA? 280 BNE EXIT ! NO COMMA 290 JSR \$FFD2 ! PRINT COMMA AS SPACE 300 JSR \$0073 ! GET NEXT BYTE 310 JMP START 320 EXIT 320 EXIT RTS 330 SYNTAX JMP \$AF08	160		LDA	#\$20		
180 LSB LDX #\$08 190 NEXT1 ASL \$14 200 BCS SETA+1 210 LDA #\$30 220 SETA BIT \$31A9 ! SETA+1 IS LDA #\$31 230 JSR \$FFD2 240 DEX 250 BNE NEXT1 260 JSR \$0079 ! GET LAST BYTE OFF BUFFER AGAIN 270 CMP #\$2C ! IS IT A COMMA? 280 BNE EXIT ! NO COMMA 290 JSR \$FFD2 ! PRINT COMMA AS SPACE 300 JSR \$0073 ! GET NEXT BYTE 310 JMP START 320 EXIT RTS 330 SYNTAX JMP \$AF08	170		JSR	\$FFD2	!	PRINT A SPACE
190 NEXT1 ASL \$14 200 BCS SETA+1 210 LDA #\$30 220 SETA BIT \$31A9 ! SETA+1 IS LDA #\$31 230 JSR \$FFD2 240 DEX 250 BNE NEXT1 260 JSR \$0079 ! GET LAST BYTE OFF BUFFER AGAIN 270 CMP #\$2C ! IS IT A COMMA? 280 BNE EXIT ! NO COMMA 290 JSR \$FFD2 ! PRINT COMMA AS SPACE 300 JSR \$0073 ! GET NEXT BYTE 310 JMP START 320 EXIT RTS 330 SYNTAX JMP \$AF08	180	LSB	LDX	#\$08		
200 BCS SETA+1 210 LDA #\$30 220 SETA BIT \$31A9 ! SETA+1 IS LDA #\$31 230 JSR \$FFD2 240 DEX 250 BNE NEXT1 260 JSR \$0079 ! GET LAST BYTE OFF BUFFER AGAIN 270 CMP #\$2C ! IS IT A COMMA? 280 BNE EXIT ! NO COMMA 290 JSR \$0073 ! GET NEXT BYTE 300 JSR \$0073 ! GET NEXT BYTE 310 JMP START 320 EXIT RTS 330 SYNTAX JMP \$AF08	190	NEXT1	ASL	\$14		
210 LDA #\$30 220 SETA BIT \$31A9 ! SETA+1 IS LDA #\$31 230 JSR \$FFD2 240 DEX 250 BNE NEXT1 260 JSR \$0079 ! GET LAST BYTE OFF BUFFER AGAIN BNE EXIT ! NO COMMA? 280 BNE EXIT ! NO COMMA 290 JSR \$FFD2 ! PRINT COMMA AS SPACE 300 JSR \$0073 ! GET NEXT BYTE 310 JMP START	200		BCS	SETA+1		
220 SETA BIT \$31A9 ! SETA+1 IS LDA #\$31 230 JSR \$FFD2 240 DEX 250 BNE NEXT1 260 JSR \$0079 ! GET LAST BYTE OFF 270 CMP #\$2C ! IS IT A COMMA? 280 BNE EXIT ! NO COMMA 290 JSR \$FFD2 ! PRINT COMMA AS SPACE 300 JSR \$0073 ! GET NEXT BYTE 310 JMP START 320 EXIT 320 EXIT JMP \$AF08	210		LDA	#\$30		
230 JSR \$FFD2 240 DEX 250 BNE NEXT1 260 JSR \$0079 ! GET LAST BYTE OFF 270 CMP #\$2C ! IS IT A COMMA? 280 BNE EXIT ! NO COMMA 290 JSR \$FFD2 ! PRINT COMMA AS SPACE 300 JSR \$0073 ! GET NEXT BYTE 310 JMP START 320 EXIT RTS 330 SYNTAX JMP \$AF08	220	SETA	BIT	\$31A9	1	SETA+1 IS LDA #\$31
240DEX250BNE NEXT1260JSR \$0079270CMP #\$2C280BNE EXIT280JSR \$FFD2290JSR \$FFD2300JSR \$0073310JMP START320 EXITRTS330 SYNTAXJMP \$AF08	230		JSR	\$FFD2		
250BNE NEXT1260JSR \$0079! GET LAST BYTE OFF BUFFER AGAIN270CMP #\$2C! IS IT A COMMA?280BNE EXIT! NO COMMA290JSR \$FFD2! PRINT COMMA AS SPACE300JSR \$0073! GET NEXT BYTE310JMP START320 EXITRTS330 SYNTAXJMP \$AF08	240		DEX			
260JSR \$0079! GET LAST BYTE OFF BUFFER AGAIN270CMP #\$2C! IS IT A COMMA?280BNE EXIT! NO COMMA290JSR \$FFD2! PRINT COMMA AS SPACE300JSR \$0073! GET NEXT BYTE310JMP START320 EXITRTS330 SYNTAXJMP \$AF08	250		BNE	NEXT1		
BUFFER AGAIN 270 CMP #\$2C ! IS IT A COMMA? 280 BNE EXIT ! NO COMMA 290 JSR \$FFD2 ! PRINT COMMA AS SPACE 300 JSR \$0073 ! GET NEXT BYTE 310 JMP START 320 EXIT RTS 330 SYNTAX JMP \$AF08	260		JSR	\$0079	ł	GET LAST BYTE OFF
270CMP #\$2C! IS IT A COMMA?280BNE EXIT! NO COMMA290JSR \$FFD2! PRINT COMMA AS SPACE300JSR \$0073! GET NEXT BYTE310JMP START320 EXITRTS330 SYNTAXJMP \$AF08						BUFFER AGAIN
280BNE EXIT! ND COMMA290JSR \$FFD2! PRINT COMMA AS SPACE300JSR \$0073! GET NEXT BYTE310JMP START320 EXITRTS330 SYNTAXJMP \$AF08	270		CMP	#\$2C	1	IS IT A COMMA?
290 JSR \$FFD2 ! PRINT COMMA AS SPACE 300 JSR \$0073 ! GET NEXT BYTE 310 JMP START 320 EXIT RTS 330 SYNTAX JMP \$AF08	280		BNE	EXIT	1	NO COMMA
300 JSR \$0073 ! GET NEXT BYTE 310 JMP START 320 EXIT RTS 330 SYNTAX JMP \$AF08	290		JSR	\$FFD2	ļ	PRINT COMMA AS SPACER
310 JMP START 320 EXIT RTS 330 SYNTAX JMP \$AF08	300		JSR	\$0073	1	GET NEXT BYTE
320 EXIT RTS 330 SYNTAX JMP \$AF08	310		JMP	START		
330 SYNTAX JMP \$AF08	320	FXIT	RTS			
	330	SYNTAX	.TMP	\$AFA8		
	000	with the	0111	770 VV		

8533	EXIT	8512	LSB
84FE	NEXT	8514	NEXT1
8504	SET	851A	SETA
84EC	START	8534	SYNTAX

LINES 10–20: On entering the first byte after the keyword is in the A register and by testing the carry and zero flags, we can check if a numeral is first.

LINES 30–150: We gather in the decimal number to convert and also print a space for presentation to move it away from the last item printed or from the border. Taking the high byte first, we check it for zero, if it is a zero we branch and just do the low byte. The x register is set to eight as there are eight bits to a byte, and we shall use it as a counter. We shift the bits of the high byte one place to the left. The leftmost bit comes off and goes to the carry flag. If the carry flag is zero, we therefore print ASCII \$30 which is zero and \$31 when the carry is set. This is repeated a further seven times for all the remaining bits.

LINES 160–250: This is a repeat of the above except for the low byte.

LINES 270–320: By calling CHRGOT, we test for a comma or if the end of the command has been reached. When the former is found, it is printed to the screen and we gather in the next byte before going to carry out the next conversion.

BIN – Binary to decimal conversion COMMAND SYNTAX BIN 8 bit binary number [,8 bit binary number,....]

Here we will convert an eight bit binary number to decimal. We supply a value that would be a high byte and one that is the low byte. For example, if you demanded 11111111 was converted, the answer would come out as 255. Only eight bit numbers are accepted but more conversions can be done by separating the items with commas.

This is essentially the reverse of the previous command. The 1s and 0s you type in will be picked up in their ASCII form. These have their rightmost bits corresponding to their numeric value, so by taking those we can build up a single byte number.

ASSEMBLY LISTING

9	*=\$85BF				
10		LDA	\$7A		
20		BNE	LOW		
30		DEC	\$7B	!	DECREASE CHARGET POINTER BY ONE
40	LOW	DEC	\$7A		
50	ANOTHER	LDX	#\$08	!	COUNTER
60	NEXT	JSR	\$0073	!	GET BYTE
70		BCC	NUMBER	!	IF NUMBER BRANCH
80	SYNTAX	JMP	\$AF08		
90	NUMBER	CMP	#\$32	!	IS > ASCII FOR 2
100		BCS	SYNTAX	!	YES
110		ROR	A	!	GET BIT 0
120		ROL	\$14	!	PUT IN \$14 MOVING 1 LEFT EVERY TIME
130		DEX			
140		BNE	NEXT	!	DONE IT 8 TIMES?
150		LDY	\$14	!	PUT ANS IN Y REG
160		LDA	#\$00	1	NO HIGH ANS
170		JSR	\$85AD	!	PRINT ANS-HEX ROUTINE
180		LDA	#\$2F		
190		JSR	\$FFD2	ļ	PRINT SLASH TO DIVIDE
					LOW ANS & HIGH ANS
200		LDA	\$14	!	NOW PUT ANS IN
					HIGH ANS REG
210		LDY	#\$00	!	SET LOW ANS REG TO 0

40

220	JSR \$85AD	! PRINT ANS-HEX ROUTINE
230	JSR \$0073	! GET NEXT BYTE OF
		INPUT BUFFER
240	CMP #\$20	! IS IT A COMMA?
250	BNE EXIT	! NO
260	JSR \$FFD2	! PRINT COMMA AS SPACER
270	JMP ANOTHER	
280 EXIT	RTS	

85C7	ANOTHER	85FB	EXIT
85C5	LOW	8509	NEXT
85D1	NUMBER	85CE	SYNTAX

LINES 10-40: This decreases CHRGET address by one.

LINES 50–140: We want to pick up eight binary digits so the x register is used as a counter. After we pick up a digit, via CHRGET, we check for a number and also if it is two or greater, in ASCII. By rotating the accumulator right, we take off bit 0 and it ends up in the carry flag. Then if we rotate location \$14 left, we move all its bits one to the left and put the carry flag state into the lowest bit. If we do this eight times, address \$14 will have a number equivalent to the 1s and 0s you typed in.

LINES 150–220: The PRINT routine in the HEX command is used twice. For this the low byte needs to be in y and the high in A. The value we want to print is in \$14 and by changing the register we load it into we can print out the states we want. A slash is printed as a separator by a visit to the KERNAL routine at \$FFD2.

LINES 230–280: Having done one conversion, we take a look to see if more are required. A comma is printed if so and then we go back to do it all again.

8 Enhancing the resident BASIC

Introduction

In the previous chapter the commands were of a toolkit nature. In this chapter they are mainly improvements to the standard 64 commands, GOTO, GOSUB, RESTORE, PRINT, INPUT, GET, PEEK, POKE and LOAD. To that end, we are supplying CGOTO, CGOSUB, PROC, DPROC, EPROC, RESET, WRITE, ENTER, INKEY\$, DEEK, DOKE and CHAIN. There are five commands which have no 64 BASIC equivalent, but which we hope will enhance your BASIC programming. They are POP, PLOT, COLOUR, LOMEM and HIMEM. The final command given is that of QUIT and exists the UTILITY.

In comparison with the toolkit commands these are shorter, but no less useful to you. No doubt you can think of existing commands which could be enhanced and even more to add. This chapter should help you on your way.

CGOTO and CGOSUB

COMMAND SYNTAX

CGOTO variable or line number CGOSUB variable or line number

A limitation of Commodore BASIC is that it does not permit the use of calculated destinations with GOTO and GOSUB. We thought it would be nice to be able to use variables and mathematical expressions, for example A+20. To allow this, we have come up with two commands CGOTO and CGOSUB, the c standing for calculated or computed – whichever you prefer.

CGOTO is the easiest of the two, not that either is complicated. The routine requires only two instructions. In the BASIC ROM routine of GOTO the first instruction gets the line number and is therefore the only thing we have missed out. So after getting the variable we only have to jump to that part of GOTO.

CGOSUB is a bit longer in that we have to copy the ROM routine for GOSUB and change the address for calling the GOTO routine as we want to use our 'computed' routine.

ASSEMBLY LISTING

9	*=\$8FAF				
10	CGOTO	JSR	\$81F5	i	GET PARAMETER
20		JMP	\$A8A3	1	GO TO ROM GOTO
30	! CGOSUB	ROUTINE			
40		LDA	#\$03		
50		JSR	\$A3FB	!	CHECK FOR ROOM ON STACK
60		LDA	\$7B	!	SAVE CHRGET ADDRESS ON STACK
70		PHA			
80		LDA	\$7A		
90		PHA			
100		LDA	\$3A	!	SAVE CURRENT LINE
110		PHA			
120		LDA	\$3B		
130		PHA			
140		LDA	#\$8D	i	MARKER FOR GOSUB ON STACK
150		Pha			
160		JSR	\$0079	!	GET LAST BYTE AGAIN
170		JSR	CGOTO		
180		JMP	\$A7AE	!	BASIC TO EXECUTE PROGRAM FURTHER

8FAF CGOTO

LINES 10–20: cGOTO: We use GET PARAMETERS simply to find the destination line number. It will evaluate any expression, and jump to the normal GOTO routine, one instruction in.

LINES 40–60: CGOSUB: These lines check if there is enough room on the stack to store the routine's information and a buffer amount for other routines. To do this the value in the accumulator is doubled and added to \$3E (62 dec). This is then compared with the stack pointer. If the stack pointer is the lesser value, then an 'OUT OF MEMORY' error is generated. In our case, the stack pointer would have to be less than \$44 (it starts at \$FF).

LINES 70–130: There are two markers we will require when the subroutine is finished. These are the present byte's address from the CHRGET routine and the line number we must return to later. The stack is used to store this information. **LINES 140–150:** Another value is put on the stack. This is used by the RETURN routine to check a GOSUB has been implemented.

LINES 160–170: Now we can go to our destination. To do this, we get the last byte collected by CHRGET again and go to our new computed GOTO routine.

LINE 180: Once the GOTO routine has been completed, in which the CHRGET has been given new values, we return to the normal flow of BASIC and the program is continued at its new address.

Procedures

COMMAND SYNTAX

PROC title – call a procedure DPROC title – define a procedure EPROC – the end of a procedure

The title is not required within quotes. If it is then the quotes will be considered as part of the name. Spaces also cannot be used as CHRGET ignores them (a space in the DPROC title will not be matched in the PROC title). On the other hand, a space in the PROC title will have no bearing on the matching. A colon is the only other character which cannot be used in a title.

You can have as many PROCS on a line as you want, but the DPROC must be on a line of its own. Everything following the DPROC to the end of the line is included as the title. EPROC follows exactly as RETURN.

64 BASIC cannot be described as a structured language. GOTOS and GOSUBS do not form the basis of a structured language.

To start you on the road to 'structured programming', we are introducing PROCEDURES. We have nothing profound to offer but by giving you an introduction we hope you will be able to take it further (IF ... THEN .. ELSE WHILE ... WEND DO ... UNTIL etc . .)

The form of procedures we have written are really no more than GOSUBS with names or variables (CGOSUB). In fact, they will be slower, but not that you would notice, than GOSUBS because of the extra code required. So what advantage will they have? Well, they can be relocated anywhere in the program without changing any directive line numbers; adding procedures from one program to another, especially if they include procedures within them, is a simple matter. If GOTO was also given the same treatment, all directive line numbers could vanish. Renumbering a program would be a simple matter of changing the line numbers rather than going through the whole program and correcting destinations. A further function they perform, and one that should not be overlooked, is that they make a program easier to follow. For instance, to see PROC PERFORM-WAIT is clearer than GOSUB 2000.
Quite simply, all we do when finding a PROC is to search through for the token DPROC and then compare the named titles. On finding it, we perform a GOSUB. The UTILITY interpreter will action the command DPROC as a REM if it encounters one. The third command of the trio is EPROC and is just a RETURN by a different name. We actually go to the RETURN routine. After the listing and description we suggest some improvements.

9	*=\$8FD2			
10		LDX #\$00		
20		DEC \$7A	!	DECREASE CHRGET ADD
30		BCS COLLECT		
40		DEC \$7B		
50	COLLECT	JSR \$0073	!	GET PROC NAME
60		BEQ NAMEEND	!	FOUND 0 OR :
70		STA \$0200,X	!	STORE IN INPUT BUFFER
80		INX		
90		BNE COLLECT	!	ENFORCED
100	NAMEEND	LDA #\$00	!	0 AT END
110		STA \$0200,X		
120		LDA #\$03		
130		JSR \$A3FB	!	CHECK STACK DEPTH
140		LDA \$7B	!	SAVE CHRGET ADDRESS
150		PHA		
130		LDA \$7A		
170		Pha		
180		LDA \$3A	!	SAVE CURRENT LINE NO
190		PHA		
200		LDA \$39		
210		PHA		
220		LDA #\$8D	!	STACKMARKER FOR GOSUB
230		Pha		
240		LDA \$2B	!	GET AND STORE
250		STA \$FB	!	ADDRESS OF 1ST
230		LDA \$2C	!	PROGRAM LINE
270		STA \$FC		
280	NEXT	LDY #\$00		
290		LDA (\$FB),Y	!	GET AND STORE LINKS
300		STA \$FD		
310		INY		
320		LDA (\$FB),Y		
330		STA \$FE	1	
340		BNE CONT	!	NOT END OF PROGRAM
350	· 2	LDX #\$11	. !	DPROC NOT FOUND

360		JMP	\$A437	!	UNDEF'D STATEMENT ERROR
370	CONT	LDY	#\$04	!	SKIP LINKS AND LINE NUMBERS
380		LDA	(\$FB),Y		
390		CMP	#\$E1	!	TOKEN OF DPROC
400		BEQ	PROC	ļ	FOUND A DPROC
410	LINE	LDA	\$FD	ļ	PUT LINKS TO
420		STA	\$FB	ļ	LINE REGISTERS
430		LDA	\$FE		
440		STA	\$FC		
450		BNE	NEXT	!	ENFORCED \$FE
					CHECKED EARLIER
430	PROC	LDX	#\$FF		
470	CHECK	INX			
480		INY			
490		LDA	(\$FB),Y	!	GET PROC TITLE
500		BEQ	ZERO		END OF LINE
510		CMP	\$0200,X	!	COMPARE FOR MATCH
520		BEQ	CHECK	!	MATCH FOUND
530		BNE	LINE	ļ	NO MATCH FIND NEXT
-			9 100 - 11		DPROC
540	ZERO	CMP	\$0200	!	COMPARE LAST BYTE
550		BNE	LINE	!	NO MATCH
560		SEC		!	PREPARE FOR SUBTRACT
5/0		LDA	\$FB	!	ADDRESS OF PROCEDURE
280		SBC	#\$01	!	DECREASE TO END UP
500		OTA	# 74	7	LAST LINE
170		JIA	⊅/A 4F0	:	PUT AS CHRGET ADD
410		CDC	⇒ru #¢00	1	IN CACE DACE CORCER
610		CTA	₩₽00 470	:	IN CHOE FRUE CRUSSED
430		TMP	#70 \$A7AF	1	BASIC TO CONT PPOG
000		0111	***/**		CHOIC IS CONTEROS

902C	CHECK	8FDA	COLLECT
9018	CONT	9020	LINE
8FE5	NAMEEND	9006	NEXT
902A	PROC	9039	ZERO

LINES 10–110: Using CHRGET, after decreasing it by one, we take the PROC title and store it at the start of the input buffer and tag a zero byte on the end for checking purposes.

LINES 120–230: Same as in CGOSUB, saving relevant details for RETURN, or in this case EPROC.

LINES 240–360: After collecting the start address of the program, we search through the program. This part gets the links and checks for the end of the program. The UNDEF'D STATEMENT error is given if the latter occurs without finding the procedure.

LINES 370–450: When inputting a BASIC line, any spaces between a line number and the first character are removed during tokenizing, (LIST inserts a space for clarity). This means that the first token in a line is the fourth byte (starting at zero, remember), so we check only this byte for the DPROC token of \$E1. If not found, the link address is placed into the line registers and the hunt continues.

LINES 460–550: Having found a DPROC token, we have to compare each character separately and as long as they match we continue checking. When we reach the end of the DPROC program line, we check the input buffer for a matching zero. When all checks succeed, we have found the required procedure.

LINES 560–630: Knowing our destination, we take the start address of the DPROC line and reduce it by one, the end of the preceding line, and store it as the CHRGET address. Finally, we jump to BASIC to evaluate.

IMPROVEMENTS?

One of the first questions that came to mind was: how could we speed up the search for the procedures? One solution to the problem is to form a table in RAM holding the start address where you first check to see if the PROC name is in it. This would involve setting aside an area of RAM: under ROM would be an ideal place, for such a table. Two characters would then have to be chosen: one to mark the end of an entry and the other, the end of the table. The make-up of the table could consist of the PROC name and its start address. How could the table be filled? When the interpreter finds the keyword PROC, the table would be searched for a match. If no match is found, then a program search, like our routine, could be instigated. On finding the DPROC with the correct name, it would be added to the table in case of another call.

There are, however, problems. Let us assume that a program containing PROCS has been RUN. This would mean the table has names and addresses within it. Before running it again, you add an extra line before the procedure. The line with the DPROC now has a different address from that in the table. Another action giving rise to the same problem is when you load another program. It may have a PROC with the same name as the previous program. Again the table may have another address. A further problem may arise in that more PROCS will be added making the table longer and longer.

Two solutions to this problem spring to mind; there are probably others. The first is to write a new RUN command, for example, PRUN,

where one of its actions would be to place the end of table marker back to the beginning – thereby effectively clearing the table. The other is to have a command that can be actioned to do just this and only this. It could be initiated in direct mode or from within a program.

A further improvement would be to allow parameters to be passed using variables which are local to the procedure. These variables could be used elsewhere in the program without losing their original values. We would envisage the PROC command to include, in say, brackets, the values or other variables to be used, for example, PROC INPUT(2,4,6) or PROC INPUT(2,A,6). The DPROC, on the other hand, will define the variables to be used. For example, DPROC INPUT(X,Y,Z). These variables may be used elsewhere but in the procedure they will start with values given in the PROC command.

What would have to be done is that when arriving at the procedure a search is carried out for the variables x,y,z in the normal variable area. If they are found, their current values would have to be transferred to a keeping area, and the new values set up. If the variable is not present, then it will have to be created. The default value of a numeric variable is zero and this will also have to be stored in the keeping area. For strings, the addresses will have to be store the original values.

The process would be the same if you wanted to incorporate GLOBAL and LOCAL commands to a BASIC extension.

Our last improvement, although we are sure you could think of more, is to allow the names of procedures to include keywords. This would be relatively simple in that all you have to do is to slightly alter the CRUNCH token routine (see Chapter 3). In that routine when it comes across a REM, for instance, it skips further crunching. All you have to do is to insert further coding to check for PROC and DPROC and follow the same path as REM.

POP – RETURN without returning

COMMAND SYNTAX

POP

There are no parameters to this command. If it is activated without a GOSUB, CGOSUB or a PROC being used, a 'RETURN WITHOUT GOSUB' error will be generated.

There are many occasions when one requires to leave a subroutine but not go back to the calling position. This is, of course, possible but leaves values on the stack; do it too often and the stack will become full and an 'OUT OF MEMORY' error will occur

POP will remove from the stack the data placed there by the last

GOSUB, or equivalent. This will mean, for example, if you called a subroutine which in turn called another and whilst in the second you called POP, then you will go back to main program when the RETURN is met, not the first subroutine. A GOTO after a POP will mean you can go anywhere from a subroutine without any worries about the stack. POP will also discharge any FOR/NEXT loops. If you happen to be in one at the time, watch out.

ASSEMBLY LISTING

- 9	*=\$8631				
10		LDA	#\$FF	ļ	RESET FOR/NEXT PTR
20		STA	\$4A		
30		JSR	\$A38A	!	SEARCH STACK FOR
					GOSUB & FOR ACTIVITY
40		TXS		!	X REGISTER TO
					STACK POINTER
50		CMP	#\$8D	ł	GOSUB MARKER ON STACK
60		BEQ	CONT		
70		LDX	#\$16	ļ	ERROR-RETURN
					WITHOUT GOSUB
80		JMP	\$A437		
90	CONT	INX		!	REMOVE GOSUB ACTIVITY
100		INX		ļ	INCREASE X AS IT WILL
					BE STACK POINTER
110		INX			
120		INX			
130		INX			
140		TXS		!	REPLACE STACK POINTER
150		RTS			

8642 CONT

LINES 10–20: By loading \$4A with \$FF, we effectively cancel any FOR/NEXT loop.

LINES 30–80: This is the ROM routine used by RETURN to look for the GOSUB marker on the stack. On return the stack pointer is in the X register and the accumulator has a value from the stack. If this is \$8D, the RETURN marker was present. An error will be produced if anything else is found.

LINES 90–150: To remove the GOSUB activity, we take the stack pointer, which is still in the X register, and increase it by five and then use it as the new stack pointer.

RESET – Selective data restorer

COMMAND SYNTAX

RESET [line number]

When no line number is present it behaves as the standard command RESTORE. With the parameter it will set the DATA pointer to the specified line.

DATA statements are extremely useful commands, and with sprites on the 64 you no doubt use them frequently. The snag comes when you want to use the same DATA statements again. RESTORE only allows you to set the pointer back to the first DATA statement, actually the start of the program, which has the same effect. To use statements again that are not at the beginning, dummy READS have to be employed to get to the desired position. To allow you greater flexibility, RESET will allow you to specify the line the next READ will start at, whether before or ahead of the present position. The RESTORE command takes the start of program address, subtracts one from it, and places it in the DATA pointer registers. RESET will take the line number, find its address, decrease it and set the pointers. Although the routine will give an error if the prescribed line number is not present, we do not check to see if it is a DATA line. This does not matter to BASIC as it will find the next DATA line when READ is sanctioned.

- 9	*=\$8611			
10		BNE RESE	T !	CHECK FOR PARAMETER
20		JMP \$A81	D !	NO - RESTORE IN
				BASIC ROM
30	RESET	JSR \$81F	5 !	GET LINE NUMBER
40		JSR \$A61	3 !	FIND BASIC LINE
50		BCS CONT	!	FOUND LINE
60		LDX #\$15	!	ILLEGAL DIRECT ERROR
70		JMP \$A43	7!	ROM ERROR ROUTINE
80	CONT	SEC	!	PREPARE FOR SUBTRACT
90		LDA \$5F	ļ	LOW ADD OF LINE
100		SBC #\$01	!	DECREASE BY ONE
110		STA \$41	!	DATA REG IN PAGE 0
120		LDA \$60	!	HIGH ADD OF LINE
130		SBC #\$00	!	IF PAGE IS CROSSED
140		STA \$42	!	DATA REGISTER
150		RTS		

LINES 10-20: If no line number, we go straight to RESTORE in the ROM.

LINES 30–70: When the line number is picked up it will be in the right location for a line search, which is immediately carried out. The carry flag set will indicate the line was found. The error given for not finding it is 'ILLEGAL DIRECT'.

LINES 80–150: Locations \$5F and \$60 will have the address of the line, and from these we subtract one and store them in the DATA pointers.

DEEK and DOKE - BASIC Addressing

It should be clear by now that addresses are stored in two locations as a low and a high byte. In the resident BASIC the only way to find the address, held in locations, is to do two PEEKS, one for each location, then multiply the high byte by 256 and add in the low byte, giving the address in decimal. To set up an address, the reverse process is used, but using POKE in place of PEEK.

The UTILITY commands are therefore obvious. We wish to read or set an address, or pair of locations, with one command. These are DEEK and DOKE.

DEEK – seeing double

COMMAND SYNTAX

DEEK (low byte location)

This returns the 16 bit value held in the given address and the following one. The rules for PEEK apply in that it must be an argument to a command (that is, a function).

9	*=\$83D7				
10		LDA \$15			
20		PHA			
30		LDA \$14			
40		PHA			
50		JSR \$AEFA	!	CHECK F	FOR (
30		JSR \$81F5	ļ	GET PAR	RAMETER
70		JSR \$AEF7	l	CHECK F	FOR)
80		LDY #\$01			
90		LDA (\$14),Y			
100		TAX			
110		DEY			
120		LDA (\$14),Y			
		-			

130		TAY			
140		PLA			
150		STA	\$14		
160		PLA			
170		STA	\$15		
180		TXA			
190		JSR	\$B391	!	A & Y IN FAC#1
200		JSR	CONVERT		
210		PLA			
220		PLA			
230		JMP	\$AD8D	!	EXIT
240	CONVERT	LDA	\$66	ļ	CHECK FOR SIGN
250		BPL	EXIT		
260		LDY	#>DATA		
270		LDA	# <data< td=""><td></td><td></td></data<>		
280		JSR	\$BA8C	ţ	CONSTANT TO FAC#2
290		JSR	\$B86A	!	ADD FAC#2 TO FAC#1
300	EXIT	RTS			
310	DATA	BYT	\$91,\$00,9	\$0(3,\$00,\$00

8401 CONVERT 840F EXIT 8410 DATA

LINES 10–40: DEEK, being a BASIC function rather than a command, is used in conjunction with other keywords. You have no doubt gathered

used in conjunction with other keywords. You have no doubt gathered that keywords use a fair number of zero page locations, notably \$14, \$15 and the FACS. We cannot take DEEK in isolation and also have to get its parameters. The latter means that we will use \$14 and \$15. We do not require to use these on exit, so we take the precaution of saving the current contents on the stack for the time being.

LINES 50–70: These not only get the parameters, but also check for the convention of them being in brackets. The ROM routines used will give the error if they are not present.

LINES 80–180: Using the address, now in \$14 and \$15, we read the contents and store them into the registers A and Y. We can restore the original values to \$14 and \$15, and do so.

