

## CLIVE EMBEREY \& BOB TURNER

# INVALUABLE UTILITIES for the 

 COMMODORE 64'The complete programmer's toolkit essential programming aids for your micro'

# Clive Emberey and Bob Turner Invaluable Utilities for the Commodore 64 

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## Introduction

This book as the title suggests, is a book of utilities for the Commodore 64. It has been written not only to provide a set of useful routines, but also to help you to begin to understand some of the more detailed workings of your 64 .

We have tried to cover a reasonable spectrum and hope that through our examples you will attack areas other than those covered here with increased confidence. Towards this end we have covered in depth the development, background and implementation of each utility.

We have made no attempt to cover programming, in either BASIC or machine code, in this book because many other texts cover this in detail. We have also assumed that most serious 64 users will be in possession of a copy of the Programmer's Reference Guide and have adopted its nomenclature throughout, particularly with reference to memory locations and KERNAL routines.

Wherever possible a utility has been implemented in both BASIC and machine code. We felt that the BASIC versions, though sometimes crude, are easier to experiment with and should also help those readers unfamiliar with machine code to appreciate the workings of the equivalent routines. They also go to prove that it is not what you know, but how you use it. In some cases it might prove beneficial to use the BASIC rather than the machine code versions. Typical circumstances might be where only one or two features are required, or when you need the full 38 K of RAM available to BASIC, or if you wish to switch to bank 2 when using graphics.

To facilitate entering the machine code it has been given in two forms: as BASIC loaders and assembler listings. The assembler listings are suitable for use with an extended monitor. For anyone not owning a monitor program, we have included Jim Butterfield's Supermon and instructions for its use in Appendices D and E. (Supermon for the 64 was first published in the January 1983 issue of Compute). However, before attempting the considerable task of typing it in and then getting it to work, ask around your friends and user groups as they may have a copy. If you do find one, you will save yourself a lot of time, effort and frustration. Jim Butterfield has also published a very complete Memory Map for the 64 in the October 1982 issue of Torpet which has since appeared in many other journals. This complements the one in the Programmer's Reference Guide as it gives the nominal entry points to
most ROM routines. A copy of this map could save you a great deal of time when disassembling ROM routines to find out how they work.

To assist in entering the BASIC code all listings have been provided in an annotated form. This, we hope, will avoid the all-too-common problems associated with deciphering the symbols for cursor keys, function keys, colours, and so on when in quotes mode. A detailed list of all mnemonics used is given in Appendix H, but you should find that most are self-explanatory. The program we wrote to generate them is included in this book in the UTILITY as the CODER command. The listings as given in the text always have a maximum line length of 40 characters in their annotated form. Where a line exceeds forty characters it is continued on the next and subsequent lines always commencing in column 1. When looking at the listings you may find it helpful to compare the rightmost characters of continued lines. When the code is typed in, replacing the mnemonics with the correct key(s), no line will exceed 80 characters on the display.

Any of the BASIC utilities intended for use from within another program have been numbered in the 60000's to allow you to merge them with your own programs (using the simple technique described in Chapter 5 or the MERGE command of the UTILITY itself).

We have chosen to put our code at 32768 ( $\$ 8000$ ), leaving the BASIC programmer with 30 K free. There is no reason why the code could not be modified and relocated elsewhere in memory (the 4 K block from $\$ C 000$ is not a bad idea) and the initialization routine adjusted to take advantage of the increased memory available. In fact, nearly all the routines were developed and tested in isolation, being enabled by a simple SYS call. They were then incorporated into the UTILITY by simply relocating and including a keyword and token to activate them. To conserve memory, common subroutines have not been duplicated. Often a pick ' $n$ ' mix approach was found useful to check out a range of extensions which relied heavily on common routines.

We have used 'hidden' RAM beneath BASIC to store data to conserve valuable user RAM and implemented a simple switching routine to access this data when necessary. Applications like setting up the function keys require access to this RAM, as does CODER. We have made extensive use of the ROM routines and RAM vectors available, but on some occasions found it more economic and faster to write our own code. The UTILITY, in the form given, occupies the same area as cartridge ROM and cannot therefore co-exist with cartridges. It was not written to run in conjunction with them and is intended as a standalone, extendable facility. As the owners of a disk unit will have received DOS 5.1. on the demo disk, the UTILITY has been written to co-exist with DOS 5.1. Some of the isolated routines will temporarily disable DOS as they make use of the same operating routine - CHRGET - but more about that later. However, a simple SYS call will restore DOS 5.1 commands.

There are many commercial utilities and BASIC extensions. These may be purchased at reasonable prices and for many applications there will be no better solution. However, if you are interested in the Commodore 64 and wish to get the most from it, you may appreciate having a range of routines which you can modify, extend and, indeed, improve upon. After all, you can pay upwards of the cost of this book for a fairly simple renumber routine.

Before we finish, we would like to leave you with two suggestions and an option:

1) Always save a program before running it
2) Always make backup copies

This is good advice for BASIC and essential where machine code is concerned.
3) It is very easy to wire a reset switch to your 64 and the necessary reset line is available at both the Serial I/O and User I/O (see the Programmer's Reference Guide, Appendix I). This is almost essential if you use machine code, but don't attempt this if you are not sure what you are doing.

Have fun (if that's the right word)!

## 1 BASIC on the 64

## Introducing basic

On powering up your 64 you will find it ready and waiting to go in BASIC, as part of the power reset sequence is to initialize BASIC and leave the user in direct mode.

The implementation of BASIC that Commodore has chosen to use for the 64 is identical to that on the VIC20 and PET microcomputers prior to the 4000 series. BASIC 2, as it is often called, differs from the later version only in its disk operating commands, the latter having a greatly improved and simplified instruction set for disk control. In producing BASIC 4 Commodore did maintain $99.99 \%$ downward compatibility and in doing so allowed users to run any program on a higher series PET. It was, therefore, a little surprising to find the BASIC on the 64 to be only V2. This may have been done to avoid a conflict of interests in so much as the new CBM micro, though in many respects far more powerful, was not quite the same.

BASIC, or to give its full title Beginners' All-purpose Symbolic Instruction Code, runs on the 64 as a high-level interpreted language. It is a subset of Microsoft's BASIC (who wrote the first implementation for the early PETs and now produce MBASIC and the MSDOS operating system for all major microcomputers). The history of BASIC is nothing to do with this book, but it is interesting to note that regardless of environment, or cost of system BASIC will usually be in there somewhere. It may only run in compiled form, or it may be syntactically different, but it is reassuring to know that a knowledge of 64 BASIC should allow you to grasp quickly other BASICs on other machines.

BASIC has its critics, particularly of CBM BASIC, who would advocate the use of Pascal, or Pilot, or Forth, or . . . Each of these languages is particularly suited to a range of tasks, but perhaps none lends itself as well as EASIC to the task of rapid development of 'untidy' and 'unstructured' programs which, most importantly, work. Arguments for and against will no doubt long continue, but as we are supplied with BASIC, let us make the most of it.

As its name imples, BASIC was developed to allow beginners to acquire programming skills rapidly. It adopted a system of naming its commands and functions to indicate the action produced. For example, if we wish to halt the execution of a program we issue the
all-too-clear command : stop. For non English-speaking countries even the use of English is no problem on the 64 as not only is it a simple matter to redefine the character set, but it is also easy to redefine the keyword table itself.

On the 64 we have 75 ( 76 if you include GO) BASIC commands, functions and operators as standard. For many applications this is perfectly adequate. Life would be simpler if more commands were available. Increasing the number of commands has the benefit of providing a more versatile programming language, but the disadvantage of slowing down the execution of the existing commands. This is true of any interpreted programming language. However, the way in which the interpreter has been implemented does allow you to add to the keywords to your heart's content, providing you understand how it works and are capable of writing the necessary machine code. For the moment the interpreter will be considered simply as a means of translating our 'meaningless entries' into something which is executable by the 6510 microprocessor at the machine code level. This it does by taking an instruction, finding the appropriate machine code routine, carrying it out and then returning to implement the next. The process is slow but very flexible and even allows us to interrupt the execution and take control should we wish to do so.

One of the best features of the 64 must be the screen editor. It allows changes to be made directly to anything appearing on the screen and, more importantly, allows you to implement these changes. The disadvantage is that the maximum length of program line or direct statement is limited to 80 characters (or two screen lines). Use of the standard abbreviations of first character and second (or third) character shifted, instead of typing the full keyword, does allow program lines, on listing, to exceed this limit. They cannot, however, be easily edited. On pressing return to acknowledge the end of the edit only two screen lines are accepted. Anything beyond this point is not included in the revised line. Still, this limitation is far outweighed by the speed at which it allows existing code to be edited and repetitive code to be entered by simply using the cursor keys, altering the line number, modifying the necessary part of the line and pressing return to enter the new line. It even allows us to write programs which can generate their own program lines, as we will see in Chapter 5.

BASIC may be used in two modes. These are direct (when a command is typed in without a line number and executed immediately) and program (a command preceded by a line number which is not executed until the program is RUN).

## Storage of basic code

If we wish to examine a program, we may do so with the LIST command. What we see has undergone many processes from the form in which it
was stored in ram. To view the code in situ we first of all need to know where to look. In the default mode on powering-up a basic program will be sorted from memory location 2049 ( $\$ 0801$ ) upwards. We can examine a program by simply Peeking out each of the locations used by
FOR I = START TO START + 200: PRINT PEEK (I) ; :NEXT I

We would see a series of decimal numbers with only the fact in common that none was less than zero or greater than 255 . We might also notice some sort of related pattern occurring, but not a great deal more. We could adopt another approach by moving an area of Ram used for program storage to the screen. This is easily accomplished, but in doing so we must also remember to set a colour at the screen location we are putting the data into for it to be visible. The resulting display is easier to decipher if the 64 is put in lower case mode by pressing the shift and logo keys together. The following line should be typed in direct mode:

$$
\begin{aligned}
& S=0: F O R \quad I=S T A R T \\
& \text { TO START }+800: \text { POKE } 1024+S, \text { PEEK }(I): \text { POKE } \\
& 55296+S, 14: S=S+1: N E X T ~ I ~
\end{aligned}
$$

If you wish to start at the beginning of a program then 2049 must be used and it assumes, as does the first example, a program to be present which occupies memory at least to START +800 (or +200 ). This time we see a series of characters and where our program has text within quotes it appears almost unchanged as do variable names, punctuation and constants. If we combine the processes, and to produce a more consistent format express the numbers in hexadecimal format, we begin to see some sort of relationship. (You had better get used to using hexadecimal notation as we use it extensively, but to help you on your way there is a table of decimal to hex conversions in Appendix B). The following program does just this and may be used to examine itself. If you wish to experiment, simply enter new lines with numbers less than 60000. Those of you with extended monitors or who jumped straight in and typed in Supermon can use the 'memory display' option.

The program displays on each line the start address and the values held in this and the next seven locations. At the right of the line the characters with ASCII (CHR\$()) codes corresponding to the byte values are printed. To avoid confusion, only those characters which are easily discernible are printed; all others are expressed by a ".". Appendix C of the PRG gives the full range of ASCII and $\mathrm{CHR} \mathrm{\$()}$ codes. If you want to display all the characters then some of the codes will have effects which will destroy the display, for example, cursor moves, clear screen, colours, and so on; so you will have to trap these. They do, however, occur in blocks and are therefore not too difficult to isolate.
As will be standard practice throughout this book, a description precedes most program listings.

## LINE ACTION

## 130 Examine selected range in groups of eight.

140 Convert current start address to low/high byte format, that is, units (0-255) and lots of 256 s .
Then convert to hex notation in two stages.
150 Get eight successive bytes from start and print two digit hex to values each time.
170
190 Convert eight bytes to ASCII characters if printable.
to
220 Else replace with a '.' and build eight character string.
230 Print string. Recycle if not end else start again.
1000 Convert start address to hex in two steps.
2000 Convert byte to two digit hex.

```
100 PRINT"MEMORY DISPLAY"
110 INPUT"DISPLAY FROM";F
120 INPUT"[10SPC]TO";T:IF T<F THEN PRINT
    "LESS THAN FROM":GOTO 120
130 FOR I=F TO T STEP 8
140 X=I:GOSUB 1000
150 FOR J=I TO I +7
160 X=PEEK(J):GOSUB 2000:PRINT X$;" ";
170 NEXT J
180 PRINT " ";
190 A$="":FOR J=1 TO I+7
200 X=PEEK(J):IF X<32 OR X>95 THEN A $=A $
+".":GOTO 220
210 A$=A$+CHR悪(X)
220 NEXT J
230 PRINT A$:NEXT I:GOTO 100
1000 MSB=INT (X/256):LSB=I -MSB*256
1010 X=MSB:GOSUB 2000:PRINT X X;
1020 X=LSB:GOSUB 2000:PRINT X$;" ";
1030 RETURN
2000 X =1=INT (X/16): X2=X-X1*16
2010 X$=CHRक(X1+48-7*(X1)9))+CHRक(X2+48-
7*(X2>9))
2020 RETURN
```

If the program is used to examine itself by entering a start of 2048 and an end of 2504 for the program as listed the following display is given:

| 0970 | 2 | , | 30 | 30 | 30 | 3A | 8 | 2 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0978 | 58 | 24 | 3B | 22 | 20 | 22 | 3B | 00 | 入郘; " |
| 0980 | 86 | 09 | 06 | 04 | 8E | 00 | A 0 | 09 |  |
| 0988 | D0 | 07 | 58 | 31 | B2 | B5 | 28 | 58 |  |
| 0970 | AD | 31 | 36 | 29 | 3A | 58 | 32 | B2 | - |
| 0998 | 58 | AB | 58 | 31 | AC | 31 | 36 | 00 | $\times . \times 1.16$. |
| 69A0 | CB | 09 | DA | 07 | 58 | 24 | B2 | C7 |  |
| 09A8 | 28 | 58 | 31 | AA | 34 | 38 | AB | 37 | ( $\times 1.48 .7$ |
| 0980 | AC | 28 | 58 | 31 | B1 | 39 | 29 | 29 | . ( $\times 1.9$ ) |
| 0988 | AA | C7 | 28 | 58 | 32 | AA | 34 | 38 | . ¢ $^{(\times 2.48}$ |
| $09 \mathrm{C0}$ | AB | 37 | AC | 28 | 58 | 32 | B1 | 39 | . $7 .(\times 2.9$ |
| $09 \mathrm{C8}$ | 29 | 29 | 00 | D1 | 09 | E4 | 07 | 8E |  |
| 0900 | 00 | 00 | 00 |  |  |  |  |  |  |

Its exact format will vary depending on how you typed the program in. A number of things are immediately apparent. All text inside quotes, all variable names, all destinations, all constants and punctuation appear unchanged. From just this information we can work out the general area of each line. Taking line 100 as an example, "MEMORY. .." is clearly visible from $\$ 0806$ to $\$ 0815$. Immediately preceding it is the value $\$ 99$ which, not surprisingly, is the tokenized value for PRINT. The two bytes before this are $\$ 64$ and $\$ 00 . \$ 64$ is the hex for 100 which is the line number. Line numbers may range from 0 up to 63999 and, like many values on the 64 , are stored in low/high byte format. The actual line number is $\$ 64+\$ 00 \star \$ 0100(100+0 \star 256)$. If we look along the hex values for line 100, we see that the byte immediately following the closing quote is a zero. This is how BASIC marks the end of a program line. The two bytes preceding the line number are $\$ 17 / \$ 08$ which is the address (low/high) of the byte immediately following this end of line zero. These two bytes are known as the link address and point to the link address (and start) of the next line. If we follow the link address through the program, the sequence runs \$0817/082D/0864....09D1 and finally 0000 . A link address of zero marks the end of the program, which in this case is \$09D1. A pointer to this address +2 is held in zero page (locations $\$ 00=\$ F F)$ at VARTAB ( $\$ 2 \mathrm{D} / 2 \mathrm{E}$ ) and marks the start of the BASIC variables. A second pointer on zero page, TXTTAB (\$2B/2C), points to the start of the program. This is the location of the first link address and in the default setting this will always be $\$ 0801$. Location $\$ 0800$ holds a zero; the byte before the start of a program must always be zero for RUN to work. This becomes of more significance when the start location of BASIC is changed.

A program can therefore be thought of as a 'linked list' of individual program lines. It is of the form:

| START | LINK | LINE |  | END |  | END PROG |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 00 | low high | low high | BASIC line | 00 | low high.... | 000000 |


|  | 00 | 17 | 08 | 64 | 00 | 99 | 22 | 4D |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 08 | 45 | 4D | 4F | 52 | 59 | 20 | 44 | 49 | EM |
| 08 | 53 | 50 | 4C | 41 | 59 | 22 | 00 | 2D | SP |
| 0818 | 08 | 6E | 00 | 85 | 22 | 44 | 49 | 53 | ...."DIS |
| 0820 | 50 | 4C | 41 | 59 | 20 | 46 | 52 | 4F | PLAY FRO |
| 0828 | 4D | 22 | 3B | 46 | 00 | 64 | 08 | 78 | M"; F |
| 0830 | 00 | 85 | 22 | 20 | 20 | 20 | 20 | 20 |  |
| 838 | 20 | 20 | 20 | 20 | 20 | 54 | 4F | 22 |  |
| 840 | 3B | 54 | 3A | 8B | 20 | 54 | B3 | 46 |  |
| 848 | 20 | A7 | 20 | 99 | 20 | 22 | 4C | 45 |  |
| 0850 | 53 | 53 | 20 | 54 | 48 | 41 | 4E | 20 | SS THAN |
| 0858 | 46 | 52 | 4F | 4D | 22 | 3A | 89 | 20 | FRO |
| 860 | 31 | 32 | 38 | 00 | 76 | 08 | 82 | 80 | 1 |
| 0868 | 81 | 20 | 49 | B2 | 46 | 20 | A4 | 20 |  |
| 0870 | 54 | 20 | A9 | 20 | 38 | 00 | 85 | 08 |  |
| 8878 | 8C | 80 | 58 | B2 | 49 | 3A | 8D | 20 |  |
| 880 | 31 | 30 | 30 | 30 | 00 | 95 | 88 | 96 | 1 |
| 888 | 00 | 81 | 20 | 4A | B2 | 49 | 20 | A4 |  |
| 890 | 20 | 49 | AA | 37 | 00 | B1 | 88 | A ${ }^{\text {d }}$ |  |
| 898 | 00 | 58 | B2 | C2 | 28 | 4A | 29 | 3A | X |
| 08A8 | 8D | 20 | 32 | 30 | 30 | 30 | 3A | 99 |  |
| 0848 | 20 | 58 | 24 | 3B | 22 | 20 | 22 | 3B |  |
| 0880 | 00 | B9 | 08 | AA | 00 | 82 | 20 | 4A |  |
| 0888 | 00 | C4 | 08 | B4 | 00 | 99 | 20 | 22 |  |
| 08C0 | 20 | 22 | 3B | 00 | DA | 08 | BE | 00 |  |
| $08 \mathrm{C8}$ | 41 | 24 | B2 | 22 | 22 | 3A | 81 | 20 |  |
| 08D0 | 4A | B2 | 49 | 20 | A4 | 20 | 49 | AA | J. |
| 0808 | 37 | 08 | 05 | 09 | C8 | 00 | 58 | B2 |  |
| 08 | C2 | 28 | 4A | 29 | 3A | 8B | 20 | 58 |  |
| 08 E 8 | B3 | 33 | 32 | 20 | B0 | 20 | 58 | B1 | . $32 . \times$ |
| 88 | 39 | 35 | 20 | A7 | 20 | 41 | 24 | B2 | 95 |
| $08 F 8$ | 41 | 24 | AA | 22 | 2E | 22 | 3A | 89 |  |
| 0908 | 20 | 32 | 32 | 30 | 08 | 14 | 09 | D2 |  |
| 0908 | 00 | 41 | 24 | B2 | 41 | 24 | AA | C7 | A |
| 0910 | 28 | 58 | 29 | 00 | 1 C | 09 | DC | 06 | (X) |
| 0918 | 82 | 20 | 4A | 00 | 2F | 09 | E6 | 00 | - |
| 0920 | 99 | 20 | 41 | 24 | 3A | 82 | 20 | 49 |  |
| 0928 | 3A | 89 | 20 | 31 | 30 | 30 | 08 | 4E | . $180 . N$ |
| 0930 | 09 | E8 | 03 | 4D | 53 | 42 | B2 | B5 |  |
| 0938 | 28 | 58 | AD | 32 | 35 | 36 | 29 | 3A | ( |
| 0940 | 4C | 53 | 42 | B2 | 49 | AB | 4D | 53 | LSB.I.MS |
| 0948 | 42 | AC | 32 | 35 | 36 | 00 | 65 | 09 | B |
| 0950 | F2 | 03 | 58 | B2 | 4D | 53 | 42 | 3A | . . $\times$.MSB |
| 0958 | 8D | 20 | 32 | 30 | 30 | 30 | 3A | 99 | 2000 |
| 0960 | 20 | 58 | 24 | 3B | 00 | 80 | 09 | FC | - |
| 0968 | 03 | 58 | B2 | 4C | 53 | 42 | 3A | 8D | $\times$ |

The link addresses are not used when a program is run, but are important during listing and editing. We can alter their values without affecting the way a program runs, but on listing some strange effects are produced.

We can now look through the display and find the start and end of a line and the associated line number. Knowing these, we can start to deduce the tokenized values for the basic keywords used. By adding lines to the program we could find out all keyword values, but to save you the effort we have produced a complete list in Appendix A. This table has been extended to include the new token values used by the utility. These values should be ignored for the moment. With a little practice, reading displays of this type becomes very easy.

From the table in Appendix A we see that all basic keywords have token values in excess of 127 (\$7F). The highest token used for standard BASIC is 203 (\$CB) for the GO command. (GO simply searches for a corresponding TO to ensure GOTO is equivalent to GO TO). When a line is entered from the keyboard it is transferred to the input buffer (buf \$0200-0258) on pressing return. The line is then tokenized in accordance with this table with keywords being processed first. If no line number is present, the BASIC interpreter immediately executes the statement(s). If one is present then the line is put in its numerically correct position and the link addresses for the whole program are recalculated and VARTAB updated. A similar process is carried out when a line is deleted. These operations are discussed in greater detail in Chapter 3.

## Variables

## General

BASIC allows three types of variable. These are real, integer and string. String and integer are distinguished from real by trailing ' $\$$ ' and ' $\%$ ' characters respectively. The default is therefore to real. Variable names may be of any length, but only the first two characters and the last character are of significance. This means $\operatorname{ABXXXX} \%$ will be considered equal to $A B Y Y Y \%$ and to $A B \%$, but different from $A B X X X X$ or $A B X X X X \$$. The last character is used to distinguish the variable's type, and if it is not one of the special characters above then the variable is treated as real. The only limitation on naming variables is that the first character must be alphabetic and the subsequent character, alphanumeric, providing that they do not form reserved keywords. For example, PEND would be treated as ' $p$ ' plus the keyword END and on running would produce a syntax error. Reserved words are any of those occurring in Appendix A with a token value exceeding 127 (bearing in mind that with the UTLITY in place the number of reserved words will be increased).

To allow each of the three types of variable (four, really, when we
include function names) to be stored in two bytes, the high bit is set or unset on each of the name bytes to give the necessary four combinations. These are:

|  | Name |  |
| :--- | :--- | :--- |
|  | Type | 1st char |
|  |  |  |
| 2nd char |  |  |
| REAL | ASCII |  |
| INTEGER | ASCII +128 | ASCII +128 |
| STRING | ASCII | ASCII+128 |
| FUNCTION | ASCII+128 | ASCII |

Where the name is only a single character, the second byte is zero or 128 as appropriate.

Each of the three types of variable may be used in multi-dimensional arrays. These subscripted variables follow the same rules as for simple variables with the addition of a ' ' following the name and type. This tells the interpreter it is dealing with an array and is handled accordingly.

## Storage of variables

Any variable created in either direct or program mode is stored after the program currently in memory. Variables are stored in the order in which they are created. Strings are stored slightly differently from numeric values mainly due to their dynamic nature and are held in two parts. The first is a pointer to the string's location and the second is the length of the string itself. Strings are stored at the top of BASIC memory (\$9FFF/40959) and grow downwards. The current lower limit of string storage is stored in FRETOP (\$33-\$34/51-52).

When searching for a variable the interpreter starts at the end of the program and searches upwards in memory for the named value according to the rules given above. If the variable does not exist, the next available space is allocated to it. Thus, if we define the more important values early on they can be accessed quicker and the time spent in 'garbage collection' reduced.

Variables are either simple or subscripted.

## Simple variables

All non-subscripted variables use seven bytes of RAM. The first two hold the name in its adjusted form. For each real variable the remaining five bytes are used in the following way: one for its exponent and the remaining four for its sign and mantissa. Integers are stored in only two bytes with the remaining three unused. Strings use one byte to indicate the length and two bytes to point to the location of the characters, which is usually at the top of memory (though not always). A function also uses seven bytes, of which the third and fourth point to
its definition (DEF FN), the next two point to the variable it uses and the last points to an initial value of the variable (zero).

The following table summarizes the storage of simple variables:

## Byte

| 1 and 2: name | 3 | 4 | 5 | 6 | 7 |
| :--- | :--- | :--- | :--- | :---: | :---: |
| REAL | exp | sign + M1 | M2 | M3 | M4 |
| INTEGER | sign+high | low | unused | unused | unused |
| STRING | length | ptr low | ptrhigh | unused | unused |
| FUNCTION | pointer to DEF FN | pointer to variable | initial value |  |  |
|  | low | high | low | high |  |

If we add the following lines of code to our memory display program, all variable types are generated (including arrays):

1 DIM A(5),B\%(5),C\$(5),D\$(1,5)
2 FOR $\mathrm{I}=0$ TO $5: \mathrm{A}(\mathrm{I})=\mathrm{I}: \mathrm{B}(\mathrm{I})=\mathrm{I}: \mathrm{C} \$(\mathrm{I})=\mathrm{CHR} \$(64+\mathrm{I}): \mathrm{D} \$(2, \mathrm{I})=\mathrm{C} \$(\mathrm{CI})$ :
NEXT
$3 \mathrm{M} \$=\mathrm{M} \$+{ }^{\prime \prime} \mathrm{STRING} 1^{\prime \prime}: \mathrm{Z} \$=\mathrm{Z} \$+{ }^{\prime \prime}$ STRING2"

We can now dump the memory associated with simple variables. The area to be displayed can be worked out from VARTAB, ARYTAB and FRETOP. To do this the program is RUN twice, the first time from 2641 (\$0A61) to 2732(\$0AAC) to display the variables and the second time from 40900 (\$9FC4) to 40960 ( $\$$ A000 - the start of BASIC ROM) to display the strings in situ:

Simple variables and string pointers


Strings in situ

| 9FC4 | 30 | 30 | 30 | 30 | 30 | 31 | 30 | 30 | 00000100 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 9FCC | 31 | 32 | 35 | 35 | 32 | 38 | 43 | 43 | $125528 C C$ |
| 9FD4 | 38 | 30 | 30 | 30 | 30 | 34 | 39 | 39 | 80000499 |
| 9FDC | 34 | 35 | 31 | 31 | 35 | 30 | 41 | 41 | $451150 A A$ |
| 9FE4 | 30 | 53 | 54 | 52 | 49 | $4 E$ | 47 | 32 | 0 STRING2 |
| 9FEC | 29 | 53 | 54 | 52 | 49 | $4 E$ | 47 | 31 | STRING1 |
| 9FF4 | 45 | 45 | 44 | 44 | 43 | 43 | 42 | 42 | EEDDCCBB |
| 9FFC | 41 | 41 | 40 | 40 | 94 | E3 | $7 B$ | E3 | AA.... |

## Real variables

Looking through the display above, it is quite easy to spot the real variables as their names are stored in unmodified ASCII. At \$0A51 we see ' $I$ ', the first non-subscripted variable to be used, with a zero second byte in its name. We can also spot F, T, X and all the others (noting that MSB and LSB are stored as MS and LS).

Real numbers are stored in binary floating point format, always to an accuracy of 31 bits. Due to the way in which they are stored in single precision form, rounding errors are introduced though these are usually not significant enough to affect the final results. Examples of this type of error are encountered all the time as in the 'X.000001'-type value. We can convert ' $I$ ' back to a decimal number quite easily.

The exponent is stored in byte 3 and is the power of two. A unit change in this doubles or halves the resulting value. Positive exponents are expressed as $129+$ EXP and negative, as 129-EXP. Therefore, the full range is from $2^{\wedge}(-129)$ to $2^{\wedge}(127)$ or in decimal, from about $10^{\wedge}(-38)$ to $10^{\wedge} 37$. The high bit of byte 4 indicates the sign and is set for negative numbers. To calculate the decimal value, we have to successively divide the mantissa starting at the right by 256 , add the result to the next on the right and so on until we reach M1, when we only divide by 128 and finally add 1 . The resulting number will lie between 1 and 1.999999 . This must finally be adjusted for its exponent and sign. The values for ' 1 ' used below are in decimal.

M4 $0 / 256=0$
M3 $\quad(\emptyset+\emptyset) / 256=0$
M2

$$
(0+16) / 256=.0625
$$

M1 $(.0625+37) / 128=0.28955$

$$
+1.00000=1.28955
$$

If this is then multiplied by the exponent of $2^{\wedge}(140-129)=2048$, the value is $2048 \star 1.28955=2640.999$ (almost 2641 ). This is the upper limit for the first memory display. A general formula may be written to convert any real variable from its floating point to decimal form:
$(-)^{\wedge}\left(\mathrm{M} 1\right.$ AND 128) $\boldsymbol{2}^{\wedge}(\mathrm{EXP}-129) \star(1+((\mathrm{M} 1$ AND 127) $+(\mathrm{M} 2+(\mathrm{M} 3+$ (M4/256)/256))/256). . ./128)

## Integer variables

These are stored in a signed high/low byte format and can range from -32768 to 32767 . The high bit of byte 3 is again used to indicate the sign. The value is easily determined from the following:
(BYTE3 AND 127) $2256+$ BYTE4 + (BYTE3 > 127) $\star 32768$

## String variables

These are the easiest of all to pick out. At \$0A59 in the display above is the variable $\mathrm{M} \$$ (its second byte is not used so is set to $\$ 80$ ). Byte 3 tells us it is seven characters long, and bytes 4 and 5 that it is located at $\$ 9$ FED. The seven bytes from $\$ 9$ FED are "STRING1" as would be expected. Strings therefore use seven plus the number of characters bytes of RAM. There is one important point to make before leaving strings. If line 3 had simply been $M \$={ }^{\prime}$ STRING1", its pointer would have pointed to the byte at which it occurred within the program itself, that is, the byte immediately following the quote. Only computed strings are stored at the top of memory which is why the line was written $\mathrm{M} \$=\mathrm{M} \$+$ "STRING1". This economizes on memory usage by only storing the string once. It does have the drawback that if another program is loaded in program mode all non-computed strings are lost.

## Subscripted variables

Arrays may be of any type, but unlike their 'simple' counterparts, only the required number of bytes are used to store the associated values. Real are stored in five bytes, integer in two and strings in three plus their length. In addition to the savings in storing the values, the array name is only stored once. Arrays are also created in the order in which they are encountered.

The area of memory used for arrays immediately follows that for simple variables. As for the latter, it, too, is recorded at two zero page locations. The start, ARYTAB, has already been mentioned when dealing with simple variables. The end, STREND, is held in $\$ 31-\$ 32 / 49-50$. For each new simple variable this whole block must be moved up seven bytes in memory. There will, of course, come a time when array storage builds up to meet that of the descending strings with the resulting 'OUT OF MEMORY' error.

Each array is preceded by a detailed header of the form shown below:

Byte

| 1 and 2 | 3 | 4 | 5 | 6 |
| :--- | :--- | :---: | :---: | :---: |
| NAME | OFFSET TO | NO. | LAST....FIRST |  |
|  | TST VALUE | DIMS | DIM $+1 \ldots$. DIM +1 |  |
| Adj. form | low high | $<256$ | low high..low high |  |

Bytes 1 and 2 hold the name in its adjusted form. Bytes 3 and 4 record
the overall memory requirement for the array (this does not include string data at the top of memory) and is the offset from its start to the next array. Byte 5 records the level of dimensioning and may not exceed 255 (a little difficult to visualize at anything more than two or three). If an undimensioned array is used, this value will default to the number of subscripts at the first occurrence. Successive pairs of bytes then hold the number of elements in each dimension (plus one for the zero subscript) in the reverse order of dimensioning. If no dimensioning has been used, these each default to $11(10+1)$. The following bytes are then used to store the data.