LINES 190–230: The calling routine will expect the result in the FAC#1 and this is all the routine at \$B391 does. Unfortunately it also stores it as a signed integer. To correct this, CONVERT is called. Having done so, we pull off the return address, but we do not require to go back to the evaluation routine, and jump back to ROM. In ROM it will check that it is numeric data and will return to the calling routine, say PRINT or DOKE.

LINES 240-300: CONVERT: If the number requires converting it will have

a negative sign. With no sign we exit the routine. Failing that we load a constant into FAC#2 which when added to the contents of FAC#1, will change it to an unsigned number. For a more detailed explanation, see MEM in Chapter 7.

DOKE – complete addressing

COMMAND SYNTAX

DOKE low byte address, value

This turns the value to a two byte number and stores it in the given address and the following one. The value has a maximum of 65535 (\$FFFF).

ASSEMBLY LISTING

9	*=\$ 83B3			
10		JSR \$81F5	ļ	GET PARAMETER
20		JSR \$AEFD	!	CHECK FOR COMMA
30		JSR \$AD8A	!	GET NEXT PARAMETER
40		LDA \$66		
50		BMI ERROR		
60		CMP #\$91		
70		BCS ERROR		
80		JSR \$BC9B	!	PUT PARAMS IN A & Y
90		LDA \$65		
100		LDX \$64		
110		LDY #\$00		
120		STA (\$14),Y		
130		INY		
140		TXA		
150		STA (\$14),Y		
160		RTS		
170	ERROR	JMP \$8248	!	ILLEGAL QUANTITY ERROR

83D4 ERROR

LINES 10–30: The first routine called is the familiar one. The address will be in \$14 and \$15 after this call. The next routine checks for a comma. The last one collects the data for storing and puts it in the FAC#1.

LINES 40–80: These check for the legality of the data and set up the FAC#1 so we can take our values off.

LINES 80–160: After getting the data from FAC#1, we store them in the addresses specified.

OUTPUT – Setting the cursor

In the standard BASIC, the normal way to set the print position is to use the cursor control codes. Although they do the job, they are not ideal. You have to remember where the current position is, they take up bytes in the program, and TAB and SPC are not much better.

A far better way would be to specify the X and Y coordinates directly. To do this, three commands are included here, PLOT, WRITE and ENTER. The first will only set the cursor, the second will set the cursor and print what you want, whilst the last is INPUT with cursor positioning. The major command as far as routines go is PLOT. It really is a subroutine for the other two.

PLOT – cursor setting

COMMAND SYNTAX

PLOT (X,Y)

The maximum value of X is 39 and of Y is 24. The top left hand corner of the screen, cursor home position, has the coordinates of 0,0.

- 9	*=\$8381				
10		JSR	\$AEFA	!	CHECK FOR (
20		JSR	\$81F5	!	GET PARAMETER
30		LDA	\$14		
40		CMP	#\$28	!	IS X > 40
50		BCC	COMMA		
30	ILLEGAL	JSR	\$B248	ļ	ILLEGAL QUANTITY ERROR
70	COMMA	PHA			
80		JSR	\$AEFD	!	CHECK FOR COMMA
90		JSR	\$81F5	!	GET PARAMETER
100		LDX	\$14		
110		CPX	#\$19	ļ	ISY > 25
120		BCS	ILLEGAL		
130		PLA		!	RETRIEVE 1ST PARAM
140		TAY			
150		CLC		!	SET NOT READ CO-OR
160		JSR	\$FFF0		

170	JSR	\$0073	!	GET	NEXT	BYTE
180	RTS					

8390 COMMA 838D ILLEGAL

LINES 10–70: The left hand bracket is checked for and the X coordinate of the command is picked up. It is then checked that it does not exceed the limit. On an occasion that it does, we go to a ROM routine whose sole purpose is to generate the 'ILLEGAL QUANTITY ERROR'. We require to use location \$14 again so the X coordinate is put on the stack for a while.

LINES 80–120: After checking for the separating comma, we get the Y coordinate. This, too, is checked for legality.

LINES 130–180: the Y coordinate was picked up in the X register and now we retrieve the X coordinate and place it in the Y register. This is the opposite to what is logical but the KERNAL routine calls for them in that order. Before calling the routine, we clear the carry flag. (If we set the carry we would read the cursor position.)

After setting the cursor we get the next byte. This is for WRITE and ENTER so that they are set up for their respective ROM routines.

WRITE and ENTER

COMMAND SYNTAX

WRITE (X,Y)[string or variable] ENTER (X,Y)[string],variable

The coordinates take the syntax of PLOT. The remainder of the commands have the same syntax as their respective standard commands, WRITE as PRINT and ENTER AS INPUT.

ASSEMBLY LISTING

9	¥=	=\$83A7				
10	ļ	WRITE	COMMAND	-PRINT		
20			JSR	\$8381	ł	PLOT ROUTINE
30			JMP	\$AAA0	!	PRINT ROUTINE IN ROM
40	!	ENTER	COMMAND	-INPUT		
50			JSR	\$8381	!	PLOT ROUTINE
60			JMP	\$ABBF	!	INPUT ROUTINE IN ROM

These simply call the previous PLOT routine and then go to their normal ROM routines.

Colour

COMMAND SYNTAX

COLOUR background[,border][,text]

The latter two parameters are optional. If they are omitted it will not affect their present values. There is no error checking on values in the command. However, only the low byte of a number is used, that is, numbers up to 255, and of that only the lower four bits have effect (15 uses four bits whilst 16 uses five). The values to be used are the same as in the *Programmer's Reference Guide* or if you prefer the key number less one, with the logo key the number plus seven. Variables can be used as parameters. If no parameters are used, the background only will be changed, and that will be to black.

ASSEMBLY LISTING

- 9	*=\$8352				
10		JSR	\$81F5	!	GET PARAMETER
20		LDA	\$14		
30		AND	#\$0F		
40		STA	\$D021	!	BACKGROUND
50		JSR	\$0079	!	GET LAST BYTE AGAIN
60		BEQ	EXIT		
70		JSR	\$AEFD	!	CHECK COMMA
80		JSR	\$81F5	!	GET PARAMETER
90		LDA	\$14		
100		AND	#\$0F		
110		STA	\$D020	!	BORDER
120		JSR	\$0079	i	GET LAST BYTE AGAIN
130		BEQ	EXIT		
140		JSR	\$AEFD	ļ	CHECK COMMA
150		JSR	\$81F5	!	GET PARAMETER
60		LDA	\$14		
170		AND	#\$0F		
180		STA	\$0286	!	TEXT
190	EXIT	RTS			

8380 EXIT

LINES 10–60: This handles the background colour. We get the first parameter of the command, and load in the low byte only. This is ANDed with \$0F which will set the top four bits to zero no matter what state they were in. The result is used to set the colour. Finally, we

check, by getting the last byte again, if the end of the command has been reached. If it has not, we continue.

LINES 70–130: This first checks for the comma. Then we can get the parameter and proceed as for the background, except to store the value in the border register.

LINES 140–190: This is the same as above except, of course, we set the text colour.

CHAIN – Passing variables

COMMAND SYNTAX

CHAIN ["filename"], [device]

The syntax for CHAIN follows that for LOAD except for the secondary address. No errors will be given for the inclusion of the secondary address, as the routine will overwrite it.

One of the problems of LOAD is that if you load a larger program, after running a smaller program, you overwrite any variables. Also LOAD, if initiated by a direct command, will perform a CLR so you will lose the variables anyway. Sometimes we wish to transfer as many of the variables as possible from one program to another, hence CHAIN.

CHAIN differs from the normal LOAD in two respects. First, it saves the data held in the variable and string areas before the load and restores it afterwards. Secondly, it automatically RUNS the program – obviously it has to be in BASIC.

CHAIN transfers the area of memory holding the variables and arrays to below the string storage. The desired program is loaded and then the data moved back down to the end of the new program. Finally all we have to do is to RUN the program.

Although a fuller, and better, explanation of the way variables are stored is given in Chapter 1, here is a reminder of areas that CHAIN cannot deal with. Defined functions are held in the program, only the pointer is in the variable area, and therefore cannot be transferred. The same applies to strings unless they are concatenated or held in arrays.

There are two listings for this command. The first is entered on the command and it will call the main CHAIN routine. Although the CHAIN routine works as designed we found that, due to the memory move routines, if there were no variables to move you ended up with a page which could contain anything. The first routine will rectify that after the main routine. This also means that CHAIN could be used as a direct command to load and run disk or tape programs.

ASSEMBLY LISTING 1

9	*=\$92B3				
10		LDA	\$32	!	END OF ARRAYS
20		CMP	\$2E	1	CHECK WITH START
					OF VARIABLES
30		BNE	ZERO+1		
40		LDA	\$31		
50		CMP	\$2D		
60		BNE	ZERO+1	!	NOT THE SAME ADDRESSES
70		LDA	#\$80		
80	ZERO	BIT	\$00A9		
90		STA	\$0C	!	\$00 FOR VARIABLES
					\$80 FOR NONE
100		JSR	\$0079	!	GET LAST CHRGET BYTE
110		JMP	\$9080	!	PERFORM CHAIN
120		LDA	\$0C	ļ	GET FLAG
130		BPL	RUN	ļ	VARIABLES
140		DEC	\$30	!	DEC ARRAY ADDS BY PAGE.
150		DEC	\$32		
160	RUN	JMP	\$A7AE		

92D3 RUN

92C1 ZER0

9	*=\$9080			
10		JSR \$E1D4	ļ	GET LOAD PARAMETERS
20		LDA #\$00	!	ENSURE RELOCATING LOAD
30		STA \$B9		
40		JSR \$B526	ļ	GARBAGE COLLECTION
50		LDA \$2D	!	START OF BLOCK TO MOVE
60		STA \$5F		
70		LDA \$2E	!	END OF RESIDENT PROG
80		STA \$60		
90		SEC		
100		LDA \$31	!	END OF BLOCK
110		STA \$5A		
120		SBC \$2F	ļ	CALC AREA OF ARRAYS
130		STA \$FD		
140		LDA \$32	1	ALSO END OF ARRAYS
150		STA \$5B		
160		SBC \$30		
170		STA \$FE		

180 190		LDA	\$33	ł	NEW END OF BLOCK
200		SRC	#401	•	NEW CHO OF DEOCK
210		STA	\$58		
220		IDA	\$34		
230		SRC	#\$00		
240		STA	450		
250		JSR	\$A3RE	ł	PERFORM MOVE
240		INA	\$37	1	SAUE END OF BASIC AREA
270		STA	\$41	•	SHVE LID OF BHOIC HIER
280		INA	\$38		
290		STA	\$422		
200		INA	\$58	I	SAUE REGINNING OF
200		LVH	*30	•	NEW BLOCK
310		STA	\$FB		
320		STA	\$37	!	SET TOP OF BASIC AREA
330		INC	\$59	!	RECTIFY PAGE
340		LDA	\$59		
350		STA	\$FC		
360		STA	\$38		
370		LDX	\$2B	!	SET LOAD ADDRESS
380		LDY	\$20		
390		LDA	#\$00	!	SET FOR LOAD
400		JSR	\$FFD5	1	KERNAL LOAD
410		BCC	STATUS	1	MAYBE GOOD LOAD
420		JMP	\$E0F9	ļ	LOAD ERROR
					DEPENDING ON A
430	STATUS	JSR	\$FFB7	!	READ I/O STATUS WORD
440		AND	#\$BF		
450		BEQ	CONT	ļ	LOAD OK
460		LDX	#\$1D		
470		JMP	\$A437	1	LOAD ERROR
480	CONT	STX	\$2D	!	SET END OF PROGRAM
490		STY	\$2E		
500		STX	\$5F	ł	SET FOR VARIABLE MOVE
510		STY	\$60		
520		LDA	\$FB	1	START OF BLOCK TO MOVE
530		STA	\$5A		
540		LDA	\$FC		
550		STA	\$5B		
560		SEC			
570		LDA	\$33	ł	END OF BLOCK
580		SBC	#501	·	
590		TAY			
600		LDA	\$34		
610		SBC	#\$99		
A. 9 A.		~~~			

620		TAX	
630		TYA	
640		SEC ! CALC AMOUNT	TO MOVE
650		SBC \$5A	
660		STA \$58	
670		TAY ! NO OF BYTES INCOMPLETE F	OF AGE
680		TXA	
690		SBC \$58	
700		TAX ! NO OF PAGES	TO MOVE
710		INX ! FOR EASIER (HECKING
720		ТҮА	
730		BEQ PAGE ! NO SEPARATE	BYTES
740		LDA \$5A ! MOVE SEPARAT BYTES FIRST	E
750		CLC	
760		ADC \$58	
770		STA \$5A	
780		BCC NOINC	
790		INC \$58	
800		CLC	
810	NOINC	LDA \$5F	
820		ADC \$58	
830		STA \$5F	
840		BCC NOINCA	
850		INC \$60	
860	NOINCA	TYA	
870		EOR #\$FF ! 1'S COMPLEME	ENT
880		TAY	
890		INY ! 2'S COMPLEME	ENT
900		DEC \$58	
910		DEC \$60	
920	PAGE	LDA (\$5A),Y	
930		STA (\$5F),Y	
940		INY	
950		BNE PAGE	
960		INC \$5B	
970		INC \$60	
980		DEX ! POINTER FOR	COMPLETION
990		BNE PAGE	
1000		SEC	
1010		LDA \$5F	
1020		STA \$31 ! NEW ARRAY E	ND
1030		SBC \$FD ! CALC ARRAY S	START
1040		STA \$2F	
1050		LDA \$60	

1060	STA \$32	
1070	SBC \$FE	
1080	STA \$30	
1090	LDA \$41	! RESET END OF BASIC
1100	STA \$37	
1110	LDA \$42	
1120	STA \$38	
1130	PLA	
1140	PLA	! REMOVE RETURN ADDRESS
1150	JSR \$A533	! RECHAIN LINES
1160	LDA #\$00	
1170	JSR \$FF90	! TURN OF KERNAL MESSAGES
1180	JSR \$FFE7	! CLALL
1190	JSR \$A677	! END OF CLR
1200	JSR \$A68E	! BACK UP TEXT POINTER
1210	JMP \$92CC	BACK TO FIRST ROUTINE
90E3 CONT	9119 NO	INC
9123 NOINCA	912C PA	GE
90D7 STATUS		

As CHAIN is just moving memory and loading, it is an amalgamation of routines previously described. Where we come across lines used elsewhere, the description will direct you there. By copying lines rather than using subroutines, we make the routine more transportable.

LISTING 1

LINES 10–110: Here we find out if there are variables to move by taking the address of the end of program away from the end of arrays address. On the result we set a flag in location \$0C to \$80 or \$00 denoting variables. We then jump to the main CHAIN routine, LISTING 2.

LINES 120–160: Having returned from the routine, we check our flag by loading and testing the sign flag. A positive result tells us that variables were transferred and no further adjustments are required. If the result was minus, then the addresses denoting the start and end of arrays are reduced by one page, the high byte of the address less one. The final action of CHAIN is to go to ROM, where the next BASIC line is executed. The main routine does the setting up for this just before it comes back to here.

LISTING 2

LINES 10-40: We use the ROM routine to get and set up the loading parameters. To ensure that the load has no secondary address, we

unset that location. The garbage collection routine at \$B526 will tidy up the variable area so that it uses the least space possible.

LINES 50–250: Although the locations from which the addresses are gathered are different, these lines are discussed in Chapter 6, Memory. Moving, lines 1190–1340 (see pages 131–136). There is, however, one extra item involved. To be able to set the start of array address, after loading, we calculate its number of bytes and store it in locations \$FD and \$FE.

LINES 260–360: The data has been moved and now we protect it by changing the pointer to the limit of BASIC. This value will be obtained from the move routine, locations \$58 and \$59, after increasing the latter by one. This is because in the move routine the high byte is decreased before checking for completion. Increasing rectifies this. The original end of BASIC pointer is stored for later use.

LINES 370–470: The loading sequence is covered in the MERGE and APPEND routine in Chapter 7, lines 550–640 (see pages 158–163).

LINES 480–990: Moving the block down is virtually identical to lines 90–550 of Memory Moving in Chapter 6.

LINES 1000–1210: With the major work done, just the clearing up remains. First we calculate the new start of arrays and set its registers. Then we restore the pointer to the end of BASIC. Six ROM routines are visited to finish off the routine. The first two rechain the lines of the BASIC program, so that the interpreter can follow them, and turn off the KERNAL messages. The call to \$FFF7 closes all open files and sets the input/output channels to their default values. The following subroutine is made halfway into CLR. This will do a RESTORE, reset CONT locations, and amend the stack point. The last two routines will do the auto-run. The former sets the CHRGET address to one byte before the program starts. The last one returns us to the calling routine to finish off.

INKEY\$ – A waiting GET

COMMAND SYNTAX

- i INKEY\$
- ii INKEY\$ ""
- iii INKEY\$ A\$ where A\$ is predefined
- iv INKEY\$ "characters"

All commands will stop and wait for a key press. The first two will wait until any key is pressed. The latter two will wait for a key press corresponding to a character within the defined string. The ASCII value of the key press will be placed in the variable sT, and will remain there until an input-output is performed on cassette, serial or RS232.

In 64 BASIC there are two commands for receiving a user input from the keyboard: INPUT and GET. The last accepts a key press without a RETURN but will not wait for one. This entails checking the input and GOTOS until the key press you want is received (see Chapter 4 on checking for function keys in BASIC). INPUT waits for a key but you also have to press RETURN, and the cursor is also in operation.

INKEY\$ will sit and wait for a key press, after emptying the keyboard buffer, and, if required will check for a particular key or keys. To allow for further checks we use the reserved variable 'ST' to store the input. Using ST is easy in that it has a predefined location in zero page.

9	*=\$904E			
10		BNE STRING	3 !	PARAMETERS PRESENT
20	ANYKEY	LDA #\$00	!	CLEAR KEY BUFFER
30		STA \$CB		
40	BYTE	JSR \$FFE4	!	GET CHARACTER
50		BEQ BYTE	ļ	NO KEY
60		STA \$90	!	ST LOCATION
70		RTS		
80	STRING	JSR \$AD9E	!	GET STRING
90		JSR \$B6A3	ļ	DISCARD UNWANTED STRING
100		CMP #\$00	ļ	NULL STRING?
110		BEQ ANYKEY	1	NULL STRING
120		STA \$FB	ļ	NO OF CHARS IN STRING
130		LDA #\$00	ļ	EMPTY KEY BUFFER
140		STA \$C6		
150	BYTEI	JSR \$FFE4	!	GET CHARACTER
160		BEQ BYTE1	!	NO KEY
170		LDY \$FB	!	GET NO. OF CHARS
180		DEY		
190	NEXT	CMP (\$22)	,Y!	CHECK STRING
200		BEQ MATCH	!	FOUND SAME CHAR
210		DEY		
220		BPL NEXT	!	CONTINUE SEARCH
230		BMI BYTE1	!	ANOTHER KEY PRESS
240	MATCH	STA \$90	1	ST LOCATION
250		RTS		

9050	ANYKEY	9054	BYTE
906C	BYTE1	907D	MATCH
9074	NEXT	905C	STRING

LINES 10–70: If there are parameters (cases ii, iii and iv of the command syntax), the zero flag will not be set and these lines are skipped over, at least for the time being. Proceeding on we set the flag for the number of characters in the keyboard buffer to zero. The KERNAL routine at \$FFE4 will return the ASCII value of key presses in the order they were placed in the buffer. If none, then the accumulator will hold zero, so we continue to call the routine until a value is returned. That value is placed in the location which the reserved variable ST uses, and we return to continue the BASIC program.

LINES 80–240: The call to the ROM routine does our string work. It finds the string, especially if it is a variable, determining its length and giving syntax errors if a non-string parameter was supplied. On returning from the routine, the number of characters will be in A and the start address in locations \$22 and \$23. If there were no characters in the string, we branch back to the previous section and wait for any key.

After clearing the buffer and getting a key press value we can check it against the string. The γ register will be loaded with the number of characters and decreased as we check the whole string. If the complete string is checked and no match is found, then the next key press is evaluated. Once a match is found, it is stored in sT and we return to carry on with your program.

LOMEM and HIMEM - Setting the area of work

COMMAND SYNTAX

LOMEM address HIMEM address

The address range that is permissible with these commands is between 1024 and 32767. 'ILLEGAL QUANTITY' errors are given outside this range. The actual start of a program will be one greater than the address given in LOMEM. Commands can be used in direct or program mode.

Changing the memory configuration is a useful, and indeed necessary, task. By raising the base of a program, you can store items such as sprite data, hires screens or even two normal screens and it will not be affected by a program.

At the other end you may wish to put a machine code routine and so to protect it at the top of memory from being overwritten by the variables, so you can set the limit of BASIC to below your routine.

LOMEM will set the lower and HIMEM the upper limit of BASIC. So that they could be used in a loader program the routine does not clear that program. Subsequent programs will be loaded to the new LOMEM address. The ideal place for these commands is at the beginning of a program before any variables are defined. Variables defined after these commands will be placed in the new area. You can use CHAIN to load the next program if there are variables you wish to transfer.

*=\$9	169					
INPU	Т		BNE	GATHER	!	PARAMETERS
			JMP	\$AF08	!	SYNTAX ERROR
GATH	ER		JSR	\$81F5	!	GET PARAMETERS
			LDA	\$15		
			CMP	#\$04	!	CHECK LOW LIMIT
			BCS	TOP	!	O.K.
ERRO	R		LDX	#\$0E	!	ILLEGAL QUANTITY
			JMP	\$A437	!	ERROR ROUTINE
TOP			CMP	#\$80	1	UPPER LIMIT
			BCS	ERROR	!	FAILED
			RTS			
! STI	ART	OF	HIMEN	1		
			JSR	INPUT		
			STA	\$38	1	SET TOP POINTER
			LDA	\$14		
			STA	\$37		
			JMP	\$A65E	!	CLR AND RETURN
! STI	ART	OF	LOMEN	1		
			JSR	INPUT		
			LDY	#\$00		
			TYA			
			STA	(\$14),	Y !	CLEAR FIRST 3 BYTES
			INY			
			STA	(\$14),	Y	
			INY			
			STA	(\$14),	Y	
			LDA	\$14		
			CLC			
			ADC	#\$01		
			STA	\$2B	!	SET START OF BASIC
			TAX			
			LDA	\$15		
			ADC	#\$00		
			STA	\$2C		
			TAY			
			TXA			
			ADC	#\$02		
			STA	\$2D	!	SET START OF VARIABLES
			TYA			
			ADC	#\$00		
	*=\$9 INPU GATH ERRO TOP ! ST ! ST	*=\$9169 INPUT GATHER ERROR TOP ! START ! START	*=\$9169 INPUT GATHER ERROR TOP ! START OF ! START OF	*=\$9169 INPUT BNE JMP GATHER JSR LDA CMP BCS ERROR LDX JMP TOP CMP BCS RTS ! START OF HIMEN JSR STA LDA STA	*=\$9169 INPUT BNE GATHER JMP \$AF08 GATHER JSR \$81F5 LDA \$15 CMP #\$04 BCS TOP ERROR LDX #\$0E JMP \$A437 TOP CMP #\$80 BCS ERROR RTS ! START OF HIMEM JSR INPUT STA \$38 LDA \$14 STA \$37 JMP \$A65E ! START OF LOMEM JSR INPUT LDY #\$00 TYA STA (\$14), INY STA (\$14), INY STA (\$14), INY STA (\$14), INY STA (\$14), INY STA (\$14), INY STA (\$14), INY STA (\$14), INY STA (\$14), INY STA \$2B TAX LDA \$15 ADC #\$00 STA \$2C TAY TXA ADC #\$02 STA \$2D TYA ADC #\$02	*=\$9169 INPUT BNE GATHER ! JMP \$AF08 ! GATHER JSR \$81F5 ! LDA \$15 CMP #\$04 ! BCS TOP ! ERROR LDX #\$0E ! JMP \$A437 ! TOP CMP #\$80 ! BCS ERROR ! BCS ERROR ! RTS ! START OF HIMEM JSR INPUT STA \$38 ! LDA \$14 STA \$37 JMP \$A65E ! ! START OF LOMEM JSR INPUT LDY #\$00 TYA STA (\$14),Y INY STA (\$14),Y INY STA (\$14),Y INY STA (\$14),Y INY STA (\$14),Y INY STA (\$14),Y LDA \$14 CLC ADC #\$01 STA \$2B ! TAX LDA \$15 ADC #\$00 STA \$2C TAY TXA ADC #\$02 STA \$2D ! TYA ADC #\$02 STA \$2D ! TYA ADC #\$00

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410	STA	\$2E				
420	JMP	\$A663	ł	CLR	AND	RETURN

9177	ERROR	916E	GATHER
9169	INPUT	917C	TOP

LINES 10–110: INPUT: This subroutine is used by both commands. It deals with the gathering and checking of addresses. First we check that there is an address. No address, then no command, and a SYNTAX ERROR is given. When the address is picked up, it is first checked for the lower limit and then for the higher.

LINES 130–170: HIMEM: After visiting the input routine, we place the address in the pointers to the limit of BASIC. We then jump to the CLR routine to finish off: this will set all the remaining relevant pointers (such as the string pointer).

LINES 190–420: LOMEM: BASIC requires that the first byte of the BASIC program area is zero (normally 2048, \$0800) and that two zeros signify the end of the program. In the new area these will be together, as there is no program, so we set those first from the address given. To set the start of the program we increase it by one, and from that we add a further two for the address to the start of the variables, or end of program if you prefer. Calling the CLR routine will set the end of variables and array pointers.

QUIT

COMMAND SYNTAX

QUIT

There are no arguments with this particular command.

QUIT disables the UTILITY and its commands, leaving you with the standard BASIC. It does not, however, reset the top of memory back to its original (\$A000). This will leave the UTILITY intact which can be reinitiated by SYS 32768.

QUIT simply restores all the vectors and pointers we changed on start up to their standard values.

*=\$9187						
	LDA	#\$76				
	STA	\$0304	ļ	TOKENISE	BASIC	TEXT
	LDA	#\$A5				
	STA	\$0305				
	*=\$9187	*=\$9187 LDA STA LDA STA	*=\$9187 LDA #\$76 STA \$0304 LDA #\$A5 STA \$0305	*=\$9187 LDA #\$76 STA \$0304 ! LDA #\$A5 STA \$0305	*=\$9187 LDA #\$76 STA \$0304 ! TOKENISE LDA #\$A5 STA \$0305	*=\$9187 LDA #\$76 STA \$0304 ! TOKENISE BASIC LDA #\$A5 STA \$0305

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50	LDA	#\$1A	
60	STA	\$0306 !	BASIC TEXT LIST
70	LDA	#\$E4	
80	STA	\$0308 !	BASIC CHAR DISPATCH
90	LDA	#\$ A 7	
100	STA	\$0307	
110	STA	\$0309	
120	LDA	#\$86	
130	STA	\$030A !	BASIC TOKEN EVALUATION
140	LDA	#\$AE	
150	STA	\$030B	
160	LDA	#\$FE	
170	STA	\$0317 !	BRK INTERRUPT
180	STA	\$0319 !	NMI INTERRUPT
190	LDA	#\$66	
200	STA	\$0316	
210	LDA	#\$47	
220	STA	\$0318	
230	SEI		
240	LDA	#\$48	
250	STA	\$028F !	KEYBOARD TABLE SETUP
260	LDA	#\$EB	
270	STA	\$0290	
280	CLI		
290	PLA		
300	PLA		
310	JMP	\$A474 !	READY FOR BASIC

9 The complete utility

Introduction

We are going to supply the complete UTILITY in the form of a Supermon listing. If you do not possess a monitor, you can find Supermon in the appendices. For the area \$80DE to \$81F4, keywords and vectors, use the M function of the monitor. You may also find it easier to use the memory dump in Chapter 6 for that area. Save to tape or disk regularly as you go.

We had thought of also giving the UTILITY in DATA statement form. This would have come to about 690 lines, of seven items of data on each, which would have been a mammoth task of programming for anyone and very prone to error.

						1						
8000	20	ØF	80	JSR	\$800F	80	3C	85	38		STA	\$38
8003	20	54	89	JSR	\$8954	80	3E	85	34		STA	\$34
8006	20	41	80	JSR	\$8941	80	40	60			RTS	
8009	20	34	80	JSR	\$8034	80	41	A9	7E		LDA	#\$7E
800C	4C	00	92	JMP	\$9200	80	43	80	16	03	STA	\$0316
800F	A9	09		LDA	#\$89	80	46	A9	61		LDA	#\$61
8011	8D	84	03	STA	\$0304	80	48	8D	18	03	STA	\$0318
8014	A9	BC		LDA	#\$BC	80	4 B	A9	89		LDA	#\$88
8016	8D	06	03	STA	\$9396	80	4D	8D	17	03	STA	\$0317
8019	A9	02		LDA	#\$02	80	50	8D	19	03	STA	\$9319
801B	8D	08	03	STA	\$0308	80	53	60			RTS	
801E	A9	29	2 Q	LDA	#\$29	80	54	78			SE1	
8020	8D	8A	03	STA	\$030A	80	55	A9	22		LDA	#\$22
8023	A9	82		LDA	#\$82	80	57	8D	8F	82	STA	\$928F
8025	8D	05	03	STA	\$0305	80	54	A9	87		LDA	#\$87
8028	8D	87	03	STA	\$0307	89	SC.	8D	98	82	STA	\$9298
802B	A9	83		LDA	#\$83	88	SF	58			CLI	-varv
802D	8D	09	03	STA	\$8389	80	60	68			RTS	
8030	8D	ØB	83	STA	\$030B	88	141	48			PHA	
8033	60			RTS		80	62	84			TYA	
8034	A9	FF		LDA	#SFF	89	43	48			PHA	
8034	85	37		STA	\$37	20	44	92			TYA	
8038	85	33		STA	\$33	20	45	40			DUA	
0004		75		LIDA		00	00	TO				
60 JA	ну	14		LUA	##/F	88	66	A7	.7F		LDA	#\$7F

8068	8D	ØD	DD	STA	\$DD0D	80C7	92			???	
806B	AC	0D	DD	LDY	\$DD0D	8008	8C	91	80	STY	\$8091
806E	10	03		BPL	\$8073	80CB	91	4D		STA	(\$4D),Y
8070	4C	72	FE	JMP	\$FE72	80CD	90	FB		BCC	\$80CA
8073	20	BC	F6	JSR	\$F6BC	80CF	85	6E		STA	\$6E
8076	20	E1	FF	JSR	\$FFE1	80D1	88			DEY	
8079	F0	03		BEQ	\$807E	80D2	60			RTS	
807B	4C	72	FE	JMP	\$FE72	80D3	8D	FF	FF	STA	\$FFFF
807E	20	15	FD	JSR	\$FD15	80D6	FF			???	
8081	20	A3	FD	JSR	\$FDA3	80D7	FF			???	
8084	20	18	E5	JSR	\$E518	80D8	00			BRK	
8087	20	54	80	JSR	\$8054	80D9	FF			???	
808A	20	41	80	JSR	\$8041	80DA	FF			???	
808D	6C	02	AØ	JMP	(\$A002)	80DB	FF			???	
8090	98			TYA		80DC	F6	FF		INC	\$FF,X
8091	87			???		80DE	Bó	F7		LDX	\$F7,Y
8092	4C	86	B2	JMP	\$B286	80E0	00			BRK	
8095	83			???		80E1	60			RTS	
8096	9F			???		80E2	00			BRK	
8097	84	EB		STY	\$EB	80E3	00			BRK	
8099	84	36		STY	\$36	80E4	D6	83		DEC	\$83,X
809B	85	BE		STA	\$BE	80E6	D6	83		DEC	\$83,X
809D	85	14		STA	\$14	80E8	00			BRK	
809F	84	51		STY	\$51	80E9	00			BRK	
80A1	83			???		80EA	00			BRK	
80A2	A6	83		LDX	\$83	80EB	68			PLA	
80A4	AE	8F	B4	LDX	\$B48F	80EC	00			BRK	
80A7	8F			???		80ED	00			BRK	
80A8	80			???		80EE	00			BRK	
80A9	83			???		80EF	40			RTI	
80AA	AC	83	51	LDY	\$5183	80F0	00			BRK	
80AD	8E	C4	89	STX	\$89C4	80F1	00			BRK	
80B0	43			???		80F2	00			BRK	
80B1	8F			???		80F3	40			RTI	
8082	A6	87		LDX	\$87	80F4	00			BRK	
80B4	92			???		80F5	40			RTI	
80B5	8B			???		80F6	4F			???	
80B6	2D	84	D1	AND	\$D184	80F7	46	C6		LSR	\$C6
80B9	8F			???		80F9	4B			???	
80 BA	3A			???		80FA	45	D9		EOR	\$D9
80BB	A9	D1		LDA	#\$D1	80FC	44			???	
80BD	A8			TAY		80FD	4F			???	
80BE	30	86		BMI	\$8046	80FE	4B			???	
80C0	Bó	91		LDX	\$91,Y	80FF	C5	54		CMP	\$54
80C2	39	8D	10	AND	\$108D,Y	8101	45	CE		EOR	\$CE
80C5	86	B5		STX	\$B5	8103	54			???	

8104	57			???		8149	54			222	
8105	CF			???		8144	CE			222	
8106	48			PHA		814B	50	52		RUC	\$819F
8107	45	D8		EOR	\$D8	814D	4F	~~		222	40171
8109	42			???		814F	03			222	
810A	49	CE		EOR	#\$CE	814F	44			222	
810C	4F			222		8150	50	52		RUC	\$0104
810D	4C	C4	43	JMP	\$4304	8152	4F	04		222	40104
8110	4F			222		8153	63			222	
8111	4C	4F	55	JMP	\$554F	8154	45	50		FOR	\$50
8114	D2			222		8154	52	00		222	400
8115	57			???		8157	4F			222	
8116	52			???		8158	63			222	
8117	49	54		EOR	#\$54	8159	50	ΔF		RUC	\$8100
8119	C5	43		CMP	\$43	815B	na	51		RNE	491AF
811B	47			222		8150	55	40		FOR	\$49 Y
811C	4F			222		815F	D4	77		222	**/ }/
811D	54			222		81.40	54			222	
811E	CF			222		81.41	52			222	
811F	43			???		81.62	41	43		FOR	(\$43.X)
8120	47			???		8164	C5	52		CMP	\$52
8121	4F			???		8166	45	53		EOR	\$53
8122	53			???		8168	45	D4		EOR	\$D4
8123	55	C2		EOR	\$C2,X	816A	43			???	877 î. j.
8125	50	4C		BVC	\$8173	816B	48			PHA	
8127	4F			???		816C	41	49		EOR	(\$49,X)
8128	D4			???		816E	CE	4C	4F	DEC	\$4F4C
8129	45	4E		EOR	\$4E	8171	4D	45	CD	EOR	\$CD45
812B	54	-		???		8174	48			PHA	
812C	45	D2		EOR	\$D2	8175	49	4D		EOR	#\$4D
812E	44			???		8177	45	CD		EOR	\$CD
8126	55	4D		EUR	\$4D,X	8179	49	4E		EOR	#\$4E
8131	DU	52		BNE	\$8185	817B	4B			???	
8133	40	4E		EUR	\$4E	817C	45	59		EOR	\$59
8133	55	CD		EUR	\$CD,X	817E	A4	4D		LDY	\$4D
8137	44	40		222		8180	45	CD		EOR	\$CD
8138	40	40		EUK	\$40	8182	41	50		EOR	(\$50,X)
0100	40	04 45		EUK	\$04 # 40	8184	50	45		BVC	\$81CB
0130	15	50		COD	⇒4U ¢E0	8186	4E	C4	54	LSR	\$54C4
0130	43	JL		DOR	⊅JZ	8189	52			222	
01/1	T/	10		CMD	#10	818A	41	~ /		222	+0/
8143	45	73		222	940	8188	40	61		LSK	\$60
8144	44			222		0100	HC			222	
8145	45	02		FOP	\$02	OIGE	JH EA			000	
8147	41	55		FOR	(\$55 Y)	0100	5A			222	
91 TI	-7 Å	00		LOK	(#00jA)	0170	HC			111	

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8191	5A	???	81BF 5	5A	???
8192	EA	NOP	81C0 5	5A	???
8193	5A	???	8101 5	5A	???
8194	5A	???	81C2 5	īΑ	???
8195	5A	???	81C3 5	5A	???
8196	5A	???	81C4 5	īΑ	???
8197	5A	???	8105 5	5A	???
8198	5A	???	81C6 E	EA .	NOP
8199	EA	NOP	8107 4	14	???
819A	5A	???	81C8 4	15 45	EOR \$45
819B	5A	???	81CA 0	B	???
819C	5A	???	81CB 0	90	BRK
819D	5A	222	81CC F	F	222
819E	54	222	81CD F	F	222
819F	5A	222	81CE F	F	222
81A0	5A	222	81CF F	F	222
81A1	EA	NOP	BIDA F	F	222
81A2	54	222	81D1 F	F	222
81A3	54	222	81D2 F	F	222
81A4	54	222	81D3 F	F	222
81A5	5A	222	81D4 F	F	222
81A6	EA	NOP	81D5 F	F	222
81A7	5A	???	81D6 F	D FF FF	SBC \$FFFF.X
81A8	5A	???	81D9 F	F	???
81A9	5A	???	81DA F	F	???
81AA	5A	???	81DB F	F	???
81AB	EA	NOP	81DC F	F	???
81AC	5A	???	81DD 7	Έ	???
81AD	5A	???	81DE F	F	???
81AE	5A	???	81DF F	F	???
81AF	5A	???	81E0 0	0	BRK
81 BØ	EA	NOP	81E1 0	0	BRK
81B1	5A	???	81E2 0	0	BRK
81 B 2	5A	???	81E3 Ø	0	BRK
81 B 3	5A	???	81E4 Ø	0	BRK
81B4	5A	???	81E5 0	0	BRK
81B5	5A	???	81E6 0	0	BRK
81B6	5A	???	81E7 0	0	BRK
81B7	EA	NOP	81E8 0	0	BRK
81 B8	5A	???	81E9 0	0	BRK
81 B 9	5A	???	81EA 0	0	BRK
81 BA	5A	???	81EB 0	8	PHP
81 BB	5A	???	81EC 0	1 00	ORA (\$00,X)
81BC	5A	???	81EE 0	0	BRK
81 BD	EA	NOP	81EF 0	0	BRK
81 BE	5A	???	81F0 0	0	BRK

81F1	00			BRK		8245	BD	00	02	LDA	\$0200,X
81F2	00			BRK		8248	38			SEC	
81F3	00			BRK		8249	F9	9E	A0	SBC	\$A09E,Y
81F4	00			BRK		824C	F0	F5		BEQ	\$8243
81F5	20	8A	AD	JSR	\$AD8A	824E	C9	80		CMP	#\$80
81F8	4C	F7	B7	JMP	\$B7F7	8250	D0	30		BNE	\$8282
81 F B	A5	01		LDA	\$01	8252	Ø 5	0B		ORA	\$0B
81 F D	29	FE		AND	#\$FE	8254	A4	71		LDY	\$71
81FF	85	01		STA	\$01	8256	E8			INX	
8201	60			RTS		8257	C8			INY	
8202	A5	01		LDA	\$01	8258	99	FB	01	STA	\$01FB.Y
8204	09	01		ORA	#\$01	825B	B9	FB	01	LDA	\$01FB.Y
8206	85	01		STA	\$01	825E	FØ	59		BEQ	\$82B9
8208	60			RTS		8269	38	- ·		SEC	
8289	Aó	7A		LDX	\$7A	8261	E9	3A		SBC	#\$ 3A
820B	AØ	94		LDY	#\$04	8263	FØ	04		BEQ	\$8269
820D	84	0F		STY	\$0F	8265	C9	49		CMP	#\$49
820F	BD	00	02	LDA	\$0200,X	8267	D0	02		BNE	\$826B
8212	10	87		BPL	\$821B	8269	85	ØF		STA	\$0F
8214	C9	FF		CMP	#\$FF	826B	38			SEC	
8216	F0	ЗE		BEQ	\$8256	826C	E9	55		SBC	#\$55
8218	E8			INX		826E	D0	9F		BNE	\$820F
8219	DØ	F4		BNE	\$820F	8270	85	08		STA	\$08
821B	C9	20		CMP	#\$20	8272	BD	00	02	LDA	\$0200,X
821D	F0	37		BEQ	\$8256	8275	F0	DF		BEQ	\$8256
821F	85	0 8		STA	\$08	8277	C5	08		CMP	\$08
8221	C9	22		CMP	#\$22	8279	F0	DB		BEQ	\$8256
8223	F0	56		BEQ	\$827B	827B	C8			INY	
8225	24	0F		BIT	\$0F	8270	99	FB	01	STA	\$01FB,Y
8227	70	2D		BVS	\$8256	827F	E8			INX	
8229	C9	ЗF		CMP	#\$3F	8280	D0	FØ		BNE	\$8272
822B	D0	04		BNE	\$8231	8282	A6	7A		LDX	\$7A
822D	A9	99		LDA	#\$99	8284	E۶	0B		INC	\$0B
822F	D0	25		BNE	\$8256	8286	C8			INY	
8231	C9	30		CMP	#\$30	8287	B9	9D	A0	LDA	\$A09D,Y
8233	90	04		BCC	\$8239	828A	10	FA		BPL	\$8286
8235	C9	3C		CMP	#\$3C	828C	B9	9E	AØ	LDA	\$A09E,Y
8237	90	1 D		BCC	\$8256	828F	D0	B4		BNE	\$8245
8239	84	71		STY	\$71	8291	AØ	FF		LDY	# \$FF
823B	A0	00		LDY	#\$00	8293	CA			DEX	
823D	84	0B		STY	\$0B	8294	C8			INY	
823F	88			DEY		8295	E8			INX	
8240	86	7A		STX	\$7A	8296	BD	00	02	LDA	\$0200,X
8242	CA			DEX		8299	38			SEC	•
8243	C8			INY		829A	F9	F6	80	SBC	\$80F6,Y
8244	E8			INX		829D	F0	F5		BEQ	\$8294

829F	C9	80		CMP	#\$80	82F7	20	47	AB	JSR	\$AB47
82A1	DØ	02		BNE	\$82A5	82FA	D0	F6		BNE	\$82F2
82A3	F0	AD		BEQ	\$8252	82FC	4C	F3	Aó	JMP	\$A6F3
82A5	A6	7A		LDX	\$7A	82FF	4C	EF	A6	JMP	\$A6EF
82A7	Еó	Ø B		INC	\$0B	8302	20	73	00	JSR	\$0073
82A9	C8			INY		8305	C9	CC		CMP	#\$CC
82AA	B9	F5	80	LDA	\$80F5,Y	8307	90	1A		BCC	\$8323
82AD	10	FA		BPL	\$82A9	8309	C9	EE		CMP	#\$EE
82AF	B9	F6	80	LDA	\$80F6,Y	830B	B Ø	16		BCS	\$8323
82B2	D0	E2		BNE	\$8296	830D	20	13	83	JSR	\$8313
82B4	BD	00	02	LDA	\$0200,X	8310	4C	EA	A7	JMP	\$A7EA
82B7	10	9B		BPL	\$8254	8313	38			SEC	
82B9	4C	89	A6	JMP	\$A689	8314	E9	CC		SBC	#\$CC
82BC	10	ЗE		BPL	\$82FC	8316	8A			ASL	
82BE	C9	FF		CMP	#\$FF	8317	A 8			TAY	
82C0	F0	ЗA		BEQ	\$82FC	8318	B9	91	80	LDA	\$8091,Y
82C2	24	ØF		BIT	\$0F	831B	48			Pha	•
82C4	30	36		BMI	\$82FC	831C	B9	90	80	LDA	\$8090,Y
8206	C9	CC		CMP	#\$CC	831F	48			PHA	·
82C8	90	0E		BCC	\$82D8	8320	4C	73	00	JMP	\$0073
82CA	38			SEC		8323	20	79	00	JSR	\$0079
82CB	E9	СВ		SBC	#\$CB	8326	4C	E7	A7	JMP	\$A7E7
82CD	AA			TAX		8329	A9	00		LDA	#\$00
82CE	A9	F6		LDA	#\$F6	832B	85	0D		STA	\$0D
82D0	85	22		STA	\$22	832D	20	73	00	JSR	\$0073
82D2	A9	80		LDA	#\$80	8330	C9	F7		CMP	#\$F7
8204	85	23		STA	\$23	8332	90	18		BCC	\$834C
8206	DU	0C		BNE	\$82E4	8334	C9	F8		CMP	#\$F8
8208	38			SEC		8336	B0	14		BCS	\$834C
8209	EY	7F		SBC	#\$7F	8338	20	30	83	JSR	\$833C
8208	AA AO	05			***	833B	60			RTS	
8200	AY	7E		LUA	#\$YE	833C	38			SEC	
02VE	80	22		514	¥ΖΖ ₩# ΔΩ	833D	E9	F6		SBC	#\$F6
0250	H7 05	H0 22		CTA	# > +10	833F	ØA			ASL	
9254	0.0	23 10		CTV	₹23 ¢10	8340	H8 D0	__	00		toore V
9754	Δ 0				#47 #455	0341	87	ED	80		\$80E3,1
8258		гг		DEY	₩₽ГГ	0344	48	54	00	PHA	ADDEA V
8259	FA	97		DEA	40252	8343	87	E4	80	LUA	\$80E4,1
82FR	60	07		TNY	₽OZFZ	0340	48	70	00	PHH MD	*0070
8250	D1	22			(422) V	8347	46	73	00	JMP	\$0073 #0070
92EE	10				(#22);1 400ED	0346	20	77	00		300/Y
82EP	20	FX		DI'L RMT	40200	0345	70	500	01	JEP	PHEOU 40155
8252	60	10			*0100	0332	70	ГЈ 1 Л	01	158	₽01LD 417
82F2	RI	22			(422) V	0333	нJ 20	14			₽14 ##0E
8255	20	<u>6</u> 0		DMT	49755	033/	27 01	9F 24	De	HNU	#701 #10001
טבו ט	50	00		0011	#OTLL	0337	ov	21	00	51A	₽U021

835C	20	79	00	JSR	\$0079	8309	A6	64		LDX	\$64
835F	F0	1F		BEQ	\$8380	83CB	A 0	00		LDY	#\$00
8361	20	FD	AE	JSR	\$AEFD	83CD	91	14		STA	(\$14),Y
8364	20	F5	81	JSR	\$81F5	83CF	C8			INY	
8367	A5	14		LDA	\$14	83D0	8A			TXA	
8369	29	0F		AND	#\$0F	83D1	91	14		STA	(\$14),Y
836B	8D	20	D0	STA	\$D020	83D3	60			RTS	
836E	20	79	00	JSR	\$0079	83D4	4C	48	B2	JMP	\$B248
8371	F0	0D		BEQ	\$8380	83D7	A5	15		LDA	\$15
8373	20	FD	AE	JSR	\$AEFD	83D9	48			Pha	
8376	20	F5	81	JSR	\$81F5	83DA	A5	14		LDA	\$14
8379	A5	14		LDA	\$14	83DC	48			PHA	
837B	29	0F		AND	#\$0F	83DD	20	FA	AE	JSR	\$AEFA
837D	8D	86	82	STA	\$0286	83E0	20	F5	81	JSR	\$81F5
8380	60			RTS		83E3	20	F7	AE	JSR	\$AEF7
8381	20	FA	AE	JSR	\$AEFA	83E6	A0	01		LDY	#\$01
8384	20	F5	81	JSR	\$81F5	83E8	B1	14		LDA	(\$14),Y
8387	A5	14		LDA	\$14	83EA	AA			TAX	
8389	C9	28		CMP	#\$28	83EB	88			DEY	
838B	90	03		BCC	\$8390	83EC	B1	14		LDA	(\$14),Y
838D	20	48	B2	JSR	\$B248	83EE	A8			TAY	
8390	48			PHA		83EF	68			PLA	
8391	20	FD	AE	JSR	\$AEFD	83F0	85	14		STA	\$14
8394	20	F5	81	JSR	\$81F5	83F2	68			PLA	
8397	Að	14		LDX	\$14	83F3	85	15		STA	\$15
8399	E0	19		CPX	#\$19	83F5	8A			TXA	
839B	B0	F0		BCS	\$838D	83F6	20	91	B3	JSR	\$B391
839D	68			PLA		83F9	20	01	84	JSR	\$8401
839E	A8			TAY		83FC	68			PLA	
839F	18			CLC		83FD	68		. –	PLA	
83A0	20	FØ	FF	JSR	\$FFF0	83FE	40	8D	AD	JMP	\$AD8D
83A3	20	73	00	JSR	\$0073	8401	A5	66		LDA	\$66
8346	60			RTS		8403	10	ØA		BPL	\$8401
83A7	20	81	83	JSR	\$8381	8405	AU	84		LDY	#\$84
8366	40	AØ	AA	JMP	\$AAA0	8407	AY	10	-	LDA	#\$10
83AD	20	81	83	JSR	\$8381	8409	20	80	BA	JSR	\$BASC
8380	40	BF	AB	JMP	\$ABBF	8400	20	6A	88	JSR	\$889 0
83B3	20	F5	81	JSR	\$81F5	840F	60	•••		RIS	
8386	20	FD	AE	JSR	\$AEFD	8410	91	66		SIA	(\$00),Y
8389	20	80	AD	JSR	\$adsa	8412	00			BKK	
83BC	A5	66		LDA	\$66	8413	00	t.		BKK	
83BE	30	14		BWI	\$83D4	8414	96			RKK	****
0360	7 با م	71		011P	#771 20207	8415	A9	FF		LDA	#\$25
0362	20	10	PC	510 100	#0304 & DCOD	8417	AU	61		LDY	#\$01
0364	20	70	DL	1 001	₽0070 445	8419	91	2B		STA	(\$2B),Y
0367	нэ	93		LUH	40J	841B	20	33	A5	JSR	\$A533

841E	A5	22		LDA	\$22	8485	38			SEC	
8420	18			CLC		8486	20	49	BC	JSR	\$BC49
8421	69	02		ADC	#\$02	8489	20	DF	BD	JSR	\$BDDF
8423	85	2D		STA	\$2D	848C	20	87	B4	JSR	\$B487
8425	A5	23		LDA	\$23	848F	20	A6	B6	JSR	\$B6A6
8427	69	00		ADC	#\$00	8492	A2	00		LDX	#\$00
8429	85	2E		STA	\$2E	8494	BD	00	01	LDA	\$0100,X
842B	4C	60	A6	JMP	\$A660	8497	9D	00	02	STA	\$0200 X
842E	20	F5	81	JSR	\$81F5	849A	F0	03		BEQ	\$849F
8431	20	FD	AE	JSR	\$AEFD	849C	E8			INX	
8434	A5	14		LDA	\$14	849D	D0	F5		BNE	\$8494
8436	85	FB		STA	\$FB	849F	60			RTS	
8438	A5	15		LDA	\$15	8440	FØ	47		BEQ	\$84E9
843A	85	FC		STA	\$FC	84A2	B0	45		BCS	\$84E9
843C	20	F5	81	JSR	\$81F5	8444	20	F5	81	JSR	\$81F5
843F	A5	14		LDA	\$14	84A7	A9	20		LDA	#\$20
8441	85	FD		STA	\$FD	8449	20	D2	FF	JSR	\$FFD2
8443	A9	4D		LDA	#\$4D	84AC	A0	02		LDY	#\$02
8445	8D	02	03	STA	\$0302	84AE	20	D1	84	JSR	\$84D1
8448	A9	84		LDA	#\$84	84B1	88			DEY	
844A	8D	83	03	STA	\$0303	84B2	D0	FA		BNE	\$84AE
844D	AD	00	02	LDA	\$0200	84B4	A4	14		LDY	\$14
8450	F0	20		BEQ	\$8472	84B6	84	15		STY	\$15
8452	A6	FB		LDX	\$FB	84B8	AØ	02		LDY	#\$02
8454	A5	FC		LDA	\$FC	84BA	20	D1	84	JSR	\$84D1
8456	20	7F	84	JSR	\$847F	84BD	88			DEY	
8459	86	63		STX	\$C6	84BE	D0	FA		BNE	\$84BA
845B	BD	00	02	LDA	\$0200,X	8400	20	79	00	JSR	\$0079
845E	9D	77	02	STA	\$0277 X	84C3	C9	2C		CMP	#\$2C
8461	CA			DEX		84C5	DØ	89		BNE	\$84D0
8462	10	F7		BPL	\$845B	84C7	20	D2	FF	JSR	\$FFD2
8464	18			CLC		84CA	20	73	89	JSR	\$0073
8465	A5	FB		LDA	\$FB	84CD	4C	A0	84	JMP	\$84A0
8467	65	FD		ADC	\$FD	84D0	60			RTS	
8469	85	FB		STA	\$FB	84D1	A2	84		LDX	#\$04
846B	90	02		BCC	\$846F	84D3	A9	00		LDA	#\$00
846D	E۵	FC		INC	\$FC	84D5	86	15		ASL	\$15
846F	4C	83	A4	JMP	\$ A483	84D7	2A			ROL	
8472	A9	83		LDA	#\$83	84D8	CA			DEX	
8474	8 D	82	03	STA	\$0302	84D9	D0	FA		BNE	\$84D5
8477	A9	A4		LDA	#\$ A4	84DB	C9	0A		CMP	#\$0A
8479	8D	03	03	STA	\$0303	84DD	90	04		BCC	\$84E3
847C	6C	02	03	JMP	(\$0302)	84DF	18			CLC	
847F	86	63		STX	\$63	84E0	69	37		ADC	#\$37
8481	85	62		STA	\$62	84E2	2C	69	30	BIT	\$3069
8483	A2	90		LDX	#\$90	84E5	20	D2	FF	JSR	\$FFD2
						1					

84E8	60			RTS	1	8550	85	62		STA	\$62
84E9	4C	88	AF	JMP	\$AF08	8552	20	7C	85	JSR	\$8570
84EC	F0	46		BEQ	\$8534	8555	A 8			TAY	
84EE	B0	44		BCS	\$8534	8556	68			PLA	
84FØ	20	F5	81	JSR	\$81F5	8557	20	AD	85	JSR	\$85AD
84F3	A9	20		LDA	#\$20	855A	20	73	8 8	JSR	\$0073
84F5	20	D2	FF	JSR	\$FFD2	855D	C9	2C		CMP	#\$2C
84F8	A5	15		LDA	\$15	855F	FØ	10		BEQ	\$8571
84FA	F0	16		BEQ	\$8512	8561	60			RTS	
84FC	A2	0 8		LDX	#\$08	8562	68			PLA	
84FE	86	15		ASL	\$15	8563	A8			TAY	
8500	B Ø	03		BCS	\$8505	8564	A9	00		LDA	#\$00
8502	A9	30		LDA	#\$30	8566	20	AD	85	JSR	\$85AD
8504	2C	A9	31	BIT	\$31A9	8569	60			RTS	
8507	20	D2	FF	JSR	\$FFD2	856A	68			PLA	
850A	CA			DEX		856B	A8			TAY	
850B	D0	F1		BNE	\$84FE	856C	A9	00		LDA	#\$00
850D	A9	20		LDA	#\$20	856E	20	AD	85	JSR	\$85AD
850F	20	D2	FF	JSR	\$FFD2	8571	A9	2C		LDA	#\$2C
8512	A2	08		LDX	#\$08	8573	20	D2	FF	JSR	\$FFD2
8514	86	14		ASL	\$14	8576	20	73	00	JSR	\$0073
8516	B0	03		BCS	\$851B	8579	4C	37	85	JMP	\$8537
8518	A9	30		LDA	#\$30	857C	A0	01		LDY	#\$01
851A	2C	A9	31	BIT	\$31A9	857E	89	62	00	LDA	\$0062,Y
851D	20	D2	FF	JSR	\$FFD2	8581	C9	30		CMP	#\$30
8520	CA			DEX		8583	90	25		BCC	\$85AA
8521	DØ	F1		BNE	\$8514	8585	C9	47		CMP	#\$47
8523	20	79	0 0	JSR	\$0079	8587	B0	21		BCS	\$85AA
8526	C9	2C		CMP	#\$ 2C	8589	C9	ЗA		CMP	#\$3A
8528	DØ	09		BNE	\$8533	858B	90	68		BCC	\$8595
852A	20	D2	FF	JSR	\$FFD2	858D	C9	41		CMP	#\$41
852D	20	73	00	JSR	\$0073	858F	90	19		BCC	\$85AA
8530	40	EC	84	JMP	\$84EC	8591	EY	37		SBC	#\$37
8533	60		. –	RTS		8373	00	63		BINE	\$83 28
8534	40	68	AF	JMP	\$AF08	8373	38	20		SEL	H+ 00
8537	85	63	~~	SIA	\$63	8376	EY	30	0.0	580	#>30
8539	20	/3	88	JSR	\$0073	8378	77	14		SIA	\$0014,T
8536	80	- ŏ∠ - ⊃C	05	51A	\$6Z	8378	88	F 0		DET	*0575
0535	20	16	80	JSK	\$837 C	0500	10	C 0		BFL	9837E ##04
0541	48	70	00	TOD	#0070	OJ7E	H0 0 4	15			##84
0542	20	10	99	JOR	700/3	0542	00	10		HOL	\$10
0343	г Ø С О	10			¥¢00¢∠ ₩¢0C	0542	00	50		DNE	40540
0540	C7				₩₽∠U 405/A	0JHJ 05AE	70	11			#0.J#10 ¢1./
0J47 0540		10		CTA	≠0J0H ¢∡2	0547		15			415
854D	20	03 72	88	JOD	+03 40072	0JH/	20 20	IJ		DTC	41J
0040	20	r J		USR	400/J	USH/				ктэ	

85AA	4C	88	AF	JMP	\$AF08	8611	D0	03		BNE	\$8616
85AD	20	91	B3	JSR	\$B391	8613	4C	1 D	A8	JMP	\$A81D
85B0	20	81	84	JSR	\$8401	8616	20	F5	81	JSR	\$81F5
85B3	20	8D	AD	JSR	\$AD8D	8619	20	13	A6	JSR	\$A613
85B6	20	DD	BD	JSR	\$BDDD	861C	BØ	05		BCS	\$8623
85B9	20	87	B4	JSR	\$B487	861E	A2	15		LDX	#\$15
85BC	4C	21	AB	JMP	\$AB21	8620	4C	37	A4	JMP	\$A437
85BF	A5	7A		LDA	\$7A	8623	38			SEC	
85C1	D0	02		BNE	\$8505	8624	A5	5F		LDA	\$5F
85C3	C6	7B		DEC	\$7B	8626	E9	01		SBC	#\$01
8505	Сð	7A		DEC	\$7A	8628	85	41		STA	\$41
85C7	A2	08		LDX	#\$08	862A	A5	60		LDA	\$60
8509	20	73	88	JSR	\$0073	862C	E9	00		SBC	#\$00
85CC	90	03		BCC	\$85D1	862E	85	42		STA	\$42
85CE	4C	0 8	AF	JMP	\$AF08	8630	60			RTS	
85D1	C9	32		CMP	#\$32	8631	A9	FF		LDA	#\$FF
85D3	BØ	F9		BCS	\$85CE	8633	85	4 A		STA	\$4A
85D5	óА			ROR		8635	20	8A	A3	JSR	\$A38A
85D6	26	14		ROL	\$14	8638	9A			TXS	
85D8	CA			DEX		8639	C9	8D		CMP	#\$8D
85D9	D0	EE		BNE	\$8509	863B	F0	05		BEQ	\$8642
85DB	A4	14		LDY	\$14	863D	A2	16		LDX	#\$16
85DD	A9	00		LDA	#\$00	863F	4C	37	A4	JMP	\$A437
85DF	20	AD	85	JSR	\$85AD	8642	E8			INX	
85E2	A9	2F		LDA	#\$2F	8643	E8			INX	
85E4	20	D2	FF	JSR	\$FFD2	8644	E8			INX	
85E7	A5	14		LDA	\$14	8645	E8			INX	
85E9	A0	00		LDY	#\$00	8646	E8			INX	
85EB	20	AD	85	JSR	\$85AD	8647	9A			TXS	
85EE	20	73	00	JSR	\$0073	8648	60			RTS	
85F1	C9	2C		CMP	#\$2C	8649	60			RTS	
85F3	D0	0 6		BNE	\$85FB	864A	4C	0 8	AF	JMP	\$AF08
85F5	20	D2	FF	JSR	\$FFD2	864D	AD	5B	80	LDA	\$805B
85F8	4C	C7	85	JMP	\$85C7	8650	C9	87		CMP	#\$87
85FB	60			RTS		8652	FØ	ØD		BEQ	\$8661
85FC	A5	9D		LDA	\$9D	8654	A9	87		LDA	#\$87
85FE	DØ	01		BNE	\$8601	8656	8D	5B	80	STA	\$805B
8966	60			RTS		8659	A9	22		LDA	#\$22
8601	20	26	85	JSR	\$B526	865B	8D	56	80	STA	\$8056
8684	38	~~		SEC		865E	20	54	80	JSR	\$8054
8600	AD	33		LDA	\$33	8661	20	79	88	JSR	\$0079
0400	E0 70	31		JDL TAY	\$ 31	8664	FØ	42		BEQ	\$8668
8484	Δ5	24			\$34	8666	20	F5	81	JSR	\$81F5
8400	F5	32		SBC	\$32	8669	20	FD	AL	JSR	SALLD
SAAF	40	AD	85	JMP	\$85AD	3000	A2	14		LUA	¥14
	.0			0 . II		SOOF	F 0	1F		RFR	⊅ 808F