If the program is again run and memory between ARYTAB and STREND displayed, the following results:

| OAAC | 41 | 00 | 25 | 08 | 01 | 08 | 86 | 00 | A. |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $0 \mathrm{AB4}$ | 00 | 00 | 00 | 00 | 81 | 00 | 00 | 00 |  |
| 0ABC | 00 | 82 | 08 | 80 | 08 | 08 | 82 | 40 |  |
| QAC4 | 08 | 00 | 00 | 83 | 00 | 00 | 08 | 00 |  |
| BACC | 83 | 20 | 88 | 08 | 80 | C2 | 80 | 13 |  |
| QAD4 | 08 | 81 | 80 | 06 | 00 | 00 | 00 | 01 |  |
| QADC | 88 | 02 | 08 | 83 | 08 | 84 | 00 | 05 |  |
| QAE4 | 43 | 80 | 19 | 00 | 01 | 00 | 06 | 01 |  |
| QAEC | FF | 9F | 01 | FD | 9 F | 01 | FB | 9F |  |
| 0AF4 | 01 | F9 | 9 F | 01 | F7 | 9F | 01 | F5 |  |
| OAFC | 9F | 44 | 80 | 3F | 00 | 02 | 08 | 06 |  |
| $0 \mathrm{B04}$ | 08 | 03 | 00 | 00 | 08 | 01 | FE | 9 F |  |
| 8B6C | 80 | 00 | 00 | 80 | 08 | 00 | 01 | FC |  |
| 0B14 | 9 F | 00 | 80 | 00 | 00 | 08 | 00 | 01 |  |
| 8B1C | FA | 9 F | 08 | 08 | 00 | 08 | 08 | 08 |  |
| 0 B 24 | 01 | F8 | 9 F | 08 | 00 | 00 | 00 | 00 |  |
| 0B2C | 08 | 01 | F6 | 9 F | 08 | 08 | 80 | 00 |  |
| 0B34 | 08 | 00 | 81 | F4 | 9 F | 08 | 00 | 08 |  |
| 083C | 08 | 01 | FF | FF | 08 | 01 | FF | FF |  |

The first array is 'A(' at \$0AAC. It occupies 37 bytes (\$0025), and has one dimension ( $\$ 01$ ) of five elements ( $\$ 06-1$ ). The six values are then held in 5 byte real format. The next array starts at $\$ 0 A A C+\$ 25=\$ 0 A D 1$. This is 'B\%(' which occupies only $19(\$ 13)$ for its six values. The values are easily read out as $0,1,2,3,4$ and 5 . The next is ' $\mathrm{C}\left({ }^{\prime}\right.$ ', and looking at the previous display we can read out its values as @, A, B, C, D and E. A little care has to be exercised here as in the loop which generated them 'D\$(' was also defined each time. This last array is the most complex of all. In its dimension statement it was defined as $\mathrm{D} \$(2,5)$. However, in its header these are reversed (last....first). The values set were assigned to $\mathrm{D} \$(1, \mathrm{I})$ and from the display we can see these occur at the second and subsequently at every third byte. This shows us that multi-dimensional arrays are
stored in the form $X(0,0) X(1,0) \ldots X(N, 0) X(0,1) X(1,1) \ldots . . X(N, N)$.
This just about concludes our section on variables, except to say that the default values are zero for numeric and null for strings.

## Link addresses and line numbers

## General

Knowing where and in what form these are stored, there is no reason why we cannot modify them from BASIC itself. This we can do using simple POKES to produce some interesting results.

## Links

If we modify a link address, the program will continue to run. It will, however, list in an unusual fashion and be difficult to edit. We can use this fact to make our programs difficult to read and modify. This we can do by hiding lines (the whole program if we wish). Hiding line 110 of the display program, as its listing was originally given, can be done by

POKE 2049,45
This simply skips the link at $\$ 0817$. We could very easily write a short routine to eliminate whole blocks of line numbers. This we leave up to you.

## Line numbers

We can change line numbers as we did link addresses. We could change the line number of 100 to any value we choose.

POKE 2051,110
will, on listing the program, give two line 110s. A little care should be taken here, because if the line number changed is the destination for a GOTO or a GOSUB, some confusion may result.

Saving modified code
BASIC'S SAVE command transfers to tape or disk a copy of the RAM between TXTTAB and VARTAB. This means any modifications are also saved. The modified code returns on loading.

## Modifying BASIC

Changing the load address
We can change the point at which BASIC programs load simply by setting the value in TXTTAB. Wherever the new start is to be, a zero must be set in the byte immediately before it. Once the new start has been set, a NEW will tidy up all other pointers.

Changing the start of BASIC is useful if using sprites or programmable characters within bank 0 (the default). The Programmer's Reference Guide recommends lowering the top of memory to make room for the
necessary data. Instead, why not move up the start of BASIC and leave yourself with far more memory to use?

## Chaining programs

Chaining in this context refers to loading one program from within another.

The question arises as to whether there is a bug in the chaining process. The answer is a qualified 'no' as there can be problems. The effect of a program LOAD is roughly equivalent to executing a GOTO the first statement of the chained program. Thus, the new program can use only real, integer and computed string variables from the first. The problem occurs when the incoming program is larger than the original. If this is the case, it will overwrite the start of the variables, causing utter confusion. Once this has happened you really need to issue a clr at the start of the new program to tidy things up. You have apparently lost all the variables anyway so there is nothing to lose.

On some micros a Chain command exists in addition to the normal load. The action of this command is to move all or only the specified variables out of the way during the loading process and then move them back and update the necessary pointers. On the 64 no such command is available. There are two solutions. The first is to ensure a larger program is never loaded from a smaller one. The second is to make the first program the largest. To do this we do not need to generate a 'large program'. All that is needed is a simple POKE and CLR sequence at the start of the first program to reserve the necessary memory. BASIC can be fooled by:

## POKE 46, (SIZE OF BIGGEST PROG)/256+8+1:CLR

In this example we have not bothered to be exact and have simply reserved to the nearest page.

## Speeding up program execution

There are many ways in which to increase the speed of a program. The speed of peripheral devices plays a major part when inputting or outputting, but the topics covered here are mainly concerned with the 64 itself.

There are commercial 'CRUNCH' or 'compactor' programs available. These traditionally remove all unnecessary spaces, REMs, combine lines not the destination of a GOTO or GOSUB and that is about all. Even these few changes can produce significant increases in execution speed. Some of this we can do from Basic itself. A short routine at the start of a program can combine lines by eliminating link addresses and line numbers, and remove all rems and all spaces not inside quotations. After each deletion, the remaining code is moved down in memory and VARTAB is updated. The end of line marker must be replaced by a ' $:$ ' to separate the last statement from the leading statement on the crunched line. There are a number of problems here. If a line number
is eliminated which is referenced by a THEN, GOTO or GOSUB, the run will fail. The second problem stems from all statements following an IF being ignored if the condition is false. The resulting compacted lines are so long that they cannot be edited. As such, a universal compactor program in BASIC is fraught with danger.

It is far more sensible to consider these points at the time of writing the program. A well-known technique is to ensure that all GOTOs have destinations as near the start of the program as possible as the interpreter starts its search there. If this cannot be done, then the destination should have its high byte greater than that of the GOTO line due to the search technique used which compares this byte first.

There are three other common methods used to optimize the code:
(i) the use of variables rather than constants;
(ii) the setting up of variables in the order of frequency of use;
(iii) and, specific to the 64, turning off the video display when not in use (see Programmer's Reference Guide, Appendix N, 'Screen Blanking').

Using these techniques, the second program below runs almost $25 \%$ faster than the first:

```
10 PRINTTI
20 POKE 49152,0
30 FOR I=0 TO 5000
40 J=J+1
5 0 ~ N E X T ~ I ~
6 0 ~ P O K E ~ 4 9 1 5 2 , 1 6 ~
70 PRINTTI
```

18 PRINTTI:POKE53265,0:P=1:FORI=0T05000:
$\mathrm{J}=\mathrm{J}+\mathrm{P}:$ NEXT: POKE53265,16:PRINTTI

The second program does leave you in $x$ and $y$ scroll mode if you are wondering just what has happened.

## Conclusion

This chapter should have given you one or two ideas to play around with. Before reading Chapter 5 , you might like to think about how to write simple renumber, delete, dump, and recover newed program routines. You might also like to think about how to overcome the chaining problem by, as the last action on leaving a program, moving all variables as high as they can go in memory. The first action of the chained program should be to move them back down to the end of this program and reset the necessary zero page pointers.

## 2 Peripherals

## Introduction

This chapter deals with some of the more common peripherals for the 64. Also included here are the keyboard and screen even though they are not quite peripherals in the same sense as a disk drive, cassette or printer.

It is not our intention to go into any of these in great detail as the subject could fill a book of its own. We have tried to look at features of more immediate use.

## Keyboard

Use of the keyboard, its ROM drive routines and RAM vectors is covered in Chapter 4. Programming of the keyboard is used extensively in Chapter 5. The following are a few useful points not directly covered elsewhere.

Keyboard as a device
The keyboard is viewed as device 0 by the 64 's operating system and is the default for input. As such it may be used like any other device and a file opened to it. This file may only be for input and any attempt to output to it will result in an error. Once opened for input the 'annoying' question mark prompt is removed from the input command display. When information is being obtained from it using input\# all warning messages, such as the double question mark for insufficient data, also disappear. The open format is the same as for any other device:

OPEN 1,0 or OPEN 1, $0,129, " Q W E R T Y "$
In the second example everything following the 0 will be ignored. Use of the keyboard in this way is highlighted in the DISK utility facility in this chapter (see below).

## Auto-repeat

See Chapter 10.
Key detection
See Chapter 4 and Appendix I.

Keyboard Buffer - KEYD (\$0277-\$0280/631-640)
The 64 provides type-ahead of up to ten characters. The buffer operates on the principle of first in/first out. However, once full no new characters will be accepted until it has been partially emptied. Characters are taken singly by a GET, up to the first RETURN on an INPUT and the buffer is emptied on an END.

The length of the buffer is determined by XMAX (\$0289/649) and as this is in RAM it may be changed. Theoretically, the buffer could be lengthened, but in practice this cannot be done as the ram immediately following it is used for other purposes. The size, however, may be decreased and is perhaps most useful when the length is set to 1 where a program requires careful, restricted input or type-ahead is to be discouraged. Setting XMAX to 0 is quite a good way to prevent unwanted user input (the stop key is still active as it is scanned by a different routine - see Chapter 10 to disable).

As the buffer is in RAM we can put data directly into it by simply pokeing the ASCII codes of the characters required. To complete the process NDX (\$C6/198) must be set to tell the system how many characters are in the buffer. This type of approach is used extensively in the BASIC utilities in Chapter 5 and we refer you there for examples of using the keyboard in this way.

One final point before leaving the keyboard which many of you may have already discovered. Pressing the 'Control' key with any other simply sets bit 6 of the ASCII code of the character low. For example, CNRLTT is the same as the del key (\$14/20).

## Cassette

The 64 does not have to use a CBM Datassette. There are, to our knowledge, two manufacturers of interfaces which allow standard cassette recorders to be used. These interfaces duplicate the part of the interface normally resident within the Datassette. It is even possible to use a standard cassette through a suitable edge connector, but do not expect a high success rate in loading back saved programs or to get anything from recordings produced on a CBM recorder (see Programmer's Reference Guide, Appendix I for connection details).

Many consider the cost of the dedicated cassette high. However, it avoids the need to adjust the volume and tone controls to ensure an accurate save (a problem on many other micros where even saving twice is not guaranteed to work). It also seems slow, but perhaps is not as slow as it at first appears. Data is transferred between the 64 and the cassette at about 300 baud (some micros offer an optional fast 1200 baud rate). When a program is saved two copies are made. On reloading the first copy is put into memory and this is then compared with the second to check for and possibly recover load errors. In our experience it has proved worth the additional expense to buy the

Datassette for peace of mind and to avoid the loss of many hours of hard work.

The speed of operation of the cassette has been chosen for reliability, but like most things on the 64 we can even change that. For many years superfast, jet, turbo, fast or whatever you care to call them, operating systems have been available for the PET, more recently for the VIC20 (ARROW) and now for the 64. The machine code listing of the original PET version has even been published. Many games now come with a high-speed load (some without the option for a normal load which has proved annoying when your cassette cannot cope with fast loads). These fast operating systems can be made to run the cassette at a higher speed than the standard operating speed of the disk drive (even this can be increased). There is no secret as to how it is done, but as many software houses pay a royalty for its use, or even sell their own versions of a high-speed loader, we have decided not to include a version of our own.

## LOAD and SAVE with cassette (see Programmer's Reference Guide, Chapter 2) <br> These are dealt with in detail in the Programmer's Reference Guide, so we will deal with them briefly here. The general syntax for SAVE is (where square brackets denote optional parameters):

SAVE['program" or string variable] [,device] [, secondary address]

If no parameters are specified, the BASIC program currently in memory will be copied to the default of cassette without a name.

The secondary address is the more interesting. A secondary address of 2 will write an end of tape marker and one of 3 appears at first sight to do exactly the same. Using either of these will prevent the tape being read beyond this point without being physically wound on. There is, however, a world of difference on loading (see below). With an address of 3 not only is the end of tape set, but an end of tape header is written which is a duplicate of the program header with a type of 5 .

The area of ram saved is that between the values held in TXTTAB and VARTAB. These pointers are automatically kept up to date by the operating system whilst a program is being edited. Should we wish to save an area of memory other than the BASIC program, we can set these up by pokeing in the appropriate values (remember low/high format). This allows us to save machine code from BASIC or even the screen itself. Data stored in memory is more economically saved this way as only single bytes are saved and not the ASCII characters which make up each number (saves at least two bytes per number between $\emptyset$ and 255). The problem is that on returning from the save, the current BASIC program and variables are lost until these pointers are restored. If you
are going to play with TXTTAB and VARTAB from BASIC, put the original values out of harm's way, say below $\$ 0800$ or above $\$ C 000$, to allow them to be recovered.
The syntax for load is identical to that for SAVE:

```
LOAD ["program" or string variable] [,device]
    [,secondary address]
```

LOAD reads the next program from tape. If a program name is specified, then the named program will be searched for and if found loaded or if an end of tape marker is found first the cassette will stop. Again it is the secondary address which is of major importance. A 1 requests the operating system to put the program at the same location from which it was saved. If no secondary address is specified, then providing the program came from an address above the current start of BAsIC it will return to its original location, but after the load TXTTAB will still hold the start of BASIC whereas VARTAB will hold the end address. The same is true when 1 is used, but in this case a load may be carried out below the start of basic. Typically, when loading machine code from basic an 'OUT OF MEMORY' error results if the code locates above $\$ 9 F F F$ due to the setting of VARTAB. A save with a secondary address of 3 ensures the code is reloaded to its original address, regardless of the syntax of the load command (extended monitors use 3).

## Tape Buffer

The tape uses a 192 byte I/O buffer, TBUFFR, which in its default setting extends from $\$ 033 \mathrm{C}-03 \mathrm{FB} / 828-1019$. TBUFFR need not reside here and may be relocated, as a pointer to its start is held in RAM at SAL (\$AC-AD/172-173). To move it, simply POKE in the new location in the usual low/high byte format (STOP/RESTORE will reset it). We have found this of use when storing sprite data blocks in bank 0 when memory is tight (\$C000 is yet again a good place to put it). Usage of the buffer is very different between program and data files. Programs only use the buffer to store their header information (see below) and the transfer of memory is direct from the I/O port without passing through the buffer. Data files, on the other hand, use the buffer initially for the header then subsequently to hold 191 byte blocks (the first byte is used as a marker). This avoids continual starting and stopping of the tape motor and by using this block system the tape is more reliable as it is allowed to pick up speed between each read/write operation. Another zero page location, BUFPNT (\$A6/166), holds the current position within the buffer.

## Tape Headers

All files are stored on tape with an initial header which is the length of the buffer. The exact format depends on the syntax of the SAVE or OPEN command (secondary address of 2 on an OPEN also writes an end of tape
marker). Each is made up of an identifier, two addresses and a file name, the format of which is given below:

## Program headers

| ID | START | END | FILE NAME (spaces to pad) |
| :--- | :--- | :--- | :--- |
| 1 | 18 | 2516 | $6566673232 \ldots .32$ |

Data headers

| ID | START | END | FILE NAME |
| :--- | :--- | :--- | :--- |
| 4 | 603 | 2523 | 6865663232 |

The ID identifies the file type and for a program may also take a value of 3 . The two bytes immediately following it are the start load address in low/high format and the next two the end address. The file name is not limited to 16 characters and in fact can be up to 187 characters. This allows machine code to be embedded in a header to add additional security to a program. When the name is printed out by LOADING ...... only the first 16 characters are displayed. The header to a data file also contains the start/end bytes but these hold the start and end of TBUFFR itself.

The last operation on completion of a save or write is to store a duplicate header. If the command had a secondary address indicating an end of tape marker, then the ID would be changed to a 5 before writing. On loading or reading to the end of a file the last operation is to get back this trailing header (which remains until the next tape operation).

## Tape directories

Tape directories as such do not exist unless you are using an improved cassette operating system such as ACOS+. There are times when it is necessary to catalogue a tape. The process is time-consuming as it is, not surprisingly, directly proportional to the length of the tape. The following program may be used to do the job. It is best left running whilst you go away to do something else.

Any header will be read with an OPEN statement. ClOSEing it immediately ceases tape operation and program execution continues. The parameters are then pulled from the buffer and stored for later use. The process is repeated for the next header. When the end of tape is reached or you stop the program, a simple GOTO 260 will display the file information. This is the file type, up to 16 characters of its name with non-alphanumeric characters replaced by a ".", and if a program its start and end addresses (in hex).

```
100 DIM F$(50),FT$(50),SA$(50),EA$(50):C
B=828
110 PRINT"[CLS]PRESS PLAY ON TAPE"
120 IF PEEK(1)<>7 GOTO 120
130 I=I +1:OPEN 1:CLOSE 1:PRINTFक(I-1)
140 FT$(I)=RIGHT$("[5SPC]"+STR$(PEEK(CB)
),4)
150 IF PEEK(CB)=4 THEN SA$(I)="[2SPC]***
*":EAक (I)="[2SPC]****":GOTO 200
160 X=PEEK(830):GOSUB 360:SA$(I)=X $
170 X=PEEK(829):GOSUB 360:SA$(I)=" $"+SA
$(I) +X$
180 X=PEEK(832):GOSUB 360:EA$(I)=X$
190 X=PEEK(831):G0SUB 360:EA$(I)=" $"+EA
$(I)+又生
200 A$="":FOR J=833 T0 848
210 X=PEEK(J):IF X<32 OR X>95 THEN A 
+".":GOTO 230
220 A$=A$+CHR$(X)
230 NEXT J
240 Fक(I)=LEFT束(" "+Aक+"[18SPC]",17)
2 5 0 ~ G O T O ~ 1 3 0 ~
260 H$="[CLS]TYPE FILENAME[9SPC]START[35
PCJEND":PRINTH$
270 FOR J=1 TO I:PRINTFT$(J);F覀(J);SA$(J
);EA$(J)
280 IF INT(J/20)<>J/20 GOTO 320
290 PRINT"PRESS RETURN FOR NEXT PAGE"
300 GET A$:IF A$<>CHR$(13) GOT0 300
310 PRINTH$
3 2 0 ~ N E X T ~ J ~
330 INPUT "REUIEW AGAIN";Y直:IF Y$="Y" GO
TO 260
340 IF Y&<>"N" GOTO 330
350 CLOSE 1:END
360 X1=INT (X/16): X2=X-X1*16
370 X$=CHRक(X1+48-7*(X1>9))+CHRक(X2+48-7
*(X2>9))
380 RETURN
```

Unfortunately，during tape I／O the internal clock variable（TI\＄）is not updated as the interrupt is used exclusively for tape timing．Had this not been the case，a read of this variable could have been used to calculate the value of the tape counter．The best suggestion we can come up with is if the file is a program then the difference in its start and end addresses could be used to determine the loading time．For a
data file bytes could be taken until the status is set to the end of file, the number of bytes read being an indication of the time. We might as well do this as the tape is running anyway. The time taken may be used to work out an approximate counter reading.

## Auto-running

Generating programs which auto-run is also discussed in Chapter 10. There are many ways to accomplish this, most of which involve fairly detailed knowledge of the operating system. The following are suggestions only for you to pursue. All but one are suitable for disk or tape.

## The stack

During load the return address is placed on the stack. As this is an area of RAM, there is no reason why we cannot load through this area and put our own address on instead. This could then go to our own machine code routine. The file type should be 3 to ensure a load to its original position. The same would apply to disk or tape if loaded with a secondary address of 1 .

BASIC warm start - \$0302
After a load in direct mode BASIC is warm-started. Again as this vector is held in ram we can load through it. The new value it then contained could jump to our machine code or straight to RUN (for BASIC programs).

IRQTMP - \$029F
This stores the current IRQ vector during tape I/O which is restored after the tape operation. Again we can do the same to this as in the above. On the first normal interrupt the action will be taken. This, of course, can only be used with tape.

CHAIN command of the UTILITY
See Chapter 8.

## Screen

The utilities in this section are confined to the text screen.
The screen on the 64 is a 40 column by 25 line memory-mapped display. Chapter 3 and Appendices B to D of the Programmer's Reference Guide cover in great detail all aspects of the screen and it is to there that we refer you. All the following utilities assume that you are familiar with or know the following.
i) The screen may be moved from its default position.
ii) There are two character sets.
iii) The screen has an associated colour map at \$D800 on.
iv) The display codes differ from the ASCII codes.
v) Commodore 'ASCII' is not true ASCII which only ranges from $\emptyset$ to 127. (Consult your printer manual.)

## Printer dump

There are two routines, both of which output the current display in standard ASCII to a printer. One is a BASIC subroutine and the other is machine code. The second is noticeably faster than the first, as would be expected.

Both routines take account of whether the 64 is in upper or lower case mode as well as checking for the location of the screen.

64 owners with Commodore printers need not concern themselves with the conversions to standard ASCII.

## BASIC printer dump subroutine

The version given here is for an RS232 printer running at 300 baud without auto-line feed. For this reason the output logical file is assigned at the start of the program. The out put file is designated ' P ' to avoid specific reference to allow for easier change to other printers. The display is centred on an 80 column display by printing 20 spaces at the start of each line.

The program first examines the lower/upper case register at 53272 by calling the subroutine at 60090 . If in lower case, LC is assigned a value of 32 (note lower case ' $a$ ' in character set 2 has a PEEK value of 1 which is standard ASCII is 97 - that is, bit 5 set). This adjustment will be applied to all letters between ' $a$ ' and ' $z$ '. The whole dump is enclosed within two loops: I for the rows and J for the columns. All screen codes are ANDed with 127 to reduce them to values in the range 0 to 127 to eliminate reversed characters. If the screen code is $<32$, we have to add 64 and the LC adjustment. If it lies between 32 and 65 , we can print it unchanged. Only if in lower case mode do we need to check for upper case letters. If we were in upper case, these would be nonprinted graphic characters. If in LC then the anded code is already in standard ASCII. If all the tests have failed, we have a graphic character so we replace it by a space to maintain the layout of printable text. Once a screen line has been processed we print it preceded by 20 spaces and recycle for all remaining 24 lines.

```
10 OPEN 129,2,0,CHR悪(6)
60000 GOSUB 60090
60010 FOR I=0 TO 24:A$="":FOR J=0 T0 39:
CH=PEEK(S+I *40+J)
60020 CH=CH AND 127
60030 IF CH<32 THEN CH=CH+64+LC:GOTO 600
7 0
60040 IF CH<65 GOTO 60070
6 0 0 5 0 ~ I F ~ C H < 9 1 ~ A N D ~ L C ~ G O T O ~ 6 0 0 7 0
60060 CH=G
60070 A$=A$+CHRक(CH):NEXT J
60080 PRINT#P,SPC(SP);A$:NEXT I:CLOSE P:
RETURN
```

$60090 \mathrm{P}=129: \mathrm{SP}=20: \mathrm{G}=32$
60100 LC=0:IF PEEK (53272) $=23$ THEN LC=32
60110 S=PEEK (648) *256:RETURN
Whenever a dump is required, simply GOSUB 60000. This could be actioned by, say, a GET statement, but should not add to the display, or if it does then only 24 lines should be printed. To improve the presentation, blank lines or a form feed should be issued at the end of the dump.

## Machine code printer dump

The logic of this routine is identical to that above and is therefore not described in detail. The differences are that it is much faster and it does not pad a line with 20 leading spaces.

The routine as written assumes logical file 2 is open to the printer at the time of calling. To change this, simply alter the byte at $\$ \mathrm{C} 001$ with a POKE. It works by changing the output device through the ChKOUT KERNAL call to that associated with file \#2 (the equivalent of a CMD from BASIC). This then allows us to use the KERNAL routine Chrout to output the data. There is a routine in ROM which could be used to do most of the conversion, but for this exercise the technique used here is adequate and easier to follow. The device need not be the printer and could be the disk or tape depending on the OPEN statement. We do not recommend you use this routine with anything other than a printer as far better screen saves follow. Once the dump is complete, the default device for output is restored to the screen before returning to basic.

The routine is used by at some point including an OPEN 2,4 or OPEN $2,2, \mathrm{CHR} \$()$ if using RS232. A simple sys 49152 will perform the dump. If your printer requires a forced line feed, make the necessary adjustment to \$C001 for a value greater than 127.

## BASIC loader for the machine code

The following must be loaded and run. Once this has been done the code remains present until overwritten by something else. Once run the machine code may be saved using an extended monitor for ease of loading later.
1 DATA $162,2,32,201,255,173,136,2$
$, 133,88,169$
2 DATA $6,133,87,173,24,208,201,21$
$, 208,6$
3 DATA $169,0,133,89,240,4,169,32$,
133,89
4 DATA $169,32,133,90,24,165,88,10$
$5,3,133$
5 DATA $91,162,4,160,0,177,87,41$,
127,24

6 DATA 201, 31, 176, 7, 24, 105, 64, 101 89, 144
7 DATA 24, 24, 201, 64, 176, 2, 144, 17, 24, 165
8 DATA 89, 240, 10, 177, 87, 24, 201, 91 176, 3
9 DATA $24,144,2,169,32,32,210,255$ 224, 1
10 DATA 208, 4, 192, 232, 240, 23, 230, 96, 201, 40
11 DATA 208, $9,169,13,32,210,255,1$ 69, 0, 133
12 DATA 96, 200, 208, 187, 230, 88, 202, 208, 182, 169
13 DATA $13,32,210,255,162,0,32,20$
1, 255, 96
14 DATA 0, 255, 255, 0, 0, 255, 255, 0, 0, 255
15 FOR I=49152 TO 49292:READ A:POKE I,A: NEXT I

Here is the assembly listing which is fully annotated to allow you to follow it:

| C000 | A202 | LDX | \#\$02 | log file to printer |
| :---: | :---: | :---: | :---: | :---: |
| C000 | 20C9FF | JSR | \$FFC9 | perform CMD2 via CHKOUT |
| C005 | AD8802 | LDA | \$0288 | screen start from HIBASE |
| C008 | 8558 | STA | \$58 | set start registers |
| COAA | A900 | LDA | \#\$00 |  |
| C.OOC | 8557 | STA | \$57 |  |
| C00E | AD1800 | LDA | \$D018 | check upper/lower case |
| C011 | C915 | CMP | \#\$15 | is it upper |
| C013 | D006 | BNE | \$C01B | no |
| C015 | A900 | LDA | \# ${ }^{\text {6 }} 8$ | set adjustment value |
| C017 | 8559 | STA | \$59 | for ASCII |
| C019 | F084 | BEQ | \$C01F | skip lower case |
| C01B | A920 | LDA | \#\$ 20 | lower case set ad.j flag |
| C010 | 8559 | STA | \$59 | as ASCII $a=97$ etc. |
| C01F | A920 | LDA | \#\$ 20 | set non-printable flag |
| C021 | 855A | STA | \$5A | to a space |
| C023 | 18 | CLC |  |  |
| C024 | A558 | LDA | \$58 | set MSB end of screen |
| C026 | 6903 | ADC | \#\$03 |  |
| C028 | 855B | STA | \$5B |  |
| C02A | A204 | LDX | \#\$04 | almost 4 pages/screen |
| co2C | A000 | LDY | \#\$08 | counter within page |


| C02E | B157 | LDA, | (\$57) , Y | get byte |
| :---: | :---: | :---: | :---: | :---: |
| C030 | 297 F | AND | \# ${ }^{\text {a }} 7 \mathrm{~F}$ | eliminate high bit 7 |
| Remember difference between screen and ASCII codes C032 18 CLC start checks |  |  |  |  |
|  |  |  |  |  |
| C033 | C91F | CMP | \# ${ }^{\text {1 }}$ F | less than a space |
| C035 | B007 | BCS | \$C03E | no go to next check |
| C037 | 18 | CLC |  |  |
| C038 | 6940 | ADC | \# $\$ 40$ | make ASCII by adding 64 |
| C03A | 6559 | ADC. | \$59 | add lower case adj. |
| C03C | 9818 | BCC | \$C056 | always taken |
| C03E | 18 | CLC |  |  |
| C03F | C940 | CMP | \#\$40 | check for upper case in |
| C041 | B002 | BCS | \$C045 | 1/c mode \& branch $) 65$ |
| C043 | 9011 | BCC | \$C056 | !-? same in both sets |
| C045 | 18 | CLC |  | check upper case |
| C046 | A559 | LDA | \$59 | if zero |
| C048 | F00A | BEQ | + C 054 | branch to avoid graphic |
| C04A | B157 | LDA | (\$57) , Y | get $1 / C$ byte again |
| C04C | 18 | CLC |  | check not gt $Z$ |
| C040 | C95B | CMP | \#\$5B |  |
| C84F | B003 | BCS |  | if so avoid graphic |
| C051 | 18 | CLC |  |  |
| C052 | 9002 | BCC | \$ C 056 | valid $A-Z$ so skip space |
| C054 | A920 | LDA | \# ${ }^{\text {20 }}$ |  |
| C056 | 20D2FF | JSR | \$FFD2 | print char |
| C059 | E001 | CPX | \#\$01 | on last prage |
| C05B | D004 | BNE | \$C061 | no - so branch |
| C05D | C6E8 | CPY | \#\$E8 | yes so check end $\$ * * E 8$ |
| C05F | F017 | BEQ | \$C078 | branch all done |
| C061 | E660 | INC | \$68 | end of screen line reg |
| C063 | C928 | CMP | \#\$28 | is it 40 dec |
| c065 | D009 | BNE | \$C070 | no so skip next bit |
| C067 | A900 | LDA | \#\$00 | output next bit |
| C069 | 20D2FF | JSR | \$FFD2 | print it |
| C06C | A900 | LDA | \#\$00 | rezero end of line reg |
| CB6E | 8560 | STA | \$60 |  |
| C070 | C8 | INY |  | continue current page |
| C071 | D9BB | BNE | \$C02E | branch if not finished |
| C073 | E658 | INC. | \$58 | inc next page register |
| C075 | CA | DEX |  |  |
| C076 | D0B6 | BNE | \$C02E | always taken |
| C078 | A900 | LDA | \#\$0D | RETURN for last line |
| C07A | 20D2FF | JSR | \$FFD2 |  |
| C07D | A200 | LDX | \#\$00 | restore screen output |
| C07F | 20C9FF | JSR | \$FFC9 |  |
| C082 | 60 | RTS |  |  |

To improve this, why not patch into the interrupt routine to, for example, dump the screen whenever a designated key is pressed rather than using the sYS command? Chapter 4 explains the interrupt in detail and Chapter 10 gives an example of its use. If you decide to do this, remember to include a routine to disable the patch. The necessary enable and disable routines can be added at the end of the code as given. The logical file will still have to be OPENED unless the appropriate KERNAL routines are called.

## Screen dumps

Three ways are given to save the screen and its associated colour map in this section. Two are in BASIC and the third is in machine code. Both BASIC programs use a sequential file to store the data, but differ in the length of file produced. The machine code saves the screen as a program file and is the most economical and by far the quickest.

A few points should be made before discussing the routines in detail. Any area of memory may be saved from BASIC by setting TXTTAB and VARTAB to its start and end addresses. The problem is that once we have changed these pointers we have temporarily lost our program. Another problem is that a LOAD will cause BASIC to warm-start, which is this case will be at the newly set TXTTAB address. The screen is an area of memory and may be loaded and saved in this way. Unfortunately, its default position is below the normal start of BASIC so a 'crash' or 'hang-up' is usually the result. Try it and see. So from a practical viewpoint we must resort to other means.

All the following routines check HIBASE for the current screen location. The resulting screens will always reload to the current screen position regardless of its location at the time of saving. The reloaded screen will be identical to that saved in both characters and colours.

Screen save using numbers
This routine firstly Peeks out the border and background colours and writes them, as numbers, to a disk file (change the OPEN command for tape). It then proceeds, writing alternate screen and colour values until finished.

To save a screen: GOSUB 60000
To load a screen: GOSUB 60050

```
60000 OPEN 2,8,2,"20:TEST,S,W"
60010 S=PEEK(648)*256:C=55296
60020 PRINT#2,PEEK(53280);",";PEEK(53281
);CHRक(13);
60030 FOR I=S TO S+999:PRINT#2,PEEK(I);"
,";PEEK(C+I-S);CHR$(13);:NEXT I :CLOSE 2
```

```
60040 RETURN
60050 OFEN 3,8,3,"TEST,S,R"
60060 INPUT#3,A,B:POKE 53280,A:POKE 5328
1,B
60070 S=PEEK(648)*256:C=55296
60080 FOR I=S TO S+999:INPUT#3,A,B:POKE
I,A:POKE C+I-S,B:NEXT I:CLOSE 3:RETURN
```

Because numbers are written as their ASCII codes three to five bytes are used for each value (spaceXXXreturn). Therefore, using this method we will generate a sequential file of between 6 and 10K, which seems rather excessive. The second method reduces the size of this file.

Screen save using characters
This time a single byte is used to store each value in the screen and colour maps. This is done by simply PeEking the value and generating the corresponding $\operatorname{CHRS}()$ character with the $\operatorname{ASC}()$ function. Zero values must be trapped as $\operatorname{ASC}(6)$ will give a syntax error. The resulting file uses only one byte for most values and the file size is therefore about 2 K . This is obviously far faster to generate and restore.

To save a screen: GOSUB 60000
To load a screen: GOSUB 60050

```
60000 OPEN 2,8,2,"a0:TEST,5,W"
60010 S=PEEK(648)*256:C=55296
60020 PRINT#2,CHR$(PEEK(53280));CHRक(PEE
K(53281));
60030 FOR I=S TO S+999:PRINT#2,CHRक(PEEK
(I));CHR#(PEEK(C+I-S));:NEXT I:CLOSE 2
60040 RETURN
60050 OPEN 3,8,3,"TEST,S,R"
60060 GET#3,A$:IF A$="" THEN A$=CHRक(0)
60070 POKE 53280,ASC(A$)
60080 GET#3,A$:IF A$="" THEN A$=CHR$(0)
60090 POKE 53281,ASC(A$)
60100 S=PEEK(648)*256:C=55296
60110 FOR I=S TO S+999:GET#3,A末:IF A$=""
    THEN A$=CHRक(0)
60120 POKE I,ASC(A$):GET#3,A$:IF A$="" T
HEN A$=CHRक(0)
60130 POKE C+I-S,ASC(A$):NEXT I:CLOSE 3:
RETURN
```


## Machine code screen save

This is by far the best method. It is very simple to use the kernal load and SAVE for both the screen and colour maps. Using these as they stand, two files would be generated - one for the colour map and one for the screen. This is no hardship, but a relocated load would be required if a screen is being restored to a different location from whence it came. This is not difficult, but perhaps is not the best way.

We have approached the problem slightly differently. Before performing the save, the screen and colour maps are combined into a 2 K block at a convenient address. This has to be out of the way of BASIC to avoid corrupting program or data areas. This could be a reserved area at the end of BASIC or even under BASIC ROM if a switch like that used in the utility is implemented to throw out and restore rom. This is possible as no BASIC ROM calls are made. For this example we have chosen to move the screen from its current position to $\$ C 400$ and the colour map to \$C800. The routine also saves the sprite pointers and if you do not wish it to do so then you will have to modify the code to move 1000 rather than its current 1024 bytes from or to each area.

All the routine does to save is to move both screen and colour maps then use the KERNAL SAVE from \$C400 to \$CC00. To restore the screen it is reloaded to $\$ C 400$ and moved back to the colour map and the current screen position.

BASIC loader for screen save
The following must be loaded and run before it can be called. Once run it may be saved using an extended monitor for ease of loading later.

```
1 DATA 32, 253, 174, 201, 76, 208, 6, 16
7, 0, 133, 87
2 DATA 240, 11, 201, 83, 240, 3, 32, 8,
175,169
3 DATA 255, 133, 87, 32, 115, 0, 32, 253
    174, 201
4 DATA 34, 240, 3, 32, 8, 175, 32, 115,
0, 165
5 DATA 122, 133,187, 165,123,133,188
, 160, 0, 177
6 DATA 122, 201, 34, 240, 8, 200, 192, 1
9, 208, 245
7 DATA 32, 8, 175, 132, 183, 152, 24, 10
1, 122, 133
8 DATA 122, 144, 2, 230, 123, 32, 115,0
    32, 253
7 DATA 174, 144, 3, 32, 8, 175, 56, 233,
    48,133
```

```
10 DATA 186, 169, 1, 133, 184, 133, 185,
    169, 0, 133
11 DATA 88, 133, 90, 133, 92, 133, 94, 1
73, 136, 2
12 DATA 133, 89, 169, 196, 133, 91, 169,
    216, 133,93
13 DATA 169, 200, 133, 95, 165, 87, 240,
    43, 160,0
14 DATA 162, 4, 177, 88, 145, 90, 177,9
2, 145,94
15 DATA 200, 208, 245, 230, 89, 230, 91,
    230, 93, 230
16 DATA 95, 202, 208, 234, 166, 94, 164,
    95, 169, 196
17 DATA 133, 91, 169, 90, 32, 216, 255,
32, 115,0
18 DATA 96, 165, 87, 32, 213, 255, 160,
0, 162, 4
19 DATA 177, 90, 145, 88, 177, 94, 145,
92, 200, 208
20 DATA 245, 230, 89, 230, 91, 230, 93,
230, 95, 202
21 DATA 208, 234, 32, 115, 0, 96, 0, 255
255, 0
22 FOR I=49152 TO 49362:READ A:POKE I,A:
NEXT I
nb"basic2"
```

To save a screen: SYS 49152,S, "filename",DEVICE
SYS 49152,S, "@0:filename",DEVICE to replace on disk
To load a screen: SYS 49152,L,"filename", DEVICE

All parameters are required and an illegal or missing parameter will produce a SYNTAX ERROR. A file name or a minimum of "'" is required, even with cassette. Remember all bytes following a sys command are ignored.

The following is the assembly listing, which should be selfexplanatory. The first part of the routine is our own version of 'GET parameters' and is a useful technique when passing parameters to a routine enabled with a SYS. CHRGET (see Chapter 3) is used to gather the necessary bytes.

## ASSEMBLY LISTING

Set up the parameters - common to both LOAD and SAVE C000 20FDAE JSR \$AEFD check for comma
C083 C94C CMP \#\$4C is next char L for LOAD
C005 D006 BNE \$C00D no then test for SAUE
C007 A900 LDA \#\$80 set flag at 57
C009 8557 STA $\$ 57$ to zero for later use
C00B F00B BEQ $\$ C 018$ always taken if LOAD
COOD C953 CMP \#\$53 is it 5 for SAVE ?

| $C 00 F$ | F003 | BEQ | $\$ C 014$ | if so continue |
| :--- | :--- | :--- | :--- | :--- |
| C011 | $2008 A F$ | JSR | $\$ A F 08$ | not $L$ or $S=S Y N T A X ~ E R R O R ~$ |

C014 A9FF LDA \#कFF set flag to FF far save
C016 8557 STA $\$ 57$

| C018 | 207308 | JSR | \$0073 | inc CHRGET (see ch.4) |
| :---: | :---: | :---: | :---: | :---: |
| C01B | 20FDAE | JSR | \$AEFD | next comma |
| C01E | C922 | CMP | \# 222 $^{\text {2 }}$ | opening quote of name |
| C020 | F003 | BEQ | 事C025 | OK so continue |
| C022 | 20884F | JSR | \$AF08 | no quote so SYNTAX ERROR |
| C025 | 207300 | JSR | \$0073 | inc CHRGET to name |
| C028 | A57A | LDA | \$7A | set FNADR low byte |
| C02A | 85BB | STA | \$BB |  |
| C02C | A57B | LDA | \$7B | do same for high |
| C02E | 85BC | STA | \$BC |  |
| C038 | A080 | LDY | \#\$80 | find length of riame |
| C032 | B17A | LDA | ( $\$ 7 \mathrm{~A}$ ) , Y | by searching for closing |
| C034 | C922 | CMP | \# $\$ 22$ | quote |

C036 F008 BEQ $\ddagger$ C040 found it so exit
C038 C8 INY
C039 C013
C03B D0F5
C03D 2008AF
CPY
BNE \$C032
JSR \#AF08
C040 84B7 STY \$B7
C042 98 TYA
C043 18 CLC

| C044 | $657 A$ | ADC $\$ 7 A$ | set CHRGET to end quote |  |
| :--- | :--- | :--- | :--- | :--- |
| C046 | $857 A$ | STA $\$ 7 A$ | low byte |  |
| C048 | 9002 | BCC $\$ C 04 C$ |  |  |
| C04A | E67B | INC $\$ 7 B$ | inc high if page crossed |  |
| C04C | 207300 | JSR $\$ 0073$ | get next byte |  |
| C04F | $20 F D A E$ | JSR | $\$ A E F D$ | comma ? |
| C052 | 9003 | BCC | $\$ C 057$ | OK |
| C054 | $2008 A F$ | JSR | $\$ A F 08$ | no then SYNTAX ERROR |
| C057 | 38 | SEC |  |  |
| C058 | E930 | SBC | $\# \$ 30$ | make byte a number 0-9 |
| C05A | $85 B A$ | STA $\$ B A$ | store current device-FA |  |
| C05C | A901 | LDA $\# \$ 01$ | set secondary add to 1 |  |


| C05E | 85B8 | STA | \$ ${ }^{\text {8 }}$ | and store at SA |
| :---: | :---: | :---: | :---: | :---: |
| C060 | 8589 | STA | \$ $\mathrm{B}^{\prime}$ | same for logical file-LA |
| Set | up pointers | to | e used in | the move |
| c0o2 | A900 | LDA | \#\$80 |  |
| C064 | 8558 | STA | \$58 |  |
| C066 | 855A | STA | \$5A |  |
| C068 | 855C | STA | \$5C |  |
| C06A | 855E | STA | \$5E |  |
| C06C | AD8802 | LDA | \$0288 | find the current screen |
| C0SF | 8559 | STA | \$59 | start from HIBASE |
| C071 | A9C4 | LDA | \#\$ C 4 |  |
| C073 | 855B | STA | \$5B |  |
| C075 | A908 | LDA | \#\$D8 |  |
| C077 | 855D | STA | \$50 |  |
| C079 | A9C8 | LDA | \#\$ C 8 |  |
| C07B | 855F | STA | \$5F |  |
| C07D | A557 | LDA | \$57 | 0 for LOAD |
| C67F | F02B | BEQ | \$C0AC | do LOAD |
| Move | screen and | calo | ur ta one | block and perform SAVE |
| C081 | A008 | LDY | \#\$00 | SAVE use Y within page |
| C083 | A204 | LDX | \#\$04 | and $X$ for page counter |
| C085 | B158 | LDA | (\$58), | read byte from screen |
| C087 | 915A | STA | (\$54) , Y | store at combined area |
| C089 | B15C | LDA | (\$5C), Y | get char colaur |
| C08B | 915 E | STA | (\$5E), Y | store in comb area+\$0400 |
| C08D | C8 | INY |  |  |
| C08E | D8F5 | BNE | \$C085 | cycle for one page |
| C090 | E659 | INC | \$59 | inc all high bytes |
| C092 | E65B | INC | \$5B |  |
| C094 | E65D | INC | \$50 |  |
| C096 | E65F | INC | \$5F |  |
| C098 | CA | DEX |  | dec $X$ and |
| C099 | D0EA | BNE | \$C085 | repeat till 4 pages done |
| C09B | A65E | LDX | \$5E | $X$ halds low end of save |
| C090 | A45F | LDY | \$5F | $Y$ the high byte |
| C09F | AGC4 | LDA | \#\$C.4 | use 5A/5B on zero page |
| C0A1 | 855B | STA | \$5B | for start of SAVE |
| COA3 | A95A | LDA | \#\$5A | A must hold offset 5 A |
| C0A5 | 2008FF | JSR | \$FFD8 | do SANE |
| C0A8 | 207300 | JSR | \$0073 | must inc CHRGET before |
| C0AB | 60 | RTS |  | returning to BASIC |
| Perfor | arm LOAD and | depl | it black | into char and colour maps |
| COAC | A557 | LDA | \$57 | read flag - $A=0$ for LOAD |
| COAE | 2005FF | JSR | \$FFD5 | do LOAD |
| COB1 | A000 | LDY | \#\$00 |  |
| C0B3 | A204 | LDX | \#\$84 |  |


| C0B5 | B15A | LDA | (\$5A), Y | reverse of SAVE |
| :---: | :---: | :---: | :---: | :---: |
| C0B7 | 9158 | STA | (\$58), Y | even if the screen |
| C0B9 | B15E | LDA | (\$5E), Y | was at a different |
| C0BB | 915C | STA | (\$5C), Y | location at the time of |
| C0BD | C8 | INY |  | SAUE it will go to the |
| C0BE | D0F5 | BNE | \$C0B5 | current pasition |
| C0C0 | E659 | INC | \$59 |  |
| C0C2 | E65B | INC | \$5B |  |
| C0C4 | E65D | INC | \$5D |  |
| C0C6 | E65F | INC | \$5F |  |
| C0C8 | CA | DEX |  |  |
| C0C9 | D0EA | BNE | \$C0B5 |  |
| C0CB | 207300 | JSR | \$0073 | must inc CHRGET before |
| CaCE | 60 | RTS |  | returning to BASIC |
| Disk |  |  |  |  |

This section deals with the 1541 disk drive though much is directly applicable to the 3040 and 4040 units.

The manual supplied with the 1541 contains all the information that most users will require. Perhaps the most difficult to master are the direct access programming commands such as BLOCK-READ, and so on. There is only one way to become proficient in their use and that is to experiment. When experimenting we suggest you use a disk containing unwanted information as disasters can happen.

We supply only one utility in this section which we like to think of as an expandable disk utility. Once direct access programming is mastered, there are all sorts of fun things you can do. To use it to its best advantage you have to know something of how the disk operates and how information is stored. To this end we give below a very short introduction and would refer you to the 1541 manual itself.

## Introduction

The 1541 is a self-contained, intelligent device. It has two processors and its disk operating system (DOS) in ROM along with an area of rAM used for input/output (buffering) operations. This differs from, say, the BBC Micro where the interface is within the micro itself and, depending on the type of interface, removes RAM from the user area. The disadvantage to this self-contained arrangement is that you cannot use non-Commodore units (there are one or two now available) as most disk manufacturers do not supply a suitable controller and DOS.

Almost the whole of the disk's capacity can be used to store data except for one or two reserved areas. The disk is divided into tracks which are further subdivided into a varying number of sectors. Tracks are numbered from 1 through to 35 whereas sectors start at zero. The following is the arrangement for both the 1541 and 4040 units (see Table 6.1, 1541 manual):

| TRACK | SECTOR | TRACK | SECTOR |
| :---: | :---: | :---: | :---: |
| $1-17$ | $\emptyset-20$ | $25-30$ | $\emptyset-17$ |
| $18-24$ | $\emptyset-18$ | $31-35$ | $\emptyset-16$ |

In order to know where to find information, the disk uses an index or directory track. This is track 18 and it has two special areas. The first is track 18 sector 0 and is the Block Availability Map (BAM). This keeps a record of all sectors in use unless direct access programming operations have been used without a block-allocate command. If this is the case then the information is there, but can be overwritten as it is empty as far as DOS is concerned. The entries in BAM are made up as follows (see Table 5.1 of manual):

## BYTES CONTENTS

000-001 Pointer to start of directory 18/01
002 Holds an 'A' for 4040 format
004-143 Four bytes for each of the 35 tracks indicating whether in use $1=$ free $0=$ allocated
144-161 Disk name plus shifted spaces to make 16 in total
162-163 Disk ID
165-166 Disk version of 2A

Each track uses four bytes in BAM. The first stores the number of free sectors on a track and is used in computing blocks free. The remaining three are used to indicate whether a particular sector is allocated (bit set low and one bit per sector). As the maximum number of sectors is 21, not all bits are used. The following is a dump of BAM from which you can pick out the information given above. All values are given in hex with the byte position within the sector given first followed by this byte's value and the next seven and at the end of the line the equivalent ASCII characters (if printable):

| START |  | TES |  |  |  |  |  |  | ASCII |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 00 | 12 | 01 | 41 | 00 | 15 | ff | $f f$ | 1 f | ..a. |
| 08 | 15 | $f f$ | $f f$ | $1 f$ | 15 | $f f$ | $f f$ | $1 f$ |  |
| 10 | 15 | ff | $f f$ | 1 f | 15 | ff | $f f$ | 1 f |  |
| 18 | 15 | $f f$ | $f f$ | $1 f$ | 15 | ff | $f f$ | $1 f$ |  |
| 20 | 15 | $f f$ | $f f$ | 1 f | 15 | ff | $f f$ | 1 f |  |
| 28 | 15 | $f f$ | $f f$ | 1 f | 00 | 00 | 00 | 00 |  |
| 30 | 00 | 00 | 00 | 00 | 00 | 00 | 00 | 00 |  |
| 90 | 42 | 4f | 4 f | 4b | 20 | 50 | 52 | $4 f$ | book pro |
| 98 | 47 | 52 | 41 | 4d | 53 | a0 | a0 | a0 | grams |
| a 0 | a0 | a0 | 31 | 31 | ab | 32 | 41 | a0 | 11 2a |

The file information starts on track 18 sector 1 and can continue throughout the remainder of the track. Each file uses 32 bytes. Therefore, one sector can hold information for eight entries. With a possible 20 sectors available, information could be held for 160 files. This is unlikely to happen as each file would have to be less then 1 K . The directory format is such that bytes $\emptyset-31$ hold file 1, 32-63 hold file 2, and so on. Each entry is divided up as follows (see Table 5.3 of the 1541 manual):

## BYTE CONTENTS

000-001 Next directory track and sector. A track of 0 indicates last sector in use. These bytes only used for the first entry.
002 The type of file $\$ 00=$ scratched or not in use.
$\$ 80=$ Deleted (scratched unclosed) $\$ 81=$ sEQuential
$\$ 82=$ Program $\$ 83=$ USer $\$ 84=$ Relative $\$ 1-4=$ unclosed
003-004 Starting track and sector of file
005-020 Name padded with shifted spaces
023 Record size of relative file
028-029 New track and sector for disk ops with replacement - @
030-031 Number of blocks file uses in low/high byte format
Below is a typical dump of the first directory sector, track 18 sector 1 , for two file entries:

| START | CONTENTS |  |  |  |  |  |  |  | ASC |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 00 | 12 | 04 | 82 | 11 | 00 | 44 | 55 | 4d | .....dum |
| 08 | 50 | 2e | 4d | 49 | 4b | a0 | a0 | a0 | p.mik |
| 10 | a 0 | a0 | a0 | ab | a0 | 00 | 00 | 00 |  |
| 18 | 00 | 00 | 00 | 00 | 00 | 00 | 0d | 00 |  |
| 20 | 00 | 00 | 82 | 11 | 03 | 44 | 55 | 4d | .....dum |
| 28 | 50 | 2 e | 4d | a0 | a 0 | a0 | a 0 | ab | p.m |
| 30 | a0 | a0 | a0 | a0 | a 0 | 00 | 00 | 00 |  |
| 38 | 00 | 00 | 00 | 00 | 00 | 00 | 01 | 00 |  |

It is worth noting that directory sectors do not follow sequentially. The same is true for file storage, as can be seen when using the disk utility.

Just to round things off, here is a dump of a BASIC program which occupies less than one block. It is in fact the loader for the UTility at the end of Chapter 9.


As it is less than one block, the linking track is zero, denoting the end. It is a straight copy of the ram and like the memory dump in Chapter 1, we can pick out the link addresses, end of program, and so on.

## Disk utility

This utility offers many of the housekeeping commands and provides a number of more interesting options. It is rather long as most of the subroutines are complete in themselves (to allow you to extract only those you want for your own programs). The listing has been left in lower case and when you are typing it in, it is easiest to put the 64 into that mode with the shift and logo keys. It makes extensive use of direct access programming so we suggest you use the information given above and the relevant sections of the 1541 manual to follow it. It has been run through the UTILITY's CODER command to produce the mnemonics. Most annotated characters are cursor moves, colours or simply capital letters.

The usual options of New, validate, SCRATCH, INITIALIZE, RENAME, COPY (within a drive) and READ DISK ERROR are all present. The directory option is unusual in that everything is input or displayed in hex notation. A much shorter way to get a directory is given in the Backup utility at the end of this chapter. The option also displays the first track and sector of a file, and if it is a program, also its load address. The listing is further split up into directory sectors and will display even scratched or Deleted files if the disk has not been validated. Two values are given for the blocks free - the usual value exclusive of erased files and another inclusive of erased files. An erased file simply has its associated BAM set to 1 (not allocated).

The trace option follows a file through displaying its associated tracks and sectors. It will also check to see if the file it is following is scratched. If this is the case, it will ask whether you wish to recover it. If your answer is 'yes', then as it traces it will also allocate each block. Providing that all the blocks found were free it has recovered the file. If an allocated block is found then the original area of the file has been
overwritten and recovery is not possible and you will be told. If the scratched file has been successfully traced, all that remains to be done is to use the MOD/DISP BLOCKS option to change its 'file type' byte (third byte in its entry) from 0 back to $\$ 81$ or $\$ 82$. The revised directory block must be rewritten to complete the process.

The MOD/DISP BLOCKS option is similar to the demonstration disk's program 'DISPLAY TRACK AND SECTOR'. The main difference is that it also allows the block to be modified and rewritten to disk. When the block is written it is also allocated. The usual options to review again and get the next track and sector are available. The subroutine called at 1680 is a little unusual and merits some explanation. Earlier in the chapter under the heading 'Keyboard' it was mentioned that a file could be opened to it. This eliminates the '?' prompt and also releases the cursor. The cursor may then be moved to the appropriate line, hex values changed and on pressing RETURN the revised values are processed and written to the disk buffer. The same 128 bytes are then redisplayed by reading from the buffer. At the end of a block the option to write the changes must be taken to change it on disk. It is also possible to recover files using this option by following a file through taking the next track and sector option (first two bytes) and always writing the changes. Unlike TRACE, it does not check to see if a file can be recovered. Files in which a READ ERROR occurs may also be reconstructed. This we discovered when the EASY SCRIPT appendices file of this book was corrupted. All we did was modify the next track and sector bytes of the preceding block to skip the corrupt block. The resulting file could be read with the loss of only 256 bytes (and was immediately saved on another disk).

The APPEND for program files is the same as that in Chapter 5 (where it is fully explained). The APPEND for sequential files (and SCRATCH) builds the command string (separated by commas) before actioning on RETURN with no input.

The final option is to modify the disk's header. This is done by simply reading BAM, moving the buffer pointer and writing in the new values. It is worth noting that whenever a byte is read, the buffer pointer is moved forward one position. So in order to write to the same position at which the read started, the pointer must be set using a ' $\overline{-}-\overline{ }{ }^{\prime}$ command.

The utility is not foolproof, but with a little attention to detail, may be used to advantage. Our last comment before the listing is to point out that when you try to allocate an already allocated block error 65 NO block occurs. This must be checked for and trapped as in MOD/DISP blocks. The locations of all the subroutines may be read from the IF statements in lines 210-340.

100 poke 53272，23：poke 53280，6：dim a\＄（10 B） t （ $\$$（5）
110 for $i=0$ to 5：read t覀（i）：rext i
120 data del，seq，prg，usr，rel，？？？
130 print＂［cls］［g＞d］［g＞i］［g＞s］［g＞k］［g＞u ］［g＞t］［g＞i］［g＞l］［g＞i］［g＞t］［g＞y］＂：print＂［ cd］［rev］［g＞n］［off］new disk＂tab（20）；＂［re v］［gンh］［off］change header＂
140 print＂［cd］［rev］［g＞v］［aff］validate d isk＂；tab（20）；＂［rev］［g＞d］［off］directory＂ 150 print＂［cd］［rev］［g＞t］［off］trace fil e＂；tab（20）；＂［rev］［g＞s］［off］scratch file （s）＂
160 print＂［cd］［reu］［g＞r］［off］rename fil e＂；tab（20）；＂［reu］［g＞e］［off］read disk er ror＂
170 print＂［cd］［rev］［g＞c］［off］copy file＂ ；tab（20）；＂［rev］［g＞a］［off］append files＂
180 print＂［cd］［rev］［g＞b］［off］backup fil
e＂；tab（20）；＂［rev］［g＞m］［off］mod／disp blo cks＂
190 print＂［cd］［rev］［g＞i］［off］initialize disk＂；tab（20）；＂［reu］［g＞x］［off］exit＂
200 gosub 2360

| 210 | ¢車＝＂n＂ | then gosub | 530：goto 130 |
| :---: | :---: | :---: | :---: |
| 220 | if f ¢＝＂u＂ | then gosub | 580：goto 130 |
| 230 |  | then gosub | 620：90to 130 |
| 240 |  | then gosub | 670：goto 130 |
| 250 | if $\mathrm{f}^{\text {車＝＂e＂}}$ | then gosut | 740：90to 130 |
| 260 |  | then gosub | 360：goto 130 |
| 270 | if $\gamma$ ¢ $=$＂c＂ | then gosub | 840：goto 130 |
| 280 | if yta＝＂d＂ | then gosut | 890：goto 130 |
| 290 | if $\mathrm{f}^{\text {¢ }}=$＂$h "$ | then gosub | 1240：goto 130 |
| 300 | if f 車＝＂x＂ | then end |  |
| 310 | if $\gamma$ ¢ $=$＂m＂ | then gosub | 1400：90to 130 |
| 320 |  | then gosub | 1740：goto 130 |
| 330 | if $\mathrm{f}_{\text {为＝＂t＂}}$ | then gosub | 1760：goto 130 |
| 340 |  | then gosub | $2130: g o t o ~ 130$ |
| 350 |  |  |  |

360 rem append
370 print＂［cls］［g＞a］［2g＞p］［g＞e］［g＞n］［g＞d ］［cd］＂：print＂［rev］［g＞p］［off］prg files［c d］＂：print＂［cd］［rev］［g＞s］［off］seq files＂ 380 gosub 2360：if $y^{\$}\langle \rangle$＂p＂and $\gamma=\langle\langle \rangle$＂s＂go to 520
390 if $\gamma \pm=" s^{\prime \prime}$ then gosub 770：return

400 rem prg files
410 print＂［cd］append prg files－sure＂：g osub 2360：if y倞〈＂y＂goto 520
420 input＂［cd］combined prg＂；fक：input＂［3s pc］first prg＂；$\times 1$ it：input＂［2spc］second prg ＂；$\times 2$ 生
430 open 3，8，3，＂0：＂＋ft＋＂，p，w＂：open 2，8，2 ，＂ $0:$＂$+\times 1 \ddagger+", p, r^{\prime \prime}$
440 get\＃2，$\gamma$ 丰
$450 \times \$=y$ 本：get\＃2，y事：if stく＞0 goto 470
460 gosub 2330：print\＃3，x ${ }^{(1)}$ ；goto 450
470 close 2：open 2， $8,2, " \theta: "+\times 2 \$+", p, r^{\prime \prime}$

490 get\＃2，x $=$ ：if st＜＞0 goto 510

510 close 2：print\＃3，chr\＄（0）；：close 3
520 return
530 rem new disk
540 print＂［cd］new disk－sure＂：gosub 236
0：if $\gamma$ ¢ $\left\rangle^{*} y^{\prime \prime}\right.$ goto 570

＝left事（f央，16）＋＂，＂＋left
560 open 15，8，15，＂n0：＂＋f象：clase 15
570 return
580 rem validate
590 print＂［cd］validate－sure＂：gosub 236
0：if $\gamma$ 事〈〉＂y＂goto 610
600 open 15，8，15，＂u＂：close 15
610 return
620 rem rename
630 print＂［cd］rename－sure＂：gosub 2360：
if yolく〉＂y＂goto 660
640 input＂［cd］old file＂；fos：input＂new fil

650 open $15,8,15, f=$ ：close 15
660 return
670 rem scratch file
680 print＂［cd］scratch－sure＂：gosub 2360 ：if $\gamma$ 卉〈〉＂y＂goto 730
690 f\＄＝＂n：print＂［cd］use＊or ？for patte rn matching＂
700 print＂hit return to delete［cd］＂


 close 15

730 return
740 rem error

close 15：print＂［cd］［g＞e］［2g＞r］［g＞o］［g＞r］

760 gosub 2360：return
770 rem append seq files
780 print＂［cd］append seq files－sure＂：9
osub 2360：if y倞〈＂＞＂goto 830
790 print＂［cd］hit return to append［cd］＂



 ，15：print\＃15，f
830 close 15：return
840 rem copy same disk
850 print＂［cd］copy－sure＂：gosub 2360：if y\＄〈〉＂y＂goto 880


870 open $15,8,15, f$ क
880 close 15：return
890 rem directory
900 print＂［cd］directory－sure＂：gosut 23
60：if $y$ 中〈〉＂＞＂goto 1230
910 open 15，8，15：open 1，8，2，＂\＃＂：t＝18：s＝0 ：$f$ क＝＂＂：bf＝0：bu＝0
920 print\＃15，＂u1＂；2；0；t；5；print\＃15，＂b－p ＂；2；144：print＂［cls］＂；tab（10）；＂［rev］＂；
930 for $i=1$ to 16：get\＃1，x $\$$ ：gosub 2290：pr
intx ${ }^{\text {o }}$ ：next i：print＂，＂；
940 print\＃15，＂t－p＂；2；162
950 for $i=1$ to 2：get\＃1，x\＄：gosub 2290：pri
ntxis；next i：print＂［off］［tlk］＂：t＝18：s＝1
960 print＂\＄blk file［13spc］type $\$ t k$ 事st $\$$ add［1 blu］＂
970 print\＃15，＂u1＂；2；0；t；s：i＝0：x＝s：gosub 2200：print＂［blk］trk 12＂；＂sct＂；x中；＂［1 b 1u］＂


$990 \mathrm{i}=\mathrm{i}+1: \mathrm{print} \mathrm{\# 15,"b-p";2;(i-1)*32+2:ge}$

1000 for $j=0$ to 5：if $x=j$ then $y=t=(j)$
1010 if $j=0$ then $x=x-128$

|  |  |
| :---: | :---: |
| 1030 get\＃1，$x \$$ ：gosub 2330：$x=a \operatorname{coc}(x \$)$ ：gosub |  |
| 1040 get\＃1， x \＄：gosub 2330： $\mathrm{x}=\mathrm{asc}(\mathrm{x} \ddagger)$ ：gosub |  |
| 1050 for $j=1$ to 16：get\＃1，x $=$ ：gosub 2290：f <br>  |  |
|  |  |
| 1060 if $\gamma \$\rangle$＂prg＂goto 1090 |  |
| 1070 open $3,8,3, f \$+", 5, p ": g e t \# 3, x \$: g o s u b$ 2330：x＝asc（x事）：gosub 2200：1末＝x |  |
| 1080 get\＃3，x\＄：gosub 2330：x＝asc（x\＄）：gosub 2200：1क＝x事＋1क：clase3 |  |
| 1090 print\＃15，＂b－p＂； 2 ；（i－1）＊32＋30 |  |
| 1100 |  |
| 1110 get\＃1，x\＄：gosub 2330：k＝asc（xt） |  |
|  |  |
| 0：bft＝bf\＄＋x ${ }^{\text {b }}$ |  |
| 1130 bu＝bu＋256＊k＋j：if $y \$\rangle " d e 1 "$ then bf＝ bf $+256 * k+j$ |  |
|  |  |
|  |  |
|  |  |
| 1150 get $\gamma$ 車：if $y \$\rangle$＂＂then gosub 2360 |  |
| 1160 if i＜8 goto 990 |  |
| 1170 if $t<>0$ goto 970 |  |
| $1180 \mathrm{bf}=664-\mathrm{bf}: x=\mathrm{bf} / 256$ ：gosub 2200： y \＄$=x$ \＄ |  |
|  |  |
| 1190 print＂［yel］＂；yक；＂blacks free［l blu ］＂ |  |
| 1200 |  |
| ：x＝bu－256：gosub 2200：y $\quad$＝$=y \$+x \$$ <br> 1210 print＂［yel］＂；y\＄；＂blocks free inc d |  |
|  |  |
| el files［l blu］＂ |  |
| 1220 |  |
| 1230 return |  |
| 1240 rem change header |  |
| 1250 print＂［cd］change header－sure＂：gos |  |
| b 2360：if $\gamma \$\rangle$＂y＂goto 1390 |  |
| 1260 open 15，8，15：open 1，8，2，＂\＃＂：t＝18：s＝ |  |
| 1270 print\＃15，＂u1＂；2；0；t；s；print\＃15，＂b－ |  |
| P＂；2；144：print＂［cd］current：［2spc］＂； |  |
| 1280 for i＝1 to 16：get\＃1，$\times$ \＄：gosub 2290：p |  |
| rintx ${ }^{\text {c }}$ ：next i ：print＂，＂； |  |
|  |  |

1300 for $i=1$ to 2：get\＃1，x $=$ ：gosub 2290：pr intxis：riext i：print
1310 input＂［4spc］name＂；f丰：f末＝1eftま（f本＋＂［ 16g＞spc］＂，16）
 $\times 1,2)$
1330 print\＃15，＂b－p＂；2；144
1340 for $i=1$ to 16：print\＃1，mid末（fis，i，1）； ：next i
1350 print\＃15，＂b－p＂；2；162
1360 prrint\＃15，＂b－p＂；2；162
1370 for $i=1$ to 2：print\＃1，mid本 $(y \neq i, 1) ;:$ next i
1380 print\＃15，＂u2＂；2；0；t；s：clase 15：clas
e 1
1390 return
1400 rem modify and display blocks
1410 print＂［cd］modify and display blocks －sure＂：gosub 2360：if y\＄〈〉＂y＂goto 1670 1420 open 15，8，15：apen 1，8，2，＂\＃＂：fま＝＂＂
1430 input＂［cd］track［cr］\＄［2c1］＂；tक：$x \neq t$
क：gosut 2250：t＝x：if $x<0$ or $x>40$ goto 143 $\square$
 sub 2250：s＝x：if $x<0$ or $x>20$ goto 1430
1450 print\＃15，＂u1＂；2；0；t；s；：print\＃15，＂b－ p＂；2；0
 gosub 2200：tnま＝x
1470 get\＃1，x $=$ gosub 2330：$x=a s c(x \neq): s r_{1}=x$ ： 90sub 2200：5n事＝x
$1480 n t \ddagger="[3 s p c][g>n][g>e][g>x][g>t]:[g>$ $t][g>r][g>a][g>c][g>k] \quad "+t n 末+"$［g＞s］［g＞e ］［g＞c］［g＞t］［g＞o］［g＞r］＂＋sno
1490 ct $\ddagger="[g>c][g>u][2 g>r][g>e][g>n][g>t$ ］：［g＞t］［g＞r］［g＞a］［g＞c］［g＞k］＂＋ti＋＂［g＞s］ ［g＞e］［g＞c］［g＞t］［g＞o］［g＞r］＂＋s解
1500 print＂［cls］＂；cta：printnt＊
1510 print\＃15，＂b－p＂；2；0
1520 for $i=0$ to $15: f=1 ": x=i * 8$ ：gosub 220 0：printx事；＂＂；for $j=0$ to 7
 ：gosub 2200：printx ${ }^{\text {ºn }}$＂；
1540 x $\ddagger=y=$ ：gosub 2290：f $\ddagger=f \$+x \$: n e x t$ j：pr intfis：next i
1550 gosub 1680：if $y 末=" y^{\prime \prime}$ goto 1500