8670	C9	11		CMP	#\$11	86C7	90	F4		BCC	\$86BD
8672	B0	1A		BCS	\$868E	8609	A9	31		LDA	#\$31
8674	63	14		DEC	\$14	86CB	85	22		STA	\$22
8676	A5	14		LDA	\$14	86CD	A9	30		LDA	#\$30
8678	8A			ASL		86CF	85	23		STA	\$23
8679	0A			ASL		86D1	20	DE	86	JSR	\$86DE
867A	8A			ASL		86D4	Еð	23		INC	\$23
867B	0A			ASL		86D6	E6	5F		INC	\$5F
867C	A 8			TAY		86D8	E8			INX	
867D	A9	A1		LDA	#\$A1	86D9	EØ	11		CPX	#\$11
867F	85	15		STA	\$15	86DB	90	F4		BCC	\$86D1
8681	A9	00		LDA	#\$00	86DD	60			RTS	
8683	85	14		STA	\$14	86DE	A0	05		LDY	#\$05
8685	A6	0A		LDX	\$0A	86E0	B9	1 C	87	LDA	\$871C,Y
8687	20	79	00	JSR	\$0079	86E3	20	D2	FF	JSR	\$FFD2
868A	C9	22		CMP	#\$22	86E6	88			DEY	
868C	F0	03		BEQ	\$8691	86E7	D0	F7		BNE	\$86E0
898E	4C	88	AF	JMP	\$AF08	86E9	A5	22		LDA	\$22
8691	20	73	99	JSR	\$0073	86EB	20	D2	FF	JSR	\$FFD2
8694	F0	0A		BEQ	\$86A0	86EE	A5	23		LDA	\$23
8696	C9	22		CMP	#\$22	86F0	20	D2	FF	JSR	\$FFD2
8698	F0	86		BEQ	\$86A0	86F3	A9	2C		LDA	#\$2C
869A	91	14		STA	(\$14),Y	86F5	20	D2	FF	JSR	\$FFD2
869C	C8			INY		86F8	A9	22		LDA	#\$22
869D	CA			DEX		86FA	20	D2	FF	JSR	\$FFD2
869E	D0	F1		BNE	\$8691	86FD	A5	5F		LDA	\$5F
86A0	A9	00		LDA	#\$00	86FF	ØA			ASL	
86A2	91	14		STA	(\$14),Y	8700	ØA			ASL	
86A4	20	73	00	JSR	\$0073	8701	ØA			ASL	
86A7	60			RTS		8702	0A			ASL	
8648	A2	00		LDX	#\$00	8703	A8			TAY	
8644	86	5F		STX	\$5F	8704	20	FB	81	JSR	\$81FB
86AC	E8			INX		8707	81	14		LDA	(\$14),Y
86AD	A9	20		LDA	#\$20	8/09	48		~~	PHA	
86AF	85	22		STA	\$22	870A	20	02	82	JSR	\$8202
86B1	A9	31		LDA	#\$31	8700	68	• •		PLA	
86B3	85	23		STA	\$23	870E	F0	96		BEQ	\$8716
86B5	A9	00		LDA	#\$00	8/10	20	02	FF	JSR	\$FFD2
86B7	85	14		STA	\$14	8713	C8			INY	
86B9	A9	A1		LDA	#\$A1	8714	DØ	EE		BNE	\$8704
86BB	85	15		STA	\$15	8716	A9	22		LDA	#\$22
86BD	20	DE	86	JSR	\$86DE	8718	20	D2	FF	JSR	\$FFD2
8908	Еó	23		INC	\$23	871B	60			RTS	
86C2	E6	5F		INC	\$5F	871C	20	59	45	JSR	\$4559
86C4	E8			INX		871F	4B			???	
8605	E0	0A		СРХ	#\$0A	8720	20	0D	A4	JSR	\$A40D
						1					

8723	CB			???		8781	68			PLA	
8724	C0	03		CPY	#\$03	8782	F0	0D		BEQ	\$8791
8726	90	04		BCC	\$872C	8784	C9	5F		CMP	#\$5F
8728	C0	07		CPY	#\$07	8786	D0	02		BNE	\$878A
872A	90	83		BCC	\$872F	8788	A9	0D		LDA	#\$0D
872C	4C	48	EB	JMP	\$EB48	878A	9D	77	02	STA	\$0277,X
872F	AD	8D	02	LDA	\$028D	878D	E8			INX	•
8732	C4	C5		CPY	\$C5	878E	C 8			INY	
8734	D0	05		BNE	\$873B	878F	D0	E7		BNE	\$8778
8736	CD	8E	02	CMP	\$028E	8791	86	63		STX	\$C6
8739	F0	F1		BEQ	\$872C	8793	A9	7F		LDA	#\$7F
873B	84	C5		STY	\$C5	8795	8D	00	DC	STA	\$DC00
873D	8D	8E	02	STA	\$028E	8798	60			RTS	
8740	C0	04		CPY	#\$04	8799	A9	48		LDA	#\$48
8742	F0	0B		BEQ	\$874F	879B	8D	56	80	STA	\$8056
8744	C0	05		CPY	#\$05	879E	A9	EB		LDA	#\$EB
8746	F0	0A		BEQ	\$8752	87A0	8D	5B	80	STA	\$8058
8748	C0	06		CPY	#\$06	87A3	20	54	80	JSR	\$8054
87 4 A	F0	89		BEQ	\$8755	87A6	60			RTS	
874C	A0	07		LDY	#\$07	87A7	20	5E	88	JSR	\$885E
874E	2C	A0	01	BIT	\$01A0	87AA	86	2B		STX	\$2B
8751	2C	A0	03	BIT	\$03A0	87AC	84	2C		STY	\$2C
8754	2C	A0	05	BIT	\$05A0	87AE	A0	00		LDY	#\$88
8757	C9	02		CMP	#\$02	87B0	98			TYA	
8759	90	07		BCC	\$8762	87B1	91	2B		STA	(\$2B),Y
875B	F0	03		BEQ	\$8760	87B3	20	10	88	JSR	\$8810
875D	A9	09		LDA	#\$09	87B6	86	2D		STX	\$2D
875F	2C	A9	0 8	BIT	\$08A9	87B8	84	2E		STY	\$2E
8762	84	BB		STY	\$BB	87BA	20	33	A5	JSR	\$A533
8764	C6	BB		DEC	\$BB	87BD	20	4D	88	JSR	\$884D
8766	18			CLC	. = =	87C0	A9	F9		LDA	#\$F9
8767	65	BB		ADC	\$88	8702	A2	87		LDX	#\$87
8/69	NA 0.4			ASL		8764	80	82	03	SIA	\$0302
876A	04			ASL		0707	OE AO	03	03	517	\$0303
8768	00			ASL		070H	H7 05	70		CTA	#>01 # 70
07/0	0H			HOL	H# A 1	9700	70 20	70		1 DA	₽/D #dEE
0700	H0 04	+ E		CTY	#PH1	0702	05	7		CTA	##FF #74
070F	04 A0	13		311	₩400	9702	0J Δ0	00		JINV	₽/H ##00
0773	Н0 0Л	14		CTV	#700 ¢1/	8704	Ri	FR			#₽00 (450) V
9775	04 <u>0</u> 0	14		TAV	414	8704	85	FD		STA	\#ED/jI \$ED
8774	Δ2	۵a			#4.00	8708	60			INY	₩1 ₩
8779	20	FR	81	JCP	\$81FR	8709	B1	FR			(\$FR) Y
877P	RI	14			(\$14) V	87DR	85	FF		STA	\$FF
8770	48	• •		PHA	\ + + 17 j l	8700	FA	24		BED	\$8883
877F	20	82	82	JSR	\$8202	87DF	C8	- 1		INY	
21 I L	~~	~~		0.011	·vavb						

87E0	B1	FB		LDA	(\$FB),Y	883F	A5	FC		LDA	\$FC
87E2	85	14		STA	\$14	8841	E9	00		SBC	#\$00
87E4	C8			INY		8843	85	FC		STA	\$FC
87E5	B1	FB		LDA	(\$FB),Y	8845	A9	00		LDA	#\$00
87E7	85	15		STA	\$15	8847	A 8			TAY	
87E9	A2	04		LDX	#\$04	8848	91	14		STA	(\$14),Y
87EB	E8			INX		884A	C8			INY	
87EC	C8			INY		884B	91	14		STA	(\$14),Y
87ED	B1	FB		LDA	(\$FB),Y	884D	A5	FD		LDA	\$FD
87EF	9D	FB	01	STA	\$01FB,X	884F	A6	FE		LDX	\$FE
87F2	D0	F7		BNE	\$87EB	8851	85	2B		STA	\$2B
87F4	8A			TXA		8853	86	2C		STX	\$2C
87F5	A8			TAY		8855	A5	FB		LDA	\$FB
87F6	20	A2	A4	JSR	\$A4A2	8857	A6	FC		LDX	\$FC
87F9	A5	FD		LDA	\$FD	8859	85	2D		STA	\$2D
87FB	A6	FE		LDX	\$FE	885B	86	2E		STX	\$2E
87FD	85	FB		STA	\$FB	885D	60			RTS	
87FF	86	FC		STX	\$FC	885E	A5	2B		LDA	\$2B
8801	DØ	C7		BNE	\$87CA	8860	85	FD		STA	\$FD
8803	A9	83		LDA	#\$83	8862	A5	2C		LDA	\$2C
8805	A2	A4		LDX	#\$ A4	8864	85	FE		STA	\$FE
8807	8D	02	03	STA	\$0302	8866	A6	2D		LDX	\$2D
880A	8E	03	83	STX	\$0303	8868	A4	2E		LDY	\$2E
880D	20	74	A4	JSR	\$A474	886A	86	FB		STX	\$FB
8810	20	D4	E1	JSR	\$E1D4	884C	84	FC		STY	\$FC
8813	A9	00		LDA	#\$00	886E	60			RTS	
8815	85	B9		STA	\$B9	886F	20	5E	88	JSR	\$885E
8817	A6	2B		LDX	\$2B	8872	8A			TXA	
8819	A4	2C		LDY	\$2C	8873	38			SEC	
881B	20	D5	FF	JSR	\$FFD5	8874	E9	02		SBC	#\$02
881E	B0	10		BCS	\$8830	8876	85	2B		STA	\$2B
8820	20	B7	FF	JSR	\$FFB7	8878	98			TYA	
8823	29	BF		AND	#\$BF	8879	E9	00		SBC	#\$88
8825	FØ	08		BEQ	\$882F	887B	85	20	~~	SIA	\$20
8827	20	38	88	JSR	\$8838	887D	20	10	88	JSR	\$8810
882A	A2	1D		LDX	#\$1D	8888	86	FB		SIX	*FB
882C	40	37	A4	JMP	\$A437	8882	84	FL	~~	511	*FL
882F	60			RIS		8884	20	40	88	JSK	\$884U
8830	48	~~	~~	PHA		8887	20	33	AD	JSR	\$A033
8831	20	38	88	JSR	\$8838	8884	00	~~		RIS OTV	****
0034	00 4C	сo	50		&F0F0	8888	80			517	キレビ
0000	Δ5	EP.	20		¢ER	8880	H3	3E		LUH	⊅JE
8832	38			SEC		0000	- 38 55	<u> </u>		SEU	#C2
883B	F9	82		SBC	#\$82	0070	CJ 05			300	₽し∠ 400
8830	85	14		STA	\$14	0072	00	DD		51A	₽ DD
0000	00	• '				0074	10			LLL	
8895	A5	FB	LDA	\$FB	88E9	E٥	60	INC	\$60		
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8897	65	49	ADC	\$49	88E8	CA		DEX			
8899	85	5F	STA	\$5F	88E9	DØ	F2	BNE	\$88DD		
889B	A5	FC	LDA	\$FC	88EB	38		SEC			
889D	69	00	ADC	#\$00	88EC	A5	2D	LDA	\$2D		
889F	85	60	STA	\$60	88EE	E5	BB	SBC	\$BB		
88A1	A5	5F	LDA	\$5F	88F0	85	2D	STA	\$2D		
88A3	65	BB	ADC	\$BB	88F2	BØ	03	BCS	\$88F7		
88A5	85	5A	STA	\$5A	88F4	63	2E	DEC	\$2E		
88A7	A5	60	LDA	\$60	88F6	38		SEC			
88A9	69	8 0	ADC	#\$00	88F7	AØ	88	LDY	#\$00		
88AB	85	5B	STA	\$5B	88F9	A5	FD	LDA	\$FD		
88AD	A5	2D	LDA	\$2D	88FB	E5	BB	SBC	\$BB		
88AF	38		SEC		88FD	85	FD	STA	\$FD		
8880	E5	5A	SBC	\$5A	88FF	91	FB	STA	(\$FB),Y		
88B2	85	58	STA	\$58	8901	85	57	STA	\$57		
88B4	A 8		TAY		8903	A5	FE	LDA	\$FE		
8885	A5	2E	LDA	\$2E	8905	E9	00	SBC	#\$00		
88B7	E5	5B	SBC	\$5B	8907	C8		INY			
8889	AA		TAX		8908	85	FE	STA	\$FE		
88BA	E8		INX		890A	85	58	STA	\$58		
88BB	98		TYA		890C	91	FB	STA	(\$FB),Y		
88BC	F0	1F	BEQ	\$88DD	890E	88		DEY	•		
88BE	A5	5A	LDA	\$5A	890F	B 1	57	LDA	(\$57),Y		
8800	18		CLC		8911	85	89	STA	\$B9		
88C1	65	58	ADC	\$58	8913	C8		INY			
88C3	85	5A	STA	\$5A	8914	B1	57	LDA	(\$57),Y		
88C5	90	03	BCC	\$88CA	8916	85	BA	STA	\$BA		
88C7	E۵	5B	INC	\$5B	8918	F0	18	BEQ	\$8932		
88C9	18		CLC		891A	88		DEY			
88CA	A5	5F	LDA	\$5F	891B	38		SEC			
88CC	65	58	ADC	\$58	891C	A5	B9	LDA	\$B9		
88CE	85	5F	STA	\$5F	891E	E5	BB	SBC	\$BB		
88D0	90	02	BCC	\$88D4	8920	AA		TAX			
88D2	Еð	60	INC	\$60	8921	91	57	STA	(\$57),Y		
88D4	98		TYA		8923	A5	BA	LDA	\$BA		
88D5	49	FF	EOR	#\$FF	8925	E9	80	SBC	#\$00		
88D7	A 8		TAY		8927	C8		INY			
88D8	C8		INY		8928	91	57	STA	(\$ 57),Y		
88D9	C6	5B	DEC	\$5B	892A	85	58	STA	\$58		
88DB	СS	60	DEC	\$60	892C	8A		TXA			
88DD	B1	5A	LDA	(\$5A),Y	892D	85	57	STA	\$57		
88DF	91	5F	STA	(\$5F),Y	892F	4C	0E 89	JMP	\$890E		
88E1	C8		INY		8932	60		RTS			
88E2	DØ	F9	BNE	\$88DD	8933	8A		TXA			
88E4	E6	5B	INC	\$5B	8934	38		SEC			
				1							

The complete utility

8935	F 5	3E		SBC	\$3E	898F	85	FD		STA	\$FD
9937	85	RR		STA	\$BB	8991	85	57		STA	\$57
8030	18	00		CI C		8993	91	FB		STA	(\$FB),Y
0034	Δ5	40		I DA	\$49	8995	A5	FE		LDA	\$FE
0020	45	RR			\$BB	8997	69	00		ADC	#\$00
0000	00	94		RCS	\$8944	8999	C8			INY	
0736	00	55		CMP	#4.55	899A	85	FE		STA	\$FE
0740	00	05			#20010	8990	85	58		STA	\$58
0742	70	17			#417	899F	91	FB		STA	(\$FB).Y
0744	HZ AC	27	~*	IMD	##17 #A427	8949	88	• -		DEY	,.
0740	40	27			420	8941	B1	57		LDA	(\$57).Y
0747	HJ				*20 4 DD	8943	85	B 9		STA	\$B9
8746	00	DD		TAV	*DD	8945	60			INY	
8740	HH	<u>م</u> ر			405	8944	B1	57		I DA	(\$57).Y
874E	HO			ADC	₽∠C ##00	8948	85	RA		STA	\$BA
8730	67 05	00		ANC	##00 #20	8966	FA	18		RFD	\$8904
8952	5	38			₽30	8940	88			DEY	+0/0/
8734	00	07		BINE	\$8730		19			rir	
8956	E4	37		LPX	\$37 400ED	0000	Δ5	80		IDA	4R9
8958	90	03		BUL	\$873U	0000	45				\$BB
895A	40	35	A4	JMP	¥A435	00007	200	00		TAY	<i>400</i>
8950	18			LLL	+ 00	0002	01	57		CTA	(457) Y
873F	AD	20		LUA	\$2D #50	0703	71				4PA
8768	80			51H	>>H ≠DD	0700	HJ ZO	DH			*DF1 #400
070Z	00	50		HUL	₽ 00 ¢E0	0707	07	00		TNIV	##00
8764	80	28		518	≫08 ¢0⊑	0707	01	57		CTA	(457) V
8766	AD			CTA	₽∠ Ε ¢50	070H	71	50		CTA	(#J//j) 450
0700	20	90		ADC	₽ JD ₩#00	0700	0.0	90		TVA	*30
070H	07	50		CTA	#700 450	07DC	01	57		CTA	457
0700	0J ^5	57		1 DA	₽J7 4ED	07DF	0J 40	J/ ۸۵	00	TMD	₽J7 ¢00∧0
070E	HJ 45				₽ГD 4/0	8904	40	HU	07	RTS	#07MU
0770	05	47		CTA	777 455	8905	20	E5	81	JSR	\$81F5
077Z	0J 55			1 DA	450 450	8908	20	FD	AF	JSR	\$AEFD
07/4	HJ	F C			₽F6 ###00	89CB	A5	14		I DA	\$14
07/0	07	20		HDC CTA	##00 #/0	89CD	85	6.2		STA	\$09
07/0	8J 20	00	^ 2	51H	₽00 ¢∧005	89CE	A5	15		LDA	\$15
07/H	20	Br	нз	JOK	PHODE	8901	85	Δſ		STA	\$CA
8970	18	~ ~			****	8903	20	E5	81	JSR	\$81E5
877E	AU	00		LDT	#>00	8904	20	FD	ΔF	JSR	
8780	AD	20		LDA	\$20 \$20	8909	Δ5	14		IDA	\$14
8982	60	RR		ADC	* 88	gong	25	RC		STA	480
8784	85	20		SIA	\$20	0000	20	55	91		\$9155
8786	70	03		BCC	*878B	8050	20 25	14			414
8788	EÓ	ZE		INC	¥∠E	0720	05	74		CTA	*17 \$RD
878A	18			ULU	450	0762	0J A5	15			415
878B	A5	FD		LDA	%FD +00	0724	HJ 05	1 J		CTA	4DE
8980	65	RR		AUC	*88	0720	67	DE		JIH	₩UL
						1					

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89E8	A5	2B		LDA	\$2B	8642	C8		INY	
89EA	85	FB		STA	\$FB	8A43	B1	FB	LDA	(\$FB).Y
89EC	A5	2C		LDA	\$2C	8A45	FØ	ED	BEQ	\$8434
89EE	85	FC		STA	\$FC	8647	6.9	22	CMP	#\$22
89FØ	A5	CA		LDA	\$CA	8649	DØ	F7	BNF	\$8442
89F2	DØ	ØF		BNF	\$8483	804R	FØ	F2	RED	\$842F
89F4	A5	C9		1 DA	\$0.70	864D	r o	8F	OMP	#48F
89F6	DØ	0R		BNF	\$8403	804F	E9	E3	RED	\$9034
89F8	AØ	02		I DY	#\$82	8451	ro on	83	CMP	#483
89FA	B1	FR		I DA	(\$FB).Y	8453	EQ.	DE	REQ	\$8034
89FC	85	6.9		STA	\$09	8455	ro	Δ7	CMP	#4407
89FF	C8			INY		8457	E9	21	RED	#+H7 \$8∆7∆
89FF	B1	FR		I DA	(\$FB).Y	8459	ro	84	OMP	#4.8A
8601	85	CA		STA	\$CA	845R	FA	10	RED	\$8070
8683	A5	6.9		I DA	\$09	8450	rø	89	CMP	#489
8495	85	14		STA	\$14	845F	FØ	20	RED	##07 #808D
8697	45	ĊΔ		IDΔ	\$CA	8041	ro	CB CB	CMP	#4CB
8489	85	15		STA	\$15	8043	DA	AR	RNF	\$8070
SAAR	20	13	<u>م</u>	JSR	\$4613	8045	C8	00	INY	POH/ 0
SARE	RA	95	N 0	BUCK	\$8015	8044	R1	FR	I DA	(4FR) V
8410	Δ2	15		IDX	#\$15	8040	60	20	CMP	#4.70
8412	4	37	۵4	IMP	\$∆437	0000	E0	20		##20 \$9045
8015	Δ5	55	N T	i na	\$5F	ON ON	0	A 4		#40405
<u>8017</u>	85	<u></u>		CTA	¢.4.1	ONZE	50	10		##HH # 0A 0D
QA10	۵5 ۵5	40			471 470	OHOL 0A70	Г0 СО	1D 9D		#0H0U
90117	85	42		CTA	₽00 ¢42	0470	C7	10		##0/ ##0/0N
	0J A0	92			##00	0H7Z	гø со	17		₽0H6 U #dE7
0010	DI				#₽00 (4CD) V	0474	C7	15		##E0 #0/0D
0411	01			CTA	(#FD/;) 4ED	0470	50	15	DEG	#0H0U
8023	C0 C0	FU		TNY	₽FD	0H70	00	ы		₽OHZ F
8624	R1	FR			(4FR) Y	967R	R1	FR		(4FR) Y
8024	85	FF		STA	\$FF	8470	60	29	CMP	#\$29
8428	DØ	83		RNF	\$8420	847F	FØ	FQ	RED	\$8474
8424	40	62	88	IMP	\$8867	8481	69	30	CMP	#\$38
8420	62	02	00	INY	+000L	8483	RØ.	83	BCS	\$8488
842F	C8			TNY		8485	88		DEY	+0//00
8A2F	60			INY		8484	DØ	Δ7	RNF	\$842F
8430	RI	FR		I DA	(\$FB).Y	8488	88		DFY	• • • • • • •
8432	DØ	RA		BNF	\$843F	8489	C9	34	CMP	#\$34
8A34	A5	FD		LDA	\$FD	8A88	RØ	A2	BCS	\$842F
8A36	85	FB		STA	\$FB	8480	68		INY	+01121
8A38	A5	FE		LDA	\$FE	8A8F	RI	FR	I DA	(\$FB).Y
8A3A	85	FC		STA	\$FC	8090	60	20	CMP	#\$29
8A3C	DØ	DF		BNE	\$8A1D	8492	FA	F9	BED	\$8A8D
8A3E	C9	22		CMP	#\$22	8494	84	49	STY	\$49
8A40	D0	0B		BNE	\$8A4D	8004	88	• *	DEY	÷ 17
	-	_				0070	00			

8A97	A2	00		LDX	#\$00	8AF3	DØ	10		BNE	\$8B05
8A99	C8			INY		8AF5	AØ	02		LDY	#\$02
8A9A	B1	FB		LDA	(\$FB),Y	8AF7	B1	FB		LDA	(\$FB),Y
8A9C	C۶	30		CMP	#\$30	8AF9	85	39		STA	\$39
8A9E	90	0A		BCC	\$8000	8AFB	C8			INY	
8008	C9	ЗA		CMP	#\$3A	8AFC	B1	FB		LDA	(\$FB),Y
8662	BØ	06		BCS	\$8000	8AFE	85	3A		STA	\$3A
8004	9D	00	02	STA	\$0200,X	8800	A2	11		LDX	#\$11
8667	E8			INX		8802	4C	37	A4	JMP	\$A437
8008	D0	EF		BNE	\$8A99	8805	C8			INY	
8000	A9	ЗA		LDA	#\$3A	8806	B1	58		LDA	(\$58),Y
8AAC	9D	00	02	STA	\$0200,X	8808	85	B7		STA	\$B7
8AAF	86	BF		STX	\$BF	880A	C8			INY	
8AB1	A9	02		LDA	#\$02	880B	B1	58		LDA	(\$58),Y
8AB3	85	7B		STA	\$7B	880D	C5	C4		CMP	\$C4
8AB5	A9	00		LDA	#\$00	880F	D0	66		BNE	\$8B17
8AB7	85	7A		STA	\$7A	8B11	A5	B7		LDA	\$B7
8AB9	20	F5	81	JSR	\$81F5	8B13	C5	CЗ		CMP	\$C3
8ABC	A5	14		LDA	\$14	8B15	F0	15		BEQ	\$8B2C
8ABE	85	СЗ		STA	\$ C3	8B17	A5	B9		LDA	\$B9
8AC0	A5	15		LDA	\$15	8B19	18			CLC	
8AC2	85	C4		STA	\$C4	8B1A	65	BC		ADC	\$BC
8AC4	C5	CA		CMP	\$CA	8B1C	85	B 9		STA	\$B9
8AC6	F0	8A		BEQ	\$8AD2	8B1E	90	02		BCC	\$8B22
8AC8	BØ	0E		BCS	\$8AD8	8820	E۵	BA		INC	\$BA
8ACA	A5	49		LDA	\$49	8B22	A5	5A		LDA	\$5A
8ACC	65	BF		ADC	\$BF	8B24	85	58		STA	\$58
8ACE	A 8			TAY		8B26	A5	5B		LDA	\$5B
8ACF	4C	55	8B	JMP	\$8B55	8B28	85	59		STA	\$59
8AD2	A5	С3		LDA	\$C3	882A	D0	BC		BNE	\$8AE8
8AD4	C5	C9		CMP	\$C9	8B2C	A6	B9		LDX	\$B9
8AD6	90	F2		BCC	\$8aca	8B2E	A5	BA		LDA	\$BA
8AD8	A5	BD		LDA	\$BD	8B30	20	7F	84	JSR	\$847F
8ada	85	B9		STA	\$B9	8B33	A5	BF		LDA	\$BF
8ADC	A5	41		LDA	\$41	8B35	85	3E		STA	\$3E
8ADE	85	58		STA	\$58	8B37	E4	3E		CPX	\$3E
8AE0	A5	BE		LDA	\$BE	8B39	F0	0B		BEQ	\$8846
8AE2	85	BA		STA	\$BA	8B3B	B0	0 6		BCS	\$8B43
8AE4	A5	42		LDA	\$42	8B3D	20	8B	88	JSR	\$888B
8AE6	85	59		STA	\$59	8840	4C	46	8B	JMP	\$8B46
8AE8	A0	00		LDY	#\$99	8B43	20	33	89	JSR	\$8933
8AEA	B1	58		LDA	(\$58),Y	8B46	A4	49		LDY	\$49
SAEC	85	5A		STA	\$5A	8848	A2	00		LDX	#\$00
SAEE	C8			INY		884A	BD	00	02	LDA	\$0200,X
BAEF	R1	58		LDA	(\$58),Y	884D	FØ	60		RED	\$8855
8AF1	85	5B		STA	\$5B	884F	71	⊦B		STA	(\$FB),Y