1560 print＂［cls］＂；cto：printnt
1570 print\＃15，＂b－p＂；2；128
1580 for $i=16$ to $31: f(\$=1 ": x=i * 8: g o s u b 22$
00：printx $\$$ ；＂＂；for $j=0$ to 7
 ：gosub 2200：printx事；＂＂；
1600x\＄＝y事：gosub 2290：f末＝f\＄＋x\＄：next jupr intfo：next i
1610 gosub 1680：if $\gamma \neq " \gamma$＂goto 1560
1620 print＂review again＂：gosub 2360：if y
報＂＂goto 1500
1630 print＂write changes＂：gosub 2360：if $y \pm=" y^{\prime \prime}$ then print\＃15，＂u2＂；2；0；t，s
1640 if $\gamma$ 事＝＂y＂then print\＃15，＂ローa＂；0；t；s ：gosub 2390
1650 print＂next $t / s ": g o s u b 2360: i f y t=" y "$ thent＝tn：s＝sn：t $\$=t n \$: s \$=s n \$: g o t o 1450$ 1660 close 1：close 15
1670 return
1680 rem any chariges
1690 print＂any changes＂：gosub 2360：if y＇〈〉＂y＂goto 1730
1700 open 9，0：input\＃9，at：close 9
$1710 \times(=1 \mathrm{eft}(\mathrm{t}(\mathrm{a}=2)$ ： gosub 2250：print\＃15，
＂b－p＂；2，x
 sub 2250：print\＃1，chr制（x）；：next i
1730 return
1740 rem intialize disk
1750 open 15，8，15，＂i0＂：clase 15：return
1760 rem trace file
1770 pririt＂［cd］trace file－sure＂：gosub
2360：if y隹く＂y＂goto 2120
1780 input＂［cd］file＂；bft
1790 open $15,8,15:$ open $1,8,2$ ，＂\＃＂：t＝18：s＝
$1: f 末="$＂
1800 print\＃15，＂ul＂；2；0；t；s：i＝0：ft＝＂＂
1810 get\＃1，xक：gosub 2330：t＝asc（x末）：get\＃1

$1820 i=i+1: p r i n t \# 15, " b-p " ; 2 ;(i-1) * 32+2: 9$

1830 for $j=0$ to $5: i f x=j$ then $y \$=t(j)$
1840 if $j=0$ then $x=x-128$
1850 next j：if $y \neq "=1$ then $y \neq t=t(5)$
1860 get\＃1，$x$ 韦：gosub 2330：$x=a s c(x+$ ）：gosub 2200：t $\$=x$ ．

1870 get\＃1，x ：gosub 2330：x＝asc（x\＄）：gosub 2200：5\＄＝x
1880 f\＄＝＂＂：for $j=1$ to 16：get\＃1，x\＄：gosub
2290：f末＝f\＄＋x $\ddagger$ ：next j

1900 if $i=8$ and $t>0$ goto 1800
1910 goto 1820
1920 if $t=0$ goto 2120
1930 print＂［cls］［rev］trace of［off］＂；bfs
；＂file type＂；y业：print：ft $\$=y \$: n o=0$
1940 if ft ${ }^{\prime}={ }^{n}$ del＂then print＂recover fil e＂：gosub 2360
1950 print\＃15，＂b－p＂；2；（i－1）＊32＋3
1960 get\＃1，x\＄：gosub 2330：t＝asc（x\＄）：get\＃1

$1970 x=t: g o s u b 2200: t \$=x \neq: x=5: g o s u b 2200$ ：s宣＝x
1980 print＂track＂；tif＂sector＂；s\＄
1990 if $\gamma$（\＄く〉＂＞＂goto 2030
2000 print\＃15，＂b－a＂；0；t；s：input\＃15，e：if e＜＞65 goto 2030
2010 print＂recovering rot on as a suppos edly＂
2020 print＂free block is allocated．＂：no＝ 1

2030 print\＃15，＂u1＂；2；0；t；s
2040 get\＃1，x\＄：gosub 2330：t＝asc（x\＄）：get\＃1

2050 if $t=0$ goto 2070
2060 goto 1970
2070 if $f t \neq{ }^{\prime \prime} d e l$＂and $y \neq " y^{\prime \prime}$ and no＝0 90 to 2090
2080 goto 2110
2090 print＂recovery ok remember to chang e＂
2100 print＂directory track and file type ＂
2110 gosub 2360
2120 close 1：close 15：return
2130 rem backup
2140 print：print＂backup file＂
2150 print：print＂［g＞t］o allow larger fil
es to be backed up＂
2160 print＂on both disk and cassette a s eparate＂

```
2170 print"utility has been provided.[3s
pc][g>f]or "
2180 print*smaller files code could be i
ncluded here."
2190 gasub 2360:return
2200 rem dec-hex
2210 x=x and 255:x1=int(x/15): x 2=x and 1
5
2220 x1क=chr$(48+x1):if x1>9 then x1$=ch
r$(55+x1)
2230 x 2$=chr$(48+x2):if x 2)9 then x 2$=ch
r本(55+x2)
2240 x$=x1$+x 2$:return
2250 rem hex-dec
2260 x ==right$("00"+xक, 2)
2270 x 1=asc(x$) -48: x 2=asc(mid$(x$, 2))-48
2280 x=16*(x1+7*(x1>9))+x2+7*(x2>9):retu
rn
2 2 9 0 ~ r e m ~ c o n v e r t ~ t o ~ a s c i i
2300 if x$=chr$(160) goto 2320
2310 if x$<" " or x$>"z" then x w="."
2320 return
2330 rem eliminate null string
2340 if x$="" then x$=chr$(0)
2350 return
2360 rem wait
2370 get yo:if y$="" goto 2370
2380 return
2 3 9 0 ~ r e m ~ e r r o r ~ o n ~ b - a ~ c h e c k
2400 input#15,en$,emक,etक,es变:if enक<"20
" or enक="65" then return
```



```
:gosub 2360:run
```

Many more options could be provided and some of the existing options could be made fully automatic. These are exercises for you to carry out.

## BACKUP

We have produced this utility to allow selective backing-up of files between disk and tape, from disk to disk and from tape to tape. Commodore provide an excellent 1541 BACKUP program on the demonstration disk, but it only backs up whole disks. The following allows selective backing-up of single files, whether they be program or sequential. It could be modified to do more than one file when going
between disk and tape, providing that the details of each file were known in advance. We wrote the program to avoid the need to produce a special program for sequential files and the use of an extended monitor to copy machine code. BASIC programs can, of course, be duplicated by a simple load and save sequence.

The program is in two parts: the machine code and the BASIC driver which uses the machine code for program files. The following describes the driver:

## LINE ACTION

100 Set top of memory to $\$ 1800 / 6144$
110 Set source device and check whether valid
120 Do same for destination device
130 Contents of disk or tape
150 Set prog or seq file, if not known use 130
170- Go to appropriate subroutine for source, destination and file type

Subroutines
250- Seq file from disk to tape so read byte/write byte until status says end of file.
Requires a file name.
290- Seq tape to disk
290 Read header and display info with sub 700 and if non-ASCII in name then offer chance to change name.
310 Final check on name.
330- Read and write bytes until status says end of file
370- Seq tape to tape. Has a limited capacity.
370 Check for non-ASCII and option to change name.
410- Read in bytes until end of file or until ram filled eliminating $\operatorname{ASC}(0)$ on the way to avoid errors.
450 Warning - only part of file read.
480 Write to destination tape
490 Pause for destination tape
510- Seq disk to disk. Same principle as for tape above
610 Pause for destination disk
630 Simple wait for any key
650 Print TAPE or DISK in set up
690 Contents of next file on tape and prompt to rewind
700 Display tape buffer in full - highlighting any non-ASCII
750 Get file name

760 Eliminate trailing spaces in file name
800 Display directory of disk
940－Prog file disk to disk．
940 Get file name and set it up in cassette buffer
950 Fill up rest of buffer with spaces
960 Set length of name register－FNLEN and enter m／c to do a relocated load
980 Delete file from destination disk
990 Do relocated save
1010－Prog disk to tape．Initial part as for disk to disk
1040 Write header created
1045 Write Ram
1060 Prog tape to tape
1100 Prog tape to disk checking as before
BASIC PROGRAM

```
100 POKE 52,24:POKE 56,24:CLR
110 PRINT"[CLS]BACKUP UTILITY":PRINT"[CD
JFROM T/D";:GOSUB 650:Fま=Y直
120 PRINT"[CD][2SPC]TO T/D";:GOSUB 650:T
$=Y$
130 PRINT"[CD]CONTENTS T/D":GOSUB 630:IF
    Y$="D" THEN GOSUB 800:GOTO 130
140 IF Y$="T" THEN GOSUB 690:GOTO 130
150 PRINT"[CD]TYPE P/S";:GOSUB 630:IF Y直
<>"P" AND Y&<>"S" GOTO 150
```



```
170 IF F&="D" AND T&="T" AND FTक="S" THE
N GOSUB 250:RUN
180 IF F&="T" AND T&="D" AND FTक="S" THE
N GOSUB 290:RUN
    190 IF F&="D" AND T$="D" AND FT क="S" THE
N GOSUB 510:RUN
195 IF F$="T" AND T$="T" AND FTक="S" THE
N GOSUB 370:RUN
200 IF F $="D" AND T 
N GOSUB 940:RLN
210 IF F$="D" AND T$="T" AND FT$="P" THE
N GOSUB 1010:RUN
220 IF F&="T" AND T&="T" AND FT&="P" THE
N GOSUB 1060:RUN
230 IF Fक="T" AND T$="D" AND FTक="P" THE
N GOSUB 1100:RUN
```

240 RUN
250 GOSUB 750：OPEN 2，8，2，N＋＋＂，S，R＂：OPEN
$1,1,1, \mathrm{~N} \$$
260 GET\＃2，Y\＄：IF ST GOTO 280
270 PRINT\＃1，Y业；：GOTO 260
280 CLOSE 2：CLOSE 1：RETURN
290 OPEN 1，1，0：GOSUB 700：IF E＝1 THEN GOS
UB 750：GOTO 310
300 GOSUB 760
310 PRINT＂［CD］FILE NAME＂；CHR\＄（34）；N\＄；CH
R末（34）；＂OK Y／N＂：GOSUB 630
320 IF Y $\ddagger=" N$＂THEN GOSUB 750
330 OPEN 3，8，3，＂20：＂＋N\＄＋＂， $5,{ }^{(1)}$
340 GET\＃1，Yま：IF ST GOTO 360
350 PRINT\＃3，Y\＄；：GOTO 340
360 CLOSE 1：CLOSE 3：RETURN
$370 \mathrm{Y}=6144$ ：OPEN 1，1，0：GOSUB 700：IF E＝1 T
HEN GOSUB 750：：GOTO 390
380 GOSUB 760
390 PRINT＂［CD］FILE NAME＂；CHR $\ddagger$（34）；N $\ddagger$ ；CH
Rも（34）；＂OK Y／N＂：GOSUB 630
400 IF Y $\ddagger=$＂N＂THEN GOSUB 750
418 GET\＃1，Y业：IF ST GOTO 460
420 IF $Y \$="$＂THEN Y $\ddagger=C H R \neq(\theta)$
430 POKE $Y$ ，ASC（Y\＄）
440 IF $Y<40959$ THEN $Y=Y+1$ ：GOTO 410
450 PRINT＂［CDJFILE TOO BIG ONLY 34 K COPI
ED＂
460 CLOSE 1
470 GOSUB 490
480 OPEN 1，1，1，N $=$ ：FOR I＝6144 TO Y：PRINT\＃ 1，CHRक（PEEK（I））；：NEXT I：CLOSE 1：RETURN 490 PRINT＂［CD］DESTINATION TAPE Y＂：GOSUB
630：IF Yまく＞＂Y＂GOTO 490
500 RETURN
510 GOSUB $750:$ OPEN 2，8，2，N\＄＋＂，S，R＂：Y＝614
4
520 GET\＃2，Y\＆：IF ST GOTO 570
530 IF $Y \ddagger="$＂THEN $Y \equiv=C H R \$(0)$
540 POKE Y，ASC（Y $\$$ ）
550 IF $Y<40959$ THEN $Y=Y+1:$ GOTO 520
560 PRINT＂［CDJFILE TOO BIG ONLY 34K COPI
ED＂
570 CLOSE2
580 GOSUB 610

```
590 OPEN 3,8,3,"20:"+N$+",S,W":FOR I=614
4 TO Y:PRINT#3,CHR$(PEEK(I));:NEXT I
600 CLOSE 3:RETURN
610 PRINT"[CD]DESTINATION DISK Y":GOSUB
630:IF Y方〈>"Y" GOTO 610
620 RETURN
630 GET Y東:IF Y$="" GOTO 630
6 4 0 ~ R E T U R N
650 GOSUB 630:IF Y&="T" THEN PRINT" TAPE
":GOTO 680
660 IF Y$="D" THEN PRINT" DISK":GOTO 680
670 GOTO 650
6 8 0 \text { RETURN}
690 PRINT"[CD]":OPEN 1:CLOSE 1:GOSUB 700
:PRINT"[CD][REV]REWIND TAPE[OFF]":RETURN
700 PRINT"TYPE FILENAME";SPC(10);"BUFFER
    START":I=PEEK(828):E=0
710 Y$=" PRG ":IF Y=4 THEN Y$=" SEQ "
720 PRINTY婁;"[REU]";:FOR I=833 T0 1019:X
=PEEK(I):IF }X<32\mathrm{ OR }X>95\mathrm{ THEN }X=63:E=
730 PRINT CHR$(X);:IF I=849 THEN PRINT"[
OFF]<";
740 NEXT I:PRINT">[REU]END[OFF]":RETURN
750 INPUT "[CD]FILENAME";N$:N$=LEFT$(N$,
16):RETURN
760 N$="":FOR I=848 TO 833 STEP -1:X=PEE
K(I)
770 IF X=32 AND N$=*" GOTO 790
780 N$=CHR$(X)+N$
790 NEXT I:RETURN
800 PRINT"[CLS]";:OPEN 1,8,0,"$0":GET#1,
Y$,Y㝖
810 I=0:GET#1,Y$,Y&,Y$, X$:IF Y$\langle>"" THEN
    I=ASC(Y$)
820 IF X$<>"" THEN I=I +ASC(X$)*256
830 PRINTRIGHT$("[2SPC]" + STR串(I),3);" n;
: I=0
840 GET#1,Y$:IF ST GOTO 930
850 IF Y$=CHR$(34) THEN I=I +1:PRINT CHR多
(34);:GOTO 840
860 IF I=0 GOTO 840
870 IF I=1 THEN PRINT Y$;:GOTO 840
880 IF I=2 THEN PRINT TAB(22);:I=I+1
890 IF I=3 AND Y$=" " GOTO 840
900 IF Y$〈\rangle"" THEN PRINT Y多;:GOTO 840
```

```
910 PRINT:GET Y主:IF Y主く>"" THEN GOSUB 63
0
920 IF ST=0 GOTO 810
930 PRINT "BLOCKS FREE":CLOSE 1:GOSUB 63
0:RETURN
940 GOSUB 750:FOR I=1 TO LEN(N&):POKE 83
2+I,ASC(MID&(N&,I,1)):NEXT I
950 FOR I=833+LEN(N*) T0 1019:POKE I,32:
NEXT I
960 POKE 183,LEN(N*):SYS 49244
970 FRINT"[CD][REU]FILE WILL BE DELETED[
OFF]":GOSUB 610
980 OPEN 15,8,15,"S0:"+N&:CLOSE 15
990 POKE 133,LEN(N*):SYS 49343
1000 RETURN
1010 GOSUB 750:FOR I=1 TO LEN(Nक):POKE 8
32+I,ASC(MID$(N$,I,1)):NEXT I
1020 FOR I=833+LEN(N*) TO 1019:POKE I,32
:NEXT I
1030 POKE 183,LEN(N利):SYS 49244
1040 SYS 49203
1045 SYS 49206
1050 RETURN
1060 SYS 49152
1070 GOSUB 490
1080 SYS 49203
1085 SYS 49206
1090 RETURN
1100 SYS 49152
1110 GOSUB 700:IF E=1 THEN GOSUB 750:GOT
0 1130
1120 GOSUB 760
1130 PRINT"[CD]FILE NAME ";CHR$(34);N$;C
HRक(34);" OK Y/N":GOSUB 630
1140 IF Y}\ddagger="N" THEN GOSUB 750
1150 PRINT"[CD][REU]FILE WILL BE DELETED
[OFF]":GOSUB 610
1160 OPEN 15,8,15,"S0:"+N$
1170 POKE 183,LEN(N$):SYS 49343
1180 RETURN
```

The following is the basic loader for the machine code and must be loaded and run before using the above program．

1 DATA 32, 44, 247, 173, 60, 3, 133, 255 169, 0, 133
2 DATA 193, 169, 24, 133, 194, 56, 173, 63, 3, 237
3 DATA 61, 3, 170, 173, 64, 3, 237, 62,
3, 168
4 DATA 24, 138, 101, 193, 133, 174, 152, 101, 194, 133
5 DATA $175,32,162,245,165,255,141$, 60, 3, 96
6 DATA 32, 183, 247, 169, 0, 133, 193, 1 69, 24, 133
7 DATA $194,56,173,63,3,237,61,3$,
170, 173
8 DATA $64,3,237,62,3,168,24,138$,
101, 193
9 DATA 133, 174, 152, 101, 194, 133, 175 32, 124, 246
10 DATA $96,169,96,133,185,169,1,1$ 41, 60, 3
11 DATA $133,184,169,8,133,186,169$, 0, 133, 195
12 DATA $133,147,169,65,133,187,169$ , 3, 133, 188
13 DATA $169,24,133,196,164,183,32$, $175,245,32$
14 DATA 213, 243, 165, 186, 32, 9, 237,
165, 185, 32
15 DATA 199, $237,32,19,238,141,61$,
3, 32, 19
16 DATA 238, $141,62,3,32,232,244,1$
65, 174, 141
17 DATA 63, 3, 56, 165, 175, 233, 24, 14 1, 64, 3
18 DATA 24, 173, 61, 3, 109, 63, 3, 141, 63,3
19 DATA $173,62,3,109,64,3,141,64$, 3,96
20 DATA $169,97,133,185,169,1,133$, 184, 169, 8
21 DATA $133,186,169,65,133,187,169$ , 3, 133, 188
22 DATA $165,185,164,183,32,213,243$ , 32, 143, 246
23 DATA $165,186,32,12,237,165,185$, 32, 185, 237

```
24 DATA 169, 0, 133, 172, 169, 24, 133,
173, 56, 173
25 DATA 63, 3, 237, 61, 3, 133, 174, 173
    64,3
26 DATA 237, 62, 3, 133, 175, 24, 169, 2
4,101,175
27 DATA 133, 175, 173, 61, 3, 32, 221, 2
37, 173, 62
28 DATA 3, 160, 0, 32, 33, 246, 96, 255,
    255,0
29 FOR I=49152 T0 49432:READ A:POKE I,A:
NEXT I
```

Once this has been run it could be saved as its machine code for later ease of loading. A detailed description of the machine code follows.

## MACHINE CODE

The machine code is called by the driver as required. It consists of four parts:
i) Read any header and do relocated load
ii) Write to tape current header and write relocated code
iii) Load from disk retaining original details but relocate
iv) Save to disk relocated code with original details

Chapter 5 of the Programmer's Reference Guide, 'The KERNAL', discusses the use of LOAD and SAVE in detail. The entry points given are for complete loads and saves (it is possible to do a relocated load, but not a relocated save using these). Unfortunately, as we are using an allpurpose BASIC driver, these entry points may overwrite it. To overcome this problem, every operation is carried out in two stages. The first is to read or write the file's details which are always stored in or taken from the cassette buffer. This avoids having to do too much moving of information. The second is to perform a relocated load or save with the correct amount of code going to or being taken from $\$ 1800$ on.

To do this we must enter the load and save routines at much later points with the parameters already set. It would consume too much space to describe these in detail, so we leave it up to you to follow them through. The only tricks are to prevent a forced load to its original address with a tape marker of 3 and to prevent a header being written with a marker of 5 (when an end of tape has been written - see Tape Headers).

| TAPE |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| Read | any tape h | header. | withou | $t$ loading |
| C000 | 202CF7 | JSR | \$F72C | read any header |
| C003 | AD3C03 | LDA | \$033C | get sec add |
| C006 | 85FF | STA | \$FF | and store for later |
| C008 | A900 | LDA | \#\$80 | set start of laad |
| COOA | 85C1 | STA | \$C1 | to $\$ 1800$ by setting |
| COOC | A918 | LDA | \#\$18 | STAL |
| COEE | 85C2 | STA | \$C2 |  |
| C010 | 38 | SEC |  | subtract MSBs |
| C011 | AD3F03 | LDA | \$033F | of start and end |
| C014 | ED3D63 | SBC | \$0330 | of original load |
| C017 | AA | TAX |  | put result in $X$ |
| C018 | AD4803 | LDA | \$0340 | do same for LSBs |
| C01B | ED3E03 | SBC | \$033E | putting answer in $Y$ |
| C01E | A8 | TAY |  |  |
| C01F | 18 | CLC |  | find overall length |
| C020 | 8A | TXA |  |  |
| C021 | 65 Cl | ADC | \$C1 | and add to STAL |
| C023 | 85AE | STA | \$AE | to give the new |
| C025 | 98 | TYA |  | end i.e. ${ }^{\text {c }} 1800$ |
| C026 | 65 C 2 | ADC | \$C2 | plus result |
| C028 | 85AF | STA | \$AF |  |
| Load | from \$1808 | 0 on |  |  |
| C02A | 20A2F5 | JSR | \$F5A2 | do the relocated load |
| C02D | A5FF | LDA | \$FF | restore the sec add |
| C02F | 8D3C03 | STA | \$033C | in case end of tape 5 |
| C032 | 60 | RTS |  | load complete |
| Write to tape in two parts the correct header and then the code from $\$ 1800$ on |  |  |  |  |
|  |  |  |  |  |
| C033 | 20B7F7 | JSR | \$F7B7 | write header in orig form |
| C036 | A908 | LDA | \#\$08 | reset STAL as it |
| C038 | 85C1 | STA | \$C1 | has been changed by |
| C03A | A918 | LDA | \#\$18 | writing the header. |
| C03C | 85 C 2 | STA | \$C2 |  |
| C03E | 38 | SEC |  | recalc the relocated end |
| C03F | AD3F83 | LDA | \$033F | MSB |
| C042 | ED3D03 | SBC | \$033D |  |
| C045 | AA | TAX |  |  |
| C046 | AD4083 | LDA | \$0340 | LSB |
| C049 | ED3E03 | SBC | \$033E |  |
| C84C | A8 | TAY |  |  |
| C04D | 18 | CLC |  |  |
| C04E | 8A | TXA |  |  |
| C04F | 65 Cl | ADC | \$ Cl | set up EAL |


| C051 | $85 A E$ | STA | \$AE |  |
| :--- | :--- | ---: | :--- | :--- |
| C053 | 98 | TYA |  |  |
| C054 | $65 C 2$ | ADC | \$C2 |  |
| C056 | $85 A F$ | STA | \$AF |  |
| Save RAM for reloading |  |  |  |  |
| C058 | $207 C F 6$ | JSR | \$F67C | save RAM |
| C05B | 68 | RTS |  | complete |
|  |  |  |  |  |
| DISK |  |  |  |  |


| Load | from disk | re | 咗 | \$1800 |
| :---: | :---: | :---: | :---: | :---: |
| c05c | A960 | LDA | \# $\$ 60$ |  |
| C05E | 85B9 | STA | \$89 | set sec add |
| C060 | A901 | LDA | \#\$01 | put type 1 in tape buffer |
| C062 | 803C03 | STA | \$833C |  |
| C065 | 85B8 | STA | \$88 | make log file 1 |
| C067 | A908 | LDA | \#\$08 | make device 8 |
| C069 | 85BA | STA | \$BA | and put in FA |
| C06B | A900 | LDA | \#\$88 |  |
| C060 | 85C3 | STA | \$C3 |  |
| C06F | 8593 | STA | \$93 | $A=0$ for load |
| C071 | A941 | LDA | \#\$41 | set pointer to file name |
| to f | le name |  |  |  |
| C073 | 85BB | STA | \$ BB $^{\text {d }}$ | in FNADR to TBUFFR + 5 |
| C075 | A903 | LDA | \#\$03 |  |
| C077 | 85BC | STA | \$BC |  |
| C079 | A918 | LDA | \#\$18 | MEMUSS set to $\$ 1800$ |
| C07B | 85C4 | STA | \$C4 | for relocated load |
| C070 | A4B7 | LDY | \$B7 | read len name set in BASIC |
| C07F | 20AFF5 | JSR | \$F5AF | print SEARCHING |
| C082 | 2005F3 | JSR | \$F3D5 | print LOADING |
| C085 | A5BA | LDA | \$BA | get current device |
| C087 | 2009ED | JSR | \$ED09 | send talk |
| C08A | A5B9 | LDA | \$B9 | get sec add |
| C08C | 20C7ED | JSR | \$EDC7 | send talk sec add |
| C08F | 2013 E | JSR | \$EE13 | receive from serial |
| C092 | 803D03 | STA | \$033D | and store LSB in TBUFFR |
| C095 | 2013 EE | JSR | \$EE13 | do same for MSB of start |
| C098 | 803E03 | STA | \$033E |  |
| C09B | 20E8F4 | JSR | \$F4E8 | do relocated laad |
| C09E | A5AE | LDA | 中AE | get end LSB |
| C0A8 | 803F93 | STA | \$033F | put in TBUFFR |
| COA3 | 38 | SEC |  | subtract relocated start |
| C8A4 | A5AF | LDA | \$AF | and end |
| C0A6 | E918 | SBC | \#\$18 | and put in appropriate |
| C0A8 | 804803 | STA | \$8348 | locations of TBUFFR |


| COAB | 18 | CLC |  |  |
| :---: | :---: | :---: | :---: | :---: |
| COAC | AD3083 | LDA | \$033D | leaving a header. |
| COAF | $603 F 03$ | ADC | \$033F | suitable for |
| C0B2 | 803F03 | STA | \$033F | a tape write |
| C0B5 | AD3E03 | LDA | \$033E |  |
| C0B8 | 604003 | ADC | \$0340 |  |
| CabB | 804003 | STA | \$0340 |  |
| C0BE | 60 | RTS |  | back to BASIC |
| Save | relacated | code | torelo | d at correct address |
| C0BF | A961 | LDA | \#\$61 | set parameters |
| CaCl | 85B9 | STA | \$ $\mathrm{B}^{\text {9 }}$ |  |
| CaC3 | A901 | LDA | \#\$01 |  |
| Cocs | 85B8 | STA | \$88 |  |
| CaC7 | A908 | LDA | \#\$08 |  |
| COC9 | 85BA | STA | 車BA |  |
| CaCB | A941 | LDA | \# \$ $41^{\text {d }}$ |  |
| COCD | 85BB | STA | \$ BB |  |
| COCF | A903 | LDA | \#\$03 |  |
| COD1 | 85BC | STA | \$BC |  |
| C0D3 | A5B9 | LDA | \$ $\mathrm{BC}^{\prime}$ |  |
| C0D5 | A4B7 | LDY | \$B7 |  |
| COD7 | 2005F3 | JSR | \$F3D5 | send sec add |
| CODA | 208FF6 | JSR | \$F68F | print SAVING |
| CODD | A5BA | LDA | \$BA | send listen |
| CODF | 200CED | JSR | \$ EDAC | device 8 |
| C0E2 | A5B9 | LDA | \$ ${ }^{\text {9 }}$ | send listen |
| C0E4 | 20B9ED | JSR | \$ EDB9 | sec add |
| C0E7 | A900 | LDA | \#\$00 | set up SAL |
| C0E9 | 85AC | STA | \$AC | with \$1800 |
| C0EB | A918 | LDA | \#\#18 |  |
| COED | 85AD | STA | \$AD |  |
| C0EF | 38 | SEC |  |  |
| COFP | AD3F03 | LDA | \$033F | calculate prag length |
| COF3 | ED3D83 | SBC | \$0330 | and put |
| COF6 | 85AE | STA | \$AE | in EAL |
| C6F8 | AD4003 | LDA | \$0340 |  |
| COFB | ED3E03 | SBC | \$033E |  |
| COFE | 85AF | STA | \$AF |  |
| C100 | 18 | CLC |  |  |
| C101 | A918 | LDA | \#\$18 | add $\$ 1890$ to |
| C103 | 65AF | ADC | \$AF | give relocated |
| C185 | 85AF | STA | \$AF | end |
| C107 | AD3D83 | LDA | \$033D |  |
| C18A | 200DED | JSR | \$EDDD | send serial deferred |
| C10D | AD3E03 | LDA | \$033E | send actual start in $A$ and $Y$ |

```
C110 A000 LDY #$00
C112 2021F6 JSR $F621 save RAM to reload
C115 60 RTS back to BASIC
```

The utility is only intended for your own files. It will not as it stands backup relative files.

As a point of interest, Supersoft's zoom monitor offers not only the option to perform relocated loads and saves, but to save in a form suitable for reloading on a PET, which eliminates an ID of 3 not used on the PET.

## 3 A token approach to BASIC

## Introduction

In this chapter we deal with the five main routines basic uses in interpreting your programs or commands. One of these, CHRGEt, picks up single bytes from the program and is called by the majority of routines in your 64. The other four routines covered are concerned with keywords - converting from ASCII to tokens, the reverse process (LIsting), and directing them to their respective routines.

Other than using sYS commands a knowledge of these routines is essential if you wish to extend existing commands or add further ones. Those of you owning a disk drive will be familiar with the program DOS 5.1 and may know that this modifies CHRGEt to trigger its commands.

## CHARGET

BASIC gets its information from the input or program lines through a routine called CHRGET (CHaracterR GET). A copy of this routine is held in the KERNAL operating system and is copied into zero page on powerup. Each time BASIC wants a character it calls this routine.
The routine is held at locations \$0073-\$008A (KERNAL is \$E3A2-\$E3B9) and is as follows:


Bytes 0073-0077
Every time the CHRGET routine is called it increases, by one, the location
from where it gets its information. After it increments the LSB, in location 007A, it checks whether the page has been crossed, that is, from $\$ F F$ to $\$ 00$. Only if it has will the msb be increased.

## Bytes 0079-007B

Here it takes the information from store and puts it into the accumulator. The store location is present before the initial entry to the routine. It is always set one byte less due to initial increment. If you were going to use the routine yourself, it would be these bytes you would change, as we shall see later.

## Bytes 007C-007F

Here it checks to see if the character is a numeral. It is testing to see if it is greater than ASCII numeral 9 (\$39). If so then the routine is left via \$008A with the carry set.

## Bytes 0080-0083

Here is a straightforward test to see if a space was picked up; if so the routine is carried out again. CHRGET cannot be left on encountering a space.

## Bytes 0084-0089

These successively subtract two numbers from the original byte and end up with the same number. You may say that is senseless but it will set two flags in the status register that help us later. These are the carry flag and the zero flag.

The carry flag. If this is clear on exit the byte will be a numeral in ASCII form. If it is set we have something else.

When subtracting two numbers in machine code the carry flag must always be set first. If the number we are subtracting from is the larger then the carry will remain set. On the other hand, if the number we are subtracting is the larger the carry will clear.

In this case we have already eliminated any byte that has a higher ASCII value than the numeral 9, in bytes 007E and 007F. It now subtracts $\$ 30$ (ASCII for digit zero) from the accumulator in preparation for setting the final flags which can be used for testing for numbers. The carry flag at this point does not matter as it is set again anyway. Bytes with ASCII values lower than numerals now range between $\$ \mathrm{D} 0$ and \$FF.

With the next subtraction the original value is restored. The carry flag will now be set or unset as BASIC requires. As numerals are the only ones less than \$D0 (the last figure subtracted) they will be the only ones to clear the carry flag.

The zero flag. This could be set in two instances. First, in bytes 007 E and 007 F where we tested our byte against $\$ 3 \mathrm{~A}$, the ASCII value for a colon. If it was a colon then the zero flag would have been set as it was equal (the carry would also have been set). Secondly, the flag would be set after the second subtraction if, and only if, the original byte was
zero (not ASCII digit zero). In that instance after the first subtraction the accumulator would hold \$D0 and subtracting the same value would set our zero flag.

A colon in a basic line signifies an end of statement and a zero, the end of a line, and hence an end of statement. Therefore, by testing the zero flag, we can quickly tell if we have reached the end of that particular instruction.

## CHRGOT

Keyword routines are entered immediately after a call to Chrget. Before using that byte it may require the accumulator for something else. To recover the CHRGET byte we can use the ChRGOT routine. This is a shortened version of CHRGET. If, instead of entering the routine at \$0073, we enter at $\$ 0079$, we miss out the instructions which update the pointer and get the original byte again.

## Wedges

If we want to patch in our own machine code routines to work alongside BASIC, one way to do this is to insert a wedge. Simply, this is a routine which diverts CHRGET to check whether it is one of our additions. From this a decision can be made whether to revert to the normal CHRGET flow or to a routine of our own.

Let us say that we put in some routines all to be actioned on the character ' $a$ '. For instance, we could have a renumber routine and a delete routine. The command for renumber might be ' $a \mathrm{R}^{\prime}$ ' and delete, ' $\left(a D^{\prime}\right.$ and could place a wedge at $\$$ cono.

The first thing we have to do is alter the Chrget routine. We want to change it after it has collected a byte but before it starts checking and manipulating it. The alteration would therefore be at \$007C with a JMP to our coding. The routine starts with six bytes of data with our changes and the original bytes. (The latter is for restoring CHRGET if you require so to do later). We can load them by using the load with the $x$ register. Our first instructions will look like this:

| C000 | 4 C | 11 | C6 | JMP | \$cal1 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| C003 | C9 | 3A |  | CMP | \# \$ $^{\text {A }}$ |
| C005 | B6 |  |  | BCS |  |
| C006 | A2 | 02 |  | LDX | \#\$02 |
| C008 | BD | 08 | C0 | LDA | \$C000, $\times$ |
| C00B | 95 | 7 C |  | STA | \$7C, $X$ |
| C:00] | CA |  |  | DEX |  |
| COBE | 10 | F8 |  | BPL | \$C008 |
| C010 | 60 |  |  | RTS |  |

C011 Our coding will start here

This is the routine to initialize our wedge and is called immediately after loading the program by using SYS 49158 (\$C006). This loads a byte $x$
places from $\$ 0000$ and stores it $x$ places from $\$ 7 C$. We decrease the counter $x$ and the branch to collect the next byte will work until we decrement it below zero, that is, no longer a positive number. If the branch fails, we go back to BASIC through the RTS.

The beginning of the ChRGET now looks like this:
0073 INC $\$ 7 A$
0075 BNE $\$ 0079$
B077 INC $\$ 7 B$
0079 LDA $\$ 0200$ (this number varies)
907C JMP $\$ C 011$

Now each time CHRGET is used it will go to our routine. As all our commands would be triggered with ' $a$ ', the first thing we would do is to check to see if it is present:

| C011 | C9 40 | CMP | \#\$40 |
| :--- | :--- | :--- | :--- |
| C013 | F0 03 | BEQ | $\$ C 018$ |

If it is, we can branch to do further checks, and if not, we continue with the next code. Here we will have to revert to the normal course of events. We have two options which we can take. First, we include the bytes of CHRGET we changed into our program - we do not want to change CHRGET itself as we want it to use again - and jump back to CHRGET at $\$ 0082$. Secondly, we can use the CHRGET routine in the KERNAL ROM, jumping in at \$E3AB, and BASIC will continue as if nothing has happened.

The first method would be like this:

```
C015 4C AB E3 JMP $E3AB
```

And the second like this (of course, the routine address will change from \$C018 to \$C01D):

| C. 013 | Fo | 87 |  | BEQ | \$C010 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| CO15 | C9 | 3 |  | CMP | \#\$3A |
| C017 | B0 | 83 |  | BCS | \$C01C |
| C019 | 4 C | 80 | 00 | JMP | \$0080 |
| Colc: | 68 |  |  | RTS |  |

We now want to find out if it is one of our routines. This we do by checking the next character without updating the CHRGET pointer (in case it isn't). The code would look like this:

| C018 | 08 | PHP |
| :--- | :--- | :--- |
| C019 | 48 | PHA |
| C01A | 98 | TYA |
| C01B | 48 | PHA |


| C01C | 84 |  | TXA |  |
| :---: | :---: | :---: | :---: | :---: |
| C010 | 48 |  | PHA |  |
| COIE | A6 | 7A | LDX | \$7A |
| C020 | E8 |  | INX |  |
| C021 | BD | $00 \quad 02$ | LDA | \$0200, x |
| C024 | C9 | 52 | CMP | \#\$52 |
| C026 | F0 | ?? | BEQ | - TO RENUMBER ROUTINE |
| C028 | C9 | 44 | CMP | \#\$ 44 |
| C02A | F0 | ?? | BEQ | - TO delete routine |
| CO 2 C | 68 |  | PLA |  |
| C02D | AA |  | TAX |  |
| C02E | 68 |  | PLA |  |
| CO2F | A8 |  | TAY |  |
| C.038 | 68 |  | PLA |  |
| C031 | 28 |  | PLP |  |
| C032 | 4. | AB E3 | JMP | \$E3AB |

## Bytes C018-C01D

Here we are preserving our registers, including the status register, on the stack in case it is not destined for our own routines.

## Bytes C01E-C023

Location \$7A has the LSB of the pointer used by CHRGET and this is one less than the next character we want. So if we load it into the $x$ register and increase it by one, we will have the position of the next byte. We can now load the byte using $x$ as a pointer.

## Bytes \$C024-\$C2B

A check is made to see if it is the letter R , signifying the renumber routine. If not we then check for the letter $D$ and if so go to delete.

## Bytes \$C2C-\$C034

It not, we restore our registers from the stack (in the reverse order we put them on). Then we continue the KERNAL routine as before.
One last point is that in the routines, such as renumber or whatever, it would be advisable to call a subroutine to remove the bytes from the stack (placed there in \$C018 to \$C01D). We do not require them, but as your routines are called the stack will become fuller and fuller, resulting in an 'OUT OF MEMORY' error.

## Keywords

A more professional approach to adding routines than altering CHRGET is the use of Keywords. This approach holds with the idea behind the BASIC language that actions can be performed by using words which are indicative of the desired action.

Keywords can be divided into two types: commands and functions.

Functions get information, for example, peek returns the contents of a location, and will always supplement a command keyword. For example, we use PRINT PEEK(xx) but never PEEK(xx)PRINT.

On the 64 it is possible to incorporate new routines actioned by keywords as they go to the relative routines through vectors held in RAM. As the vectors are in RAM we can change them to go to routines of our own.

The other items we will have to add are three tables of data. One will have the new keywords in ASCII code. This will be used when LISTing and tokenizing. The end of a word is indicated by adding $\$ 80(128)$ to its last letter. To signify the end of the table, a zero is used. If we had a table of two keywords, say END and NOT, it would look like this:

$$
\begin{array}{lllllll}
45 & 4 \mathrm{E} & \mathrm{C} 4 & 4 \mathrm{E} & 4 \mathrm{~F} & \mathrm{D} 4 & 00
\end{array}
$$

The other two tables will have the addresses of our routines, one for command keywords and one for functions. These will hold the address of each routine less 1 . The reason for this is that we will put them on the stack for an RTS instruction. The program counter will add 1 when it takes them from the stack, thus getting the correct address for the routine. The table will have two bytes for each routine, the LSB and then the msb. For example, if we had a routine at \$C4DF, on the table it will look like this: DE C4.

There are four vectors we will have to change (three if not adding functions).

ADD OF VECTOR ADD OF ROM DESCRIPTION

| $\$ 0304 / 5$ | \$A57C | Tokenize BASIC |
| :--- | :--- | :--- |
| $\$ 0306 / 7$ | \$A71A | Print Tokens(LIST) |
| $\$ 0308 / 9$ | \$A7E4 | Token Dispatch - Command words |
| $\$ 030 A / B$ | \$AE86 | Token Dispatch - Function words |
| Tokenize bASIC Text |  |  |

The object of this subroutine is to take an input line, check it for keywords, tokenize them and condense the line. It does this by taking every byte from the input buffer, not using CHRGET, and then checking through the keyword table for a match. If the letters do not make a keyword, it stores them as variables, meaning that variables cannot have keywords in them.

There are two ways we could approach the problem of incorporating our own keywords and tokens. First, we could copy the baSIC from ROM into RAM, and alter the tokenize routine within BASIC so that if it cannot find a match it jumps to a routine to check through our table of keywords. This would mean we would not have to change the vector but would lose the RAM area under the BASIC ROM which is useful for storing hires screens, data tables, and so on. The second way would be to change the vector to a routine of our own. Here we have a copy of
the rom routine altered slightly to be able to search our table as well. We would only use the rom routine when we finished tokenizing a whole line.
We shall describe the second method, which we think is the better in the long run. A description follows the code.

| 10 |  | LDX | \$7A |
| :---: | :---: | :---: | :---: |
| 20 |  | LDY | \#\$84 |
| 30 |  | STY | \$ FF |
| 40 | ANOTHER | LDA | \$0200, X |
| 50 |  | BFL | SPACE |
| 60 |  | CMP | \#\#FF |
| 70 |  | BEQ | Store |
| 80 |  | INX |  |
| 90 |  | BNE | ANOTHER |
| 180 | SPACE | CMP | \#\$20 |
| 110 |  | BNE | STORE |
| 120 |  | STA | \$08 |
| 130 |  | CMP | \#\$22 |
| 140 |  | BEQ | QUOTE |
| 150 |  | BIT | क 0 F |
| 160 |  | BUS | STORE |
| 170 |  | CMP | \# ${ }^{\text {S }}$ 3F |
| 180 |  | BNE | NUMBER |
| 190 |  | LDA | \#\$99 |
| 200 |  | BNE | STORE |
| 210 | NUMBER | CMP | \#\$30 |
| 220 |  | BCC | CONT |
| 230 |  | CMP | \#\#3C |
| 240 |  | BCC | STORE |
| 250 | CONT | STY | \$71 |
| 260 |  | LDY | \#\$00 |
| 270 |  | STY | 棟 |
| 280 |  | DEY |  |
| 290 |  | STX | \$7A |
| 300 |  | DEX |  |
| 310 | NEXT LETTER | INY |  |
| 320 |  | INX |  |
| 330 | CONT 1 | LDA | \$0200, X |
| 340 |  | SEC |  |
| 350 |  | SBC | \$A09E, Y |
| 360 |  | BEQ | NEXT LETTER |
| 370 |  | CMP | \#\$80 |
| 390 |  | BNE | NEXT WORD |
| 400 | STORE A | ORA | \$0B |
| 410 | FOUND | LDY | \$71 |

420 STORE INX
430 INY

440 STA $\$ 01 \mathrm{FB}, \mathrm{Y}$
450 LDA $\$ 01 \mathrm{FB}, \mathrm{Y}$
460 BEQ END
470
480
490
500
510
SEC
SBC \# ${ }^{\text {\# }}$ 3A
BEQ COLON
CMP \# 49
BNE DATA
520 COLON STA S0F
530 DATA SEC
540
550
560
570 LINE
580
590
600
618 QUOTE
620
630
SBC \#\$55
BNE ANOTHER
STA $\$ 08$
LDA \$8200, X
BEQ STORE
CMP $\$ 08$
BEQ STORE
INY
STA $\$ 01 \mathrm{FB}, \mathrm{Y}$
INX
640 BNE LINE
650 NEXT WORD LDX \$7A
660 INC $\$ 0 B$
670 FIND INY
680 LDA \$A89D, Y
690
700
710
720
730
740 NEXT
750
760 NEXT B
770
780
790
800
810
820
830 NEXT NEW
840
850 NEXT A
860
870
BPL FIND
LDA \$AB9E,Y
BNE CONT 1
LDY \#\$FF
DEX
INY
INX
LDA $\$ 0200, x$
SEC
SBC \$START OF OUR WORD TABLE, $X$
BEQ NEXT
CMP \#\$ 80
BNE NEXT NEW
BEQ STORE A
LDX \$7A
INC \$0B
INY
LDA $\$$ START OF OUR WORD TABLE $-1, X$
BPL NEXT A

880
890
900
910
920 END
LDA GSTART OF OUR WORD TABLE, $X$
BNE NEXT B
LDA $\$ 0200, X$
BPL FOUND
JMP \$A609

LINES 10-30: Initialization. Location \$7A will have a value the same as the position within the input buffer. In immediate mode this will be 0 , the start of the buffer. If inputting a program line, it would be after the line number, which has already been taken care of. Now $x$ will be our pointer to the original contents of the buffer and $\gamma$ will be the pointer to our new buffer set up. $\gamma$ is stored in $\$ 0 F$ just to initialize that location. The value does not matter as long as it was not over $\$ 3 \mathrm{~F}$, as we shall see.
LINES 40-50: We load in a byte and check to see if it is under $\$ 80$ (128). If it is, we branch off to line 100.

LINES 60-90: Values of $\$ 80$ and over arrive here and we check to see if it is ' PI ' ( $\$ \mathrm{FF}$ ). If it is, we branch further into the routine to store it. If not, we branch back to get another byte. This means we cannot use the first character of a keyword as a shifted letter because it would be greater than $\$ 80$.

LINES 100-140: If a space is found we branch off to store it, otherwise we store the accumulator in a register, in case the following check succeeds, for comparison later. Now we check to see if the byte is a quote. Items between quotes do not require tokenizing and therefore we branch and continually store them until we come across another quote or the end of the line.

LINES 150-160: Here we are checking to see if location \$0F has bit 6 set or not. Amongst other things, the BIT instruction takes bit 6 of $\$ 0 \mathrm{~F}$ and places it in bit 6 of the status register, the overflow register. This bit in $\$ 0 F$ will only be set in this routine if we tokenize DATA later in the routine. It means that after DATA all characters will be stored and do not go through the keyword table. Bit 6 can be unset by a colon if outside quotes and for this reason colons have to be in quotes within DATA statements. A colon outside quotes will mean that information after is tokenized. BASIC instructions can be placed at the end of a data line and will be actioned as the line is encountered. BASIC differentiates between data and rem. On rem it will go to the next line whereas with data it will search through it for another BASIC instruction.
LINES 170-200: The '?' is the shortened version of the keyword PRINT. It is the only keyword which can be shortened to a single character. These lines check for the question mark and if found place the token for PRINT, $\$ 99$, in the accumulator and go to store it.
LINES 210-240: Here we find out if the byte is a numeral, colon or
semicolon. If it is off we go to store it, if not then we continue the routine.

LINES 250-300: We now set up for the search through the table of keywords. We store our $y$ register, which, if you remember, is the pointer to the 'new look' buffer. Location sob will be our counter to the number of keywords we encounter; it will not hold the token value but helps determine it. We store $x$ in $\$ 7 \mathrm{~A}$. This is part of ChRGEt, but we are only using it as a store. We decrease both $x$ and $y$ as the first part of the next section will increment them.

LINES 310-390: This section of the routine explains why on the Commodore 64 we can use shortened keywords using shifted letters. There are two things to remember here. First, the last letter of the keyword in the table is the value of the letter plus $\$ 80(128)$ which when loaded will set the negative flag. Secondly, the value of a shifted letter (not logo shift) is also the value of the letter plus $\$ 80$.

Back in the routine we increase the registers and load in our byte again; it will later load the next byte. We set the carry for subtraction and subtract the value of a letter, from the keyword table, from our byte. If we are left with zero then we have a match and we go back to get the next letter from the buffer. If it fails, we check to see if we have the value $\$ 80$ left. This would indicate we have a match by either the input letter being the shifted letter or we have reached the end of the keyword in the table and have also matched. Failing this second check, it branches off to find the end of that word so we can check the next for a match.

As the second letter can be shifted to give our match this explains not only shortened keywords, but also why some require at least three. As an example, take the keywords CLR and ClOSE. CLR comes before ClOSE in the table so that $C$ and shift $L$ will match with the former before it gets to Close, which will therefore need two standard letters before a shifted letter. We can also explain why there is no shortened version for InPUT. INPUT\# Comes before INPUT and so any shortened version will always match with INPUT\#.

LINES 400-460: Back to the routine now. We have in the accumulator the value $\$ 80$ and have found a match in the table. Here we perform the logical OR of the accumulator and location 50b. Later on we will find out that every time we pass through a keyword that does not match, we increase the value of $\$ 0 B$. As we started at $0, \$ 0 B$ will have the number of the match word, the first word in the table being 0 . The instruction OR forces bits into the accumulator if they are not set. In this case it has the same effect as adding but saves bytes doing it this way. The accumulator always has $\$ 80$ and if you OR it with a value of one you get $\$ 81$. This is how we arrive at the token value. The keyword GO is the last word in the table and that is the 76th, giving a value in $\$ 9 B$ of $\$ 4 B$ (75), and oring it
with $\$ 80$ gives a value of $\$ C B$ which is the token value of GO.
Having got our token value, we load back into the y register from $\$ 71$ - the pointer to the new buffer layout. We have now reached the point where we store a lot of the characters into the new buffer layout. This is the point where many of the earlier branches arrive. First, we increase both buffer pointers. The base location to store is \$01FB indexed with Y . As we started with Y equal to $\$ 04$ and increase it immediately, our first location will be $\$ 01 \mathrm{FB}+5$, giving $\$ 0200$. We will not overwrite anything in the original line we have not checked for two reasons. First, if there was a line number at the beginning it has been dealt with and is no longer needed. Secondly, the routine only shortens and never lengthens.

Once the byte has been stored, it is loaded back in and checked for zero which signifies the end of the line.

LINES 470-550: This section is going to test if the byte we have stored is either a colon or one of the keywords data or rem. If it is a colon, it unsets bit 6 of location $\$ 0$ which we discussed earlier, then goes to get more bytes. DATA will set the 6 th bit of that location before getting more bytes.

REM is slightly different in that nothing after it requires tokenizing. We set the location of $\$ a b$ to zero, as that is the result of the subtractions, which we will not actually need for checking but which stops us from branching out of the next section for any reason other than the end of line.

LINES 560-640: These lines are used only in two instances: on encountering a quote or encountering rem. All this does is to move bytes from their original to their new position in the input buffer, until we reach the end of the line or, in the case of a quote, we find a closing quote. Finding either of these, we branch back to the normal store lines of $420-460$. The branch in line 640 is enforced as $x$ can never be zero as before we get here we have looked at a minimum of two bytes so the least $x$ can be is 1 and the maximum $\$ 58$ (the maximum input line length).

LINES 650-710: This is the section we come to if we did not find a keyword match. All this does is to search for the next character in the table that has a value greater than $\$ 80$, and return with its position. The pointer will be increased before we start another match. It also checks for a 0 , signifying the end of the table.

So far the routine is the same as in ROM but now we change the course of events. In the ROM routine, when it finds the end of the table it assumes the characters are variables and stores them as such. We, on the other hand, want to see if it is one of our keywords, so on getting zero we have to search our table. The $y$ register is loaded at the beginning of the section because either way we want to get back the
first character of this particular check from the original input buffer line up.

LINES 720-910: This is a repeat of the checking of the standard keywords except it will have the address of our keyword table. We will now be checking for our keywords.

LINE 920: When we have found and stored the end of line zero, we get here. The routine now jumps off to end or to the original rom routine. There it will reset the CHRGET pointers to their initial setting of \$01FF and continue the normal flow of BASIC to either store the line or carry out its instructions if in direct mode.

## Print tokens

This routine is part of the LIST routine in BASIC. It takes the token value, finds the keyword and prints it to the screen, or other device. The rom print token routine is not a subroutine by itself but an integral part of LIST, but thankfully it is vector-started. The vector points to the next instruction in the ROM routine. What we would have to do is to change the vector to a routine of our own, PRINT keywords of either the standard ones or our own, and then jump back to the LIST routine at an appropriate point. The coding for such a routine is as follows:

10
20
30
40
50
60
70
80
90
100
110
120
130
148
150
160
170
180
190
200
210
220
230 START
240
250 NEXT WORD

BFL ROM 1
CMP \#まFF
BEQ ROM 1
BIT \$0F
BMI ROM 1
CMF \#\#CC
BCC CBM TOKEN
SEC:
SBC $\ddagger \mathrm{CB}$
TAX
LDA \# LSB START OF OUR KEYWORD TABLE
STA $\$ 22$
LDA \#\$MSB START OF OUR KEYWORD TABLE
STA $\$ 23$
ENVE START
CBM TOKENS SEC
SBC \#\$7F
TAX
LDA 9 9E
STA $\$ 22$
LDA \#\$A0
STA $\$ 23$
STY $\$ 49$
LDY \#\#FF
DEX

| 260 | BEQ WORD FOUND |
| :--- | :--- |
| 270 NEXT CHAR | INY |
| 280 | LDA (\$22), Y |
| 290 | BPL NEXT CHAR |
| 300 | BMI NEXT WORD |
| 310 WORD FOUND | INY |
| 320 | LDA (\$22), Y |
| 330 | BN1 ROM 2 |
| 340 | JSR \$AB47 |
| 350 | ENE WORD FOLND |
| 360 ROM 1 | JSR \$A6F3 |
| 370 ROM 2 | JSR \$A6EF |

LINE 10: This tests the negative flag. A value of 980 (128) or over is signalled as negative. As all tokens are $\$ 80$ or over, this branch will succeed; values under $\$ 80$ go back to LISt unchanged.
LINES 20: $\mathbf{~ S F F}$ is the value of ' Pl '. If it is that value we again return to LIST.
LINES 40-50: What we are doing here is putting bit 7 of location $\$ 0 F$ into the negative flag, although it does other things which are of no concern to us. Location \$0F is the flag used by the LIST routine to signal if it is listing in quotes or not. If bit 7 of SOF is 1, then it is listing in quotes and we do not want to print tokens but the ASCII of the bytes. Therefore, if the negative flag is set, we branch to go back to rom.
LINES 60-70: Here we find out whether it is one of the standard tokens or one of ours. It will branch off if it is standard.

## LINES 80-150 OUR TOKENS

80-100: Here we subtract a number that is one less than our first token value. The result is then transferred to the $x$ register to act as a counter. The value of $x$ is one greater than the position in the table (starting at $\emptyset$ ) but $x$ will be decreased before we start the search.

110-140: We store the start address of our table in what will be our search registers.
150: This is enforced as the last figure in the accumulator, the MSB of our table, will not be zero. We are hardly likely to have a keyword table in the zero page which has many important bASIC locations.

LINES 160-220 CBM TOKENS: This is a duplicate of lines 80-140 except it is for the standard tokens and keyword table.

LINES 230-240: So far we have not used or altered the y register but we store it here in location $\$ 49$ for the LIST routine as that is where it will expect to find it later. We initialise r with $\$$ FF but will increase it before our search so it will start a zero.
It may be worth a note here that we will not alter the values of the search registers, $\$ 22$ and $\$ 23$, as the 64 's keyword table is not longer than

256 bytes and it is unlikely that ours would be. Therefore, incrementing y through its 256 range ( $0-\mathrm{FF}$ ) will serve our purpose. It also saves bytes and time.

LINES 250-260: Every time we read a word from the table we will come here and decrease the $x$ register. If $x$ is zero, then we have found the position one byte before the keyword we want. In that case we branch off to PRINT the keyword.

LINES 270-300: Here we increase our table pointer, the y register, and then load in the next character from the table. Remember that the last character of a keyword is its ASCII form plus $\$ 80$ (128) and this is what we look for. This will set the negative flag in the status register.

The first check is to see if the negative flag is unset signifying a branch back to get the next character. If the negative is set, then the end of the word is found and we branch back to test $x$ to see if we have come far enough. One of these two branches must work as a byte is determined as either negative or positive.

LINES 310-350: We have found our word and now have to print it out. First we increase our pointer to pick up the first character. We load it and test to see if it is the last character. If it is we go to the bASIC ROM to have it printed through the LIST routine. Failing this, we go to another ROM routine to have the character printed. We will return with the same character in the accumulator. As a keyword will not have a byte of zero value, the branch in line 350 is enforced, to get another character to print.

LINE 360-ROM 1: The character was not a token at the beginning so we go back to the LIST routine to have it printed and continue with the listing.

LINE 370-ROM 2: We have here the last character of the keyword in the accumulator. Now we go back to the LIST routine where it will be turned into the proper ASCII value, printed and the listing continued.

## BASIC token dispatch

This is the routine that BASIC uses on finding a token to get the address for the routine. It deals only with command keywords, such as PRINT. It is a subroutine in itself. What we need to do is put in a routine that it goes to first, through the vector.

10
20
30
40
50
60
70

JSR 18073
CMP \#\#CC
BCC ROM
CMP \#HIGHEST COMMAND TOKEN VALUE BCS ROM
JSR DISPATCH
JMP まATEA

| 80 | DISPATCH |
| :--- | :--- |
| 90 | SEC |
| 100 | SBC \#\$CC |
| 110 | ASL |
| 120 | TAY |
| 130 | LDA START OF OUR UECTOR TABLE $+1, Y$ |
| 140 | PHA |
| 150 | LDA START OF OUR UECTOR TABLE,Y |
| 160 | PHA |
| 170 ROM | JMP $\$ 0073$ |
| 180 | JSR $\$ 0079$ |

LINE 10: Get the token from the input buffer or program line through the CHRGET routine.

LINES 20-30: We check to see if it is one of our tokens. If it is not, we branch off to the normal routine in ROM.

LINES 40-50: Now we find out if it is a command or a function token of ours. If it is a function vector, then it is a 'SYNTAXERROR' so we branch to rom to print it.

LINES 60-70: Here we go to our subroutine for dispatch. When the keyword routine has been completed, the program flow will come back here where we shall jump back to BASIC for continuation.

LINES 80-150: We subtract our lowest token value from the value we have. This will give us values of $\emptyset$ upwards. Now as each routine has a two byte address, we must double our 'new' token value to get its proper place in the vector table. The instruction ast does just this by shifting all bits one place left and putting a zero in bit $\emptyset$. This new value is transferred to the $y$ register as a pointer in the table. What we are going to do is to put a new return address on top of the stack (the program counter expects the LSB on top with the MSB underneath). Therefore, we take the second byte of the table first, put it on the stack, and then the first. Remember our vector table is made up of LSB then the msb.

We now Imp to the ChrGet routine to pick up the next byte. The rTS at the end of Chrget will now be to our keyword routine as we have just put its address on the stack. We came to these lines ( $80-150$ ) by a ISR command so its return address was originally on the top of the stack. We then put another address on top of that which was pulled off in CHRGET leaving our original return address once more at the top. At the end of the keyword routine this address will be pulled off and we will return to line $7 \emptyset$ of this routine.

LINES 170-180: Here we go to the normal dispatch routine. This is not the address normally found in the Token Dispatch Vector because we
will miss out the first instruction which is to get the next byte. We go to CHRGOT first not to get the byte we have already got but to set the flags that the ROM routine wants to test.

BASIC function dispatch
This is the routine that will find the routine addresses of function keywords.

| 10 |  | LDA | \#00 |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 20 |  | STA | \$0D |  |  |  |
| 30 |  | JSR | \$0073 |  |  |  |
| 48 |  | CMP | \#\$LOWEST | FUNCTION | TOKEN | Value |
| 50 |  | BCC | ROM |  |  |  |
| 60 |  | CMP | \#\$HIGHEST | T TOKEN VA | Lue |  |
| 70 |  | BCS | ROM |  |  |  |
| 80 |  | JSR | DISPATCH |  |  |  |
| 90 |  | RTS |  |  |  |  |
| 108 | DI SPATCH | SEC |  |  |  |  |
| 110 |  | SBC | \# ${ }^{\text {L }}$ OWEST | FUNCTION | TOKEN | VALUE |
| 120 |  | ASL |  |  |  |  |
| 130 |  | TAY |  |  |  |  |
| 140 |  | LDA | \#\$START O | OUR FUNCT | VECT | TABLE+1,Y |
| 150 |  | PHA |  |  |  |  |
| 160 |  | LDA | \#\$START | QUR FUNCT | VECT | TABLE, Y |
| 170 |  | PHA |  |  |  |  |
| 180 |  | JMP | \$0673 |  |  |  |
| 190 | ROM | JSR | \$0079 |  |  |  |
| 200 |  | JMP | \$AE8D |  |  |  |

This is basically the same as the previous routine. The return addresses to rOM are different, as will be the table address. The first two lines load a location which BASIC uses to decide whether to accept numeric or string data, the latter value would be $\$ 80$.

The other difference is that on return from the function routine we will arrive back at line 90 . The previous routine went back to BASIC for another command, but here we RTS as functions will be performed as part of a command routine and therefore we go back to it.

## 4 Keyboard revisited making use of the wasted keys

On the far right of your keyboard there are four keys that do not really do much, at least at the moment. They are, of course, the function keys. In this chapter we are going to show you how to make use of them. First we thought it a good idea to describe the rom routine in the 64 which services the keyboard. In doing so we will also come across the locations that appertain to the keys.

## The hardware interrupt vector

Every $1 / 60$ th second the computer hands control to an interrupt system. When the microprocessor receives an interrupt signal it will not do anything until the present instruction has been completed. The processor will then save the program counter and the status register. The program counter is then loaded with the contents of locations \$FFFE and $\$$ FFFF. This will start a routine at $\$$ FF48 which saves the register contents on the stack before doing an indirect jump to the vector at $\$ 0314$ and $\$ 0315$.

The interrupt routine found at this vector points to address \$EA31. This KERNAL routine performs several housekeeping operations such as the update of the system clock, but it also scans the keyboard. The key that you press is picked up by the Complex Interface Adapter \#1, and in particular the Data Port b within that chip. From this the value of the key pressed, and shift keys if used, is calculated and stored.

There seems to be some doubt from what we have read about which location the current key value is stored in. The current value is stored in $\$ C B$ (203) and the last in \$C5 (197). This to the BASIC programmer does not make a lot of difference unless the key buffer is full when the key value is not logged except in \$CB. The shift, logo and CTRL keys have the same system, with the current location being \$028D (653) and the last press in \$028E (654).

Having stored your current input it will check to see if it is the same as the last key press. Its next action will depend on whether it is the same, or not. If the same, it will see if it is a repeat function such as the cursor keys or if location $\$ 02 \mathrm{BA}(650)$ has been set for all keys to repeat. Failing this, the value will not be placed into the keyboard buffer. Where the key values are processed it does so by looking up a table to obtain the ASCII code for your key press. This value is placed in the keyboard buffer and its counter updated.

The keyboard buffer is situated at \$0277-\$0280, a size of ten characters. It operates on the system that the first character in will also be the first out. The pointer for the number of characters in the buffer at a particular time is $\$ C 6$ (198). The size of the buffer can be reduced from its initial value of ten by setting register $\$ 0289$ (649).

Earlier we said that every key has a value. These are from 0 to 64 , the latter being no key press. A table of these, and the shift key values are given in the appendices.

Here is a summary of keyboard locations:

| $\$ C B$ | 203 | Current key press. |
| :--- | :--- | :--- |
| $\$ C 5$ | 197 | Last key pressed. |
| $\$ 028 \mathrm{D}$ | 653 | Current shift etc. |
| $\$ 028 \mathrm{E}$ | 654 | Last shift etc. |
| $\$ 028 \mathrm{~A}$ | 650 | Repeat flag: $\$ 80$ all, $\$ 00$ normal. |
| $\$ 0277-\$ 0280$ | $631-640$ | Keyboard buffer. |
| $\$ 0289$ | 649 | Size of keyboard buffer. |
| $\$ C 6$ | 198 | No of chars in buffer. |

## The Function Keys

These keys have values and ASCII codes like any other key. They are:

| Value |  |  |
| :---: | :---: | :---: |
| Function key | (\$CB and \$C5) | ASCII |
| F1 | 4 | 133 |
| F2 | 4 | 137 |
| F3 | 5 | 134 |
| F4 | 5 | 138 |
| F5 | 6 | 135 |
| F6 | 6 | 139 |
| F7 | 3 | 136 |
| F8 | 3 | 140 |

Knowing these values and the locations mentioned earlier, we can make use of the function keys.

## Function keys within a BASIC program

One of the most used BASIC statements for evaluating a key press is the GET function. This function returns the ASCII code for the first key in the keyboard buffer, or the latest key if the buffer is empty. It will not wait for a key press. A BASIC routine could look like this:

100 GET A\$
110 IF A $\$=$ " $[F 1]$ " THEN 1000: REM ACTION ON F1 PRESS
120 REM ACTION IF ANY OTHER KEY PRESSED

This routine will not stop and wait for a key press. It will only branch off to line 1000 if key F1 is pressed at the same time as the GEt statement is actioned or the next character in the keyboard buffer is the ASCII for F1.

We could adapt this so that it will wait until a key is pressed - any key.

$$
100 \text { GET A\$ :IF A } \$==^{\prime \prime} \text { 'THEN } 100
$$

Here line 100 will be repeated until one key is pressed or there is a value in the key buffer that has not been read.

The next thing we could add is a line to clear the input buffer before we GET a character. The easiest way is to set the register for the number of buffer characters to zero.

## 90 POKE 198, 0

At the moment the routine actions on any key. If we wanted it to action on only two keys, say F1 and F7, we would have to alter line 120 to:

```
120 IF A$ <> "[F7]" THEN 90
130 REM ACTION ON F7 PRESSED
```

Now the routine will wait until a key is pressed. Once a key is pressed it goes to 110 to see if it was F 1 and branches if so. Failing that it goes to line 120 where we look to see if it was not $\mathrm{F7}$. On $\mathrm{F7}$ the program will continue its flow. Now lines $90-120$ will keep repeating until either F 1 or 57 is pressed.

The only other alteration we could do is to rid ourselves of the graphic characters in the quotes that represent the function keys. This would make it easier for someone else to read and on a nonCommodore printer the graphic character would not print. This we can do by using the CHRS ( function when checking As. Line 110 would now look like:

## 110 IF A\$ = CHR\$(133) THEN 1000: REM <br> ACTION ON F1 PRESSED

In the GET statement all eight function keys can be tested in the same way, either by changing the character in the quotes or changing the CHR\$ value.

Another way of testing for the keys is by examining the key press registers set in the interrupt routines. From a BASIC programmer's viewpoint it does not really matter whether you test the current or the last key register. The snag with this method is that without checking the shift register only four of the function keys can be detected. On the other hand, by checking the shift register with all its combinations you can have up to 32 function key combinations. Here is a routine that tests for function key F 1 :

90 POKE 198,0:REM CLEAR KEY BUFFER
100 IF PEEK $(203)=4$ THEN 1000 : REM F1 VALUE.
110 REM PROG CONTINUES IF NOT F1.
This will not wait for F 1 . To wait, line 110 will have to be changed to:
110 GOTO 90
We have now set up a loop and the only exit is F 1 being pressed. Now if we wanted to test for $F 2$, the shift flag would have to be introduced. Line 100 could look like this:

## 100 IF PEEK(203)=4 AND PEEK(648)=1 THEN <br> 1000 : REM ACTION ON F2 PRESSED

If you wanted to go to line 1000 on any key, or no key, apart from F2, then the equals sign should be replaced by greater than and less than signs.

## Programming the keys in immediate mode

## Our interrupt routine

The routines that follow will allow you to program the function keys with commands or phrases to be actioned as if you typed them in full, but using only one keystroke.

Most of the routines we have seen to do this operation change the vector address of the Hardware Interrupt Routine in $\$ 0314$ and $\$ 0315$. They alter it to point to their routine, which when finished will return direct to the normal interrupt routine. This course of action has drawbacks. First, it adds to the length of the interrupt, especially if the user's routine has to be completely followed through. Secondly, it means that you have to set up your own registers for checking to see if it was the same action as the last time or not, to avoid auto-repeat. A further drawback is that if you want to use the data assigned to function keys within quotes, it is more difficult to suppress the graphic character that is generated in the quotes mode along with your phrase.