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8B51	C8			INY		8BA6	85	FE		STA	\$FE
8852	E8			INX		88A8	D0	05		BNE	\$8BAF
8853	D0	F5		BNE	\$8B4A	8BAA	68			PLA	
8855	B1	FB		LDA	(\$FB),Y	8BAB	68			PLA	
8B57	C9	2C		CMP	#\$2C	8BAC	4C	74	A4	JMP	\$A474
8859	F0	04		BEQ	\$885F	8BAF	63			INY	
8B5B	88			DEY		88BØ	C8			INY	
885C	4C	2F	8A	JMP	\$8A2F	8BB1	C8			INY	
885F	4C	8D	8A	JMP	\$8A8D	8BB2	B 1	FB		LDA	(\$FB),Y
8B62	A0	00		LDY	#\$00	8BB4	D0	9C		BNE	\$8BC2
8B64	B1	41		LDA	(\$4 1),Y	88B6	A5	FD		LDA	\$FD
8B66	85	5A		STA	\$5A	8888	85	FB		STA	\$FB
8B68	C8			INY		8BBA	A5	FE		LDA	\$FE
8B69	B1	41		LDA	(\$41),Y	8BBC	85	FC		STA	\$FC
8B6B	85	5B		STA	\$5B	8BBE	A2	00		LDX	#\$00
886D	D0	05		BNE	\$8B74	8BC0	F0	DB		BEQ	\$8B9D
886F	68			PLA		8BC2	C9	FF		CMP	#\$FF
8B70	68			PLA		8BC4	D0	04		BNE	\$8BCA
8B71	4C	74	A4	JMP	\$A474	8BC3	85	3E		STA	\$3E
8B74	C8			INY		8BC8	F0	21		BEQ	\$8BEB
8B75	A5	BD		LDA	\$BD	8BCA	C9	22		CMP	#\$22
8B77	91	41		STA	(\$41),Y	8BCC	D9	09		BNE	\$8BD7
8B79	C8			INY		8BCE	E8			INX	
887A	A5	BE		LDA	\$BE	8BCF	E0	02		CPX	#\$02
8B7C	91	41		STA	(\$41),Y	8BD1	D0	0 8		BNE	\$8BDB
887E	18			CLC		8BD3	A2	00		LDX	#\$00
887F	A5	BD		LDA	\$BD	8BD5	F0	DA		BEQ	\$8BB1
8881	65	BC		ADC	\$BC	8BD7	E0	01		СРХ	#\$01
8883	85	BD		STA	\$BD	8BD9	D0	D6		BNE	\$8BB1
8885	90	02		BCC	\$8887	8BDB	85	3E		STA	\$3E
8887	Eð	BE		INC	\$BE	8BDD	C9	CØ			#\$0050
8888	A5	56		LDA	\$0A	880F	90	02		BUU	\$88E3
8888	85	41		51A	\$41 450	SRF1	EY	60		SBC	#>00 #+ /0
8880	AD	28			* 4 2 * 4 2	8BE3		00			400CD
8885	83	42		DAIC	342 20072	88E2	80	04		CMD	₽0DED ##21
8871	00			DINE	7000£	BBE1	U7 D0			DCC	#₽∠1 40001
8873	AZ	99			#700 #70	ODE7	05			CTA	4 2D ⊅0001
8873	HJ OF			CTA		ODED	0.0	30		CTV	\$30 \$40
8877	80	20		JIN	₽FD 400	ODED	04	47		CTV	#7/ #20
8877	HJ OF			CTA			00 00	30		1 67	₽3C #d:01
8878	80	rt aa		JIN	₽FC #±00	0052	HZ CO	01			#701
0070	H0 D4	00			#700 (45P) V	ODEA	- LØ - D1	C D			(GEP) V
0075	DI			CTA	\#F₽/j! ¢FN			25		CMP	1760/91 42F
0001	60	rυ			₽ĽV	0000	- LJ	30			#JE doden
0043	UQ D1	ED		1 04	(4EP) V	OPTA	50	03			≠00FV
00H4	DI	гр		L.UT1	141 107 91	OBLH	сð			11/1/	

8BFB	DØ	F۵		BNE	\$8BF3	8C59	F0	03		BEQ	\$8C5E
8BFD	86	3E		STX	\$3E	8C5B	C 8			INY	
8BFF	EØ	02		CPX	#\$02	8050	91	FB		STA	(\$FB),Y
8001	B0	12		BCS	\$8C15	8C5E	A5	3F		LDA	\$3F
8003	CA			DEX		8698	F0	03		BEQ	\$8C65
8004	A5	ЗD		LDA	\$3D	8C62	C 8			INY	
8003	C9	20		CMP	#\$20	8693	91	FB		STA	(\$FB),Y
8038	FØ	8		BEQ	\$8C12	8C65	A9	47		LDA	#\$47
8C0A	A9	00		LDA	#\$00	8C67	C8			INY	
8C0C	85	40		STA	\$40	8C68	91	FB		STA	(\$FB),Y
8C0E	85	ЗF		STA	\$3F	8C6A	A9	ЗE		LDA	#\$3E
8C10	FØ	15		BEQ	\$8C27	8C6C	C 8			INY	
8C12	4C	33	8D	JMP	\$8D33	8C4D	91	FB		STA	(\$FB),Y
8015	A9	00		LDA	#\$00	8C6F	A5	ЗD		LDA	\$3D
8017	20	7F	84	JSR	\$847F	8C71	C8			INY	
8C1A	A2	00		LDX	#\$00	8C72	91	FB		STA	(\$FB),Y
8010	BD	00	02	LDA	\$0200,X	8C74	A9	5D		LDA	#\$5D
8C1F	85	ЗF		STA	\$3F	8C76	C8			INY	
8C21	E8			INX		8C77	91	FB		STA	(\$FB),Y
8C22	BD	00	02	LDA	\$0200,X	8079	A6	30		LDX	\$30
8C25	85	40		STA	\$40	8C7B	4C	B 1	8B	JMP	\$8BB1
8C27	A5	ЗD		LDA	\$3D	8C7E	85	3D		STA	\$3D
8C29	C9	61		CMP	#\$61	8C80	A9	50		LDA	#\$50
8C2B	90	51		BCC	\$8C7E	8C82	85	62		STA	\$62
8C2D	C9	7B		CMP	#\$7B	8C84	A9	A3		LDA	#\$A3
8C2F	B0	4D		BCS	\$8C7E	8088	85	63		STA	\$63
8C31	38			SEC		8038	A2	51		LDX	#\$51
8C32	E9	20		SBC	#\$20	8C8A	A0	00		LDY	#\$00
8C34	85	ЗD		STA	\$3D	8C8C	20	FB	81	JSR	\$81FB
8C36	A2	07		LDX	#\$07	8C8F	B1	62		LDA	(\$62),Y
8C38	A5	40		LDA	\$40	8C91	48			Pha	
8C3A	D0	66		BNE	\$8C42	8092	20	82	82	JSR	\$8202
8C3C	CA			DEX		8095	68			PLA	
8C3D	A5	3F		LDA	\$3F	8096	C5	ЗD		CMP	\$3D
8C3F	D0	61		BNE	\$8C42	8078	F0	20		BEQ	\$8CBA
8C41	CA			DEX		8C9A	C8			INY	
8C42	E4	3E		CPX	\$3E	8C9B	C8			INY	
8C44	F0	6B		BEQ	\$8051	8090	C8			INY	
8046	BØ	96		BCS	\$8C4E	8C9D	CA			DEX	
8C48	20	8B	88	JSR	\$888B	8C9E	10	EC		BPL	\$8C8C
8C4B	4C	51	8C	JMP	\$8C51	8CA0	A5	3D		LDA	\$3D
8C4E	20	33	89	JSR	\$8933	8CA2	C9	1B		CMP	#\$1B
8C51	A4	49		LDY	\$49	8CA4	90	05		BCC	\$8CAB
8053	A9	5B		LDA	#\$5B	8CA6	A2	0D		LDX	#\$0D
8C55	91	FB		STA	(\$FB),Y	8CA8	4C	37	A4	JMP	\$A437
8C57	A5	40		LDA	\$40	8CAB	69	40		ADC	#\$40

8CAD 8D 48 A	4 STA \$A448	8D08 C8	INY	
8CB0 A9 43	LDA #\$43	8D09 91 FB	STA (\$F	·B),Y
8CB2 85 62	STA \$62	8D08 84 49	STY \$49	, '
8CB4 A9 A4	LDA #\$A4	8000 A0 00	LDY #\$0	0
8CB6 85 63	STA \$63	8D0F 20 FB	81 JSR \$81	FB
8CB8 D0 12	BNE \$8CCC	8D12 C8	INY	
8CBA C8	INY	8D13 B1 62	LDA (\$6	2).Y
8CBB 20 FB 8	1 JSR \$81FB	8D15 84 C2	STY \$C2	2
8CBE B1 62	LDA (\$62),Y	8D17 A4 49	LDY \$49	•
8CC0 48	PHA	8D19 C8	INY	
8001 08	INY	8D1A 91 FB	STA (\$F	Έ),Υ
8CC2 B1 62	LDA (\$62),Y	8D1C 84 49	STY \$49	1
8CC4 85 63	STA \$63	8D1E A4 C2	LDY \$C2	
8CC6 20 02 8	2 JSR \$8202	8D20 C4 C1	CPY \$C1	
8009 68	PLA	8D22 D0 EE	BNE \$8D	12
8CCA 85 62	STA \$62	8D24 20 02	82 JSR \$82	02
800 0A 3338	LDY #\$00	8D27 A4 49	LDY \$49	r
8CCE 20 FB 8	1 JSR \$81FB	8D29 A9 5D	LDA #\$5	D
8CD1 B1 62	LDA (\$62),Y	8D2B C8	INY	
8CD3 85 C1	STA \$C1	8D2C 91 FB	STA (\$F	B),Y
8CD5 20 02 8	2 JSR \$8202	8D2E A6 3C	LDX \$3C	
8CD8 A5 C1	LDA \$C1	8D30 4C B1	8B JMP \$8B	B1
8CDA 18	CLC	8D33 A4 49	LDY \$49	1
8CDB 69 04	ADC #\$04	8D35 A6 3C	LDX \$3C	
8CDD AA	TAX	8D37 4C B1	8B JMP \$8B	B1
8CDE A5 40	LDA \$40	803A A5 90	LDA \$9D	
8CE0 D0 06	BNE \$8CE8	8D3C F0 21	BEQ \$8D	5F
8CE2 CA	DEX	8D3E 78	SEI	
8CE3 A5 3F	LDA \$3F	8D3F A9 FF	LDA #\$F	F
8CE5 D0 01	BNE \$8CE8	8D41 8D 47	8E STA \$8E	47
8CE7 CA	DEX	8D44 A9 FF	LDA #\$F	F
8CE8 E4 3E	CPX \$3E	8D46 8D 46	BE STA \$8E	46
8CEA FØ ØB	BEQ \$8CF7	8D49 AE 08	03 LDX \$03	0 8
8CEC 80 06	BCS \$8CF4	8D4C 8E 50	8E STX \$8E	50
8CEE 20 88 88	3 JSR \$888B	8D4F A9 79	LDA #\$7	9
8CF1 4C F7 80	C JMP \$8CF7	8D51 8D 08	03 STA \$03	08
8CF4 20 33 89	7 JSR \$8933	8D54 AE 09	03 LDX \$03	89
8CF7 A4 49	LDY \$49	8D57 8E 51	BE STX \$8E	51
8CF9 A9 58	LDA #\$58	805A A9 80	LDA #\$8	D
8CFB 91 FB	STA (\$FB),Y	8D5C 8D 09	03 STA \$03	99
8CFD A5 3F	LDA \$3F	8D5F 58	CLI	
8CFF F0 03	BEQ \$8D04z	8D60 60	RTS	
8D01 C8	INY	8D61 78	SEI	
8D02 91 FB	STA (\$FB),Y	8D62 A5 9D	LDA \$9D	
8D04 A5 40	LDA \$40	8D64 F0 F9	BEQ \$8D	5F
8D06 F0 03	BEQ \$8D0B	8D66 A9 00	LDA #\$01	8

8068	8D	46	8E	STA	\$8E46	8DD4	20	D2	FF	JSR	\$FFD2
8D6B	AD	50	8E	LDA	\$8E50	8DD7	A9	12		LDA	#\$12
8D6E	8D	88	03	STA	\$0308	8DD9	20	D2	FF	JSR	\$FFD2
8D71	AD	51	8E	LDA	\$8E51	8DDC	A5	3A		LDA	\$3A
8D74	8D	09	03	STA	\$0309	8DDE	8D	4F	8E	STA	\$8E4F
8D77	58			CLI		8DE1	A6	39		LDX	\$39
8D78	60			RTS		8DE3	8E	4E	8E	STX	\$8E4E
8D79	8D	49	8E	STA	\$8E49	8DE6	20	CD	BD	JSR	\$BDCD
8D7C	08			PHP		8DE9	A9	92		LDA	#\$92
8D7D	8E	4 A	8E	STX	\$8E4A	8DEB	20	D2	FF	JSR	\$FFD2
8D80	8C	4B	8E	STY	\$8E4B	8DEE	18			CLC	
8D83	A5	9D		LDA	\$9D	8DEF	B0	96		BCS	\$8D87
8D85	F0	0D		BEQ	\$8D94	8DF1	AE	4D	8E	LDX	\$8E4D
8D87	AD	49	8E	LDA	\$8E49	8DF4	AC	4C	8E	LDY	\$8E4C
808A	AC	4B	8E	LDY	\$8E4B	8DF7	20	F0	FF	JSR	\$FFF0
8D8D	AE	4 A	8E	LDX	\$8E4A	8DFA	20	E4	FF	JSR	\$FFE4
8D90	28			PLP		8DFD	F0	25		BEQ	\$8E24
8D91	6C	50	8E	JMP	(\$8E50)	8DFF	C9	2F		CMP	#\$2F
8D94	A5	39		LDA	\$39	8E01	90	21		BCC	\$8E24
8D96	AD	46	8E	LDA	\$8E46	8E03	C9	ЗA		CMP	#\$3A
8D99	F0	EC		BEQ	\$8D87	8E05	B0	1 D		BCS	\$8E24
8D9B	38			SEC		8E07	E9	30		SBC	#\$30
8D9C	20	F0	FF	JSR	\$FFF0	8E09	D0	07		BNE	\$8E12
8D9F	8E	4D	8E	STX	\$8E4D	8E0B	A9	FF		LDA	#\$FF
8DA2	80	4C	8E	STY	\$8E4C	8E0D	8D	47	8E	STA	\$8E47
8DA5	18			CLC		8E10	D0	12		BNE	\$8E24
8DA6	A2	66		LDX	#\$00	8E12	AA			TAX	
8048	AU	18		LUT	#\$18 #FFF6	8E13	38	~~		SEC	****
8044	20	- 10 0	FF	100		8E14	- AY	99			##00
80AU	AZ	01			##0F ##00	BE10				RUL	
80AF	87	20		LUH	#720 #EED2	0017		50			#0E14
0001	20	υz	ГГ		PFFUZ	00010	00	10	or		40E10
0004		го			¢QDAE	OCIN		40	OE		70E40 #±00
0000	10	гo			₽00AI	OCIC	on	47	or	CTA	40547
0007	10	99			#400	9E22	FA	14	0C	BED	49538
0000		10			###18	8524		47	QE		49547
ODDA	20	E0	FF	JSP	\$FFF0	8527	FQ	AF		BED	48F38
ODDC	70	12			#412	8529	01	20		CMP	#4.20
8001	20	n2	FF	JSR	\$FFD2	BE2B	EA	1.4		BED	##20 \$8F43
80001		ΔF	8F		\$8F4F	8E2D	20	F4	FF	JSR	
8007	AF	4F	8F	I DX	\$8E4E	SE 30	FA	FR	•••	RED	\$8F2D
8004	20	CD	BD	JSR	\$BDCD	8532	on o	20		CMP	#\$28
8000	Δ9	92	50	I DA	#\$92	8F34	FP	an		RED	\$8F42
8005	20	02	FF	JSR	\$FFD2	BE34	50	00 65		RNC	480FD
8002	- <u>6</u> 9	20	••	I DA	#\$20	8538	ΔF	49	8F		\$8F49
0002	/	20		2011		0230		-10		LUX	₽0 240

8E3B	A0	FF		LDY	#\$FF	8E8A	20	A2	BB	JSR	\$BBA2
8E3D	88			DEY		8E8D	20	DD	BD	JSR	\$BDDD
8E3E	DØ	FD)	BNE	\$8E3D	8E90	20	1 E	AB	JSR	\$AB1E
8E40	CA	1		DEX		8E93	A9	FF		LDA	#\$FF
8E41	D0	F8	3	BNE	\$8E3B	8E95	D0	7A		BNE	\$8F11
8E43	38	ł		SEC		8E97	29	7F		AND	#\$7F
8E44	F0	A9	•	BEQ	\$8DEF	8E99	20	D2	FF	JSR	\$FFD2
8E46	FF			???		8E9C	A9	24		LDA	#\$24
8E47	FF			???		8E9E	20	D2	FF	JSR	\$FFD2
8E48	00			BRK		8EA1	20	33	8F	JSR	\$8F33
8E49	FF			???		8EA4	A9	22		LDA	#\$22
8E 4A	09	0B	ł.	ORA	#\$0B	8EA6	20	D2	FF	JSR	\$FFD2
8E4C	00			BRK		8EA9	A0	00		LDY	#\$00
8E4D	00			BRK		8EAB	B1	BB		LDA	(\$BB),Y
8E4E	00			BRK		8EAD	AA			TAX	·
8E4F	00			BRK		8EAE	F0	15		BEQ	\$8EC5
8E50	02			???		8EB0	C8			INY	
8E51	83			???		8EB1	B1	BB		LDA	(\$BB),Y
8E52	A5	9D		LDA	\$9D	8EB3	85	22		STA	\$22
8E54	62	80		CMP	#\$80	8EB5	C8			INY	
8E56	F0	01		BEQ	\$8E59	8EB6	B1	BB		LDA	(\$BB),Y
8E58	60			RTS		8EB8	85	23		STA	\$23
8E59	A5	2D		LDA	\$2D	8EBA	A0	00		LDY	#\$00
8E5B	85	BB		STA	\$BB	8EBC	B1	22		LDA	(\$22),Y
8E5D	A5	2E		LDA	\$2E	8EBE	20	D2	FF	JSR	\$FFD2
8E5F	85	BC		STA	\$BC	8EC1	C8			INY	
8E61	A5	BC		LDA	\$BC	8EC2	CA			DEX	
8E63	C5	30		CMP	\$30	8EC3	D0	F7		BNE	\$8EBC
8E65	DØ	07		BNE	\$8E6E	8EC5	A9	22		LDA	#\$22
8E67	A5	RR		LDA	\$BB	8EC7	20	D2	FF	JSR	\$FFD2
8E69	05	21		UMP	\$2F	8ECA	F0	45		BEQ	\$8F11
SE9R	00	61		BNE	\$8F9F	8ECC	D0	43		BNE	\$8F11
8EOU	00	~ ~		RIS	***	8ECE	BØ	91		BCS	\$8E61
OF 70	AU Di	00		LDT	#\$UU	8ED0	29	7F		AND	#\$7F
0570	CO	DD		CMD	(\$86),T	8ED2	20	D2	FF	JSR	\$FFD2
0574	C7	50		DCC	# ₽ 00	8ED5	C8			INY	
9E74	20	DP	FF	JCD	₽0ED0	8ED9	B1	BB		LDA	(\$BB),Y
8F79	ra.	02		INY	₽11 DZ	8ED8	C9	7		CMP	#\$/
SE74	RI	RR			(\$88) Y	8EDA	BN	12		BCS	\$8EEE
SE7C	C0	75		CMP	#475	8EDC	20	02	FF -	JSR	\$FFD2
SE7E	RØ	17		BUS	##/1 \$8F97	SEDF	20	33	8F	JSR	\$8F33
SER C	20	D2	FF	JCD	\$FED2	8EE2	A7	46		LDA	#\$46
8F83	20	32	8F	JCD	48F33	BEE4	20	02	F F	JSR	\$FFD2
SESX	Δ5	RP	51		400 JJ	8EE7	A9	4E		LDA	#\$4E
8F88	Δ4	RC			\$ ₿€	8EE9	20	D2	FF	JSR	\$FFD2
0200	717	50			* 00	8EEC	D0	23		BNE	\$8F11

8EEE	29	7F		AND	#\$7F	8F51	A2	15		LDX	#\$15
8EF0	20	D2	FF	JSR	\$FFD2	8F53	4C	37	A4	JMP	\$A437
8EF3	A9	25		LDA	#\$25	8F56	A5	5F		LDA	\$5F
8EF5	20	D2	FF	JSR	\$FFD2	8F58	85	FB		STA	\$FB
8EF8	20	33	8F	JSR	\$8F33	8F5A	A5	60		LDA	\$60
8EFB	A0	00		LDY	#\$00	8F5C	85	FC		STA	\$FC
8EFD	B1	BB		LDA	(\$BB),Y	8F5E	20	79	00	JSR	\$0079
8EFF	85	62		STA	\$62	8F61	C9	2C		CMP	#\$2C
8F01	C8			INY		8F63	D0	E1		BNE	\$8F46
8F02	B1	BB		LDA	(\$BB),Y	8F65	20	73	00	JSR	\$0073
8F04	85	63		STA	\$63	8F68	D0	0F		BNE	\$8F79
8F06	A2	90		LDX	#\$90	8F6A	38			SEC	
8F08	20	44	BC	JSR	\$BC44	8F6B	A5	2D		LDA	\$2D
8F0B	20	DD	BD	JSR	\$BDDD	8F6D	E9	02		SBC	#\$02
8F0E	20	1E	AB	JSR	\$AB1E	8F6F	85	5F		STA	\$5F
8F11	A9	0D		LDA	#\$0D	8F71	A5	2E		LDA	\$2E
8F13	20	D2	FF	JSR	\$FFD2	8F73	E9	80		SBC	#\$00
8F16	18			CLC		8F75	85	60		STA	\$60
8F17	A5	BB		LDA	\$BB	8F77	D0	18		BNE	\$8F91
8F19	69	05		ADC	#\$05	8F79	B 0	СВ		BCS	\$8F46
8F1B	85	BB		STA	\$BB	8F7B	20	F5	81	JSR	\$81F5
8F1D	90	02		BCC	\$8F21	8F7E	E6	14		INC	\$14
8F1F	E٥	BC		INC	\$BC	8F80	20	13	A6	JSR	\$A613
8F21	20	E4	FF	JSR	\$FFE4	8F83	A5	FC		LDA	\$FC
8F24	20	E1	FF	JSR	\$FFE1	8F85	C5	60		CMP	\$60
8F27	D0	01		BNE	\$8F2A	8F87	90	0 8		BCC	\$8F91
8F29	60			RTS		8F89	DØ	BB		BNE	\$8F46
8F2A	A5	CB		LDA	\$CB	8F8B	A5	FB		LDA	\$FB
8F2C	C9	40		CMP	#\$40	8F8D	C5	5F		CMP	\$5F
8F2E	D0	F1		BNE	\$8F21	8F8F	B0	85		BCS	\$8F46
8F30	38			SEC		8F91	AØ	00		LDY	#\$00
8F31	B0	9B		BCS	\$8ECE	8F93	A5	5F		LDA	\$5F
8F33	A9	ЗD		LDA	#\$3D	8F95	91	FB		SIA	(\$FB),1
8F35	20	D2	FF	JSR	\$FFD2	8F97	C8			INY	+ (0
8F38	18			CLC		8F98	A5	60		LDA	\$60 (***
8F39	A5	BB		LDA	\$BB	8F9A	91	FB		SIA	(\$+8),1
8F3B	69	02		ADC	#\$02	8F9C	C8			INY	(+ FB))(
8F3D	85	BB		STA	\$BB	8F9D	B1	FB		LDA	(\$FB),Y
8F3F	90	02		BCC	\$8F43	8F9F	AA			TAX	
8F41	E6	BC		INC	\$BC	8FA0	C8			INY	
8F43	60			RTS		8FA1	B1	FB		LDA	(\$FB),Y
8F44	90	03		BCC	\$8F49	8FA3	20	7F	84	JSR	\$8476
8F46	4C	0 8	AF	JMP	\$AF08	8FA6	68			PLA	1
8F49	20	F5	81	JSR	\$81F5	8FA7	68	_		PLA	
8F4C	20	13	A6	JSR	\$A613	8FA8	A2	FF		LDX	#\$FF
8F4F	B0	05		BCS	\$8F56	8FAA	A9	01		LDA	#\$01
						1					

.

8FAC	4C	7D	92	JMP	\$927D	9008	B1	FB		LDA	(\$FB),Y
8FAF	20	F5	81	JSR	\$81F5	900A	85	FD		STA	\$FD
8FB2	4C	A3	A8	JMP	\$A8A3	900C	C8			INY	
8FB5	A9	03		LDA	#\$03	900D	B1	FB		LDA	(\$FB),Y
8FB7	20	FB	A3	JSR	\$A3FB	900F	85	FE		STA	\$FE
8FBA	A5	7B		LDA	\$7B	9011	DØ	05		BNE	\$9018
8FBC	48			PHA		9013	A2	11		LDX	#\$11
8FBD	A5	7A		LDA	\$7A	9015	4C	37	A4	JMP	\$A437
8FBF	48			PHA		9018	A0	04		LDY	#\$84
8FC0	A5	ЗA		LDA	\$3A	901A	B1	FB		LDA	(\$FB),Y
8FC2	48			PHA		901C	C9	E1		CMP	#\$E1
8FC3	A5	3B		LDA	\$3B	901E	FØ	0A		BEQ	\$902A
8FC5	48			PHA		9828	A5	FD		LDA	\$FD
8FC6	A9	8D		LDA	#\$8D	9022	85	FB		STA	\$FB
8FC8	48			PHA		9824	A5	FE		LDA	\$FE
8FC9	20	79	00	JSR	\$0079	9026	85	FC		STA	\$FC
8FCC	20	AF	8F	JSR	\$8FAF	9828	DØ	DC		BNE	\$9886
8FCF	4C	AE	A7	JMP	\$A7AE	902A	A2	FF		LDX	#\$FF
8FD2	A2	00		LDX	#\$00	902C	E8			INX	
8FD4	63	7A		DEC	\$7A	902D	C 8			INY	
8FD6	BØ	02		BCS	\$8FDA	902E	B1	FB		LDA	(\$FB).Y
8FD8	63	7B		DEC	\$7B	9030	FØ	07		BEQ	\$9039
8FDA	20	73	99	JSR	\$0073	9032	DD	00	82	CMP	\$0200.X
8FDD	F0	06		BEQ	\$8FE5	9035	FØ	F5		BEQ	\$902C
8FDF	9D	00	82	STA	\$0200,X	9837	DØ	E7		BNE	\$9828
8FE2	E8			INX		9039	DD	00	02	CMP	\$0200.X
8FE3	D0	F5		BNE	\$8FDA	903C	DØ	E2		BNE	\$9820
8FE5	A9	00		LDA	#\$00	903E	38			SEC	
8FE7	9D	00	02	STA	\$0200,X	903F	A5	FB		LDA	\$FB
8FEA	A9	83		LDA	#\$03	9041	E۶	01		SBC	#\$01
8FEC	20	FB	A3	JSR	\$A3FB	9843	85	7A		STA	\$7A
8FEF	A5	7B		LDA	\$7B	9045	A5	FC		LDA	\$FC
8FF1	48			Pha		9047	E9	00		SBC	#\$80
8FF2	A5	7A		LDA	\$7A	9849	85	7B		STA	\$7B
8FF4	48			Pha		904B	4C	AE	A7	JMP	\$A7AE
8FF5	A5	ЗA		LDA	\$3A	904E	DØ	0C		BNE	\$905C
8FF7	48			Pha		9050	A9	00		LDA	#\$00
8FF8	A5	39		LDA	\$39	9052	85	СВ		STA	\$CB
8FFA	48			Pha		9854	20	E4	FF	JSR	\$FFE4
8FFB	A9	8D		LDA	#\$8D	9057	FØ	FB		BEQ	\$9054
8FFD	48			Pha		9059	85	90		STA	\$98
8FFE	A5	2B		LDA	\$2B	905B	60			RTS	
9000	85	FB		STA	\$FB	905C	20	9E	AD	JSR	\$AD9E
9002	A5	2C		LDA	\$2C	905F	20	A3	B6	JSR	\$B6A3
9004	85	FC		STA	\$FC	9062	C9	00		CMP	#\$80
9006	A0	00		LDY	#\$00	9064	FØ	EA		BEQ	\$9050
						I					

9066	85	FB		STA	\$FB	9001	E٥	59		INC	\$59
9068	A9	00		LDA	#\$00	90C3	A5	59		LDA	\$59
906A	85	Сð		STA	\$C6	9005	85	FC		STA	\$FC
906C	20	E4	FF	JSR	\$FFE4	9007	85	38		STA	\$38
906F	F0	FB		BEQ	\$906C	9009	Aó	2B		LDX	\$2B
9071	A4	FB		LDY	\$FB	90CB	A4	2C		LDY	\$2C
9073	88			DEY		90CD	A9	00		LDA	#\$00
9074	D1	22		CMP	(\$22),Y	90CF	20	D5	FF	JSR	\$FFD5
9076	F0	05		BEQ	\$907D	90D2	90	03		BCC	\$90D7
9078	88			DEY		90D4	4C	F9	E0	JMP	\$E0F9
9079	10	F9		BPL	\$9074	90D7	20	B7	FF	JSR	\$FFB7
907B	30	EF		BMI	\$906C	90DA	29	BF		AND	#\$BF
907D	85	90		STA	\$90	90DC	F0	05		BEQ	\$90E3
907F	60			RTS		90DE	A2	1 D		LDX	#\$1D
9080	20	D4	E1	JSR	\$E1D4	90E0	4C	37	A4	JMP	\$A437
9083	A9	00		LDA	#\$00	90E3	86	2D		STX	\$2D
9085	85	B9		STA	\$B9	90E5	84	2E		STY	\$2E
9087	20	26	B5	JSR	\$B526	90E7	86	5F		STX	\$5F
908A	A5	2D		LDA	\$2D	90E9	84	60		STY	\$60
908C	85	5F		STA	\$5F	90EB	A5	FB		LDA	\$FB
908E	A5	2E		LDA	\$2E	90ED	85	5A		STA	\$5A
9090	85	60		STA	\$60	90EF	A5	FC		LDA	\$FC
9092	38			SEC		90F1	85	5B		STA	\$5B
9093	A5	31		LDA	\$31	90F3	38			SEC	
9095	85	5A		STA	\$5A	90F4	A5	33		LDA	\$33
9097	E5	2F		SBC	\$2F	90F6	E9	01		SBC	#\$01
9099	85	FD		STA	\$FD	90F8	A8			TAY	
909B	A5	32		LDA	\$32	90F9	A5	34		LDA	\$34
909D	85	5B		STA	\$5B	90FB	E9	00		SBC	#\$00
909F	E5	30		SBC	\$30	90FD	AA			TAX	
90A1	85	FE		STA	\$FE	90FE	98			TYA	
90A3	A5	33		LDA	\$33	90FF	38			SEC	
90A5	38	~ ~		SEC		9100	E5	5A		SBC	\$5A
70A6	EY	01		SBC	#\$01	9102	85	58		STA	\$58
90A8	82	28		SIA	\$58	9104	A8			TAY	
90AA	AD EO	34		LDA	\$34	9105	8A			TXA	
OBAC	E7 05	50		SBL	#>00 #F0	9106	E5	5B		SBC	\$5B
SUBC	20	DC	A2	TCD	#J7 #A00E	9108	AA			TAX	
9083	Δ5	27	H 0	IDA	₽H3DF ¢07	9109	E8			INX	
9085	85	41		CTA	#37 #/1	910A	98			TYA	
9087	Δ5	20		IDA	₽41 ¢00	910B	F0	1F		BEQ	\$912C
90 R9	85	42		CTA	#30 \$42	910D	A5	5A		LDA	\$5A
90RR	Δ5	74 50		INA	*72 \$50	910F	18			CLC	
988D	85	FR		CTA	+J0 4CD	9110	65	58		ADC	\$58
90 RF	85	27		CTV	+r⊡ ⊈07	9112	85	5A		STA	\$5A
/ 001	00	57		JIH	401	9114	70	63		BCC	29213