So how are we going to achieve this desirable routine of making the function keys really useful? Earlier we described the interrupt routine and how your key presses are interpreted. What we did not say was that there is a vectored jump within it. This occurs after the value from the Data Port is put into the current key registers but before it is actioned. The vector is held in addresses $\$ 028 \mathrm{~F}$ and $\$ 0290$ ( 655 and 656 ) and is known as the 'Keyboard Table Setup Vector'. If we change the address in this vector to point to a routine of our own we can process the data first. If the data concerns us we can process it jumping back to the normal interrupt routine at a point which misses out the normal key press routine. When the data does not concern us, control will be handed back to the normal flow of things.

The use of the vectors by Commodore has allowed us an easy way to program the keys. This cannot be said of the values that have been assigned to the function keys. It would have been easier if F 1 had a value of 1 and F , of 3 , but this is not the case.

We are going to have 16 programmed function keys. To get this number, you have to use the keys in conjunction with the shift and logo keys as follows:
KEYS F1, F3, F5, F7 - THE KEY ONLY
KEYS F2, F4, F6, F8 - THE KEY + SHIFT
KEYS F9, F11, F13, F15 - THE KEY + LOGO
KEYS F10, F12, F14, F16 - THE KEY + SHIFT + LOGO
This gives us keys in the range of 1 to 16 , but for the routine it is easier to use 0 to 15 . We shall load the data into the keyboard buffer so we are limited to ten characters. We also require a marker for the end of data for each key, which will be a zero, meaning a maximum 11 bytes storage for each. It is easier, and quicker, to use 16 bytes per key. This wastes five bytes but as we are going to store the data in the ram under the basic rom this is unimportant. This will mean the value of the key needs to be multiplied by 16 to get the start of its data. Multiplying by 16 for the low numbers we are using, 0 to 15 , simply involves moving the four lower bits to the four higher bits and filling the lower ones with zeros, four ASL instructions will achieve this. Sixteen bytes of data for the 16 keys will take one page, 256 bytes, exactly.

To summarise:
i) Find out if the key is a function key, yes - continue, no - go to interrupt.
ii) Calculate key number less 1 .
iii) Multiply key number by 16 for table position.
iv) Get data off the data table and store in the key buffer.

ASSEMBLY LISTING

| 9 | $*=\$ 8722$ |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 10 |  | LDY | \$CB | ! | CURRENT | KEY PRESS |
| 20 |  | CPY | \#\$83 | ! | IS IT A | FUNCTION KEY |
| 30 |  | BCC | NORTMAL | ! | NO |  |
| 40 |  | CPY | \#\$07 | ! | IS IT A | FUNCTION KEY |
| 50 |  | BCC | CONT | ! | YES |  |
| 60 | NORMAL | JMP | \$EB48 | ! | NORMAL IN KEY ROUT | NTERRUPT INE |
| 70 | CONT | LDA | \$028D | ! | CURRENT | SHIFT PRESS |
| 80 |  | CPY | \$C5 | ! | IS CURREI | ENT KEY=LAST |
| 90 |  | BNE | CONT2 | ! | NO |  |
| 180 |  | CMP | \$028E | ! | IS CURREN | NT SHIFT=LAST |
| 110 |  | BEQ | NORMAL | ! | KEY AND | SHIFT AS LAST |


| 128 | CONT 2 | STY | \$0:5 | ! | STORE CURRENT KEY <br> IN LAST REGISTER |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 130 |  | STA | \$028E | ! | STORE CURRENT |
|  |  |  |  |  | SHIFT IN LAST REG |
| 140 |  | CPY | \#\$04 | ! | IS IT F1 |
| 150 |  | BEQ | F1+1 | ! | YES |
| 160 |  | CPY | \#\$05 | ! | IS IT F3 |
| 170 |  | BEQ | F3+1 | ! | YES |
| 180 |  | CPY | \#\$06 | ! | IS IT F5 |
| 190 |  | BEQ | F5+1 | $!$ | YES |
| 200 |  | LDY | \# 507 | ! | IT IS F7 |
| 210 | F1 | BIT |  | ! | UALUE FOR F1 |
| 220 | F3 | BIT | \$0340 | ! | UALUE FOR F3 |
| 230 | F5 | BIT | \$05A8 | ! | UALUE FOR F5 |
| 240 |  | CMP | \#\$02 | ! | WHAT SHIFT |
| 250 |  | BCC | NOCHANGE |  | NONE OR SHIFT UALUES CORRECT |
| 260 |  | BEQ | CBM +1 | ! | LOGO KEY |
| 270 |  | LDA | \#\#89 | $!$ | UALUE FOR SHIFT+LOGO |
| 280 | CBM | BIT | \$08A9 | ! | VALUE FOR LOGO |
| 290 | NaCHANGE | STY | \$BB |  |  |
| 300 |  | DEC | \$ BB | $!$ | ONLY WANT NO'S 0-15 |
| 310 |  | CLC |  |  |  |
| 320 |  | ADC | \$BB | $!$ | GET FINAL UALUES |
| 338 |  | ASL | A | ! | MULTIPLY UALUE BY 16 |
| 340 |  | ASL | A |  |  |
| 350 |  | ASL | A |  |  |
| 360 |  | ASL | A |  |  |
| 370 |  | LDY | \#\# A1 | ! | HIGH ADDR KEY TABLE |
| 380 |  | STY | \$15 |  |  |
| 390 |  | LDY | \#\$80 |  |  |
| 400 |  | STY | \$14 | ! | LOU ADDRESS |
| 410 |  | TAY |  | ! | TRANSFER TO Y AS POINTER |
| 420 |  | LDX | \# 000 | ! | COUNTER KEY BUFFER |
| 430 | NEXT | JSR | \$81FB | ! | SWITCH OFF BASIC |
| 440 |  | LDA | (\$14), Y | ! | GET BYTE OF DATA |
| 450 |  | PHA |  | ! | Stare TEMP |
| 460 |  | JSR | \$8202 | ! | SUITTCH ON BASIC |
| 470 |  | PLA |  | ! | get back data |
| 480 |  | BEQ | \$EXIT | ! | END OF DATA |
| 490 |  | CMP | \# $\ddagger 5 \mathrm{~F}$ |  | ARROW FOR RETURN |
| 508 |  | BNE | \$STORE | ! | NO |
| 510 |  | LDA | \#\$0D | ! | LOAD CODE FOR RETURN |
| 520 |  | STA | \$0277, X | ! | STORE IN KEYBGARD |
|  |  |  |  |  | BUFFER |


| 530 | INX | ! INCREASE COUNTER |
| :--- | :--- | :--- |
| 540 | INY | ! INCREASE POINTER |
| 550 | BNE NEXT | ! FORCED-GET NEXT DATA |
| 560 EXIT | STX $\$ C 6$ | ! NO OF CHARS IN KEY |
|  |  | BUFFER |
| 570 | LDA \#\$7F | ! RESET CIA DATA PORT |
| 580 | STA $\$ 0 C 00$ |  |
| 590 | RTS |  |


| 875 F CE 1 | 872 F CONT |
| :--- | :--- |
| 873 B CONT2 | 8791 EXIT |
| 874 F F1 | 8751 F3 |
| 8754 F5 | 8778 NEXT |
| 8762 NOCHANGE | 872 C NORTMAL |
| 878A STORE |  |

LINES 10-60: What we do here is to get into the y register the value of the current key press and see if it is a function key or not. Function keys have values from 3 to 6 inclusive. Line 60 has the normal address of the Keyboard Table Setup Vector and if we do not find a function key this is where we direct the flow.

LINES 70-110: This part of the routine checks to see if the last kEy is the same as the current KEY. If it is, then off to the standard routine to avoid auto-repeat. At this point we have the current key value in the $y$ register and the current shift value in the accumulator.

LINES 120-130: This is part of the housekeeping. We copy the current values we have obtained into the last key registers. This is not only for our routine but also for the normal key interpreting routine.

LINES 140-230: We now take our key value, find which key it is and give a number corresponding to the number on the key itself. The BIT commands will not alter any data at all except for the status flags (which we are not testing here). They allow us to 'hide' an instruction within the address, in these cases loading the y register, saving bytes and branch instructions. For instance, the bIt address in line 210 is \$01a0 which is stored in memory as A001, which is the code for LDY \#\$01.

LINES 240-280: We now do the same for the shift value. If there is no shift or just the standard shift, there will be no need for any alterations so they would branch off in line 250 . The logo key requires the value of 8 ( $1+8$ giving key 9 and so on) and both shifts 9 . We again do this using the BIT function.

LINES 290-360: Here we subtract one from the key value and then add the result to the shift value, ending up with a value between 0 and 15.

This total will be in the accumulator which is then increased 16 times by the four ast instructions. We now end up with a value between $\square$ and 15 which will be the pointer to the data for that particular key.

LINES 370-420: The start position of the data table is put in registers $\$ 14$ and $\$ 15$. We also transfer the pointer in the accumulator to the $y$ register. Lastly, we initialize the $x$ register to zero to use as a counter to the number of characters we put in the keyboard buffer.

LINES 430-550: At last we can get our data and use it. Earlier we said that we were going to put our data in the ram under the basic rom. To read it back, we have to 'remove' the rom to access it. This we do by a call to an earlier routine in the Utility which you will come to later. Now we pick up a byte of data and put it on the stack for temporary safe keeping, as we require the accumulator for re-enabling the BASIC ROM. With the rom back, and having recovered the byte, we have two checks before storing it. The first in line 480 is to see if the byte is zero, signifying that all the relevant data has been collected and we can finish up. The second is a check for the 'left arrow', which signifies the user wants a return to be included (more of this in programming the keys). If this succeeds, we will change the byte to the ASCII code for return.

The data is stored in the keyboard buffer starting at the beginning and working upwards - it will be removed in the same order. We do not need to check for overflow as we are only allowed ten characters to be programmed (see next section). Therefore, the zero, which is not stored, cannot be later than the eleventh byte.

Having stored our byte, the two registers are increased by one and we branch back to get a further byte. The branch is enforced as we will not increase $Y$ enough to return it to a zero. The highest value $Y$ will achieve is $\$ F B$ ( 251 dec ).

LINES 550-590: The end is near. Having stored all our data, the $x$ register will hold a number equal to the total number of characters we put into the buffer. This is put into the register denoting how many characters are in the buffer and the operating system will only take that many off. The following two instructions are again housekeeping in that we reset the data port for collection of the next press. A return follows, but didn't we come by a JMP? This is true, but the whole key routine is entered by a JSR where the vectored jump is found. We do not now need the use of the normal key interpreting routine so we can go straight back to the main interrupt.

## Key

COMMAND SNYTAX
KEY
Displays the current data assigned to the keys in a form which can be amended.

KEY[number between 1 and 15], "[data]"
Assign data to a particular key. If a return is required, type $\mathrm{a}^{\prime \prime} \longleftarrow$ ' to signify this. Quotes cannot be used as data. A typical command could look like this:

$$
\text { KEY 7, "LIST } \leftarrow "
$$

Here is a full list of the key numbers and how to achieve them:

$$
\begin{array}{lll}
\text { KEY } & 1 & \text { - F1 ONLY } \\
\text { KEY } & 2 & \text { - F1 + SHIFT } \\
\text { KEY } & 3 & \text { - F3 ONLY } \\
\text { KEY } & 4 & - \text { F3 + SHIFT } \\
\text { KEY } & 5 & \text { - F5 ONLY } \\
\text { KEY } & 6 & \text { - F5 + SHIFT } \\
\text { KEY } & 7 & \text { - F7 ONLY } \\
\text { KEY } & 8 & \text { - F8 + SHIFT } \\
\text { KEY } & 9 & \text { - F1 + LOGO KEY } \\
\text { KEY } & 10 & \text { - F1 + SHIFT + LOGO KEY } \\
\text { KEY } & 11 & \text { - F3 + LOGO KEY } \\
\text { KEY } & 12 & \text { - F3 + SHIFT + LOGO KEY } \\
\text { KEY } & 13 & \text { - F5 + LOGO KEY } \\
\text { KEY } & 14 & \text {-F5 + SHIFT + LOGO KEY } \\
\text { KEY } & 15 & \text { - F7 + LOGO KEY } \\
\text { KEY } & 16 & \text { - F7 + SHIFT + LOGO KEY }
\end{array}
$$

KEY $0 . .$. will generate a SYNTAX ERROR. We had thought about using this as a way of turning off the key routines, but decided on a separate command. This makes it more of a conscious decision rather than a typing error. The command will be off, which is discussed later.

We have seen that we can make use of the four 'mystery' keys by getting data output on their use and in fact having 16 keys when used with the shift and logo keys. Now we have a routine to program the data, in which the user can decide what data to apply. This operation is acted upon through the keyword KEY.

KEY will perform three functions. It will 'switch' on the keys if they are off. This is performed in both of the following options. The choices are to program a key or to display the data applied to all the keys, which can then be amended on the display.

As we have said, there are two routines included in this. There is one routine to program individual keys and one routine to display the data
assigned to all keys. The latter is very similar to the interrupt routine discussed earlier except that the data goes to the screen rather than a buffer. The former in many ways is the reverse: we take data from a buffer - the input buffer - and put it in a tabie.

ASSEMBLY LISTING

|  | *=\$864D |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 10 |  | LDA | \$805B | ! | CHECK IF INTERRUPT SET FOR KEYS |
| 20 |  | CMP | \# $\$ 87$ |  |  |
| 30 |  | BEQ | START | $!$ | YES |
| 40 |  | LDA | \#\$87 |  |  |
| 50 |  | STA | \$885B |  |  |
| 60 |  | LDA | \# $\$ 22$ |  |  |
| 70 |  | STA | \$8056 |  |  |
| 80 |  | JSR | \$8054 | ! | SET INTERRUPT |
| 90 | START | JSR | \$0079 | ! | GET LAST BYTE AGAIN |
| 100 |  | BEQ | DI SPLAY |  |  |
| 110 |  | JSR | \$81F5 | ! | GET PARAMETER |
| 120 |  | JSR | \$AEFD | ! | CHECK FOR COMMA |
| 130 |  | LDA | \$14 |  |  |
| 140 |  | BEQ | SYNTAX | ! | NO KEY0 |
| 150 |  | CMP | \#\$11 | ! | HIGHEST KEY IS 16 |
| 160 |  | BCS | SYNTAX |  |  |
| 170 |  | DEC: | \$14 |  |  |
| 180 |  | LDA | \$14 | $!$ | set to calculate POINTER |
| 198 |  | ASL | $A$ | ! | CALCULATING POINTER |
| 208 |  | ASL | A |  |  |
| 210 |  | ASL | $A$ |  |  |
| 220 |  | ASL | A |  |  |
| 230 |  | TAY |  |  |  |
| 240 |  | L[AA | \#\# ${ }^{\text {A }}$ | ! | H1 ADD FOR KEY TABLE |
| 250 |  | STA | \$15 |  |  |
| 268 |  | LDA | \# $0^{0} 0$ | ! | SET LO ADD FOR KEYS |
| 270 |  | STA | \$14 |  |  |
| 280 |  | LDX | \$8A | $!$ | COUNTER MAX NO OF CHARS |
| 290 |  | JSF | \$0079 | ! | GET LAST BYTE AGAIN |
| 300 |  | CMP | \# $\mathbf{\$ 2}^{2}$ | ! | IS IT A QuOTE |
| 310 |  | BEQ | CONT 2 | ! | YES |
| 320 | SINTAX | JMP | \$AF88 | ! | PRINT SYNTAX ERROR |
| 330 | CONT2 | JSR | \$0073 | ! | GET NEXT BYTE |
| 340 |  | BEQ | 2ER0 | ! | END OF DATA INPUT |
| 350 |  | CMP | \# ${ }^{\text {2 }} 2$ |  | IS IT A QUOTE |


| 360 |  | BEQ | ZER0 |  | END OF DATA INPUT |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 370 |  | STA | (\$14), Y |  | STORE DATA IN TABLE |
| 380 |  | INY |  |  | Inc table pointer |
| 396 |  | DEX |  |  | dex char count |
| 400 |  | BNE | CONT2 | ! | IF ZERO MAX NO CHARS REMAINDER I GNORED |
| 410 | ZERO | LDA | \#\#00 | ! | END OF WORD MARKER |
| 420 |  | STA | (\$14), Y |  |  |
| 430 |  | JSR | \$8073 | ! | GET NEXT BYTE |
| 448 |  | RTS |  |  | FINISHED |
| 450 | display | LDX | \#\$00 |  | SET COUNTER |
| 460 |  | STX | \$5F |  |  |
| 470 |  | INX |  |  |  |
| 480 |  | LDA | \#\$20 |  | space as no ten's DIGIT |
| 490 |  | STA | \$22 |  |  |
| 500 |  | LDA | \#\$31 | ! | ASCII FOR ONE |
| 510 |  | STA | \$23 |  |  |
| 520 |  | LDA | \# 00 | ! | lo byte of data table |
| 530 |  | STA | \$14 |  |  |
| 540 |  | LDA | \#\#A1 | ! | hI byte of data table |
| 550 |  | STA | \$15 |  |  |
| 560 | PD1 | JSR | PRINT |  |  |
| 570 |  | INC. | \$23 | $!$ | INCREASE NUMERAL |
| 580 |  | INC | \$5F | ! | INCREASE KEY COUNT |
| 590 |  | INX |  |  |  |
| 600 |  | CPX | \# 0 8A | $!$ | have we done keysi-9 |
| 610 |  | BCC | PD1 | ! | No |
| 620 |  | LDA | \#\#31 |  | now have a ten digit |
| 630 |  | STA | \$22 |  |  |
| 640 |  | LDA | \# 30 |  |  |
| 650 |  | STA | \$23 |  |  |
| 660 | PD2 | JSR | FRINT |  |  |
| 676 |  | INC | \$23 |  |  |
| 680 |  | INC | \$5F |  |  |
| 690 |  | INX |  |  |  |
| 700 |  | CPX | \#\#11 | $!$ | have we done 16 |
| 710 |  | BCC | PD2 |  | NO |
| 720 |  | RTS |  | ! | YES |
| 730 | PRINT | LDY | \#\$05 | ! | COUNTER |
| 740 | NEXTA | LDA | PDATA, Y |  | PRINT " KEY |
| 750 |  | JSR | \$FFD2 |  |  |
| 760 |  | DEY |  |  |  |
| 770 |  | BNE | nexta |  |  |


| 780 |  | LDA |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 790 |  | JSR | \$FFD2 | ! | PRINT TEN'S |
|  |  |  |  |  | NUMERAL OR SPACE |
| 800 |  | LDA | \$23 | ! | PRINT LOW NUMERAL |
| 810 |  | JSR | \$FFD2 |  |  |
| 820 |  | LDA | \#\$2C |  |  |
| 830 |  | JSR | \$FFD2 | $!$ | PRINT COMMA |
| 840 |  | LDA | \#\$22 |  |  |
| 350 |  | JSR | \$FFD2 | ! | PRINT QUOTE |
| 860 |  | LDA | \$5F | ! | CALC TABLE POINTER |
| 870 |  | ASL | A |  |  |
| 880 |  | ASL | A |  |  |
| 890 |  | ASL | A |  |  |
| 900 |  | ASL | $A$ |  |  |
| 910 | CONT | TAY |  | ! | PUT POINTER IN Y |
| 920 | NEXT | JSR | \$81FB | ! | SWITCH OFF BASIC |
| 930 |  | LDA | (\$14), Y | $!$ | GET CHAR OFF TABLE |
| 940 |  | PHA |  | $!$ | TEMP STORE |
| 950 |  | JSR | \$8202 | ! | SWITCH ON BASIC |
| 960 |  | PLA |  | ! | RETRIEUE CHAR |
| 970 |  | BEQ | ExIT | ! | FOUND END OF WORD |
| 980 |  | JSR | \$FFD2 | ! | PRINT CHAR |
| 990 |  | INY |  |  |  |
| 1080 |  | BNE | NEXT | 1 | ENFORCED |
| 1810 | EXIT | LDA | \# ${ }^{\text {2 }} 2$ |  |  |
| 1020 |  | JSR | \$FFD2 |  | PRINT A QUOTE |
| 1038 |  | RTS |  |  |  |
| 1040 | PDATA | BYT | \$20, 'Y, |  | , $K$, \$20, $\$ 0 \mathrm{D}$ |
| 8691 | Cont 2 |  | 8648 DIS | PL |  |
| 8716 | EXIT |  | 8704 NEX |  |  |
| 86E0 | NEXTA |  | $86 \mathrm{BD} \mathrm{PD1}$ |  |  |
| 86D1 | PD2 |  | 871 C PDA | TA |  |
| 860E | PRINT |  | 8661 STA | RT |  |
| 868 E | SINTAX |  | 86A0 2ER |  |  |

LINES 10-80: Earlier in the UTILITY, a routine will exist that is used when the extension is initialized or when stop/Restore is used. This sets the Keyboard Table Setup Vector to where we want it to point to. These addresses can be changed by the off command. Here we look to see if the high byte of the address is pointing to our interrupt routine. If not, we change the address in the setting routine to point to our interrupt routine and then call the setting routine to initialize.
LINES 90-270: A call first to the CHRGOT routine to get the byte after the

KEY token. This is necessary as we have used the accumulator and so overwritten the byte. If the byte has set the zero flag, then there are no further parameters and a display of the key data is required. The program in that case branches to the display which starts at line 450.

Knowing we have got some parameters, off we go to our 'GET PARAMETER' routine (Chapter 6) and to a ROM coding to see if the byte after the key number is a comma. This coding will not only update the CHRGET address but will generate a SYNTAX ERROR if a comma is not found.

The parameter we want is now held in location $\$ 14$ - the key number. This value is put in the accumulator and checked for two things. If it is zero or greater than sixteen, it is out of bounds, so an error message is required, and therefore we branch off to get this printed. As in the interrupt routine it is easier to work in numbers $\emptyset$ to 15 rather than 1 to 16 so we decrease the value in $\$ 14$ by one and then reload back into the accumulator.

To get the pointer to the required position in the data table, the number is multiplied by 16 with the ast instructions. The $y$ register will be the pointer so the value is transferred to it. Next we load two registers with the address of the data table start. Now we are in a position to get, and store, the data.

LINES 280-440: The data generated by using the function keys will be placed in the keyboard buffer. This buffer is only ten characters in length so we have to limit the input to that number. This is achieved by setting the $x$ register to ten (\$0A). We said earlier that the comma check updates the Chrget address so a call to the ChRGOT routine will get the next byte we want. This should be a quote; if not a syntax error is generated (remember that ChRGET skips spaces).

Now to get the data and store it. To get the data we make use of CHRGET. If the zero flag is set, the end of the command has been reached with either a colon or a zero placed by the BASIC input routine. The second quote is checked which also signifies the end of data input. If any of these are found, we branch off to end the routine at line 410. We can now store our data in the table under the BASIC ROM. We do not have to disable the ROM as you cannot store data in ROM so it is automatically stored in the RAM underneath. We increase the y register which points to the table position. We decrease $x$ which checks for overflow of data. If $x$ reaches zero, the maximum number of characters has been stored. The flow only branches back to get the next byte if $x$ is greater than zero.

To finish off, we store a zero after the last byte of data. This will help when retrieving the data to signify all data has been gathered.

We do another visit to CHRGET to get the next byte as BASIC expects this. This will cause a syntax error if you have input more than ten characters of data though the first ten bytes will have been logged.

The RTS returns us to BASIC for further operations.

## DISPLAYING THE KEY DATA

LINES 450-550: These instructions set up the registers used in the display itself. The $x$ register is again used as a counter. Location $\$ 5 \mathrm{~F}$ will have the value of the key number less one and will be used to calculate the pointer for data collection. Locations $\$ 22$ and $\$ 23$ hold the ASCII values of the key number. As keys up to and including 9 have only one digit location $\$ 22$ is loaded with the ASCII for a space character. $\$ 23$ starts with the ASCII for 1 and will be incremented. Finally, we load up the address of the start of the data table into registers $\$ 14$ and $\$ 15$.

LINES 560-610: Call the coding to print key data for keys 1 to 9. After calling, the ASCII value in $\$ 23$ is increased along with the key number register $\$ 55$. Register $x$ is also increased and checked to see if it has reached \$0A (ten). If so, we would have to reset the ASCII numbers before printing further data. If $x$ has not reached this value, we branch back to call the print coding for the next key.

LINES 620-720: First we reset locations $\$ 22$ and $\$ 23$. Key numbers from 10 to 16 have to be displayed so we have two digit numbers, the first always being one. Therefore, $\$ 22$ is loaded with $\$ 31$, the ASCII for one. The other is initialized to zero in ASCII format. We now continue to print out the key data, incrementing $\$ 5 \mathrm{~F}$, and x each time, until x reaches the value of 17 just after being incremented. This value of $x$ signifies we have finished the display so we exit from the routine and hand control back to BASIC.

## The Print Routine to Display the Key Data

This part of the command is entered 16 times in total to print the data to the screen. The value for calculating the pointer, held in $\$ 5 \mathrm{~F}$, is set before these lines are implemented, as are the ASCII values of the key number. We use the KERNAL routine at \$FFD2 to print a character to the screen. The data is printed out in the same format as it was entered. This is done so that it can be changed, just like normal screen editing, if required.

LINES 730-770: The start of every key display line will be the same. These lines will print this from the area of data at the end of the routine (Line 140). We start with a return so it starts on a new line, then a space to give better clarity if the border and screen are different colours, especially if the border and text colours are the same. key is printed next, followed by a space for presentation.

LINES 780-850: The key number is printed, followed by a comma and the first set of quotes.

LINES 860-910: The key number, less one, is taken from $\$ 5 \mathrm{~F}$ and
increased 16 times with the now familiar four ASL instructions. The result is transferred to the $\gamma$ register for the data pointer.

LINES 920-1030: Get the key data. First we switch off BASIC to get the data underneath. After returning the BASIC, we print the data as long as it is not the 'end of data' zero. Printing finished, we update the pointer and go back to get the next byte. When the zero is found we exit and print the closing quote. Then it's back to the main key display routine.

## OFF - Turn off the keys

COMMAND SYNTAX
OFF
There are no parameters with this command.
If you want to use the function keys within a program simply as keys, you will want to be able to disable the programming they have been given. The command that enables you to do this is OfF. All we do is to alter the addresses in the routine that sets the Keyboard Table Setup Vector back to its normal address. Once changed, we call the routine to change them in the BASIC work area. Do not forget that they can be re-enabled with any KEY command.

ASSEMBLY LISTING

* $=$ \$ 8799

| 10 | LDA $\# \$ 48$ |  |
| :--- | :--- | :--- |
| 20 | STA $\$ 8056$ | ! CHANGE LON ADD IN |
|  |  | SETTING ROUTINE |
| 30 | LDA \#\$EB |  |
| 40 | STA $\$ 805 B$ | ! CHANGE HIGH ADD IN |
|  |  | SETTING ROUTINE <br> 50 |
| 60 | JSR $\$ 8054$ | ! CALL SETTING ROUTINE |

Stand alone programmable function keys
Perhaps this chapter would have been better located between Chapters 6 and 7. It was difficult to decide on its position as it also uses information from both Chapters 2 and 3, but will not work as it stands.

To provide programmable function keys without using the keyword enable routine, the 'get parameter' and 'switch off BASIC' routines have to be copied from Chapter 6 . The whole routine may then be relocated and the actions of KEY and OFF performed using SYS commands.

## 5 Utilities in basic

## General

This chapter includes many of the utilities in the form of BASIC subroutines and programs. You do not really need any of that which follows if you load-up the UTility each time. In time we suspect that the simple routines contained here will not only prove useful, but will also give you plenty of ideas of your own.

A number of the utilities require that you generate an ASCII file of a program on tape or disk. This produces a file in the same format as would be received at a printer or the screen itself. The resulting sequential file contains the program in 'un-tokenized' form. To do this, output must be directed to the desired device with an OPEN and CMD sequence. For a tape this is:

OPEN 1,1,1,"PROGRAM": CMD1:LIST[from - tol
PRINT\#1:CLOSE 1
and for a disk:
OPEN 2,8,2,"PROGRAM,S,W":CMD2:LIST[from - to] PRINT\#2:CLOSE 2

Most of the utilities given here are in the form of subroutines and have been numbered in the 60000 s to allow them to be easily added on to your own programs as and when appropriate. They may be included in whole, or in part, by a suitable merge or append technique. You may wish to combine a number of them together to form useful modules which in the future may save many hours of repetitive work. This you can easily do by using the mini-renumber and merge programs given. Many of the routines can be extended, but they have deliberately been kept as short as possible. Always try to adopt a 'house' format to simplify the creation of future programs. This may only be a simple line numbering sequence where: the working part of your program lies between lines 100 and 9999; the specific subroutines lie between 10000 and 50000 ; and your library routines are from 50000 on.

The information upon which much of the following is based is contained in Chapter 1 and we refer you to that chapter. The utilities that follow are arranged in alphabetical order.

## Keyword - APPEND

## Append 1

Function: To append two BASIC programs in memory (nose-to-tail)
In the past, whenever you have loaded a program, it has erased the one currently in memory. This need not be the case. BASIC can start at any address in memory and need not always be the default of 2049 ( $\$ 0801$ ). The pointer (TXTTAB) to tell the 64 where BASIC begins is held in RAM and can therefore be changed. It is even possible to have two BASIC programs resident in memory concurrently by changing the necessary zero page pointers, though only one could be running at any time. We can manipulate these pointers to allow us to append one program to another.

With a program in memory change TXTTAB to point to its end (VAR-TAB-2) by:

$$
\mathrm{A}=\mathrm{PEEK}(45): \text { POKE } 43, \mathrm{~A}-2: \text { POKE } 44, \operatorname{PEEK}(46)+\mathrm{A}<2
$$

The program to be appended will now be loaded at the end of the existing one. Resetting the start of BASIC will make the 64 see both programs as one by:

POKE 43,1:POKE 44,8 (assumes original start was 2049/\$0801)
The resulting program may then be edited or saved in the usual way. Many texts say the appended program should have line numbers higher than the original. This is not essential, but some confusion can result if this is not so. Try appending when the second program does not have higher line numbers and see.

The combined program will run correctly until a coto or gosub references a line which occurs twice. By virtue of the way these commands work, the branch will always be taken to the first occurrence of a line.

## Append 2

Function: To append two programs on disk (BASIC or machine code)
Program files on disk store an image of the memory which the program occupied. The first two bytes record the load address and the last byte is a zero to mark the end of file. They can, however, be read and written in a sequential manner. This allows us to append files in much the same way as we did above, but this time performing the operation solely on disk. The following program will append two programs which will load at the address of the first:

## LINE ACTION

130 Open up＇Program＇files for read and write．
140 Read first program and make a byte by byte copy
TO in the combined file．Skip the terminating zero byte
180 and jump to read the second program．
200 Read the load address and discard it．
210－Copy the remainder through to produce the combined file．

```
100 INPUT"FRONT PROGRAM";F毒
110 INPUT"[2SPC]END PROGRAM";E#
120 INPUT"FINAL PROGRAM";R直
130 OFEN 2,8,2,Fक+",P,R":OPEN 3,8,3,Rक+"
,P,W"
140 GET#2,Aま
150 B&=A&:GET#2,A&:IF ST AND 64 GOTO 180
160 IF A =="" THEN A ==CHRक(0)
170 FRINT#3,B$;:GOTO 150
180 CLOSE 2
190 OPEN 2,8,2,E$+",P,R"
200 GET#2,A$:GET#2,A直
210 GET#2,A&:IF ST AND 64 GOTO 240
220 IF A$="" THEN A ==CHRक(0)
230 PRINT#3,A$;:GOTO 210
240 PRINT#3,CHR$(婁;
250 CLOSE 3:CLOSE 2
```

Append 3
Function：To reopen an existing closed sequential file on disk and continue writing data from the previous end of file．

This is a standard disk command which is not made clear in the disk manual．Its format is：

OPEN 2，8，2，＂TEST，A＂

## Subroutine keyword－AUTO NUMBER

Function：To automatically generate line numbers as code is entered．
Initiation：RUN 60000

This allows the start line and increment to be set．The line number is printed，followed by any characters typed．When RETURN is pressed the program enters the line，resets the line number variables（as an edit destroys all variables）and reruns itself by forcing two RETURNS into the keyboard buffer．As written，the program will not accept any line not followed by BASIC code（equivalent of delete line）．

## LINE ACTION

60010 Position cursor to 3rd line down and print line number in black．
60020 Generate a flashing cursor－not normally present on a GET
60040 Watch out for null lines
60060 Print line to－GOTO 60010，reset variables，restart program and move to home．
60070 Set ndx for two characters in keyboard buffer．Put two returns in k／b buffer（KEYD）．On END KEYD will be emptied and the returns will enter the line and execute line from 60060

```
60000 INPUT "START[4SPC]";LN: INPUT "INC
REMENT "; I%
60010 B$="n:PRINT CHR$(147);CHR$(17);CHR
$(17);CHRक(17);CHR$(144);LN;CHR$(154);
60020 POKE 204,0:POKE 207,0
60030 GET A$:IF A$="" GOTO 60020
60040 PRINT Aक;:IF Bक="" AND ASC(Aま)=13
GOTO 60010
60050 B$=A生:IF ASC(A末)<>13 GOTO 60020
60060 PRINT "LN=";LN+I%;":I%=";I%;":GOTO
    60010";CHRक(19)
60070 POKE 198,2:POKE 631,13:POKE 632,13
: END
```

The version below is a little more flexible．It will not only delete an existing line if RETURN is pressed after its number，but also allows you to change the printed line number to any value．Subsequent line numbers will increment from the new value until it is again changed． The main difference is the addition of the code to evaluate the current line number（（60060 and 60070$)$ ．This is done by reading from the start of the fourth screen line until a non－numeric code is encountered and reassigning the line number variable＇ $\mathrm{L} N$＇．

```
60000 INPUT "START[4SPC]";LN: INPUT "INC
REMENT";I%
60010 PRINT "[CLS][3CD][BLK]";MID&(STR&(
LN),2);"[L BLU]";
60020 POKE 204,0:POKE 207,0
60030 GET A$:IF A$="" GOTO 60020
60040 IF ASC(A$)<>13 THEN PRINT A$;:GOTO
    6 0 0 2 0
60050 PRINT:B$=" ":I=1143
60060 I=I +1 :IF PEEK(I)>47 AND PEEK(I) <58
    THEN B手=Bま+CHRक(PEEK(I)):GOTO 60060
```

```
60070 LN=UAL(B&):PRINT "LN=";LN+I%;":I%=
";I%;":GOTO 60010[HOM]";
60080 POKE 198,3:POKE 631,13:POKE 632,13
:POKE633,13:END
```

Program keyword－BASES
Function：To convert hex to decimal，binary to decimal and vice versa
This program contains four useful inter－base conversion subroutines． The hex to decimal is most useful if you wish to use hex rather than decimal values in the DATA statements for a machine code BASIC loader． The Programmer＇s Reference Guide，Chapter 3 uses binary patterns for the sprite data in the＇BALLOON＇program but pictorial data is also enlightening when setting up user－defined characters and makes for easier editing．

No explanation is given as the program is easy to follow．

```
100 PRINT"1 HEX/DEC":PRINT"2 DEC/HEX":PR
INT"3 BIM/DEC":PRINT"4 DEC/BIN"
110 PRINT:INPUT"SELECT ";N
120 ON N GOSUB 150,240,330,400
130 GOT0100
140 ON N GOSUB 150,240
150 PRINT: INPUT "HEX[4SPC]";A害
160 IF LEN(Aま)<4 THEN A$=LEFT$("0000"+A$
,4-LEN(A事))+A事
170 A=ASC(A末)-48
180 B=ASC(MID本(A* , 2,1))-48
190 C=ASC(MID*(A*,3,1))-48
200 D=ASC(MID*(A青,4,1))-48
210E=256*(16*(A+7*(A)9))+B+7*(B>9))+16*
(C+7*(C)9))+D+7*(D)99)
220 PRINT:PRINT"覀 ";A申;" = D";E:PRINT
230 RETURN
240 PRINT:INPUT"DEC[4SPC]";G:A=INT(G/256
):B=G-A*256:IF G<0 OR G>65535 GOTO 240
250 C=INT (A/16):D=A-16*C
260 C = =CHR&(48+C):IF C>9 THEN C = =CHR& (C+
55.
```



```
55)
280 E=INT(B/16):F=B-16*E
290 Eq=CHR多(48+E):IF E>9 THEN Eq=CHR京(E+
55.
```

```
300 F&=CHR本(48+F):IF F>9 THEN F$=CHR末(F+
55)
```



```
    $ ";A#:PRINT
320 RETURN
330 PRINT:INPUT"BIN[4SPC]";A末
340 A=0:Aक=RIGHTक\"0000000000000000"+A立,
16)
350 FOR I=16 TO 1 STEP - 1
360 B$=MID$(\hat{A}\ddagger,I,1):IF B$="1" THEN A=A+2
*(16-I)
370 NEXT I
380 PRINT:PRINT"B ";LEFT妻(A争,8);" ";RIGH
T本(Aま, 8);" = D";A:PRINT
390 RETURN
400 PRINT:INPUT"DEC[4SPC]";A:IF A>65535
OR A<0 GOTO 400
410 B立="":D=A:FOR I=15 TO 0 STEP -1
420 B=INT (A/2^I):IF B=1 THEN B B=B$+" 1":G
OTO 440
430 B$=Bま+"0"
440 A=A-B*2^1:NEXT I
450 PRINT:PRINT"D";D;"= B ";LEFT我(B#,8);
" ";RIGHT&(Bも,8):PRINT
460 RETURN
```


## Program keyword－DATALINES

Function：To generate BASIC data statements for machine code programs．

Once again the keyboard buffer is used to generate program lines．This time there are more variables in use than would conveniently fit on a single assign line so they have been put＇out of the way＇in the cassette buffer．Only variables in the normal BASIC variable storage area are lost by an edit．The resulting data values are generated to the nearest ten bytes．

## LINE ACTION

60000－Data input．
60060－POKE values to tBuFFr．
60090 Recycle from here．Re－read next line number，
60100 step，
60110 start address，
60120 end address for current line，

60130 and end address of program. If finished stop program.
60140- Print line number, DATA, the values and GOTO 60090.
60210- Increment line number, address, and set up k/b ready for end.

```
60000 INPUT"START ADDRESS";B
60010 INPUT"END ADDRESS[2SPC]";E
60020 F=B:L=F+10
60030 INPUT"START LINE[ 3SPC]";S
60040 INPUT"LINE INC[5SPC]";T
60050 PRINT"[4CD]"
60060 POKE831,INT(E/256)
60070 POKE832,E-INT(E/256)*256
60080 POKE828,T:GOT060160
60090 S=PEEK(826)*256+PEEK(827)
60100 T=PEEK(828)
60110 L=PEEK(829)*256+PEEK(830)
60120 E=PEEK(831)*256+PEEK(832)
60130 IFL>=EGOT060270
60140 F=L+1:L=L+10
60150 PRINT"[CU][145PC]"
60160 PRINTS;
60170 PRINT"DATA";
60180 FORP=FTOL:PRINTPEEK(P);"[CL],";:NE
XTP
60190 PRINT"[CL][3SPC]"
60200 PRINT"GOT060090[4CU]";
60210 POKE198,2:POKE631,13:POKE632,13
60220 S=S+T
6 0 2 3 0 ~ P O K E 8 2 6 , I N T ~ ( S / 2 5 6 ) ~
60240 POKE827,S-INT (S/256)*256
60250 POKE829,INT(L/256)
60260 POKE830,L-INT (L/256)*256:END
6 0 2 7 0 ~ S T O P
```


## Subroutine keyword - DELETE

Function: To remove unwanted program lines en masse
Two delete routines follow. Both use the link address and line number storage at the start of a BASIC line during execution to perform the deletion. The first deletes line numbers as they are encountered whereas the second only deletes one line as the final step in the process. The first line of each routine reads tхTTAB to find out the current start of BASIC.

Delete 1
This routine deletes lines using the all-too-familiar keyboard sequence and as such requires no explanation.

```
60000 TX=PEEK(43)+PEEK(44)*256
60010 INPUT"DELETE FROM";LL:M=256:INPUT"
[7SPC]TO[2SPC]";UL
60020 IF PEEK(TX+2) +PEEK(TX + 3)*M<LLTHENT
X=PEEK(TX) +PEEK(TX+1) *M:GOTO 60020
60030 POKE 828,UL-INT (UL/M)*256:POKE 829
,UL/M:GOTO 60050
60040 M=256:TX=PEEK(830)+PEEK(831)*M:UL=
PEEK(828)+PEEK(829)*M
60050 IF PEEK(TX+2)+PEEK(TX+3)*M)UL OR P
EEK(TX)+PEEK(TX+1)*M=0 THEN END
60060 PRINT "[CLS][3CD]";PEEK(TX+2)+PEEK
(TX+3)*M:PRINT" GOTO 60040[HOM]"
60070 POKE830,TX-INT (TX/M) *M:POKE831,TX/
M:POKE198,2:POKE631,13:POKE632,13:END
```


## Delete 2

```
This is, perhaps, a more refined way to carry out the task. It takes fullest advantage of the way programs are stored in RAM and in particular the function of link addresses. The routine scans the line numbers until the start of the block to be removed is found. It records the address of this link address and then continues to scan for the end line number for the delete. Once a line number equal or greater is found, this link address is substituted at the start of the block link. One very large line has thus been created in memory. A simple keyboard program is then used to remove the start line and all others go with it. This is without doubt a lot faster than the first method, but has the disadvantage that you cannot see the lines as they go.
```

```
60000 TX=PEEK(43)+PEEK(44)*256
60010 INPUT"DELETE FROM";LL:INPUT"[7SPC]
TO[2SPC]";UL
60020 L=PEEK(TX+2) +PEEK(TX+3)*256
60030 IF L<LL THEN TX=PEEK(TX)+PEEK(TX+1
)*256:GOTO 60020
60040 IF L=0 THEN PRINT"LOWER LIMIT";LL;
"NOT FOUND":END
60050 LL=L:D=TX
60060 L=PEEK(TX+2) +PEEK(TX+3)*256
60070 IF L=0 THEN FRINT"UPPER LIMIT";UL;
"NOT FOUND":END
60080 IF L<UL THEN TX=PEEK(TX)+PEEK(TX+1
)*256:GOT0 60060
60090 POKE D,PEEK(TX):POKE D+1,PEEK(TX+1
)
60100 PRINT"[CLS][3CD]";LL;"[HOM]":POKE
198,1:POKE 631,13:END
```


## Subroutine keyword - DUMP

Function: To display the current values of all simple numeric, string and function variables

Initiation: Туре сото 60000
This routine will display the values of all simple variables in use at the time of calling. The variables will be displayed in the order in which they were created by the program. The routine will not handle arrays nor will it work if editing has been carried out prior to its being called (simply because all variable pointers will be reset to the end of program). It also displays the values of the variables it uses $-s v, v \$$, and so on. As these are the last variables to be created they will be the final ones to be displayed. Output may be directed to a printer by a simple:

OPEN 4,4:CMD 4:GOTO 60000
The display may be stopped by holding down any key and will resume on the release of that key. Pressing the stop key will 'break' into the program and allow you to use the cursor keys to move up and change values. If you resume program execution with a GOTO, then the amended values will be used. A simple CONT would re-enter the dump subroutine at the break and dump any remaining variables.

The routine makes extensive use of the information contained in Chapter 1 on the storage of BASIC variables. Remember the first two bytes are the variable name adjusted for its type. The following is a description of the routine:

## LINE ACTION

60030 Read the current value of vartab.
60040 Do the same for ARYtab.
60050 Default values.
60055 If equal then no simple variables, edit used, or finished. If not equal more variables exist so continue.
60060 Read the seven bytes used for variable.
60070 Determine the type from the two name bytes and Gото the to appropriate subroutine, these being real, integer, string or func60100 tion. The name bytes must be changed back to their unmodified ASCII values by the subtraction of 128 , as necessary, and ' $\%$ ' or '\$' suffixes printed where required.
60105 Pause if key held down ( $64=$ no key at sfDx. Note this is the current key not LSTX as in the Programmer's Reference Guide).
60110 Increment 7 bytes to next variable and recycle.
61000 Subroutine to convert 5 floating point binary bytes to decimal.
61500 Subroutine to convert the 2 of the 5 bytes used to a signed integer.

62000 Subroutine to read string length and location then find and build string．
62005 Avoids the single pass through for next if null string．
62020 Surround a string with quotes－required for changing its value on a break．
62500 Subroutine to detect a function and simply acknowledge the fact as its current value will be picked up by one of the other routines．

```
6 0 0 0 0 ~ : ~
60010 :REM DUMP VARIABLES
60020 :
60030 SU=PEEK(45)+PEEK(46)*256 :REM STAR
T OF VARIABLES
60040 SA=PEEK(47)+PEEK(48)*256 :REM STAR
T OF ARRAYS
60050 \&="":\U未="":\=0:UV=0:REM DEFAULTS
60055 IF SA=SU THEN END: :REM NO SIMPLE
UARIABLES OR EDIT USED
60060 FOR U=0 TO 6:V(U)=PEEK(SU+U):NEXT
U:REM READ UARIABLE NAME AND UALUE
60070 IF U(0)<128 AND U(1)<128 THEN GOSU
B 61000:REM REAL
60080 IF U(0)>128 AND U(1)>127 THEN GOSU
B 61500:REM INTEGER
60090 IF U(0)<128 AND U(1)>127 THEN GOSU
B 62000:REM STRING
60100 IF U(0)>128 AND U(1)<128 THEN GOSU
B 62500:REM FUNC:TION
6 0 1 0 1 ~ I F ~ P E E K ( 2 0 3 ) \lll 6 4 ~ G O T O ~ 6 0 1 0 1 ~
60110 SU=SU+7:GOT0 60040:REM INCREMENT C
OUNTER AND DO NEXT
61000 U$=CHR$(V(0))+CHRक(U(1)):REM REAL
NAME
6 1 0 1 8 U = ( - 1 ) ` ( V ( 3 ) A N D 1 2 8 ) * 2 \wedge ( V ( 2 ) - 1 2 9 )
61020 UV=(1+(<U(3) AND127)+(U(4)+(U(5)+U(
6)/256)/256)/256)/128)
61030 U=U*VU:PRINT U*;"=";V:RETURN
61500 U$=CHRま(V(0)-128)+CHR末(U(1)-128)+"
%":REM INTEGER NAME
61510 U=(V(2)AND 127)*256+U(3)+(U(2)>127)
*32768
61520 PRINTU&;"=";U:RETURN
62000 U$=CHRक(U(0))+CHRक(V(1)-128)+"$":R
EM STRING NAME
62005 IF U(2)=0 GOTO 62020
```



```
\(U(3)+V(4) * 256+U-1)):\) NEXT \(U\)
```



```
: RETURN
62500 RETURN: REM FUNCTION PICKED UP BY 0
THER ROUTINES
```

An alternative approach might be to use the technique in RENUMBER (see below). Namely, print a line which reads: PRINT the variable name and cото the point at which program execution should be resumed. If we get the cursor movements right and poke a return into the keyboard buffer, a dump could be performed. To tidy up, we should really clear the line which says 'PRINT and GOTO' with more cursor movements and spaces, and so on.
An obvious extension would be to include arrays. The logic involved in determining and printing the values of subscripted variables is identical to the above and, with care, the same subroutines could be used. The tricky bit is deciphering the array header to determine the number of dimensions and the size of each dimension. If you do decide to try this, do remember integer array values are stored in only two bytes and string pointers in three bytes, unlike their simple variable counterparts. You must also check that arrays do exist by examining strend and comparing it with ARYTAB +1 . Array headers have also been covered in Chapter 1. Including arrays will greatly increase the size of DUMP and in applications where memory is tight, prove impracticable. It is also difficult, so do not worry if your efforts are not rewarded immediately as a simple error in the logic can cause some very unexpected results.

## Program keyword - LISTER

Function: To produce dated, paged and neatly formatted listings
The version given below has been written for an RS232 printer operating at 300 baud, 1 stop bit and no parity (see Programmer's Reference Guide, Chapter 6: 'Input/Output Guide'). The printer used also required a carriage return/line feed sequence to be generated at the end of each line. Therefore, the logical file number used has to be greater than 127, in this case \#129. When using any RS232 device, it is advisable to OPEN-up the file at the start of the program to allocate the input and output buffers. For other printers, the OPEN and PRINT\# statements below will have to be amended to suit.

If your printer does not support the CBM special characters, the program to be listed should first be run through CODER before generating the ASCII file. With a cassette, the OPEN command to read sequential data on line 210 should read OPEN $2,1,0$, AS (I).
The listing produced is ideal for permanent record, though as the
process takes a little time it is not recommended for intermediate listings．The final listing will have all text inset to column 7 and any wrap－around lines will also be inset．Specifying a line width less than the maximum available has the benefit of allowing space for comments （can save a lot of time in the future）．A brief description follows：

## LINE ACTION

100 See above．
110－Set parameters．
160－Allocate files to be listed to array A\＄（）．
210 See above．
220 This line is included to get any leading returns．The number of these will depend on exactly how the ASCII file was generated． Once a cmD has been issued all returns normally sent to the screen will go to the file．Typically this will be two for the list．If zeros appear on your output then you will have to adjust the program or the way you generate the file．
230 Create bottom margin．
250 Build one line into string as．
260 Same problem as 220 at end of file．Assume a line number of zero is the end．
270 Reset line for text to start in col 7.
290 If length＜max then print it．
300－Else split it and print first part．Recycle each time adding 6 leading spaces to continuations．
340 Print blank lines to next top of form before next program．

```
100 OPEN 129,2,0,CHR$(6)
110 PRINT"LISTER UTILITY":PRINT
120 INPUT"DATE[10SPC]";D$
130 INPUT"LINES/PAGE[4SPC]";LP:IF LP=0 T
HEN LP=66
140 INPUT"MAX CHARS/LINE";CP
150 INPUT"NO.OF PROGS[3SPC]";N:DIM A$(N)
160 PRINT:FOR I=1 TO N
170 INPUT"PROGRAM[7SPC]";A古(I)
180 NEXT I
190 I=0
200 I=I+1:LC=0:IF I >N THEN END
210 Z=1:OPEN 2,8,2,A隹(I)+",S,R"
220 Z=1:GET#2,A*:GET#2,A*:GOSUB 320
230 IF LC>=LP-8 THEN FOR J=1 TO LP-LC:PR
INT#129,"":NEXT J:GOSUB 320
240 J=0:B申=" "
250 J=J+1:GET#2,A末:IF A$〈>CHR$(13) THEN
B車=B多+A支:GOTO 250
```

```
260 IF UAL(B丰)=0 THEN GOSUB 340:GOTO200
270 L$=STR&(VAL(B&)):B&=MID$(L&+"[6SPC]"
,2,6)+MID*(B*,LEN(L京))
280 L=LEN(Bも)
290 IFL<=CPTHEN PRINT#129,B#:LC=LC+1:GOT
0230
300 L$=LEFT$(B*,CP):PRINT#129,Lक:LC=LC+1
310 B&="[7SPC]"+MIDक(B*,CP+1):GOT0 280
320 PRINT#129,"PROGRAM ";A主(I);" LISTED
ON ";DF;" LISTING PAGE";Z:LC=2:Z=Z+1
330 PRINT#129," ":RETURN
340 FOR J=1 TO LP-LC:PRINT#129,"":NEXTJ:
CLOSE2:RETURN
```


## Subroutine keyword－MERGE

Function：To merge two BASIC programs
In all the following where line numbers are common to both the program in memory and the merging program those of the latter will take precedence．

## Merge 1

Where a program is less than 22 screen lines when listed，it may be merged very easily indeed．Simply load the short program and list it． Type NEW and move the cursor to the line below the last line of the list． LOAD the main program and then move up and simply press RETURN on all lines to be included in the final program．

This is the reason for having short keyword routines，to allow the above technique to be used on many of them．

## Merge 2

The following subroutine will merge programs of any length．The program（or part of）to be merged must be stored as an ASCII file on disk or tape．The program resident in memory must，of course，include the merge subroutine．

Initiation：RUN 60000
The resulting program will be an amalgamation of the two programs and unlike APPEND the lines will be in the correct numerical sequence．At the completion of the merge an＇OUT OF DATA＇or＇SYNTAX ERROR＇will be displayed depending on how the ASCII file was generated and which program had the highest line number，but who cares，as the result is exactly what we wanted．The program may then be saved in the normal way（after deleting lines 60000－if they are no longer needed）．The version given is for disk and the necessary changes for cassette have
been included in the description below，but should be only too familiar by now．

The program uses the keyboard programming technique for the most part．There is one problem and that is that whenever an edit is performed all OPEN files are closed．So in theory only one line may be read from the file．Any further attempts to obtain input will result in a ＇FILE NOT OPEN＇error．The solution is simple．BASIC is made to believe a file is open even though an edit has been carried out by pokeing the necessary values into the zero page file registers for current logical file （LA），secondary address（SA）and device number（FA）．

## LINE ACTION

60010 For tape OPEN 1，1，0，F\＄
60020 Get bytes until numeric code．This overcomes the problem in LISTER and perhaps should also be used in that program．
60030 Set file parameters by poking into LA，SA and FA．For tape use 2,0 and 1 （ $0=$ read 1 ＝cassette）．
60050 As the first numeric character has been found，mustn＇t forget to print it．－B\＄
60060－As all other programs using keyboard．
60080 Set up k／b buffer on END to enter printed line and GOTO 60030，the cycle repeating until all done．

```
60000 INPUT"PROGRAM ";F直
60010 OPEN2,8,2,Fक+",5,R"
60020 GET#2, B手:IF UAL(B&) <1 GOTO 60020
60030 POKE 184,2:POKE 185,2:POKE 186,8:P
OKE152,1
60040 PRINT"[CLS][3CD]";
60050 PRINTB害;:B牛=""
60060 GET#2,A$:PRINTA$;:IF A$〈>CHR$(13)
GOTO 60060
60070 PRINT"GOTO 60030[HOM]"
60080 POKE 198,2:POKE 631,13:POKE632,13:
END
```


## Merge 3 （tape only）

This is the cleverest tape merge we have seen．It was originally worked out by J．Butterfield and B．Templeton for the PET and all we have done is to modify it for the 64.

Again，the program to be merged must be on tape in ASCII format． The statements may be typed in direct mode or，as in this case，be a subroutine．In direct mode the contents of the quotes should be typed after performing the necessary cursor moves and RETURN pressed at the end．Line 60030 is needed only in program mode．

## Initiation: RUN 60000

The key to this is the POKE 153,1 (DFLTN) which changes the default input device after each line has been merged from the usual default of $\emptyset$ (the keyboard) back to the cassette (1).

```
60000 INPUT "PROGRAM ";F叓
60010 POKE 19,1:OPEN 1,1,0,F$
60020 PRINT"[CLS][3CD]POKE 153,1:POKE 19
8,1:POKE 631,13:PRINT CHRक(19)"
60030 POKE 198,1:POKE 631,13:PRINT "[HOM
]":END
```

The most common problem with merge is if a program line is in excess of 80 characters when listed (possible if abbreviations have been used). The merge will be unsuccessful as the cursor movements will be incorrect and also BASIC's input buffer will overflow.

## Program keyword - OLD

Function: To recover newed programs
The command NeW does not actually erase the program in memory, it simply changes the first line's link address to 0000 (2049 and 2050) and therefore fools BASIC into thinking that there isn't a program present. In addition, all variable pointers are reset to the end of the program, which in this case is the start of BASIC itself (action of CLR). The following uses these facts to recover the program by resetting the necessary pointers.

To use OLD, the start of BASIC (TXTTAB) must be set above the end of the Newed program and TXTTAB-1 set to zero by:

POKE 43, 01 :POKE 44, no. of pages:POKE (no. of pages) $\star 256,0$ :NEW
OLD may then be loaded and run. The erased program will be recovered and you are back in business. As a point of interest OLD is still present higher in memory and will remain so until overwritten by variable data or a larger program.

The program works by hunting from the input value of TXTTAB +4 (ignore first three zeros) for three consecutive zero bytes which mark the end of the erased program. En route the first link is changed to point to the second line. Once found, тхПтAB is changed to point to the specified start and VARTAB, to the end address. A CLR then tidies up and the original program is listed.

```
60000 INPUT"TXTTAB(2049)";TX:MX=256
60010 POKE 828,TX-INT(TX/MX) *MX:POKE 829
    ,TX/MX
60020 X=TX+4+J:IF PEEK(X)<>0 THEN J=J+1:
GOTO 60020
60030 POKE TX,X+1-INT((X+1)/MX)*MX:POKE
TX+1,(X+1)/256:TX=X+1
60040 X=PEEK(TX) +PEEK(TX+1)*MX:IF X<>0 T
HEN TX=X:GOTO 60040
60050 TX=TX+1:POKE 830,TX-INT (TX/MX) *MX:
POKE 831,TX/MX
60060 POKE 43,PEEK(828):POKE 44,PEEK(829
):POKE 45,PEEK(830):POKE 46,PEEK(831)
60070 CLR:LIST
```


## Subroutine keyword - PLOT

Function: To position the cursor to a specified screen location
There are many ways of positioning the cursor. The most common way is to include the necessary control characters inside quotation marks. This can be expensive on memory if a lot of cursor movement is used. The movement is also relative to the current location and not absolutely fixed to some reference point unless a clear screen or home cursor is first issued. Many micros have $\operatorname{TAB}(x, y), \operatorname{POS}(x, y)$ or $H T A B x$ and VTAB y functions within their BASICS to position the cursor. The following are just two ways of doing this on the 64 with its unmodified BASIC.

## Plot 1

This uses a simple subroutine into which are passed the line and column position. First, two strings are defined - preferably at the start of the program as they remain unchanged throughout the run for speed of access. They are:

$$
1 \mathrm{Y} \$="[24 C D]^{\prime \prime}: \mathrm{X} \$={ }^{\prime}[40 C R]^{\prime \prime}
$$

and have been allocated line number 1. To position to any location, the $x$ and $y$ coordinates are passed to the subroutine which simply homes the cursor then prints the appropriate number of cursor downs and rights.

In the example below, lines 100 to 130 have been included for demonstration purposes.

```
1 Y$="[23CD]":X$="[40CR]"
100 INPUT"COLUMN";X
110 INPUT"[3SPC]ROW";
120 GOSUB 1000:PRINTX;",";Y
130 GOTO 100
1000 PRINT"[HOM]";RIGHTक(Y生,Y);RIGHT&(X ( 
,X);:RETURN
```

The top left of the screen is considered as ' $0, \eta^{\prime}$ '. The semicolon at the end of the print in 60000 is included to hold the cursor at the set location.

The idea of holding frequently used character patterns, control characters, and so on as string variables can reduce memory usage and also makes for easier-to-read code.

## Plot 2

This second plot routine uses the same zero page locations as the kernal function Plot see Programmer's Reference Guide. These are PNT (209/210), PNTR (211) and TBLX (214). If you look at the memory map in Chapter 5 of the PRG or Appendix K of this book, you'll notice locations from 200 to 245 all relate to the screen in some way or other. We are not going to run through them all, but try experimenting with them and see what happens. If in trouble, turn off the 64.

Let us look at the three locations we are going to use to accomplish plot in a little more detail.

PNT: contains the address of the start of the current line in low/high format. With the screen at its default start (1024-2023), this will hold a value $1024+40$ *row where row is in the range $0-24$. Unusual results are produced if this does not correspond to the start of a physical screen line.

PNTR: holds the offset from the address held in PNT. It is the absolute screen column ( $0-39$ ) when PNT holds the start of line address.

TBLX: holds the current physical screen row.
Using only PNT and PNTR, we can position the cursor to any $X, Y$ location. The next print would occur at the specified point. However, when the cursor returns after the print, it reappears at or below the line it was on before PNT and PNTR were set. This is difficult to put into words and much easier to see. For example, if an input took place on line 23 and the cursor was then moved to line 10, column 10 and a PRINT took place without a semicolon, the cursor would reappear at the start of line 24 and not 11 as might be expected. To avoid this, we simply also set tblx and all will be well. The routine below has the same effect as the first PLOT routine given. Again, lines 100 to 130 are included for demonstration only.

```
100 INPUT"COLUMN";}
110 INPUT"[3SPC]ROW";Y
120 GOSUB 1000:PRINTX;",";Y
130 GOTO 100
1000 POKE 214,Y:Y=1024+Y*40
1010 POKE 209,Y-INT(Y/256)*256:POKE 210,
INT(Y/256)
1020 POKE 211,X
1030 RETURN
```


## Subroutine keyword－PRINT USING

PRINT USING is a very powerful output formatting command available in some BASIC languages．It allows numbers to be right or left aligned to a specified number of decimal places，or to be expressed in exponential format and much more．There are equally as many possibilities for strings．A routine to duplicate all the facilities would be very long，so here we have only considered the problem of formatting numbers．

Very quickly everybody picks up on the idea of：

$$
\mathrm{X}=\mathrm{INT}\left(\mathrm{X} \star 10^{\wedge} \mathrm{W}+.5\right) / 10^{\wedge} \mathrm{W}
$$

to get numbers to a set number of decimal places，where $W$ is the number of places．Unfortunately，due to the way numbers are stored， this is not guaranteed to produce the expected result．By way of a trivial example，try printing ．01＊649 and 649＊． 01 and see the difference． The result of any calculation is very much dependent on the order in which it was evaluated．To overcome the problem we have to resort to strings as these are the only type of variable we can fully format．

The following routine will format numbers not in scientific notation and will avoid the xx．x00001 type occurrence by not using any division． The value returned is right aligned to $w$ decimal places and padded with leading spaces to a set width of $L$ ．The variable transferred is in $x$ and the string $x \$$ is returned．

```
1 INPUT "X";X :INPUT"W";W:INPUT"L";L: GO
SUB 60000:PRINTX传
2 GOTO 1
60000 X$=STR&(INT(X*10^W+.5)):LE=LEN(X$)
60010 S2$=".000000000000000":S2$="[31SPC
]"
60020 IF LE<W+2 THEN X 聿=LEFT$(X 
(S2#,1,W+2-LE) +RIGHT$(X$,LE-1)
60030 IF LE)=W+2 THEN X$=LEFT$(X$,LE-W)+
"." + RIGHT生(X生,W)
60040 X $=RIGHT$(S2%+X$,L):RETURN
```

To illustrate its use，the following display

| COL1 | COL2 COL3 | COL4 |
| :---: | :---: | :---: |
| 99.000 | 100.00 .999 | 9.51456 |
| 100.091 | 98.22 .010 | 11.00000 |

would be produced by the program lines；where C1－4 represent the values to be printed in cols 1－4（set elsewhere within your own pro－ gram）．

```
1 C1=99.00001:C2=100.00123:C3=.99888:C4=
9.514569:REM EXAMPLES
10 L=10:W=3:X=C1:GOSUB 60000:PRINT X 招;
20 W=2:X=C2:GOSUB 60000:PRINT X悉;
30 L=5:W=3:X=C3:GOSUB 50000:PRINT X直;
40 L=14:W=5:X=C4:GOSUB 60000:PRINT X韦
```

Where numbers are very large or very small，simply raise them to an appropriate power of ten prior to calling the routine and head the output ${ }^{*} 10^{\circ} \mathrm{N}^{\prime}$ ．

## Subroutine keyword－RENUMBER

Function：To renumber a specified section of a program
It is not possible to write a full renumber program in BASIC which does not use ASCII disk files（somebody will no doubt wish to disprove this statement）．There are many problems，the biggest of which is in renumbering GOTOS，GOSUBS and THENS line destinations．It is relatively easy，albeit slow，to hunt these out by their token values．The problem arises in correcting destinations which are held in ASCII form．For example，GOTO100 is held as 137494848 （ $\$ 89 \$ 31 \$ 30 \$ 30$ in hex）．If during the renumbering process the destination changes by a magnitude of ten or more（the overall length changes），we have to move all code from the byte following the reference up or down in memory，recalculating link addresses as we go．If all references are entered as five figures as standard，this problem is eliminated，for example，GOTO00100．Entering line numbers in this way is rather tedious and is considered imprac－ tical．Machine code renumber programs use the＇crunch tokens＇ routine and the necessary memory moves are performed as part of this routine when a line is added or removed．See renum in Chapter 7.
The program below only renumbers the lines．It will renumber all or only a set block．The new line numbers need not even be in sequence with the rest of the program，though problems will arise if they are referenced．The user will have to manually change all gotos，etc．This subroutine is really intended to allow you to put together a number of the shorter routines in this chapter．

```
60000 TX=PEEK(43)+PEEK(44)*256:MX=256
60010 INPUT"RENUMBER FROM";LL:INPUT"[9SP
C]TO[2SPC]";UL
60020 INPUT"[5SPC]NEW LINE"; S:INPUT"[9SP
C]STEP";I
60030 IF PEEK(TX+2)+PEEK(TX+3)*MX<LL THE
N TX=PEEK(TX)+PEEK(TX+1)*MX:GOTO 60030
60040 S=S+J*I:IF TX=0 OR PEEK(TX+2)+PEEK
<TX+3)*MX>UL THEN END
60050 POKE TX+2,S-INT(S/MX)*MX:POKE TX+3
,S/MX:TX=PEEK(TX) +PEEK(TX+1)*MX
60060 J=J+1:GOT0 60030
```


## Subroutine keyword - SQUASH

Function: To increase the speed of execution of BASIC programs
Many 'crunch' or 'compactor' programs are available, both commercially and in various journals. Their function is to increase the speed of execution of a BASIC program by the removal of redundant code.

There are many reasons why code is slower than it need be. Much of this code is useful at the time of developing the program, but is not required at run-time. Some examples have been given at the end of Chapter 1, but there are many more. Listing the more obvious:

Line numbers: When they are the reference for a goto, gosub or then they are held in ASCII form. The shorter they are (that is, the lower the number), the quicker they are converted to numeric form. Therefore, a renumbering with an increment of 1 is advantageous.
REM: These are ignored at run-time and need only be retained if they are a destination. rems also use valuable memory.
Spaces: Including spaces in a program makes for easier reading, but is unnecessary and wasteful (this is true only outside quotes).
Variable names: One-character names use less space and are found quicker.
Destinations: See Chapter 1 (page 23)
Screen: See Chapter 2 (page 32)
Print: Semicolons separating print lists are sometimes superfluous. They must be retained after a numeric variable and at the end of a PRINT list if a carriage return is to be inhibited.
Line length: Short lines use an extra five bytes each time (link=2 line $=2$ end $=1$ ) and also take time in working out the next line's details. Lines which are not destinations can be strung together, taking due care of the logic of any If statements. Lines may be of any length, but are difficult to edit or generate once they exceed 80 characters (even if all the possible abbreviations are used, there is a limit to BASIC's input buffer).

FOR/NEXT loops: A surprising increase in speed is gained by omitting the variable on the NEXT statement. This eliminates the look-up operation for the variable name. Try timing:

$$
\begin{gathered}
\mathrm{FORI}=1 \mathrm{TO} 255: \mathrm{FOR} \mathrm{~J}=1 \mathrm{TO} 255: \text { NEXTJ:NEXTI } \\
\text { and FOR......NEXT:NEXT }
\end{gathered}
$$

Operating system: Once spaces have been eliminated, Chrget itself may be modified to get rid of the test for spaces. (See Chapter 3).