9114	FA	58	2	TNC	458	0171	A5	15		1 54	#15
9118	18		•		400	0172	- HJ - CO	01			₽1J ##04
9119	45	5F			\$5F	0175	D0	07			#704 40170
911B	65	58		ADC	\$58	9177	Δ2	05 0F			₽7176 ₩¢0E
911D	85	5F	•	STA	\$5F	0170		27	~4		##8C #A407
911F	20	R2		BCC	\$9123	0170	-40 -00	00	H4	CMD	₽H437 #±00
9121	FA	20	•	INC	4/120	9170	D0	57			#700 40177
9123	98	00		TYA	400	0100	20	F 7		DUD	₽71//
9124	49	FF		FOR	#\$FF	9191	20	40	01	TCD	₫01 ∡0
9124	48	• •		TAY		9184	85	30	/1	STA	\$29
9127	C8			INY		9184	Δ5	14			4 1 4
9128	0.0	58		DEC	\$58	9188	85	37		STA	\$37
912A	60	60		DEC	\$60	9184	40	63	A 4	.TMP	\$4663
912C	B1	54		LDA	(\$5A).Y	918D	20	69	91	JSR	\$9169
912E	91	5F		STA	(\$5E).Y	9198	ÂÅ	ññ		I DY	#\$99
9130	C8			INY		9192	98	••		TYA	
9131	DØ	F9		BNE	\$912C	9193	91	14		STA	(\$14).Y
9133	E6	5B		INC	\$5B	9195	C8			INY	,.
9135	E٥	60		INC	\$60	9196	91	14		STA	(\$14).Y
9137	CA			DEX		9198	C8			INY	···· , ·
9138	D0	F2		BNE	\$912C	9199	91	14		STA	(\$14),Y
913A	38			SEC		919B	A5	14		LDA	\$14
913B	A5	5F		LDA	\$5F	919D	18			CLC	
91 3D	85	31		STA	\$31	919E	69	01		ADC	#\$01
913F	E5	FD		SBC	\$FD	91A0	85	2B		STA	\$2B
9141	85	2F		STA	\$2F	91A2	AA			TAX	
9143	A5	60		LDA	\$60	91A3	A5	15		LDA	\$15
9145	85	32		STA	\$32	91A5	69	00		ADC	#\$00
9147	E5	FE		SBC	\$FE	91A7	85	2C		STA	\$20
9149	85	30		STA	\$30	91A9	A 8			TAY	
914B	A5	41		LDA	\$41	91AA	8A			TXA	
914D	85	37		STA	\$37	91AB	69	02		ADC	#\$02
914F	A5	42		LDA	\$42	91AD	85	2D		STA	\$2D
9151	85	38		STA	\$38	91AF	98			TYA	
9153	68			PLA		91B0	69	88		ADC	#\$00
9154	68	~~		PLA		91B2	85	2E		SIA	\$2E
9100	20	33	A5	JSR	\$A533	9184	40	63	A6	JMP	\$A663
7128	AY OO	00		LDA	#\$88	91B7	AY	70		LDA	#\$/0
913A	20	70		JSK	\$FF70	A1 BA	80	04	03	514	\$0304 Hear
9130	20	E/		JSK	\$FFE7	ALBC.	AY	AD		LUA	#\$AD
7100	20	//	H0	JSR	\$A677	AIRE	80	60	03	514	\$0300 #++A
7103	20		H0 02	JSK	#HOUL	7101	HY on	18	a 2	LUA CTA	#\$1H
7100 0120	46 60	UF 03	7 L		₽72UF	7163	80	00	03	514	₽0300 ##E4
7107 Q120	10	03 00	AE		₽710E	7160	H7 on	E4	a -2	CTA	#₽C4 40000
710D 012E	70	00	P1F 01	TCD	#MF80 #0155	7168	οU ΛΟ	00 7	03	JIH	₽0300 #4A7
7 1 OC	20	гJ	01	JOK	\$01L2	TICR	H7	H7		LUH	# 71-1 /

91 CD	8D	87	83	STA	\$0307	9239	8D	11	86	STA	\$0611
9100	80	A 9	03	STA	\$8389	9230	80	DA	95	STA	\$0506
9103	49	84		I DA	#\$86	923E	8D	24	84	STA	\$9626
9105	8D	8A	83	STA	\$030A	9242	A9	R 5		LDA	#\$05
9108	A9	AF	00	LDA	#\$AE	9244	8D	C1	D9	STA	\$D9C1
91DA	8D	0B	03	STA	\$030B	9247	8D	11	DA	STA	\$DA11
9100	49	FF		I DA	#\$FF	9244	8D	DA	D9	STA	\$D9D6
91DF	8D	17	83	STA	\$0317	924D	80	26	DA	STA	\$DA26
91E2	8D	19	83	STA	\$0319	9250	A9	0D		LDA	#\$9D
91E5	A9	66		LDA	#\$66	9252	20	D2	FF	JSR	\$FFD2
91E7	8D	16	03	STA	\$0316	9255	4C	74	A4	JMP	\$A474
91EA	A9	47		LDA	#\$47	9258	18		•••	CLC	
91EC	8D	18	03	STA	\$0318	9259	20	F0	FF	JSR	\$FFF0
91EF	78			SEI		925C	A9	2A		LDA	#\$2A
91F0	A9	48		LDA	#\$48	925E	A2	16		LDX	#\$16
91F2	8D	8F	8 2	STA	\$028F	9260	20	D2	FF	JSR	\$FFD2
91F5	A9	EB		LDA	#\$EB	9263	CA			DEX	
91F7	8D	90	02	STA	\$0290	9264	DØ	FA		BNE	\$9268
91FA	58			CLI		9266	60			RTS	
91 F B	68			PLA		9267	2A			ROL	
91 F C	68			PLA		9268	20	59	54	JSR	\$5459
91FD	4C	74	A4	JMP	\$A474	926B	49	40		EOR	#\$4C
9200	20	63	A6	JSR	\$A663	926D	49	54		EOR	#\$54
9203	A9	93		LDA	#\$93	926F	55	20		EOR	\$20,X
9205	20	D2	FF	JSR	\$FFD2	9271	43			???	
9208	A9	00		LDA	#\$00	9272	49	53		EOR	#\$53
920A	8D	20	DU	SIA	\$D020	9274	41	42		EOR	(\$42,X)
920D	80	21	DN	514	\$D021	9276	20	20	4E	JSR	\$4E20
9210	AY	85			#\$00 #000/	9279	43			???	
9212	80	86	92	516	₩ #0A	927A	50	20		BAC	\$9290
9215	AZ	8A			#>0H #++00	9270	2A	~ ~	~~	RUL	+
9217	60 00	87	00		#707 40050	927D	AD	02	83	LDA	\$0302
9219	20	28	72		₽72J0 ##00	9280	80	97	92	SIA	\$9297
9210	HZ A	00			##0C ##00	9283	AD	03	03	LUA	\$0303 #0000
921E	10	07			₩₽U7	9286	80	90	ΥZ	514	\$7276 #+0/
9220	10	E 0	FF	TCD	45550	9289	A7	96		LUA	#7470 #0000
7221	20	15	FF		₽FFF0 #dt15	9288	80	9Z	03	518	¥030∠ #+00
7224	PL	47	02		##13 #9247 Y	928E	AY	72		LUH	#\$72 #0000
7220	20	07	72	JCP	\$FED2	9298	80	03	03	518	¥0303
7227	20 CA	UL	r r		#11 DZ	9293	40	80	A4		⊅H480 #≠00
7220	10	57			\$977A	9290	AY OD	83		LUH	##03 #0000
7220 9775	70	00			#*##	7278	80	02	62	518	₽030£ #4∧4
0721	<u>20</u>	00			#\$09	7278	HY	- A4	0 -2		##H4 #0000
0222	70	50	92	JCP	\$9258	7270	80	83	03 AE	51A	70303 48500
0721	20	- 00 - C1	05	STA	\$9501	7240	20	చచ	HJ	0.2K	⊅HJ 33
1200	00	01	00	0111	+0001	72H3	18			LLL	

92A4	A5	22		LDA	\$22	928E	C5	2D		CMP	\$2D
92A6	69	02		ADC	#\$02	9200	DØ	03		BNE	\$9205
92A8	85	2D		STA	\$2D	92C2	A9	80		LDA	#\$80
92AA	A5	23		LDA	\$23	92C4	2C	A9	00	BIT	\$00A9
92AC	69	00		ADC	#\$00	9207	85	0C		STA	\$0C
92AE	85	2E		STA	\$2E	9209	20	79	00	JSR	\$8879
92BØ	20	60	A6	JSR	\$A660	92CC	4C	80	90	JMP	\$9080
92B3	4C	74	A4	JMP	\$ A474	92CF	A5	9C		LDA	\$0C
92B6	A5	32		LDA	\$32	92D1	10	94		BPL	\$92D7
92B8	C5	2E		CMP	\$2E	92D3	63	30		DEC	\$30
92BA	D0	09		BNE	\$9205	92D5	63	32		DEC	\$32
92BC	A5	31		LDA	\$31	92D7	4C	AE	A7	JMP	\$ A7AE

Loading the utility

Once the UTILITY and the data for CODER have been set up, a loader program something like the following should be used.

10 A=A+1:IF A=1 THEN LOAD"UTILITY DATA",8,1

20 IF A=2 THEN LOAD"UTILITY",8,1

30 SYS32768

10 Bits 'n' pieces

General

This chapter is a collection of snippets of information we have found out since acquiring our 64s about 18 months ago. No detailed code here, just the bare facts and a few ideas.

AUTO-REPEATS and INTERRUPTS

We have seen two articles on the subject of providing a repeat on all keys. Both articles were based upon the same idea used on the pre-8000 series PETS. In essence the normal IRQ service routine is patched to include additional code by changing the vector at CINV from its default of \$EA31. The additional routine simple scans SFDX – if a key is being pressed then it is reset to no key (\$40) – and ends with a JUMP to \$EA31 to process the normal interrupt. This will then detect a key as being pressed and enter the appropriate character in the keyboard buffer. Alternatively, for a repeat on all keys, simply POKE 650,128. To disable the repeat, POKE 650,1. (For a full description of the IRQ service routine, see Chapter 4.)

The second method is obviously far easier, but the first does allow a selective auto-repeat to be implemented.

Whilst on the subject of the hardware interrupt (see Chapter 4 for its implementation in the KEY commands), here is a short example to demonstrate what can be done. The following program patches IRQ to scan for function keys 1 and 3. These keys are used to increment the border and background colours respectively. The routine only takes the appropriate action once every 60 interrupts (about a second). If you remove the interrupt counter from \$C01F to \$C02D, the effects produced are quite unusual, but the routine becomes of little practical use as it is too fast to exercise selective control.

To enable: To disable:	SYS 49152 SYS 49170					
C000 78 C001 A91 C003 8D1 C003 A9C	SEI F LDA 403 STA 20 LDA	#\$1F \$0314 #\$C0	ENABLE ENTRY SET CINV TO	POINT	то	\$C01F

C008	8D1503	STA	\$0315	
C00B	A900	LDA	#\$00	SET IRQ COUNTER TO ZERO
C00D	8D00C1	STA	\$C100	
C010	58	CLI		
C011	60	RTS		
C012	78	SEI		DISABLE ENTRY
C013	A931	LDA	#\$31	RESTORE CINV TO \$EA31
C015	8D1403	STA	\$0314	
C018	A9EA	LDA	#\$EA	
C01A	8D1503	STA	\$0315	
C01D	58	CLI		
C01E	30	RTS		
C01F	EE00C1	INC	\$C100	NEW IRQ ENRTY
CØ22	AD00C1	LDA	\$C100	
C025	C93C	CMP	#\$ 3C	60 INTERRUPTS ???
C027	D02A	BNE	\$C053	NO - SKIP KEY SCAN
C029	A900	LDA	#\$00	YES - SO RESET COUNTER
C02B	3D00C1	STA	\$C100	
C02E	A5CB	LDA	\$CB	SFDX - CURRENT KEY PRESS
C030	18	CLC		
C031	C904	CMP	#\$04	F1 ???
C033	D00F	BNE	\$C044	
C035	18	CLC		
C036	AD20D0	LDA	\$D020	BDR COLOUR
C039	290F	AND	#\$0F	BITS 0-3 ONLY (0-15 DEC)
C03B	6901	ADC	#\$01	INCREMENT IT
C03D	8D20D0	STA	\$D020	
C040	A900	LDA	#\$00	ENSURE SKIP TAKEN
C042	F00F	BEQ	\$C053	SKIP BKD COLOUR
C044	C905	CMP	#\$05	BKD COLOUR
C046	D00B	BNE	\$C053	
CØ48	18	CLC		
C049	AD21D0	LDA	\$D021	
C04C	290F	AND	#\$0F	
C04E	6901	ADC	#\$01	
C050	8D21D0	STA	\$D021	
C053	4031EA	JMP	\$EA31	CONTINUE NORMAL IRQ
The BAS	sic loader:			

1 DATA 120, 169, 31, 141, 20, 3, 169, 19 2, 141, 21, 3 2 DATA 169, 0, 141, 0, 193, 88, 96, 120, 169, 49 3 DATA 141, 20, 3, 169, 234, 141, 21, 3, 88, 96

```
4 DATA 238, 0, 193, 173, 0, 193, 201, 60
, 208, 42
5 DATA 169, 0, 141, 0, 193, 165, 203, 24
, 201, 4
6 DATA 208, 15, 24, 173, 32, 208, 41, 15
, 105, 1
7 DATA 141, 32, 208, 169, 0, 240, 15, 20
1, 5, 208
8 DATA 11, 24, 173, 33, 208, 41, 15, 105
, 1, 141
9 DATA 33, 208, 76, 49, 234, 237, 61, 3,
170, 173
10 FOR I=49152 TO 49238:READ A:POKE I,A:
NEXT I
```

The IRQ vector can be used to great advantage. One common use is to provide interrupt driven music (see *The Companion to the Commodore 64* pub. by Pan/PCN) and as in the UTILITY to make the function keys programmable (Chapter 4).

Simple program protection

Some BASICs include commands to 'unlist' or generate protected files; the 64's, however, does not. In order to protect our software we have to resort to programming tricks.

There are many ways to afford a program some degree of protection from unauthorized change. Most of these are well known and do little to prevent the experienced user from gaining access. Chapter 1 showed how the link addresses could be modified to make program lines invisible and list out of sequence. Another way to hide areas of code is to end lines with a REM"[DEL][DEL]....." sequence. On listing, the 'deletes' will erase characters to their left. Most other techniques require a program to be RUN.

Once a program has been run we can destroy some of the vectors from \$0300 to \$0333. These include the PRINT TOKENS LINK, IQPLOP, which could be directed to, say, print 'SYNTAX ERROR' at \$AF08, the SAVE vector at ISAVE to prevent saving and also disable the RUN/STOP at ISTOP. We could also put a specified value somewhere in memory which, if not found, will cause the program to crash, erase itself or even perform a cold boot of the system. Unfortunately (or perhaps fortunately) any BASIC program can be loaded without being run. To produce programs which auto-run on loading requires knowledge of both machine code and the operating system of the 64 (see CHAIN in Chapter 8). Nearly all commercial software uses a number of levels of protection, one of which is usually auto-running. We have covered a number of ways to accomplish this for your own software in Chapter 2, but purposely leave out many of the techniques used by commercial software houses. (Remember it is illegal to reproduce commercial software.)

Commodore Computing International Volume 2 No.II has an article on program protection. It contains the usual:

DISABLE RUN/STOP	POKE 808,251
DISABLE LIST	POKE 774,131:POKE 775,164
DISABLE SAVE	POKE 818,131:POKE 819,164

The first simply bypasses the test by jumping to the end of the routine (RTS). The latter two jump to 'ready for BASIC'. Similar changes may be made to RUN/RESTORE at NMINV (see Chapter 6). The article does give a program to generate auto-run programs from your own code. If you are interested then, as the program is copyrighted, we suggest you get hold of a copy of the magazine. A further tape protection idea was given in *Commodore User* Volume 1 No.10.

Specified input

One of the most difficult and time-consuming tasks in producing software for use by others is in making it 'crash-proof'. BASIC does not allow the programmer to specify which keys are valid during input. The results of incorrect entries in type, size or number can spoil a wellthought-out, pleasing display or even crash the program. The way round this problem is to write your own input routine.

Commodore Horizons magazine of February 1984 published a very good machine code 'Keyscan' input program written by Adrian Warman which does just about everything you could ask for. We see little point in re-inventing the wheel, so we suggest that you read that article. However, we have approached the problem from a different angle and produced a simple routine entirely in BASIC. This is intended to be called when input is required. The type of input expected is set using the variable 'F' which is set to Ø for a real number, 1 for an integer and 2 for a string. Strings may contain commas and quotes. Editing an input may only be carried by using the DELETE key. The returned value may, if required, be converted to a number by a simple VALO. If the routine is to be used more than once, A\$ must be emptied by: A\$="""

- 60000- Generate a flashing cursor.
- 60030 If it is a delete check chars are there to be removed.
- 60040 'Return' marks end of input and the resulting string A\$ is passed back to the main program.
- 60050 Real numbers.
- 60060 Integer only.
- 60070 String.
- 60090 Update the display and wait for next char.

- 60110– Real numbers may begin with + or and may contain only numerals and a decimal point.
- 60170– As for real but may not have a decimal point and must lie in the range given.
- 60230– Allows any of the standard alphanumerics. To provide for lower case mode where uppercase characters have their high bit set some graphics are permitted (128+32 to 128+64).

```
60000 POKE 204,0:POKE 207,0
60010 GET Y$:IF Y$="" GOTO 60000
60020 A=LEN(A$)
60030 IF Y$=CHR$(20) AND A>0 THEN A$=LEF
T$(A$,A-1):GOTO 60090
60040 IF Y$=CHR$(13) GOTO 60100
60050 IF F=0 THEN GOSUB 60110:GOTO 60090
60060 IF F=1 THEN GOSUB 60170:GOTO 60090
60070 IF F=2 THEN GOSUB 60230:GOTO 60090
60080 GOTO 60000
60090 PRINT Y$;:GOTO 60000
60100 RETURN
60110 REM REAL
60120 IF Y$="+" OR Y$="-" AND A=0 THEN
A$=Y$:GOTO 60160
60130 IF Y$>"/" AND Y$<":" THEN A$=A$+Y$
:GOTO 60160
60140 IF Y$="." THEN A$=A$+Y$:GOTO 60160
60150 Y$=""
60160 RETURN
60170 REM INTEGER
60180 IF Y$="+" OR Y$="-" AND A=0 THEN
A$=Y$:GOTO 60220
60190 IF Y$>"/" AND Y$<":" THEN A$=A$+Y$
:GOTO 60220
60200 IF VAL(A$)>32767 OR VAL(A$)<-32768
 THEN A$=LEFT$(A$,A):GOTO 60220
60210 Y$=""
60220 RETURN
60230 REM STRING
60240 IF (ASC(Y$) AND 127)(32 OR (ASC(Y$
) AND 127)>95 THEN Y$="":GOTO 60260
60250 A$=A$+Y$
60260 RETURN
```

The above is intended only as a starting point. Obvious improvements would be to allow the use of the cursor keys by manipulating the string with LEFT\$, MID\$ and RIGHT\$. The maximum length of the string field could also be set to prevent overwriting an existing display.

Invisible characters

Most readers will no doubt be aware that character data may be directly POKEd to the screen. They will also have discovered that on occasions no effect is apparent. A screen character (normal mode) takes on the colour set in the corresponding location of the colour map (\$D800 on). If no character has been printed at this location nor a colour set since the last clear screen the adopted colour is that of the background. To see if the character is there, simply move the cursor to the location.

We can use this to good effect by making displays change quite quickly from BASIC with just pokes to the colour map. It is important to remember that on an INPUT even though the character cannot be seen it is still there and active.

PET-64/64-PET

Commodore has maintained almost complete compatibility in the storage of programs on tape and to a lesser degree on disk. A tape prepared on any machine may be read by another. The 1541 disk drive uses an identical format to the 4040 unit. It can also read 2040 and 3040 formatted disks, but will corrupt these disks if it writes to them. You may also find that you get write problems on 4040 formatted disks. This compatibility does not mean to say that a program written for one machine will work directly on another.

A BASIC program saved on a PET can be loaded and run directly on the 64, whereas the reverse is not true. A word of warning to cassette users that the secondary address of 3 available on the 64 is not recognised by the PET. This is due to the different start address of BASIC on the two machines (\$0401 on PET). Loading a program with a secondary address of zero will not allow it to load below the current start of BASIC. This means that on the 64 the load will be forced to \$0801. A 64 program will normally have a start greater than \$0401 and will go in above the start of BASIC and is not directly usable.

There are two ways to overcome the problem. The first is raise the start of BASIC on the PET to \$0801 by (BASIC 2 & 4)

POKE 41,08:POKE 2048,0:NEW

before loading. All pokes to the screen will have to be adjusted for the PET screen which starts at \$8000 (32768) and all pokes to the colour map removed. Defining the start of the screen with a variable and using offsets from this simplifies the conversion. The easiest way is to avoid using anything other than PRINT for output.

The second technique involves configuring the 64 to look like a PET – BASIC at \$0401 and the screen in bank 2 at \$8000. The *Programmer's Reference Guide* (Chapter 3, 'Screen Memory') tells us how to relocate the screen by setting bits 7 to 4 of \$D018 (53272), remembering to tell the screen editor where it has gone by setting HIBASE (\$0288/648). The start of BASIC is lowered by setting TXTTAB and the top of memory set using MEMTOP and string storage with FRETOP. The following program if run will carry out the necessary changes. SYNTAX ERRORS may result until the screen is cleared due to invisible characters.

```
10 POKE 51,000:POKE 55,000
20 POKE 52,128:POKE 56,128
30 POKE 56578,PEEK(56578) OR 3:POKE56576
,(PEEK(56576)AND252) OR 1
40 POKE 53272,(PEEK(53272)AND15)OR0:POKE
 648,128:POKE 1024,0:POKE 44,4:NEW
```

There are some very sophisticated PET emulators on the market and even cross-assemblers for machine code and cross-compilers for BASIC. If a lot of your work is in an area where portability is important, it might well be worthwhile pursuing the matter.

Load and run

Pressing SHIFT and RUN/STOP will load and run the first program on tape providing it is in BASIC. This may also be performed by:

POKE 631,131:POKE 198,1

The advantage of the second is that it can be used from within a program to avoid the problems associated with chaining if variables are not to be retained. Less well known is its use with disk. The format is:

```
LOAD"PROG",8:[Press SHIFT & RUN/STOP]
```

Disk bugs

When using sequential data files problems may be encountered if the same logical file number is used for both read and write operations. Typically, error 63 FILE EXISTS is reported. The only way to be sure is to use different numbers for input and output.

A less annoying feature is that null strings written to a data file are ignored on reading back. One way to overcome the problem is to always default null strings to a set value which is recognized on reading back. Alternatively, GET# may be used to pick up returns and commas (a bit laborious).

Appendices

APPENDIX A: Storage of BASIC text

Standard CBM 64 tokens

hex	dec		hex	dec		hex	dec		hex	dec	
\$20	32	sp	\$40	64	@	\$80	128	END	\$A6	166	SPC(
\$21	33	!	\$41	65	А	\$81	129	FOR	\$A7	167	THEN
\$22	34	"	\$42	66	В	\$82	130	NEXT	\$A8	168	NOT
\$23	35	#	\$43	67	С	\$83	131	DATA	\$A9	169	STEP
\$24	36	\$	\$44	68	D	\$84	132	INPUT#	\$AA	170	+ add
\$25	37	%	\$45	69	E	\$85	133	INPUT	\$AB	171	– minus
\$26	38	&	\$46	70	F	\$86	134	DIM	\$AC	172	* multi
\$27	39	'	\$47	71	G	\$87	135	READ	\$AD	173	/ div
\$28	40	(\$48	72	Н	\$88	136	LET	\$AE	174	↑ power
\$29	41)	\$49	73	I	\$89	137	GOTO	\$AF	175	AND
\$2A	42	*	\$4A	74	J	\$8A	138	RUN	\$BØ	176	OR
\$2B	43	+	\$4B	75	Κ	\$8B	139	IF	\$B1	177	>gt
\$2C	44	,	\$4C	76	L	\$8C	140	RESTORE	\$B2	178	= eq
\$2D	45	-	\$4D	77	Μ	\$8D	141	GOSUB	\$B3	179	<lt< td=""></lt<>
\$2E	46	•	\$4E	78	Ν	\$8E	142	RETURN	\$B4	180	SGN
\$2F	47	/	\$4F	79	Ο	\$8F	143	REM	\$B5	181	INT
\$30	48	0	\$50	80	Р	\$90	144	STOP	\$B6	182	ABS
\$31	49	1	\$51	81	Q	\$91	145	ON	\$B7	183	USR
\$32	50	2	\$52	82	R	\$92	146	WAIT	\$B8	184	FRE
\$33	51	3	\$53	83	S	\$93	147	LOAD	\$B9	185	POS
\$34	52	4	\$54	84	Т	\$94	148	SAVE	\$BA	186	SQR
\$35	53	5	\$55	85	U	\$95	149	VERIFY	\$BB	187	RND
\$36	54	6	\$56	86	V	\$96	150	DEF	\$BC	188	log
\$37	55	7	\$57	87	W	\$97	151	POKE	\$BD	189	EXP
\$38	56	8	\$58	88	Х	\$98	152	PRINT#	\$BE	190	COS
\$39	57	9	\$59	89	Y	\$99	153	PRINT	\$BF	191	SIN
\$3A	58	:	\$5A	90	Z	\$9A	154	CONT	\$C0	192	TAN
\$3B	59	;	\$5B	91	[\$9B	155	LIST	\$C1	193	ATN
\$3C	60	<	\$5C	92	£	\$9C	156	CLR	\$C2	194	PEEK
\$3D	61	=	\$5D	93	1	\$9D	157	CMD	\$C3	195	LEN
\$3E	62	>	\$5E	94	ſ	\$9E	158	SYS	\$C4	196	STR\$
\$3F	63	?	\$5F	95	←	\$9F	159	OPEN	\$C5	197	VAL
						\$A0	160	CLOSE	\$C6	198	ASC
						\$A1	161	GET	\$C7	199	CHR\$
						\$A2	162	NEW	\$C8	200	LEFT\$
						\$A3	163	TAB(\$C9	201	RIGHT\$
						\$A4	164	ТО	\$CA	202	MID\$
						\$A5	165	FN	\$CB	203	GO

Extended basic tokens

hex	dec		hex	dec	
\$CC	204	OFF	\$E6	230	RESET
\$CD	205	KEY	\$E7	231	CHAIN
\$CE	206	DOKE	\$E8	232	LOMEM
\$CF	207	TEN	\$E9	233	HIMEM
\$D0	208	TWO	\$EA	234	INKEY\$
\$D1	209	HEX	\$EB	235	MEM
\$D2	210	BIN	\$EC	236	APPEND
\$D3	211	OLD	\$ED	237	TROFF
\$D4	212	COLOUR	\$EE	238	unused
\$D5	213	WRITE	\$EF	239	unused
\$D6	214	CGOTO	\$FØ	240	unused
\$D7	215	CGOSUB	\$F1	241	unused
\$D8	216	PLOT	\$F2	242	unused
\$D9	217	ENTER	\$F3	243	unused
\$DA	218	DUMP	\$F4	244	unused
\$DB	219	RENUM	\$F5	245	unused
\$DC	220	DELETE	\$F6	246	unused
\$DD	221	MERGE	\$F7	247	DEEK
\$DE	222	CODER	\$F8	248	unused
\$DF	223	AUTO	\$F9	249	unused
\$E0	224	PROC	\$FA	250	unused
\$E1	225	DPROC	\$FB	251	unused
\$E2	256	EPROC	\$FC	252	unused
\$E3	227	POP	\$FD	253	unused
\$E4	228	QUIT	\$FE	254	unused
\$E5	229	TRACE			

APPENDIX	B: Hex to	decimal and	decimal to	hey convertor
	D. HEX IU	uecimai anu	uecimai to	nex converter

hex	deci	mal	hex	deci	mal	hex	deci	mal	her	dec	imal
	low	high		low	high	nen	low	high	IIC A	low	high
Ф00 ¢01	1	0	\$29	41	10496	\$52	82	20992	\$7B	123	31488
501 ¢00	1	256	\$2A	42	10752	\$53	83	21248	\$7C	124	31744
\$02 \$02	2	512	\$2B	43	11008	\$54	84	21504	\$7D	125	32000
\$03 ¢04	3	/68	\$2C	44	11264	\$55	85	21760	\$7E	126	32256
504 ¢05	4	1024	\$2D	45	11520	\$56	86	22016	\$7F	127	32512
\$Ø5	5	1280	\$2E	46	11776	\$57	87	22272	\$80	128	32768
\$Ø6	6	1536	\$2F	47	12032	\$58	88	22528	\$81	129	33024
\$0/	/	1792	\$30	48	12288	\$59	89	22784	\$82	130	33280
\$08	8	2048	\$31	49	12544	\$5A	90	23040	\$83	131	33536
\$09	9	2304	\$32	50	12800	\$5B	91	23296	\$84	132	33792
\$0A	10	2560	\$33	51	13056	\$5C	92	23552	\$85	133	34048
\$ØB	11	2816	\$34	52	13312	\$5D	93	23808	\$86	134	34304
\$0C	12	3072	\$35	53	13568	\$5E	94	24064	\$87	135	34560
\$0D	13	3328	\$36	54	13824	\$5F	95	24320	\$88	136	34816
\$ØE	14	3584	\$37	55	14080	\$60	96	24576	\$89	137	35072
\$ØF	15	3840	\$38	56	14336	\$61	97	24832	\$8A	138	35328
\$10	16	4096	\$39	57	14592	\$62	98	25088	\$8B	139	35584
\$11	17	4352	\$3A	58	14848	\$63	99	25344	\$8C	140	35840
\$12	18	4608	\$3B	59	15104	\$64	100	25600	\$8D	141	36096
\$13	19	4864	\$3C	60	15360	\$65	101	25856	\$8E	142	36352
\$14	20	5120	\$3D	61	15616	\$66	102	26112	\$8F	143	36608
\$15	21	5376	\$3E	62	15872	\$67	103	26368	\$90	144	36864
\$16	22	5632	\$3F	63	16128	\$68	104	26624	\$91	145	37120
\$17	23	5888	\$40	64	16384	\$69	105	26880	\$92	146	37376
\$18	24	6144	\$41	65	16640	\$6A	106	27136	\$93	147	37632
\$19	25	6400	\$42	66	16896	\$6B	107	27392	\$94	148	37888
\$1A	26	6656	\$43	67	17152	\$6C	108	27648	\$95	149	38144
\$1B	27	6912	\$44	68	17408	\$6D	109	279 0 4	\$96	150	38400
\$1C	28	7168	\$45	69	17664	\$6E	110	28160	\$97	151	38656
\$1D	29	7424	\$46	70	17920	\$6F	111	28416	\$98	152	38912
\$1E	30	7680	\$47	71	18176	\$70	112	28672	\$99	153	39168
\$1F	31	7936	\$48	72	18432	\$71	113	28928	\$9A	154	39424
\$20	32	8192	\$49	73	18688	\$72	114	29184	\$9B	155	39680
\$21	33	8448	\$4A	74	18944	\$73	115	29440	\$9C	156	39936
\$22	34	8704	\$4B	75	19200	\$74	116	29696	\$9D	157	40192
\$23	35	8960	\$4C	76	19456	\$75	117	29952	\$9E	158	40448
\$24	36	9216	\$4D	77	19712	\$76	118	30208	\$9F	159	40704
\$25	37	9472	\$4E	78	19968	\$77	119	30464	\$AØ	160	40960
\$26	38	9728	\$4F	79	20224	\$78	120	30720	\$A1	161	41216
\$27	39	9984	\$50	80	20480	\$79	121	30976	\$A2	162	41472
\$28	40	10240	\$51	81	20736	\$7A	122	31232	\$A3	163	41728

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hex	decir	lecimal hex		decimal		hex	decii	mal	al hex		decimal	
	low	high		low	high		low	high		low	high	
\$A4	164	41984	\$BB	187	47872	\$D2	210	53760	\$E9	233	59648	
\$A5	165	42240	\$BC	188	48128	\$D3	211	54016	\$EA	234	59904	
\$A6	166	42496	\$BD	189	48384	\$D4	212	54272	\$EB	235	60160	
\$A7	167	42752	\$BE	190	48640	\$D5	213	54528	\$EC	236	60416	
\$A8	168	43008	\$BF	191	48896	\$D6	214	54784	\$ED	237	60672	
\$A9	169	43264	\$CØ	192	49152	\$D7	215	55040	\$EE	238	60928	
\$AA	170	43520	\$C1	193	49408	\$D8	216	55296	\$EF	239	61184	
\$AB	171	43776	\$C2	194	49664	\$D9	217	55552	\$FØ	240	61440	
\$AC	172	44032	\$C3	195	49920	\$DA	218	55808	\$F1	241	61696	
\$AD	173	44288	\$C4	196	50176	\$DB	219	56064	\$F2	242	61952	
\$AE	174	44544	\$C5	197	50432	\$DC	220	56320	\$F3	243	62208	
\$AF	175	44800	\$C6	198	50688	\$DD	221	56576	\$FA	244	62464	
\$BØ	176	45056	\$C7	199	50944	\$DE	222	56832	\$F5	245	6272 0	
\$B1	177	45312	\$C8	200	51200	\$DF	223	57088	\$F6	246	62976	
\$B2	178	45568	\$C9	201	51456	\$EØ	224	57344	\$F7	247	63232	
\$B3	179	45824	\$CA	202	51712	\$E1	225	57600	\$F8	248	63488	
\$B4	180	46080	\$CB	203	51968	\$E2	226	57856	\$F9	249	63744	
\$B5	181	46336	\$CC	204	52224	\$E 3	227	58112	\$FA	250	64000	
\$B6	182	46592	\$CD	205	52480	\$E4	228	58368	\$FB	251	64256	
\$B7	183	46848	\$CE	206	52736	\$E5	229	58624	\$FC	252	64512	
\$B8	184	47104	\$CF	207	52992	\$E6	230	58880	\$FD	253	64768	
\$B9	185	47360	\$DØ	208	53248	\$E7	231	59136	\$FE	254	65024	
\$BA	186	47616	\$D1	209	53504	\$E8	232	59392	\$FF	255	65280	

APPENDIX C: Machine code mnemonics and hex values

6510 OP-CODES

The tables below are a quick reference guide only and for more detailed information a 6502 assembler book should be consulted.

The tables should be read row then column. If in doubt, remember LDA immediate mode is \$A9. The following abbreviations have been used:

# – immediate mode	A=accumulator
\$ – absolute address	X=X index register
Z – zero page	Y=Y index register
I – indirect address	

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0	1	2	4	5	6
Ø BRK	ORA (1,X)			ORA Z	ASL Z
1 BPL	ORA (I),Y			ORA Z,X	ASL Z,X
2 JSR	AND (I,X)		BIT Z	AND Z	ROLZ
3 BMI	AND (I),Y			AND Z,X	ROL Z,X
4 RTI	EOR (I,X)			EOR Z	LSR Z
5 BVC	EOR (I),Y			EOR Z,X	LSR Z,X
6 RTS	ADC (I,X)			ADC Z	ROR Z
7 BVS	ADC (I),Y			ADC Z,X	ROR Z,X
8	STA (1,X)		STY Z	STA Z	STX Z
9 BCC	STA (I),Y		STY Z,X	STA Z,X	STX Z,Y
A LDY#	LDA (1,X)	LDX #	LDY Z	LDA Z	LDX Z
B BCS	LDA (I),Y		LDY Z,X	LDA Z,X	LDX Z,Y
C CPY#	CMP (1,X)		CPY Z	CMP Z	DEC Z
d bne	CMP (I),Y			CMP Z,X	DEC Z,X
E CPX #	SBC (1,X)		CPX Z	SBC Z	INC Z
F BEQ	SBC (I),Y			SBC Z,X	INC Z,X

8	0	Δ	C		г
<u> </u>	J	~	C	U	E
0 PHP	ORA #	ASL A		ORA \$	ASL \$
1 CLC	ORA\$,Y			ORA \$,X	ASL \$,X
2 PLP	AND #	ROLA	BIT \$	AND \$	ROL\$
3 SEC	AND \$,Y			AND \$,X	ROL\$,X
4 PHA	EOR #	LSR A	JMP \$	EOR \$	LSR \$
5 CLI	EOR \$,Y			EOR \$,X	LSR \$,X
6 PLA	ADC #	ROR A	JMP I	ADC \$	ROR \$
7 SEI	ADC \$,Y			ADC \$,X	ROR \$,X
8 DEY		TXA	STY \$	STA \$	STX \$
9 TYA	STA \$,Y	TXS		STA \$,X	
A TAY	LDA #	TAX	LDY \$	LDA \$	LDX \$
B CLV	LDA \$,Y	TSX	LDY \$,X	LDA \$,X	LDX \$,Y
C INY	CMP #	DEX	CPY \$	CMP \$	DEC \$
D CLD	CMP \$,Y			CMP \$,X	DEC \$,X
E INX	SBC #	NOP	CPX \$	SBC \$	INC \$
F SED	SBC \$,Y			SBC \$,X	INC \$,X

APPENDIX D: BASIC loader for Supermon

This is the BASIC program to produce Jim Butterfield's Supermon monitor. Type it in and SAVE it before running. There are lots of numbers so it is easy to make a mistake. We have put in some checksums to help isolate errors. Once loaded correctly you will be able to save the machine code version from the monitor itself. Instructions for using Supermon are given in Appendix E.

Supermon normally loads to the top of BASIC memory. We have modified it to sit at \$C000 on to allow you to enter the code for the UTILITY. Once you have it up and running, you can use Supermon to modify itself to sit anywhere. There are addresses which require changing so we have included a relocater program after the loader.

The Loader

```
10 A=49152:C=0
20 READB: IFB=-1THEN40
30 POKEA.B:A=A+1:C=C+B:GOT020
40 IFC=27914THEN60
50 PRINT DATA ERROR IN 1000 - 1300":END
60 C=0
70 READB: IFB=-1THEN90
80 POKEA.B:A=A+1:C=C+B:GOT070
90 IFC=26078THEN110
100 PRINT DATA ERROR IN 1310 - 1600":END
110 C=0
120 READB: IFB=-1THEN140
130 POKEA, B:A=A+1:C=C+B:GOT0120
140 IFC=26897THEN160
150 PRINT DATA ERROR IN 1610 - 1900":END
160 C=0
170 READB: IFB=-1THEN190
180 POKEA, B:A=A+1:C=C+B:GOT0170
190 IFC=28055THEN210
200 PRINT DATA ERROR IN 1910 - 2200":END
210 C=0
220 READB: IFB=-1THEN240
230 POKEA, B:A=A+1:C=C+B:GOT0220
240 IFC=25343THEN260
250 PRINT DATA ERROR IN 2210 - 2500":END
260 C=0
270 READB: IFB=-1THEN290
280 POKEA.B:A=A+1:C=C+B:GOT0270
290 IFC=25432THEN310
300 PRINT DATA ERROR IN 2510 - 2800":END
310 C=0
320 READB: IFB=-1THEN340
330 POKEA.B:A=A+1:C=C+B:GOT0320
340 IFC=27324THEN360
350 PRINT DATA ERROR IN 2810 - 3100":END
```

```
360 C=0
370 READB: IFB=-1THEN390
380 POKEA.B:A=A+1:C=C+B:GOT0370
390 IFC=25335THEN410
400 PRINT DATA ERROR IN 3110 - 3400":END
410 C=0
420 READB: IFB=-1THEN440
430 POKEA, B:A=A+1:C=C+B:GOT0420
440 IFC=28057THEN460
450 PRINT DATA ERROR IN 3410 - 3700":END
460 C=0
470 READB: IFB=-1THEN490
480 POKEA.B:A=A+1:C=C+B:GOT0470
490 IFC=20514THEN510
500 PRINT DATA ERROR IN 3710 - 4000 :END
510 C=0
520 READB: IFB=-1THEN540
530 POKEA, B:A=A+1:C=C+B:GOT0520
540 IFC=22061THEN560
550 PRINT DATA ERROR IN 4010 - 4290":END
560 PRINT"DATA CORRECT":PRINT
570 PRINT SYS49152 TO USE"
580 END
1000 DATA76,233,192,255,0,0,255
1010 DATA255,0,0,255,255,0,0
1020 DATA255,255,0,0,255,255,0
1030 DATA0,255,255,0,0,255,255
1040 DATA0,0,255,255,0,0,255
1050 DATA255,0,0,255,255,0,0
1060 DATA255,255,0,0,255,255,0
1070 DATA0,255,255,0,0,255,255
1080 DATA0,0,255,255,0,0,255
1090 DATA255,0,0,255,255,0,0
1100 DATA255,255,0,0,255,255,0
1110 DATA0,255,255,0,0,255,255
1120 DATA0,0,255,255,0,0,255
1130 DATA255,0,0,255,255,0,0
1140 DATA255,255,0,0,255,255,0
1150 DATA0,255,255,0,0,255,255
1160 DATA0,0,255,255,0,0,255
1170 DATA255,0,0,255,255,0,0
1180 DATA255,255,128,0,255,255.0
1190 DATA0,255,255,0,0,255,255
1200 DATA0,0,255,255,0,0,255
1210 DATA255,0,0,255,255,0,0
1220 DATA255,255,0,0,255,255,0
```

DATA0,255,255,0,0,255,255 1230 1240 DATA0.0.255.255.0.0.255 1250 DATA255.0.0.255,255,0,0 1260 DATA255,255,0,0,255,255,0 1270 DATA0,255,255,0,0,255,255 1280 DATA0,0,255,255,0,0,255 1290 DATA255,0,0,255,255,0,0 1300 DATA255,255,0,0,255,255,0,-1 1310 DATA0,255,255,0,0,255,255 1320 DATA0,0,255,255,0,0,255 1330 DATA255.0.169.160.133.56,173 1340 DATA230,200,141,22,3,173,231 1350 DATA200,141,23,3,169,128,32 1360 DATA144.255.0.216.104.141.62 1370 DATA2,104,141,61,2,104,141 1380 DATA60,2,104,141,59,2,104 1390 DATA170,104,168,56,138,233,2 1400 DATA141,58,2,152,233,0,141 1410 DATA57,2,186,142,63,2,32 1420 DATA87,198,162,66,169,42,32 1430 DATA87,195,169,82,208,52,230 1440 DATA193,208,6,230,194,208,2 1450 DATA230,38,96,32,207,255,201 1460 DATA13,208,248,104,104,169,154 1470 DATA32,210,255,169,0,133,38 1480 DATA162,13,169,46,32,87,195 1490 DATA169,159,32,210,255,32,62 1500 DATA193,201,46,240,249,201,32 1510 DATA240,245,162,14,221,183,200 1520 DATA208,12,138,10,170,189,199 1530 DATA200,72,189,198,200,72,96 1540 DATA202,16,236,76,237,195,165 1550 DATA193,141,58,2,165,194,141 1560 DATA57,2,96,169,8,133,29 1570 DATA160,0,32,84,198,177,193 1580 DATA32,72,195,32,51,193,198 1590 DATA29,208,241,96,32,136,195 1600 DATA144,11,162,0,129,193,193,-1 1610 DATA193,240,3,76,237,195,32 1620 DATA51,193,198,29,96,169,59 1630 DATA133,193,169,2,133,194,169 1640 DATA5,96,152,72,32,87,198 1650 DATA104,162,46,76,87,195,169 1660 DATA154,32,210,255,162,0,189 1670 DATA234,200,32,210,255,232,224 1680 DATA22,208,245,160,59,32,194

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1690 DATA193,173,57,2,32,72,195
1700 DATA173,58,2,32,72,195,32
1710 DATA183,193,32,141,193,240,92
1720 DATA32,62,193,32,121,195,144
1730 DATA51,32,105,195,32,62,193
1740 DATA32,121,195,144,40,32,105
1750 DATA195,169,154,32,210,255,32
1760 DATA225,255,240,60,166,38,208
1770 DATA56,165,195,197,193,165,196
1780 DATA229,194,144,46,160,58,32
1790 DATA194,193,32,65,195,32,139
1800 DATA193,240,224,76,237,195,32
1810 DATA121,195,144,3,32,128,193
1820 DATA32,183,193,208,7,32,121
1830 DATA195,144,235,169,8,133,29
1840 DATA32,62,193,32,161,193,208
1850 DATA248,76,71,193,32,207,255
1860 DATA201,13,240,12,201,32,208
1870 DATA209,32,121,195,144,3,32
1880 DATA128,193,169,154,32,210,255
1890 DATA174,63,2,154,120,173,57
1900 DATA2,72,173,58,2,72,173,-1
1910 DATA59,2,72,173,60,2,174
1920 DATA61,2,172,62,2,64,169
1930 DATA154,32,210,255,174,63,2
1940 DATA154,108,2,160,160,1,132
1950 DATA186,132,185,136,132,183,132
1960 DATA144,132,147,169,64,133,187
1970 DATA169,2,133,188,32,207,255
1980 DATA201,32,240,249,201,13,240
1990 DATA56,201,34,208,20,32,207
2000 DATA255,201,34,240,16,201,13
2010 DATA240,41,145,187,230,183,200
2020 DATA192,16,208,236,76,237,195
2030 DATA32,207,255,201,13,240,22
2040 DATA201,44,208,220,32,136,195
2050 DATA41,15,240,233,201,3,240
2060 DATA229,133,186,32,207,255,201
2070 DATA13,96,108,48,3,108,50
2080 DATA3,32,150,194,208,212,169
2090 DATA154,32,210,255,169,0,32
2100 DATA239,194,165,144,41,16,208
2110 DATA196.76.71.193.32.150.194
2120 DATA201,44,208,186,32,121,195
2130 DATA32,105,195,32,207,255,201
2140 DATA44,208,173,32,121,195,165
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2150 DATA193,133,174,165,194,133,175 2160 DATA32,105,195,32,207,255,201 2170 DATA13,208,152,169,154,32,210 2180 DATA255, 32, 242, 194, 76, 71, 193 2190 DATA165,194,32,72,195,165,193 2200 DATA72,74,74,74,74,32,96,-1 2210 DATA195,170,104,41,15,32,96 2220 DATA195,72,138,32,210,255,104 2230 DATA76,210,255,9,48,201,58 2240 DATA144,2,105,6,96,162,2 2250 DATA181,192,72,181,194,149,192 2260 DATA104,149,194,202,208,243.96 2270 DATA32,136,195,144,2,133,194 2280 DATA32,136,195,144,2,133,193 2290 DATA96,169,0,133,42,32,62 2300 DATA193,201,32,208,9,32,62 2310 DATA193,201,32,208,14,24,96 2320 DATA32,175,195,10,10,10,10 2330 DATA133,42,32,62,193,32,175 2340 DATA195,5,42,56,96,201,58 2350 DATA144,2,105,8,41,15,96 2360 DATA162,2,44,162,0,180,193 2370 DATA208,8,180,194,208,2,230 2380 DATA38,214,194,214,193,96,32 2390 DATA62,193,201,32,240,249,96 2400 DATA169,0,141,0,1,32,204 2410 DATA195,32,143,195,32,124,195 2420 DATA144,9,96,32,62,193,32 2430 DATA121,195,176,222,174,63,2 2440 DATA154,169,154,32,210,255,169 2450 DATA63,32,210,255,76,71,193 2460 DATA32,84,198,202,208,250,96 2470 DATA230,195,208,2,230,196,96 2480 DATA162,2,181,192,72,181,39 2490 DATA149,192,104,149,39,202,208 2500 DATA243,96,165,195,164,196,56,-1 2510 DATA233,2,176,14,136,144,11 2520 DATA165,40,164,41,76,51,196 2530 DATA165,195,164,196,56,229,193 2540 DATA133,30,152,229,194,168,5 2550 DATA30,96,32,212,195,32,105 2560 DATA195,32,229,195,32,12,196 2570 DATA32,229,195,32,47,196,32 2580 DATA105,195,144,21,166,38,208 2590 DATA100,32,40,196,144,95,161 2600 DATA193,129,195,32,5,196,32

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2610 DATA51,193,208,235,32,40,196
2620 DATA24,165,30,101,195,133,195
2630 DATA152,101,196,133,196,32,12
2640 DATA196,166,38,208,61,161,193
2650 DATA129,195,32,40,196,176,52
2660 DATA32,184,195,32,187,195,76
2670 DATA125,196,32,212,195,32,105
2680 DATA195,32,229,195,32,105,195
2690 DATA32,62,193,32,136,195,144
2700 DATA20,133,29,166,38,208,17
2710 DATA32,47,196,144,12,165,29
2720 DATA129,193,32,51,193,208,238
2730 DATA76,237,195,76,71,193,32
2740 DATA212,195,32,105,195,32,229
2750 DATA195,32,105,195,32,62,193
2760 DATA162,0,32,62,193,201,39
2770 DATA208,20,32,62,193,157,16
2780 DATA2,232,32,207,255,201,13
2790 DATA240,34,224,32,208,241,240
2800 DATA28,142,0,1,32,143,195,-1
2810 DATA144,198,157,16,2,232,32
2820 DATA207,255,201,13,240,9,32
2830 DATA136,195,144,182,224,32,208
2840 DATA236,134,28,169,154,32,210
2850 DATA255,32,87,198,162,0,160
2860 DATA0,177,193,221,16,2,208
2870 DATA12,200,232,228,28,208,243
2880 DATA32,65,195,32,84,198,32
2890 DATA51,193,166,38,208,141,32
2900 DATA47,196,176,221,76,71,193
2910 DATA32,212,195,133,32,165,194
2920 DATA133,33,162,0,134,40,169
2930 DATA147,32,210,255,q69,154,32
2940 DATA210,255,169,22,133,29,32
2950 DATA106,197,32,202,197,133,193
2960 DATA132,194,198,29,208,242,169
2970 DATA145,32,210,255,76,71,193
2980 DATA160,44,32,194,193,32,84
2990 DATA198,32,65,195,32,84,198
3000 DATA162,0,161,193,32,217,197
3010 DATA72,32,31,198,104,32,53
3020 DATA198,162,6,224,3,208,18
3030 DATA164,31,240,14,165,42,201
3040 DATA232,177,193,176,28,32,194
3050 DATA197,136,208,242,6,42,144
3060 DATA14,189,42,200,32,165,198
```

3070 DATA189,48,200,240,3,32,165 3080 DATA198.202.208.213.96.32.205 3090 DATA197,170,232,208,1,200,152 3100 DATA32,194,197,138,134,28,32,-1 3110 DATA72,195,166,28,96,165,31 DATA56,164,194,170,16,1,136 3120 DATA101,193,144,1,200,96,168 3130 DATA74,144,11,74,176,23,201 3140 DATA34,240,19,41,7,9,128 3150 DATA74,170,189,217,199,176,4 3160 3170 DATA74,74,74,74,41,15,208 3180 DATA4,160,128,169,0,170,189 3190 DATA29,200,133,42,41,3,133 3200 DATA31,152,41,143,170,152,160 3210 DATA3,224,138,240,11,74,144 3220 DATA8,74,74,9,32,136,208 3230 DATA250.200.136.208.242.96.177 3240 DATA193,32,194,197,162,1,32 DATA254,195,196,31,200,144,241 3250 3260 DATA162,3,192,4,144,242,96 3270 DATA168,185,55,200,133,40,185 DATA119,200,133,41,169,0,160 3280 3290 DATA5,6,41,38,40,42,136 DATA208,248,105,63,32,210,255 3300 DATA202,208,236,169,32,44,169 3310 3320 DATA13,76,210,255,32,212,195 DATA32,105,195,32,229,195,32 3330 DATA105,195,162,0,134,40,169 3340 3350 DATA154,32,210,255,32,87,198 3360 DATA32,114,197,32,202,197,133 DATA193,132,194,32,225,255,240 3370 3380 DATA5, 32, 47, 196, 176, 233, 76 3390 DATA71,193,32,212,195,169,3 3400 DATA133,29,32,62,193,32,161,-1 DATA193,208,248,165,32,133,193 3410 3420 DATA165,33,133,194,76,70,197 3430 DATA197,40,240,3,32,210,255 3440 DATA96,32,212,195,32,105,195 3450 DATA142,17,2,162,3,32,204 3460 DATA195,72,202,208,249,162,3 3470 DATA104,56,233,63,160,5,74 3480 DATA110,17,2,110,16,2,136 3490 DATA208,246,202,208,237,162,2 3500 DATA32,207,255,201,13,240,30 3510 DATA201,32,240,245,32,208,199 3520 DATA176,15,32,156,195,164,193

```
3530 DATA132,194,133,193,169,48,157
3540 DATA16,2,232,157,16,2,232
3550 DATA208,219,134,40,162,0,134
3560 DATA38,240,4,230,38,240,117
3570 DATA162,0,134,29,165,38,32
3580 DATA217,197,166,42,134,41,170
3590 DATA188,55,200,189,119,200,32
3600 DATA185,199,208,227,162,6,224
3610 DATA3,208,25,164,31,240,21
3620 DATA165,42,201,232,169,48,176
3630 DATA33,32,191,199,208,204,32
3640 DATA193,199,208,199,136,208,235
3650 DATA6,42,144,11,188,48,200
3660 DATA189,42,200,32,185,199,208
3670 DATA181,202,208,209,240,10,32
3680 DATA184,199,208,171,32,184,199
3690 DATA208,166,165,40,197,29,208
3700 DATA160,32,105,195,164,31,240,-1
3710 DATA40,165,41,201,157,208,26
3720 DATA32,28,196,144,10,152,208
3730 DATA4,165,30,16,10,76,237
3740 DATA195,200,208,250,165,30,16
3750 DATA246,164,31,208,3,185,194
3760 DATA0,145,193,136,208,248,165
3770 DATA38,145,193,32,202,197,133
3780 DATA193,132,194,169,154,32,210
3790 DATA255,160,65,32,194,193,32
3800 DATA84,198,32,65,195,32,84
3810 DATA198,169,159,32,210,255,76
3820 DATA176,198,168,32,191,199,208
3830 DATA17,152,240,14,134,28,166
3840 DATA29,221,16,2,8,232,134
3850 DATA29,166,28,40,96,201,48
3860 DATA144,3,201,71,96,56,96
3870 DATA64,2,69,3,208,8,64
3880 DATA9,48,34,69,51,208,8
3890 DATA64,9,64,2,69,51,208
3900 DATA8, 64, 9, 64, 2, 69, 179
3910 DATA208,8,64,9,0,34,68
3920 DATA51,208,140,68,0,17,34
3930 DATA68,51,208,140,68,154,16
3940 DATA34,68,51,208,8,64,9
3950 DATA16,34,68,51,208,8,64
3960 DATA9,98,19,120,169,0,33
3970 DATA129,130,0,0,89,77,145
3980 DATA146,134,74,133,157,44,41
```

```
3990 DATA44,35,40,36,89,0,88
4000 DATA36,36,0,28,138,28,35,-1
4010 DATA93,139,27,161,157,138,29
4020 DATA35,157,139,29,161,0,41
4030 DATA25,174,105,168,25,35,36
4040 DATA83,27,35,36,83,25,161
4050 DATA0,26,91,91,165,105,36
4060 DATA36,174,174,168,173,41,0
4070 DATA124,0,21,156,109,156,165
4080 DATA105,41,83,132,19,52,17
4090 DATA165,105,35,160,216,98,90
4100 DATA72,38,98,148,136,84,68
4110 DATA200,84,104,68,232,148,0
4120 DATA180,8,132,116,180,40,110
4130 DATA116,244,204,74,114,242,164
4140 DATA138,0,170,162,162,116,116
4150 DATA116,114,68,104,178,50,178
4160 DATA0, 34, 0, 26, 26, 38, 38
4170 DATA114,114,136,200,196,202,38
4180 DATA72,68,68,162,200,58,59
4190 DATA82,77,71,88,76,83,84
4200 DATA70,72,68,80,44,65,66
4210 DATA194,53,194,204,193,247,193
4220 DATA86,194,137,194,244,194,12
4230 DATA195,62,196,146,196,192,196
4240 DATA56,197,91,198,138,198,172
4250 DATA198,70,193,255,192,237,192
4260 DATA13,32,32,32,80,67,32
4270 DATA32,83,82,32,65,67,32
4280 DATA88,82,32,89,82,32,83
4290 DATA80,108,239,255,254,221,222,-1
```

Relocation

If you need to relocate Supermon to another area of memory, the following program will assist. First load in SUPERMON and use the TRANS-FER option to copy it to the desired location. Now run the program below and it changes the appropriate locations so that it will run at its new address.

```
10 PRINT"[CLS][REV]SUPERMON RELOCATION":
PRINT:PRINT
20 INPUT"NEW START ADDRESS - DECIMAL";AD
30 IFAD>40960THENPOKEAD+234,160:GOTO90
40 PRINT"DO YOU WANT IT PROTECTED FROM BASIC"
50 PRINT"Y OR N"
```
```
60 GETA$: IFA$<>"Y"ANDA$<>"N"THEN60
70 IF AS="N"THEN POKE AD,160:GOT090
80 POKEAD+234, INT((AD/256+.5))
90 READ OF: IF OF=-1 THEN130
100 READ DA:DA=DA+AD:D2=INT(DA/256):D1=D
A-(D2*256)
110 POKE OF+AD, D1: POKE OF+AD+1, D2
120 GOT090
130 END
140 CO=0
150 READ A: IF A=-1 THEN170
160 CO=CO+A:GOT0150
170 IF CO=451387THENPRINT "[REV]DATA OK"
: END
180 PRINT"[REV]DATA INCORRECT"
190 REM ***** START OF DATA *****
200 DATA1,233,238,2278,244,2279,294,1623
,301,855,341,855,349,318,370,2247
210 DATA374,2246,382,1005,402,1620,407,8
40,410,307,418,904,431,1005,434,307
220 DATA453,1623,459,855,469,2282,482,45
0,488,840,494,840,497,439,500,397
230 DATA505,318,508,889,513,873,516,318,
519,889,524,873,553,450,556,833,559,395
240 DATA564,1005,567,889,572,384,575,439
,580,889,589,318,592,417,597,327,611,889
250 DATA616,384,719,1005,733,904,758,662
,770,751,779,327,782,662,789,889
260 DATA792,873,802,889,813,873,828,754,
831,327,836,840,846,864,853,864
270 DATA890,904,897,904,909,318,925,943,
934,318,937,943,973,318,986,972
280 DATA989,911,992,892,998,318,1001,889
,1020,327,1023,1620,1069,1075
290 DATA1088,980,1091,873,1094,997,1097,
1036,1100,997,1103,1071,1106,873
300 DATA1115,1064,1124,1029,1127,307,113
2,1064,1147,1036,1158,1064
310 DATA1163,952,1166,955,1169,1149,1172
,980,1175,873,1178,997,1181,873
320 DATA1184,318,1187,904,1198,1071,1207
,307,1212,1005,1215,327,1218,980
330 DATA1221,873,1224,997,1227,873,1230,
318,1235,318,1242,318,1265,911
340 DATA1281,904,1297,1623,1317,833,1320
,1620,1323,307,1335,327,1338,980
```

```
350 DATA1365,1386,1368,1482,1384,327,138
9,450,1392,1620,1395,833,1398,1620
360 DATA1405,1497,1409,1567,1413,1589,14
34,1474,1444,2090,1447,1701
370 DATA1450,2096,1455,1701,1462,1485,14
71,1474,1477,840,1515,2009,1533,2077
380 DATA1570,1474,1575,1022,1591,2103,15
96,2167,1629,980,1632,873,1635,997
390 DATA1638,873,1650,1623,1653,1394,165
6,1482,1668,1071,1673,327,1676,980
400 DATA1683,318,1686,417,1699,1350,1710
,980,1713,873,1721,972,1762,2000
410 DATA1767,924,1866,985,1814,2103,1817
,2167,1820,1977,1843,1983,1848,1985
420 DATA1860,2096,1863,2090,1866,1977,18
76,1976,1881,1976,1892,873,1905,1052
430 DATA1917,1005,1943,1482,1957,450,196
0,1620,1963,833,1966,1620
440 DATA1974,1712,1978,1983,1330,1071
450 DATA362,2231,916,318,1806,1497
460 DATA2246,578,2248,565,2250,460
470 DATA2252,503,2254,598,2256,649
480 DATA2258,756,2260,780,2262,1086
490 DATA2264,1170,2266,1216,2268,1336
500 DATA2270,1627,2272,1674,2274,1708
510 DATA2276,326,2278,255,2280,237,-1
```

APPENDIX E: Instructions for the use of Supermon

TO USE SUPERMON (relocated version)

LOAD "SUPERMON", DEVICE, 1 NEW SYS49152

GENERAL NOTE

On entering Supermon it will save the stack which is restored on exit. It further changes the BREAK vector so when a BRK is met in a program Supermon is called.

All values are entered in hex. Only in ASSEMBLER mode do they have to be prefixed with '\$' and then only for the operand.

Once Supermon has been loaded, it is resident until the 64 is either turned off or a program loaded which uses memory from scoon.

INSTRUCTIONS

A – ASSEMBLER – Allows simple assembly of machine code. A \$START OPCODE OPERAND.

For example, A 8000 LDA #\$0A

Supermon will prompt with the next address.

Entering <RETURN> after address will exit assembler mode.

Branches are written with the destination address and not its displacement, that is, BEQ \$C456.

 \boldsymbol{D} – DISASSEMBLER – Disassembles 22 instructions from the address specified.

D \$START

for example, D 8000

Hex values may be changed by overtyping and on *<*RETURN> the same locations will again be disassembled.

Typing D on the bottom line will disassemble the next 22 instructions. Typing SPACE> <return> will exit disassembler mode.

F – FILL MEMORY – Fill an area of memory with a specified value.

F \$FROM \$TO BYTE

for example, F 5000 6000 FF

Useful to set up defaults prior to assembly, in particular to fill with NOPs (\$EA).

G – GO RUN – execute machine code.

G – Starts execution at address currently in the Program Counter Register (PC).

G \$START – Starts execution at specified address.

H – HUNT – search memory for specified bytes.

H \$START \$TO DATA

for example, H 5000 6000 'READ – Hunts for ASCII string "READ"

H 5000 6000 A9 0A - Hunts for LDA #\$0A

A maximum of 32 bytes may be set.

L – LOAD – Loads a program at its secondary address, leaving BASIC pointers unchanged.

L "filename" ,DEVICE – Device in hex. 08 disk 01 tape

M – MEMORY DISPLAY – Displays hex values.

M \$FROM \$TO

for example, M 0801, 0891

11 Bytes may be overtyped to change.

P – PRINT DISASSEMBLY – Output hard copy of disassembled listing.

If in Supermon then exit (see below) and set up printer as for normal listing. Re-enter Supermon with subsequent output being directed to the specified device.

For example OPEN4,4:CMD4:SYS49152 P \$FROM \$TO When complete, exit Supermon and close printer channel.

R – REGISTER DISPLAY – Displays current register values. This displays the PC, IRQ, Status Register (SR), A, X, Y and Stack Pointer (SP).

Values can be overtyped to change. This is of particular use in debugging operations where any of the registers may be altered and program execution continued with a GO command.

S – SAVE – Saves an area of memory to tape or disk.

S "filename", DEVICE, START, SEND

Saves from the start up to, but not including the end address.

For example, S ''NAME'',08,5000,6001 – Saves from \$5000 to \$6000, but not the byte at \$6001.

 \mathbf{T} – TRANSFER – Transfers an area of memory to another leaving the original intact.

T \$FROM \$TO \$START

for example, T 5000 6000 1000

You can also use MEMORY DISPLAY this way.

This option may be used in conjunction with the relocator to generate versions of Supermon for use at other locations.

X - EXIT SUPERMON

Х

The stack saved when Supermon was entered will be restored. A CLR from BASIC should tidy up any stack problems.

COPYING SUPERMON

Use the save command as normal with the following addresses: SUPERMON \$C000 \$C000

APPENDIX F: Extended BASIC memory map

The following gives the main entry points for the UTILITY:

Address	Description
\$8000	Initialize Extended BASIC
\$800F	Set up Keyword Vectors
\$8034	Set Top of Memory
\$8041	Set NMI and BRK Vectors
\$8054	Set Keyboard Table Set-up Vector
\$8061	NMI Routine
\$807E	BRK Routine
400/ L	

Address	Description
\$8090	Keyword Vector Table
\$80F6	Keyword Table – Command Keywords
\$81C7	Keyword Table – Function Keywords
\$81F5	Routine – GET PARAMETER
\$81FB	Switch off BASIC
\$8202	Switch on BASIC
\$8209	CRUNCH Tokens
\$82BC	PRINT Tokens
\$8302	Token DISPATCH – Command Keywords
\$8329	Token DISPATCH – Function Keywords
\$8352	Perform COLOUR
\$8381	Perform PLOT
\$83A7	Perform WRITE
\$83AD	Perform ENTER
\$83B3	Perform DOKE
\$83D7	Perform DEEK
\$8401	Routine – Convert to Positive
\$8415	Perform OLD
\$842E	Perform AUTO
\$847F	Routine – Convert to ASCII
\$84AØ	Perform TEN
\$84EC	Perform TWO
\$8537	Perform HEX
\$85BF	Perform BIN
\$85FC	Perform MEM
\$8611	Perform RESET
\$8631	Perform POP
\$864D	Perform KEY
\$8722	KEY Interrupt Routine
\$8799	Perform OFF
\$8/A/	Perform MERGE
1000¢	Perform APPEND
\$000D ¢0022	Routine to close up memory and rechain
\$0933 ¢00CF	Routine to open up memory and rechain
\$09C5 ¢0D02	Perform KENUM
\$0D93 ¢0D2A	Perform CODEK
	Periorm DLIAD
\$0E32 ¢QE44	
Φ0Γ44 \$8ΕΔΕ	
	Perform CCOSUR
₽0ΓD Ο €8ΕΓ\Ο	Perform PPOC
\$0FD2 \$004E	
\$904E	renorm inkey\$

Address	Description
\$9080	CHAIN routine
\$9169	HIMEM/LOMEM routine
\$9181	Perform HIMEM
\$918D	Perform LOMEM
\$91B7	Perform QUIT
\$9200	Start up message
\$927D	Completion of DELETE
\$92B6	Perform CHAIN
\$92DA-	
\$9FFF	unused (expansion for sound/graphic/disk)

Appendix G: Reading an assembler listing

The machine code routines in this book have been presented in two formats. The first was generated using Supermon which is given in Appendix D and instructions for its use in Appendix E. The second was produced using Supersoft's MIKRO assembler cartridge. This appendix deals with listings generated using MIKRO as we feel they require some explanation.

PSEUDO-OPS

These are instructions to the assembler and are not directly executable op-codes.

★=\$C000

This tells the assembler to start its assembly at address \$cooo.

WOR, BYT, and TXT

These instructions reserve bytes in memory. Both WOR and BYT may be followed by any number of arguments separated by commas up to the limit of two screen lines. WOR reserves two bytes and is used to store absolute addresses in low/high format.

WOR \$C000,\$0100

Puts the four bytes \$00,\$C0,\$00,\$01 in four consecutive addresses. BYT reserves single bytes.

BYT \$A9,\$FF

TXT is followed by a quoted string and places the hex values of the ASCII codes sequentially in memory.

TXT ''ABCD''

Puts \$41,\$42,\$43 and \$44 into memory.

Each of these directives allocates bytes from the address at which it appears.

LABELS

Labels are used to identify an absolute address in memory. They are normally used as the destination for branches and jumps. They may also be used as operands.

LDA #\$00 BYTE=#\$FF BEQ ZERO or LDX BYTE

• • • • • •

ZERO RTS

The absolute value of an address may be divided into low/high format by the use of '#<' and '#>' operators.

★=\$C000 START LDA #<START LDX #>START

This loads A with \$00 and X with \$00. Simple numerical calculations may be performed.

STORE BYT \$00,\$FF LDA STORE LDX STORE+1

This loads A with the value held in STORE which has been set to \$00 and X with the byte from the next location, that is, \$FF.

ADDRESS TABLES

Where labels have been used their values, starting at an arbitrary address, have been given. This is useful to determine the hex values for all branches and jumps.

LINE NUMBERS

The assembler code is entered exactly as one would type in a BASIC program. The same editing rules apply to a MIKRO program as to a BASIC program. Generally, we have retained these line numbers in the listings given for clarity and to aid description.

COMMENTS

An exclamation mark is used in the same way as a REM from BASIC and

prefixs comment statements. It tells MIKRO to ignore anything which follows it.

This is by no means the definitive MIKRO manual. We have limited ourselves to using only a few of the options available to allow easier conversion to other assemblers.

Appendix H: Mnemonics generated by CODER

The code	es generated	d are:			
[BLK]	– BLACK	[GR1]	– GRAY1	[DEL]	– DELETE
[WHT]	– WHITE	[GR2]	– GRAYS2	[INS]	– INSERT
[RED]	– RED	[LT GRN]	– LIGHT GREEN	[REV]	– REVERSE
-					ON
[CYN]	– CYAN	[LT BLU]	– LIGHT BLUE	[OFF]	– REVERSE
-					OFF
[PUR]	– PURPLE	[GR3]	– GRAY3	[SPC]	– SPACE
[GRN]	– GREEN	[CLS]	– CLEAR SCREEN	[G>SPC]	– SHIFTED
					SPACE
[BLU]	– BLUE	[HOM]	– HOME CURSOR	[G>?]	– GRAPHIC
					WITH
					SHIFT
[YEL]	– YELLOW	[CU]	– CURSOR UP	[G]</td <td>– GRAPHIC</td>	– GRAPHIC
					WITH LOGO
[ORG]	– ORANGE	[CD]	- CURSOR DOWN	[CTRL?]	– CONTROL
					WITH
					LETTER
[BRN]	– BROWN	[CR]	– CURSOR RIGHT	[F?]	– FUNCTION
					KEYS
[LT RED]	– LIGHT RED) [CL]	– CURSOR LEFT	[PI]	PI
•					3.1416

MULTIPLE CHARACTERS are coded as [10CD]

SPACES

Single, unshifted spaces are not coded. We thought it unnecessary as it detracted from the legibility of the listing.

SPECIAL CODES

The following is an extract from the *Programmer's Reference Guide*, page 74:

There are some other characters that can be PRINTed for special functions, although they are not easily available from the keyboard. In order to get these into quotes, you must leave empty spaces for them in the line, hit <RETURN> or <SHIFT><RETURN>, and go back to the spaces with the cursor controls. Now you must hit <CTRL> <RVS/ON>, to start typing reversed characters, and type the keys shown below:

Function	Type	Appears As
SHIFT RETURN	SHIFTM	
switch to lower case	N	Ν
switch to upper case	SHIFT N	7
disable case switching keys	Н	A
enable case switching keys	1	

Functions 1 and 3 of the above are achieved as stated. CODER replaces them with:

[CRG>M] – SHIFT RETURN [CRG>N] – SWITCH TO UPPER CASE

The other three can be achieved far more easily. Whilst PRINTING in quotes mode, press <CTRL> and the appropriate letter.

Appendix I: Key codes

The following are the values assigned to keys in locations SFDX and LSTX (\$CB/203 & \$C5/197):

dee	c key	dec key	dec key
Ø	INST/DEL	22 T	44 .
1	RETURN	23 X	45 :
2	CRSR R/L	24 7	46 <i>@</i>
3	F7	25 Y	47 ,
4	F1	26 G	48 £
5	F3	27 8	49 *
6	F5	28 B	50 ;
7	CRSR U/D	29 H	51 CLR/HOME
8	3	30 U	52 None
9	W	31 V	53 =
10	А	32 9	54 ^
11	4	33	55 /
12	Z	34 J	56 1
13	S	35 Ø	57 ←
14	E	36 M	58 None
15	None	37 K	59 2
16	5	38 O	60 SPACE
17 [.]	R	39 N	61 None
18	D	40 +	62 Q
19	6	41 P	63 RUN/STOP
20	С	42 L	64 No key press
21	F	43 –	

The following are the values of the shift registers SHFLAG and LSTSHF (\$028D/653 and \$028E/654):

de	ec key pattern
0	NO SHIFTS
1	SHIFT
2	LOGO
3	SHIFT AND LOGO
4	CTRL
5	CTRL AND SHIFT
6	CTRL AND LOGO
7	CTRL, SHIFT AND LOGO

Appendix J: Summary of UTILITY commands

APPEND "program name", device

As for merge except that the appended program is tagged on the end of the memory program. Line numbers are not altered. Peculiar listings can be the result. Use RENUM after an append.

AUTO first line number, increment

Automatic line numbering when entering code.

BIN 8 bit binary number[,....

Prints out decimal conversion of binary number in two forms. The first as a low byte conversion and then, separated by an oblique, the high byte conversion (low \star 256). The binary number must be of eight bits.

CGOTO variable, calculation or line number Line numbers can be mathematical equations.

CGOSUB variable, calculation or line number Line numbers can be mathematical equations.

CHAIN ["filename"][,device]

Will load and run a BASIC program. It also transfers most variables from one program to another.

CODER

Will replace non-standard ASCII and graphic codes with mnemonics. See Appendix H for full list.

COLOUR screen[,border][,text]

Values over 15 can be input, but only the lower four bits will be considered. Border and text parameters are optional.

DEEK(address)

Two byte PEEK. Returns memory location held in address and address+1.

DELETE first line to be deleted, [last line to be deleted]

Deletes lines in the range specified. No last line parameter, it will delete to the end of program.

DPROC name Start of procedure called 'name'.

DOKE address, value

Two byte POKE. Stores value (0–65535) in address and address+1.

DUMP

Displays the values of all simple variables currently in use.

ENTER (x,y) Same as INPUT, but first sets cursor position as in PLOT.

EPROC

End of a procedure.

HEX hexadecimal number[,hex number][,...

Prints decimal conversion of hex input. The hex input can be of either two or four characters, but does not require a prefix of '\$'.

HIMEM address

Will set the top of memory to the given address, within the range of 1024 to 32767.

INKEY\$ [string or string variable]

Will wait for a key press. With no parameter, it will wait for any key. With a parameter, it will wait for a key to correspond to any character in the string. The ASCII value of the key press is placed in reserved variable 'ST'.

KEY 1 to 16,"data"

Loads function keys with data. Maximum of ten characters per key is permissible. Inputs over ten characters will generate a SYNTAX ERROR, but the first ten characters of data will be assigned to the particular key. To generate a return in the data, use " \leftarrow ".

for example, KEY 7, "LIST←"

KEY

Will display the data assigned to all 16 keys in the format they were first entered. This will allow you to overtype the displayed data to amend the key assignations.

To obtain keys: KEY 1–8 as marked on keys KEYS 9,11,13,15 key with logo KEYS 10,12,14,16, key with shift and logo Note: any KEY command will enable the keys if they have previously been disabled.

LOMEM address

Will set the start of BASIC to the given address, within the range of 1024 to 32767.

MEM

Display amount of memory free as an unsigned number.

MERGE ["program name"],[device]

Merges a stored program with that currently in memory according to their line numbers. Lines numbers of the merging program take precedence. If no program name and/or device then the command will default to tape. With no name then first program on tape will be merged.

OFF

Disable the function keys.

OLD

Restores a BASIC program after a NEW or system reset have been actioned. This will not work if an edit has been carried out before OLD is actioned.

PLOT (x,y)

Sets the cursor column and row position. x-0 to 39 and y-0 to 24. 0,0 is the top left hand corner of the screen (cursor home).

POP

Rectifies stack on leaving a subroutine before a RETURN has been called.

PROC name

Calls a procedure called 'name'.

QUIT

Disables the utility and its commands, but protects the area that it uses. The UTILITY can be initialized again by SYS 32768.

RENUM first line number to be changed or 0, increment, new start line number

If first parameter is 0, the whole program will be renumbered, otherwise, from designated line to the end of program. Renumbers the following tokens: GOTO, GO TO, GOSUB, IF THEN, RUN, ON GOTO, ON GOSUB and RESET. It will not renumber CGOTO or CGOSUB.

RESET [line number]

Restore DATA pointer to specific line or start of program.

TEN decimal number[,.....

Prints hex conversion of a decimal number.

TRACE

A diagnostic to follow the execution of a BASIC program as it runs.

TROFF

Disables TRACE function.

TWO decimal number[,.... Prints binary conversion of decimal number.

WRITE (x,y)

Same as PRINT, but sets cursor position first as in PLOT.

Note: All commands performing number conversions will do more than one conversion if the values are separated by commas.

ERRORS

These are particular to the UTILITY.

CODER

STRING TOO LONG – more than 254 bytes have been generated by the mnemonics for one program line.

OUT OF DATA - found a character not handled by CORDER.

RENUM

ILLEGAL DIRECT – line at which to start renumbering does not exist. UNDEF'D STATEMENT – no destination found for a GOTO or GOSUB directive.

APPENDIX K: 64 low memory map

The following is the first few pages of the memory map in the *Programmer's Reference Guide* (PRG), Chapter 5. It is included to avoid continual reference to the PRG to look up label addresses. Some of the descriptions have been changed through personal experience or preference to those in J. Butterfield's map.

LABEL	hex	decimal	Description
D6510	0000	0	6510 direction register
R6510	0001	1	6510 I/O, memory and tape
ADRAY1	00030004	3-4	Float to fixed vector
ADRAY2	0005-0006	56	Fixed to float vector
CHARAC	0007	7	Search character
ENDCHR	0008	8	End of quote flag
TRMPOS	0009	9	Save screen last TAB
VERCK	000A	10	Flag: LOAD=0 VERIFY=1
COUNT	000B	11	Ptr input buffer/#subscripts
DIMFLG	000C	12	Default DIM to 10 flag
VALTYP	000D	13	DATA type:string=255 numeric=0
INTFLG	000E	14	:integer=128 float=0

LABEL	hex	decimal	Description
GARBFL	000F	15	DATA scan/LIST quote/garbage
			collection flag
SUBFLG	0010	16	subscript/user fn call
INPFLG	0011	17	\$00=INPUT \$40=GET \$80=READ
TANFLG	0012	18	TAN sign/comparison
	0013	19	current I/O prompt
LINNUM	0014-0015	2021	integer value
TEMPPT	0016	22	pointer: temp string stack
LASTPT	00170018	23–24	last temp string address
TEMPST	0019-0021	25-33	stack for temp strings
INDEX	0022-0025	34-37	utility pointer area
RESHO	0026002A	38-42	product area for multiply
TXTTAB	002B002C	4344	pointer start of BASIC
VARTAB	002D-002E	4546	, pointer start of variables
ARYTAB	002F-0030	4748	pointer start of arrays
STREND	0031-0032	49–50	pointer end of arrays
FRETOP	00330034	51-52	pointer bottom of strings
FRESPC	0035-0036	53–54	utility string pointer
MEMSIZ	00370038	55-56	pointer highest address used by BASIC
CURLIN	0039-003A	57–58	current BASIC line number
OLDLIN	003B003C	59–60	previous basic line number
OLDTXT	003D003E	6162	BASIC statement for CONT
DATLIN	003F0040	6364	current DATA line
DATPTR	00410042	6566	current DATA address
INPPTR	00430044	67–68	INPUT vector
VARNAM	00450046	6970	pointer current variable name
VARPNT	00470048	71–72	pointer current variable data
FORPNT	0049004A	73–74	pointer variable for FOR/NEXT
	004B004C	75–76	' Y-save/op-save/BASIC pointer save
	004D	77	comparison symbol accumulator
	004E-0050	78-83	misc work area
	00540056	84-86	jump vectors for functions
	00570060	87-96	misc numeric work area
FACEXP	0061	97	FPACC#1:exponent
FACHO	00620065	98–101	FPACC#1:mantissa
FACSGN	0066	102	FPACC#1:sign
SGNFLG	00 67	103	pointer series evaluation constant
BITS	0068	104	FPACC#1:overflow digit
ARGEXP	00 69	105	FPACC#2:exponent
ARGHO	006A006D	106109	FPACC#2:mantissa
ARGSGN	006E	110	FPACC#2:sign
ARISGN	006F	111	sign comparison result
FACOV	0070	112	FPACC#1:low order rounding
FRUEPT	0072-0072	113–114	pointer cassette buffer

LABEL	hex	decimal	Description
CHRGET	0073-008A	115–138	subroutine: get next byte of BASIC
CHRGOT	0079	121	entry point to get same byte
TXTPTR	007A-007B	122–123	pointer current byte of BASIC
RNDX	008B008F	139–143	RND seed value
STATUS	0090	144	KERNAL I/O status ST
STKEY	0091	145	switch:STOP and RVS keys
SVXT	0092	146	timing constant for tape
VERCK	0093	147	LOAD=0 VERIFY=1
C3PO	0094	148	serial output: deferred char flag
BSOUR	0095	149	serial output deferred char
SYNO	0096	150	tape EOT received
	0097	151	register save
LDTND	0098	152	how many open files#
DFLTN	0099	153	input device (default=0)
DFLTO	009A	154	output device (default=3)
PRTY	009B	155	tape char parity
DPSW	009C	156	tape byte received flag
MSGFLG	009D	157	BASIC mode flag \$00=program
			\$80=direct
PTR1	009E	158	tape pass 1 error log
PTR2	009F	159	pass 2 error log
TIME	00A0-00A2	160–162	real-time jiffy clock
	00A3	163	serial bit count/EO1 flag
	00A4	164	cycle count
CNTDN	00A5	165	tape sync countdown/bit count
BUFPNT	00A6	166	pointer tape I/O buffer
INBIT	00A7	167	RS232 input bits
			tape wrt ldr/rd count
BITCI	00A8	168	RS232 input bit count
			tape wrt new byte/rd error
RINONE	00A9	169	RS232 start bit flag
RIDATA	00AA	170	RS232 input byte buffer
			tape scan/counter/ldr
RIPRIY	ØØAB	171	RS232 input parity
			tape wrt ldr length'rd checksum
SAL	00AC00AD	172–173	pointer tape buffer/scrn scroll
EAL	00AE-00AF	174-175	tape end address/end program
СМРО	00B0-00B1	176–177	tape timing constants
TAPE1	00B200B3	178–179	pointer start of tape buffer
BILLS	00B4	180	RS232 out bit count/tape enabled=1
NXIBII	00B2	181	KS232 next bit to send/tape EOT
KODATA	00B6	182	KS232 out byte buffer/rd char error
FNLEN	00B7	183	Length current file name
LA	00B8	184	Current logical file number

LABEL	hex	decimal	Description
SA	00B9	185	Current secondary address
FA	00BA	186	Current device number
FNADR	00BB00BC	186–187	Ptr current file name address
ROPRTY	00BD	189	RS232 out parity/tape rd input char
FSBLK	00BE	190	tape #blocks left to wrt/rd
MYCH	00BF	191	Serial word buffer
CAS1	00C0	192	Tape motor control
STAL	00C1-00C2	193–194	I/O start address
MEMUSS	00C300C4	195–196	KERNAL setup ptr/tape temp address
LSTX	00C5	197	Last key pressed
NDX	00C 6	198	#characters in k/b queue
RVS	00 C7	199	RVS char print flag 1=yes Ø=no
INDX	00C 8	200	Ptr end of line for INPUT
LXSP	00C9-00CA	201–202	Cursor row, col at start of INPUT
SFDX	00CB	203	Current key pressed 64=no key
BLNSW	00CC	204	0=blink cursor
BLNCT	00CD	205	Cursor countdown timer
GDBLN	00CE	206	Character at cursor pos
BLNON	00CF	207	Cursor blink flag on/off
CRSW	00D0	208	Flag: INPUT from screen or
			GET from keyboard
PNT	00D1-00D2	209–210	Ptr current start of screen line add
PNTR	00D3	211	Cursor col on above line
QTSW	00D4	212	Flag: Ø=cursor in edit mode else in
			quote mode
LNMX	00D5	213	Physical screen line length
TBLX	00D6	214	Current row where cursor lives
	00D7	215	Last inkey/checksum/buffer temp data
INSRT	00D8	216	<pre>#inserts outstanding</pre>
LDTB1	00D9-00F2	217–242	Screen line link table
USER	00F3-00F4	243–244	Ptr screen colour
KEYTAB	00F5-00F6	245–246	K/b decode table vector
RIBUF	00F7-00F8	247–248	RS232 input buffer ptr
ROBUF	00F900FA	249–250	RS232 output buffer ptr
FREKZP	00FB00Fe	251–254	Free zero page area
BASZPT	00FF	255	BASIC temp data area
	0100-010A	256–266	Float to ASCII work area
	0100-013E	256–318	Tape error log
	0100-01FF	256–511	Processor stack
BUF	0200-0258	512600	System input buffer
LAT	02590262	601-610	Logical file table
FAT	0263026C	611–620	Device number table
SAT	026D0276	621-630	Secondary address table
KEYD	0277-0280	631-640	Keyboard buffer

LABEL	hex	decimal	Description
MEMSTR	0281-0282	641-642	Start of BASIC memory
MEMSIZ	02820283	643–644	Top of BASIC memory
TIMOUT	Ø285	645	Serial bus time out flag
COLOR	Ø286	646	Current character colour
GDCOL	0287	647	Background colour under cursor
HIBASE	Ø288	648	Start of screen memory: page number
XMAX	Ø289	649	Size of k/b buffer
RPTFLG	Ø28A	650	Flag: 128=repeat all keys
KOUNT	Ø28B	651	Repeat speed counter
DELAY	028D	653	Flag: shift/ctrl/logo key
LSTSHF	Ø28E	654	Last shift pattern
KEYLOG	028F-0290	655-656	K/b table setup ptr
MODE	0291	657	Flag: Ø=disable shift keys 128=enable
AUTODN	Ø292	658	Ø=scroll down enable
M51CTR	029 3	659	RS232 control register
M51CDR	0294	660	RS232 command register
M51AJB	0295-0296	661662	RS232 non-standard baud rate
RSSTAT	0 297	663	RS232 status register
BITNUM	Ø 298	664	RS232 bits left to send
BAUDOF	0299-029A	665-666	RS232 Baud rate
RIDBE	Ø29B	667	RS232 index to end of input buffer
RIDBS	029C	668	RS232 page number of start
			of input buffer
RODBS	029D	669	RS232 page number of start
			of output buffer
RODBE	029E	670	RS232 index to end of output buffer
IRQTMP	029F02A0	671-672	IRQ save during tape I/O
ENABL	02A1	673	RS232 enable/CIA 2 (NMI) interrupt control
	0 2A2	674	CIA 1 timer A control log during
			tape I/O
	Ø2A3	675	CIA 1 interrupt log tape read
	Ø2A4	676	CIA 1 Timer A enable log tape read
	02A5	677	Screen line marker
	Ø2A6	678	PAL/NTSC flag Ø=NTSC 1=PAL
	02A7-02FF	679–767	Unused
	02C0-02FE	704766	Block 11 for sprites
ERROR	0300-0301	768–769	Vector: BASIC error message (\$E3B8)
MAIN	03020303	77 0 –771	Vector: BASIC warm start(\$A483)
CRNCH	0304-0305	772–773	Vector:Crunch BASIC tokens(\$A57C)
QPLOP	0306-0307	774–775	Vector:Print BASIC tokens(\$A71A)
GONE	03080309	776–777	Vector: Start new BASIC line(\$A7E4)
EVAL	030A030B	778–779	Vector: BASIC token evaluate(\$AE86)
SAREG	030C	780	Save A register

LABEL	hex	decimal	Description
SXREG	030D	781	Save X register
SYREG	030E	782	Save Y register
SPREG	030F	783	Save status register
USRPOK	0310	784	USR function jump instrn (\$4C)
USRADD	0311-0312	785–786	USR address low/high form(\$B248)
	Ø 313	787	Unused
CINV	03140315	788–789	Vector:Hardware IRQ(\$EA31)
CBINV	0316-0317	790791	Vector:BRK interrupt(\$FE66)
NMINV	0318-0319	792–793	Vector:NMI(\$FE47)
IOPEN	031A031B	794795	Vector:KERNAL OPEN(\$F34A)
ICLOSE	031C031D	796–797	Vector: KERNAL CLOSE(\$F291)
ICHKIN	031E-031F	798–799	Vector:KERNAL CHKIN(\$F20E)
ICKOUT	0320-0321	800-801	Vector: KERNAL CHKOUT(\$F250)
ICLRCH	0322-0323	802803	Vector: KERNAL CLRCHN(\$F333)
IBASIN	03240325	804805	Vector: KERNAL CHRIN(\$F157)
IBSOUT	03260327	806807	Vector: KERNAL CHROUT(\$F1CA)
ISTOP	03280329	808-809	Vector: KERNAL STOP(\$F6ED)
IGETIN	032A032B	810-811	Vector: KERNAL GETIN\$F13E)
ICLALL	032C032D	812-813	Vector: KERNAL CLALL(\$F32F)
USRCMD	032E032F	814815	Vector:Warm start(\$FE66)
ILOAD	0330-0331	816-817	Vector: KERNAL LOAD(\$F4A5)
ISAVE	0 332 0 333	818-819	Vector: KERNAL SAVE(\$F5ED)
	0334033B	820-827	Unused
TBUFFR	033C03FB	8281019	Tape I/O buffer
	03FC-03FF	10201023	Unused
	0340037E	832894	Block 13 sprite data
	038003BE	896–958	Block 14 sprite data
	03C003FE	960-1022	Block 15 sprite data
VICSGN	040007FF	1024-2047	Screen memory
	040007E7	1024-2023	Visible memory
	07F807FF	20402047	Sprite block data pointers 0–7
	0800	2048	Start of BASIC (TXTTAB-1)



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