The subroutine below will remove all unnecessary spaces, semicolons and rems. Renumbering is left up to you. Once again, an ASCII file must first be generated of the program. The program is based on Merge 2 (see page 110) and only the differences from that program are described below:

## LINE ACTION

60030 L is set to 1 to account for $\mathrm{B} \$$ in first line.
60090 As Merge 2 line 60060, returning to k/b bit once a return found.
60100 Once a rem found, ignore all chars except return.
60110 Flag to indicate in or out of quote mode.
60120 Ignore spaces out of quotes.
60130 Keep spaces in quotes.
60140 Semicolons out of quotes require careful checking and this is carried out at 60210 on.
60150 Semicolons in quotes - keep.
60160 If not an ' $\mathrm{M}^{\prime}$, don't look for REM.
60170 Else see if preceding two chars were 'RE'.
60180 If they were, replace by a ' $:$ ' and set RE to ignore everything following (see 60100).
60190 Build line to be printed.
60210 Handle the semicolon when out of quotes and eliminate if possible. Do this by getting next byte and if the list continues check for a preceding string or opening quote. Finally, re-enter the main body of the program where appropriate.

Initiation: RUN 60000

```
60000 INPUT"FROGRAM ";Fま
60010 OPEN2,8,2,F$+",S,R"
60020 GET#2,Bま:IF UAL(B$)<1 GOTO 60020
60030 L=1
60040 POKE 184,2:POKE 185,2:POKE 186,8:P
OKE152,1
60050 PRINT"[CLS][3CD]";
60060 GOSUB 60090
60070 PRINT"GOTO 60040[HOM]"
```

```
60080 POKE 198,2:POKE 631,13:POKE632,13:
END
60090 GET#2,A$:IF A$=CHR$(13) GOTO 60200
60100 IF RE=1 GOTO 60090
60105 IF A&="T" AND Q=0 AND RIGHTま(B本,4)
="PRIN" THEN P=-1
60106 IF A$=":" AND Q=0 THEN P=0
60110 IF A$=CHR&(34) THEN Q=NOT(Q)
60120 IF A$=" " AND Q=0 GOTO 60090
6 0 1 3 0 ~ I F ~ A \$ = " ~ " ~ A N D ~ Q = - 1 ~ G O T O ~ 6 0 1 9 0
60140 IF Aक=";" AND Q=0 AND F=-1 GOTO S0
210
60150 IF A$=";" AND Q=-1 GOTO 60190
60160 IF A$<> "M" GOTO 60190
60170 IF MIDま(B斗,L-1,2)<>"RE" GOTO 60190
60180 Bま=LEFT&(B卉,L-2)+":":RE=1:GOTO 600
90
60190 B==Bま+A方:L=L+1:GOT0 60090
60200 PRINTB急:RETURN
```



```
=L+1:F=0:GOTO 60190
```



```
60200
60230 IFRIGHT本(B立,1)="韦"ORRIGHTक(B*,1)=C
HR$(34) ORC = =CHR$(34)THENA$=C $:GOT060100
60240 A$=A末+C末:L=L+1:GOT0 60190
```


## Conclusion

We hope that this chapter has given you food for thought．By way of a project，why not write a routine to recover as much of the data as possible after an edit or New has been performed？

## 6 Routines old and new

## Introduction

In Chapter 4 we gave listings in machine code to make use of the function keys. These are actioned by keywords. At the present time, bASIC will not understand these. All the functions of the UTILITY, the remainder of which are in the following two chapters, require some sort of 'driving mechanism'. That is, routines which will not only recognize the keywords, but will action them. Those routines are the PRINT tokens, DISPATCH BASIC CHARS and BASIC EVALUATION. In Chapter 3 these were fully discussed, so we are only supplying in this chapter the coding that is particular to the UTility.

To initialize the utility we need to change the addresses in certain locations. These fall into three categories. First, we have to change the vector addresses so that BASIC will go to our token routines; secondly, we need to protect the UTILITY from being overwritten by programs and strings; and lastly we need to retain its operation during a Non-Maskable-Interrupt, that is when run/STOP ReStore is pressed.

There are certain subroutines which will be used by more than one command, so we include them in this chapter. These deal with getting parameters, the switching in and out of the BASIC ROM and memory moving.

That has dealt with the new, and now for the old. A few of the resident ROM routines are useful. Many of them will be covered when describing our new commands. The later part of this chapter describes some more.

## Initialization

When you start up the UTILity with sys 32768 these instructions will be the first to be actioned. They will set up and protect the UTIIITY. At the end of the four subroutines we return control back to you, with a screen message, and the utility in operation.

ASSEMBLY LISTING



We feel that this listing up to 430 is fairly self-explanatory, especially with a memory map. The remaining lines are dealt with in the next section.

## BRK and NMI routines

These are included in the listing of the previous section, lines 430 to 620. When either of these are initiated, it will be to these lines they will come. The majority of these routines are copies of the equivalent ROM routines, plus a couple of directions to our set up routines to keep the utility in service.

NMI
The NmI is initiated by the use of the restore key (although there are means to initiate it through the cartridge slot). Not only does it tell the microprocessor it has been actioned, but it also sets a flag in the CIA \#2. The processor will not action it immediately, but will wait until the present instruction is complete. The processor then saves the program counter and the status register on the stack. It will load the address stored at $\$$ FFFA and $\$$ FFFB into the program counter. This is normally $\$$ FF43. At this address it sets the interrupt flag, so that the other interrupt does not interfere, and then jumps to the vector address that we have changed. Note that the routine has so far not stored the $A, X$ and $Y$ registers.

The NMI in the UTILITY will end up at our routine, which is a series of subroutines. After saving the processor registers on the stack, it clears the NMI flag in the Interrupt Control Register of the CIA\#2 chip, which deals with inputs and outputs of the computer. It then loads that location back into $y$ and if the Nmi flag is still clear then it jumps ahead, missing out for the time being an RS232 reset. The following routine checks the STOP and RVS flags at location \$91. A call to the KERNAL routine to check for stop follows. If on exit the accumulator is zero, then the STOP was initiated and we go to the BRK routine. Finally, we jump to the routine to reset the RS232 locations.

BRK
The first subroutine resets the KERNAL set up vector from \$0314 to \$0333 to their default values from a list held in the kernal rom itself. This will reset two we have changed, the BRK and NMI vectors. The following routine will service the two CIA interface chips, by restoring them to their setup levels.

The routine at $\$$ E518 performs the remaining functions of a BRK. It restores the output device to the screen and the input device to the keyboard. The video chip is next for the restoration treatment. The screen and character set are returned to their default positions, and sprite graphics turned off. After this it is the keyboard's turn, with the buffer, delays and set-up vector all returned to default values. The routine finishes off by resetting the input/output flags, clearing the screen, setting the colours and putting the cursor in the home position.

We now put in two calls ourselves so we can reset the NMI, BRK and keyboard vectors to those we require. Finally, there is the indirect jump of a002 which sets the stack pointer to its start, prints 'ready' and gives control back to the user.

## Routine vectors and keywords

There is sufficient space, using the existing token system, for 51 further keywords. These will be split up into an area for command keywords and an area for function keywords. In the UTILITY we are supplying 34 commands and 1 function. Between the last command keyword vector and that of the function keyword there is space for a further nine commands (token values 238 to 246 (\$EE-\$F6)). Seven extra functions could be added within the space available. The vector table is positioned at $\$ 8090$ to $\$ 80 F 5$.

The keyword table is exactly 255 bytes long. Out of that our keywords use up 155 plus a zero byte to mark the end of the table. The amount of space available to you if wish to extend it is 58 bytes for command keywords and 41 bytes for functions. Remember that the last letter of each word has $\$ 80$ (128) added to it. In our table, the space
between our last command, TROFF, and the only function, DEEK, has been filled with bytes $\$ 5 \mathrm{~A}$ and sea to make up the nine unused token values.

MEMORY DUMP

| 809 | 8 | 87 | 4C | 86 | B2 | 83 |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| :8098 | EB | 84 | 36 | 85 | BE | 85 | 14 | 84 |
| : 80A0 | 51 | 83 | A6 | 83 | AE | 8F | B4 8 | 8 F |
| : 80A8 | 80 | 83 | AC | 83 | 51 | 8 E | C. 48 |  |
| : 80B0 | 43 | 8F | A6 | 87 | 92 | 8B | 2 D | 84 |
| :8088 | D1 | 8F | 3A | A9 | D1 | A8 | 30 | 86 |
| $80 \mathrm{C0}$ | B6 | 91 | 39 | 8D | 10 | 86 | B5 | 2 |
| 80 CB | 8C | 91 | 80 | , |  |  | FB | 85' |
| 0 DO | 6E | 88 | 60 | 8 D | FF | F | FF |  |
| :8008 | 00 | FF | FF | FF | F6 | FF | B6 | F |
| :80E0 | 00 | 60 | 00 | 00 | D6 | 83 | D6 | 8 |
| 89E8 | 00 | 00 | 00 | 68 | 00 | 00 | 00 | 40 |
| 80 Fb | 00 | 00 | 00 | 40 | 00 | 40 | 4F | 46 |
| 80F8 | C6 | 4B | 45 | D9 | 44 | 4F | 4B |  |
| 8100 | 54 | 45 | CE | 54 | 57 | CF | 48 | 45 |
| 08 | D8 | 42 | 49 | CE | 4F | 4C | C4 | 4 |
| 110 | 4F | 4C | 4F | 55 | D2 | 57 | 52 |  |
| 8118 | 54 | C5 | 43 | 47 | 4F | 54 | CF | 4 |
| : 8120 | 47 | 4F | 53 | 55 | C2 | 50 | 4C | 4 F |
| :8128 | D4 | 45 | 4E | 54 | 45 | D2 | 44 | 55 |
| 8130 | 4D | D0 | 52 | 45 | 4E | 55 | CD | 44 M |
| :8138 | 45 | 4C | 45 | 54 | C5 | 4D | 45 | $52^{\prime}$ ELE |
| : 8140 | 47 | C5 | 43 | 4F | 44 | 45 | D2 | 41 |
| :8148 | 55 | 54 | CF | 50 | 52 | 4F | C3 | 44 |
| : 8150 | 50 | 52 | 4F | C3 | 45 | 50 | 52 | 4 F |
| :8158 | C3 | 50 | 4F | D0 | 51 | 55 | 49 | D4 cP |
| . 8160 | 54 | 52 | 41 | 43 | C5 | 52 | 45 | $53^{\prime}$ T |
| : 8168 | 45 | D4 | 43 | 48 | 41 | 49 | CE | 4C'EtCH |
| :8170 | 4F | 4D | 45 | CD | 48 | 49 | 4D | 45' OMEmHIME |
| :8178 | CD | 49 | 4E | 4B | 45 | 59 | A4 | 4D'mINKEY |
| :8180 | 45 | CD | 41 | 50 | 50 | 45 | 4E | C4'EmAPPEN |
| :8188 | 54 | 52 | 4F | 46 | C6 | 5A | 5A | $5 A^{\prime}$ TROF 5222 |
| :8190 | 5A | 5A | EA | 5A | 5A | 5A | 5A 5 | 5A' 22.22222 |
| :8198 | 5A | EA | 5A | 5A | 5A | 5A | 5A 5 | $5 A^{\prime} 2.222222$ |
| :81A0 | 5A | EA | 5A | 5A | 5 A | 5A | EA 5 | $5 A^{\prime} 2.2222 .2$ |
| :81A8 | 5A | 5A | 5A | EA | 5A | 5A | 5A 5 | $5 A^{\prime} 222.2222$ |
| : 81B0 | EA | 5A | 5A | 54 | 5 A | 5A | 5A E | EA'. 227222 |
| : 81 B8 | 5A | 5A | 5A | 5A | 5A | EA | 5A | $5 A^{\prime} 22722$ |
| :81C0 | 5A | 5A | 5A | 5A | 5 5 | 5A | EA | 44 |
| :81C8 | 45 | 45 | CB | 00 | FF | FF | FF F | FF'EEK |

```
.:81D0 FF FF FF FF FF FF FD FF'........
.:81D8 FF FF FF FF FF 7F FF FF,........
.:81E000 00 00 00 00 00 00 00%........
.:81E8 00 00 00 08 01 00 00 00%........
.:81F0 00 00 00 00 00 20 8A AD'........
```

This has been produced in upper case mode and as such the end shifted letter of each command is printed in lower case. If putting it into your computer in a way other than the dump, remember that they are shifted. The last letter in location $\$ 817 \mathrm{E}$ is a shifted $\$$, giving the keyword inkeys.

## Getting parameters and controlling BASIC

ASSEMBLY LISTING


## Parameters

Lines 10 and 20 hold the only two instructions that we need to incorporate, but they do a lot of work in getting our numeric parameters. Let us look at the instructions one at a time.

> JSR \$AD8A

The first action of this is to call the evaluate expression routine at \$AD9E.
This is a complex routine which deals not only with numeric data, but also with strings. After setting the CHRGET pointer back one place, it proceeds to start picking up data after the command keyword. It will then go through checking to see whether a mathematical operator or a function keyword (such as Peek), a variable or simply a number has been obtained. From the information obtained it will (after calculating if necessary) store the result or findings in the FAC\#1. For numbers up to SFFFF, the relevant numbers will be in locations $\$ 64$ and $\$ 65$ of this accumulator.

We now return to our original subroutine at \$AD8A, where we check to see if the data received was numeric or not. The evaluate expression
will set a flag in the zero page location s0D. The value of $\$ F F$ indicates string data, whilst zero designates numeric data. If this subroutine finds the former, a 'TYPE MISMATCH' error is generated and the command, and program, is terminated.

## JMP \$B7F7

We have our numeric parameter. This routine will do two checks and then transfer our data. The checks are to make sure that neither a negative number nor one over 65535 (\$FFFF) was given. In either case, failure will result in the 'Illegal quantity' error. The data is now transferred from $\$ 64$ and $\$ 65$ to locations $\$ 14$ and $\$ 15$. The reason for this is that the FAC\#1 is used for many applications. The RTS at the end of this routine will return us to the place that called our complete GET PARAMETERS routine, that will most likely be a command routine.

The BASIC switch
As we said, when dealing with the function keys, the area of ram under the BASIC ROM is a useful place for hiding data, or indeed routines which do not use the BASIC interpreter. To use this area, BASIC must be 'removed'. We have no trouble writing to the RAM as the computer, through its decoding logic, will select it when the processor sends a write signal. When reading, the ROM has priority unless we tell the electronics that it is not there. The main difference between the 6510 processor in the 64 and the normal 6502 is that the former has input/ output ports. The user can control these using locations $\$ 0000$ and $\$ 0001$. The first deals with the direction of the data, that is, whether the ports, of which there are six, are going to be input or output. The second location deals with the data itself, one bit for each port, either a one or a zero which gives a switching mode. Three of the ports are connected to the cassette port. The other three control three ROMS: BASIC, KERNAL and the Character rom. A zero will switch all of these off. The one we are concerned with, BASIC, uses bit $\emptyset$ of the data register and so by changing this bit, making sure not to disturb the others, we can remove or replace as required.

Lines 30 to 60 perform the switching-out of BASIC. We load the register and set bit $\emptyset$ to zero. The AND instruction will do this without changing any other bit. After placing the result back, the ROM is no longer present as far as the computer is concerned.

Lines 70 to 100 reverse the process by using the ORA code which will only affect the bits according to the data with the instruction.

> To switch off BASIC - JSR $\$ 81 \mathrm{FB}$
> To switch in BASIC - JSR $\$ 8202$

## Dealing with the keywords

In Chapter 3 the routines that BASIC uses to deal with keywords and
tokens were fully described. Below are the listings to use with the UTILITY, which require no further explanation.

## ASSEMBLY LISTING - CRUNCH TOKENS

| 9 | * $=\$ 8289$ |  | 410 | STORE | INX |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 10 |  | LDX $\$ 7 \mathrm{~A}$ | 420 |  | INY |  |
| 20 |  | LDY \#\#04 | 430 |  | STA | \$01FB, Y |
| 30 |  | STY 08 F | 440 |  | LDA | \$01FB, Y |
| 40 | ANOTHER | LDA $\$ 8200$, X | 450 |  | BEQ | EXIT |
| 50 |  | BPL SPACE | 460 |  | SEC |  |
| 60 |  | CMP \#\#FF | 470 |  | SBC | \#\$3A |
| 70 |  | BEQ STORE | 480 |  | BEQ | COLON |
| 80 |  | INX | 490 |  | CMP | \#\$49 |
| 90 |  | BNE ANOTHER | 500 |  | BNE | data |
| 100 | SFACE | CMF \# 20 | 510 | COLON | STA | \$0F |
| 110 |  | BNE STORE | 520 | data | SEC: |  |
| 120 |  | STA $\$ 08$ | 530 |  | SBC | \#\$55 |
| 130 |  | CMP \#\$ 22 | 540 |  | BNE | ANOTHER |
| 140 |  | BEQ QuOTE | 550 |  | STA | \$08 |
| 150 |  | BIT \$0F | 568 |  | LDA | \$0200, X |
| 160 |  | BUS STORE | 570 | LINE | BEQ | Stare |
| 170 |  | CMP \#\#3F | 580 |  | CMP | \$08 |
| 180 |  | BNE NUMBER | 590 |  | BEQ | Store |
| 190 |  | LDA \#\$99 | 600 | Quote | INY |  |
| 200 |  | BNE STORE | 610 |  | STA | \$01FB, Y |
| 210 | NUMBER | CMP \#\$ 30 | 620 |  | INX |  |
| 228 |  | BCC: CONT | 630 |  | BNE | LINE |
| 230 |  | CMP \#\#3C | 640 | NEXTWORD | LDX | \$7A |
| 240 |  | BCC: STORE | 650 |  | INC | \$0B |
| 250 | CONT | STY \$71 | 660 | FIND | INY |  |
| 260 |  | LDY \#\#80 | 670 |  | LDA | \$A09D, Y |
| 270 |  | STY \$0B | 680 |  | BPL | FIND |
| 280 |  | DEY | 690 |  | LDA | \$A09E,Y |
| 290 |  | STX \$7A | 700 |  | BNE | CONT1 |
| 300 |  | DEX | 710 |  | LDY | \#कFF |
| 310 | NEXTLETTER | INY | 720 |  | DEX |  |
| 320 |  | INX | 730 | NEXT | INY |  |
| 330 | CONT 1 | LDA $\$ 0200, X$ | 740 |  | INX |  |
| 340 |  | SEC | 750 | NEXTB | LDA | \$0200, $\times$ |
| 350 |  | SBC \$A09E,Y | 760 |  | SEC |  |
| 360 |  | BEQ NEXTLETTER | 770 |  | SBC | \$80F6, ${ }^{\text {\% }}$ |
| 370 |  | CMP \#\$80 | 780 |  | BEQ | NEXT |
| 380 |  | BNE NEXTWORD | 790 |  | CMP | \#\$80 |
| 390 | STOREA | ORA \$0B | 808 |  | BNE | NEXTNEW |
| 400 | FGUND | LDY \$71 | 810 |  | BEQ | STOREA |


| 870 | LDA | \$80F6, Y | 820 | NEXTNEW | LDX $\ddagger 7 A$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 880 | BNE | NEXTE | 830 |  | INC \$0B |
| 890 | LDA | \$0200, $X$ | 840 | NEXTA | INY |
| 908 | BPL | FOUND | 850 |  | LDA \$80F5,Y |
| 910 EXIT | JMP | \$A609 | 860 |  | BPL NEXTA |
| 820F ANOTHER |  | 8269 C. |  |  |  |
| 8239 CONT |  | 8245 CO |  |  |  |
| 826 B DATA |  | $82 \mathrm{B9}$ EX |  |  |  |
| 8286 FIND |  | 8254 F0 |  |  |  |
| 8275 LINE |  | 8294 NE |  |  |  |
| 82A9 NEXTA |  | 8296 NE |  |  |  |
| 8243 NEXTLETTER |  | 8245 NE | TNEW |  |  |
| 8282 NEXTWORD |  | 8231 NU | BER |  |  |
| 827B QUOTE |  | 821 B SP |  |  |  |
| 8256 STORE |  | 8252 ST | EA |  |  |

ASSEMBLY LISTING - PRINT TOKENS
$*=\$ 82 \mathrm{BC}$

10
20
30
40
50
60
78
80
90
100
110
120
130
140
150
160 CBMTOKEN
170
180
190
200
210
220
230 START
240
250 NEXTWORD

BPL ROM1
CMP \#कFF
BEQ ROM1
BIT $50 F$
BMI ROM1
CMP \#\$CC
BCC CBATTOKEN
SEC
SBC: \#\#CB
TAX
LDA \#\#F6
STA $\$ 22$
LDA \#ま80
STA $\$ 23$
BNE START
SEC
SBC \#\#7F
TAX
LDA \#\$9E
STA $\$ 22$
LDA \#\#AB
STA $\$ 23$
STY $\$ 49$
LDY \#\$FF
DEX

| 260 | BEQ WORDFOUND |
| :--- | :--- |
| 270 NEXTCHAR | INY |
| 280 | LDA (\$22),Y |
| 290 | BPL NEXTCHAR |
| 300 | BMI NEXTWORD |
| 318 WORDFOUND | INY |
| 320 | LDA (\$22),Y |
| 330 | BMI ROM2 |
| 340 | JSR \$AB47 |
| 350 | BNE WORDFOUND |
| 360 ROM1 | JMP \$A6F3 |
| 370 ROH2 | JMP \$A6EF |


| 82D8 CBMTOKEN | 82EB NEXTCHAR |
| :--- | :--- |
| 82E8 NEXTWORD | 82FC ROM1 |
| 82FF ROM2 | $82 E 4$ START |
| 82F2 WORDFOUND |  |

ASSEMBLY LISTING - DISPATCH AND EVALUATION

| $9 \%$ \% 8302 |  |  |  |
| :---: | :---: | :---: | :---: |
| 10 |  | JSR | \$8873 |
| 20 |  | CMF | \#\$CC |
| 30 |  | BCC | ROM3 |
| 40 |  | CMP | \#SEE |
| 50 |  | BCS | ROM3 |
| 60 |  | JSR | DISPATCH |
| 70 |  | JMP | \$A7EA |
| 80 | DISPATCH | SEC |  |
| 90 |  | SBC | \#\$ CC |
| 108 |  | ASL | A |
| 110 |  | TAY |  |
| 120 |  | LDA | 8091, Y |
| 130 |  | PHA |  |
| 140 |  | LDA | \$8090, Y |
| 150 |  | PHA |  |
| 160 |  | JMP | \$0073 |
| 170 | ROM3 | JSR | \$0079 |
| 180 |  | JMF | \$A7E7 |
| 190 |  | LDA | \#\$00 |
| 280 |  | STA | \$0D |
| 210 |  | JSR | \$0073 |
| 220 |  | CMP | \#\$F7 |
| 230 |  | BCC | ROM4 |
| 240 |  | CMP | \#\$F8 |


| 250 | BCS ROM4 |
| :--- | :--- |
| 260 | JSR DISPATCHI |
| 270 | RTS |
| 280 | DISPATCHI |
| 290 | SEC |
| 300 | SBC \#\$F6 |
| 310 | ASL A |
| 320 | TAY |
| 330 | LDA $\$ 80 E 5, Y$ |
| 340 | PHA |
| 350 | LDA $\$ 80 E 4, Y$ |
| 360 | PHA |
| 370 | ROM4 |
| 380 | JMP $\$ 0873$ |
|  | JSR $\$ 0079$ |
|  | JMF $\$ A E 8 D$ |
| 8313 DISPATCH |  |
| 8323 ROM3 | $833 C$ DISPATCH1 |

## The start up message

This is the final subroutine called during the initialization of the UTILITY. It performs a CLR, to set all the variable addresses, changes the screen and text colours, and finally puts a message on the screen indicating that the utility is in operation.

ASSEMBLY LISTING
*=\$9208

| 10 | JSR \$A663 |  | CLR |
| :---: | :---: | :---: | :---: |
| 20 | LDA \#\#93 | ! | CLEAR SCREEN |
| 30 | JSR \$FFD2 |  |  |
| 40 | LDA \#\#80 | ! | SET COLOURS TO |
| 50 | STA \$0020 | ! | BORDER |
| 60 | STA \$D021 | ! | BACKGROUND |
| 70 | LDA \#\$05 |  |  |
| 80 | STA 8286 | ! | GREEN TEXT |
| 90 | LDX \#\$0A |  |  |
| 100 | LDY \#\$09 |  |  |
| 110 | JSR STARS |  |  |
| 120 | LDX \#\$0C |  |  |
| 130 | LDY \#\$09 |  |  |
| 140 | CLC |  |  |



## Memory moving

RENUMBer and CODER, described in Chapter 7, both require some manipulation of memory in the form of either gaining space or removing unnecessary bytes. This section deals with the two subroutines, CLOSE and OPEN, which perform these operations. CLOSE is self-contained whilst OPEN uses a ROM routine for the actual moving of memory. In the BASIC interpreter there are routines to both open and close up a BASIC program, used when you insert or delete lines, but we can only really use the opening routine. It is a subroutine on its own whereas the closing-up is integral with the inputting of a basic line. We
have written coding that is virtually identical to the one in ROM as it is efficient enough.

Having moved the program about, all the link addresses, from the line the move started, will now be wrong by the amount of the move. There is a subroutine in the interpreter which changes the link addresses but we have not used it. The reason for this is one of speed as during the course of using CODER or RENUM, these subroutines may be called several times and would prove to be very slow.

The ROM routine for rechaining the lines goes through the whole program, byte by byte, to calculate the link addresses and store them. It has been done this way as it is a multi-purpose routine, catering for the lengthening and shortening of code. What we have done is to write separate routines for each direction of movement and place them immediately after the moving instructions. These will only rechain from the program line in which the alteration occurred. In addition, we only need to look at the link addressess as we know by how much they have changed so we can subtract or add as required.

To set the scene, as they say, here are the locations that need to be set before calling these subroutines:
\$FB and \$FC- The address of the start of the current BASIC line.
\$49- The number of the current position on that line. This will be where the replacement code will start.
\$FD and \$FE- The address of the next BASIC line, that is, the link address of the line in \$FB and \$FC.
\$3E- The number of bytes in the original code to changed.
register- The number of bytes in the replacement code.
ASSEMBLY LISTING
$9 *=\$ 888 B$

| 10 | STX $\$ C 2$ |  |
| :--- | :--- | :--- |
| 20 | LDA $\$ 3 E$ | ! FIND HOW MANY |
| 30 | SEC |  |
| 40 | SBC $\$ C .2$ |  |
| 50 | STA $\$ B B$ |  |
| 60 | CLC |  |
| 70 | LDA $\$ F B$ |  |
| 80 | ADC $\$ 49$ | ! FIND START OF |
| 90 |  | BLOCK TO MONE |
| 180 | STA $\$ 5 F$ |  |
| 110 | ADA $\$ F C$ |  |
| 120 | STA $\$ 60$ |  |
| 130 | LDA $\$ 5 F$ |  |


| 148 |  | ADC | \$BB |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 150 |  | STA | \$5A | ! | START + AMOUNT OF REDUCTION |
| 160 |  | LDA | \$60 |  |  |
| 170 |  | ADC: | \# $\$ 00$ |  |  |
| 180 |  | STA | \$5B |  |  |
| 190 |  | LDA | \$20 | ! | END OF PROG |
| 200 |  | SEC |  |  |  |
| 210 |  | SBC | \$5A | ! | CALCULATE TOTAL AMOLNT TO MONE |
| 220 |  | STA | \$58 |  |  |
| 238 |  | TAY |  | $!$ | NO OF BYTES OF INCOMPLETE PAGE |
| 240 |  | LDA | \$2E |  |  |
| 250 |  | SBC | \$5B |  |  |
| 260 |  | TAX |  | ! | NO OF PAGES TO MOVE |
| 270 |  | INX |  | ! | FOR EASIER CHECKING |
| 280 |  | TYA |  |  |  |
| 290 |  | BEQ | PAGE | ! | NO SEPARATE BYTES |
| 300 |  | LDA | \$5A | ! | MOVE SEPARATE BYTES FIRST |
| 310 |  | CLC |  |  |  |
| 320 |  | ADC | \$58 |  |  |
| 330 |  | STA | \$5A |  |  |
| 340 |  | BCC | NOINC. |  |  |
| 350 |  | INC | \$5B |  |  |
| 360 |  | CLC |  |  |  |
| 370 | NOINC | LDA | \$5F |  |  |
| 380 |  | ADC | \$58 |  |  |
| 390 |  | STA | \$5F |  |  |
| 400 |  | BCC | NOINCA |  |  |
| 410 |  | INC | \$68 |  |  |
| 420 | NOINCA | TYA |  |  |  |
| 430 |  | EOR | \#\$FF |  |  |
| 440 |  | TAY |  |  |  |
| 450 |  | INY |  |  |  |
| 468 |  | DEC | \$5B |  |  |
| 470 |  | DEC | \$60 |  |  |
| 480 | PAGE | LDA | ( $\ddagger 5 A$ ), Y |  |  |
| 490 |  | STA | (\$5F), $Y$ |  |  |
| 500 |  | INY |  |  |  |
| 518 |  | BNE | PAGE |  |  |
| 520 |  | INC: | \$5B |  |  |
| 530 |  | INC | \$68 |  |  |
| 548 |  | DEX |  | 1 | POINTER - COMPLETION |


| 550 |  | BNE | PAGE |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 568 |  | SEC: |  |  |  |
| 570 |  | LDA | \$2D | ! | SET END OF PROG |
| 580 |  | SBC | \$BB |  |  |
| 590 |  | STA | \$2D |  |  |
| 608 |  | BCS | RECHAIN |  |  |
| 610 |  | DEC | \$2E |  |  |
| 628 |  | SEC: |  |  |  |
| 630 | RECHAIN | LDY | \#\$00 |  |  |
| 640 |  | LDA | कFD | $!$ | GET LINK |
| 650 |  | SBC | \$BB | ! | CALC NEU ADDRESS |
| 660 |  | STA | FFD |  |  |
| 670 |  | STA | (解B), Y | ! | STORE IN LINE |
| 680 |  | STA | \$57 |  |  |
| 690 |  | LDA | \$FE |  |  |
| 710 |  | SBC | \#\$00 |  |  |
| 710 |  | INY |  |  |  |
| 720 |  | STA | \$FE |  |  |
| 730 |  | STA | \$58 |  |  |
| 740 |  | STA | ( $\ddagger \mathrm{FB}$ ) , Y |  |  |
| 750 | NEXT1 | DEY |  |  |  |
| 760 |  | LDA | (\$57), Y | ! | GET LINKS |
| 770 |  | STA | \$89 | ! | STORE THEM |
| 780 |  | INY |  |  |  |
| 790 |  | LDA | (\$57), Y |  |  |
| 800 |  | STA | \$EA |  |  |
| 810 |  | BEQ | EXIT | $!$ | COMPLETED RECHAINING |
| 820 |  | DEY |  |  |  |
| 830 |  | SEC |  |  |  |
| 840 |  | LDA | \$89 |  |  |
| 850 |  | SBC | \$BB | ! | CALC NEW LINK ADDS |
| 860 |  | TAX |  | ! | TEMP STORE |
| 870 |  | STA | (\$57), Y |  |  |
| 880 |  | LDA | \$BA |  |  |
| 890 |  | SBC | \# $\$ 00$ |  |  |
| 900 |  | INY |  |  |  |
| 910 |  | STA | (\$57), Y |  |  |
| 920 |  | STA | \$58 |  |  |
| 930 |  | TXA |  |  |  |
| 948 |  | STA | \$57 |  |  |
| 950 |  | JMP | NEXT 1 | $!$ | GET NEXT LINE |
| 968 | EXIT | RTS |  |  |  |
| 970 |  | TXA |  |  |  |
| 988 |  | SEC |  |  |  |
| 990 |  | SBC | \$3E | ! | CALCULATE NO OF SPACES REQUIRED |


| 1000 |  | STA | \$BB |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1010 |  | CLC |  |  |  |
| 1020 |  | LDA | \$49 |  |  |
| 1830 |  | ADC: | \$BB |  |  |
| 1040 |  | BCS | ERROR | ! | >255 CHARS IN LINE |
| 1050 |  | CMP | \$FE |  |  |
| 1060 |  | BCC | CONT | ! | ONLY 254 ALLOWED <br> -NEED END MARKER |
| 1070 | ERROR | LDX | \# ${ }^{\text {1 }} 17$ |  |  |
| 1080 |  | JMP | \$4437 | ! | ERROR STRING TOO LONG |
| 1090 | CONT | LDA | \$20 |  |  |
| 1180 |  | ADC. | \$BB | ! | ENOUGH MEMORY? |
| 1110 |  | TAX |  |  |  |
| 1120 |  | L[AA | \$2E |  |  |
| 1130 |  | ADC | \#\$00 |  |  |
| 1140 |  | CMF | \$38 |  |  |
| 1150 |  | BNE | CONT2 | ! | ENOUGH MEMORY |
| 1160 |  | CPX | \$37 |  |  |
| 1170 |  | BCC | CONT2 |  |  |
| 1180 |  | JMF' | \$4435 | ! | ERROR OUT OF MEMORY |
| 1190 | CONT2 | CLC |  |  |  |
| 1208 |  | LDA | \$20 | $!$ | SET ADDS FOR MOVE |
| 1210 |  | STA | \$5A |  |  |
| 1220 |  | ADC | \$BB |  |  |
| 1230 |  | STA | \$58 |  |  |
| 1240 |  | LDA | \$2E |  |  |
| 1250 |  | STA | \$5B |  |  |
| 1260 |  | ADC: | \#\# 00 |  |  |
| 1270 |  | STA | \$59 |  |  |
| 1280 |  | LDA | \$FB |  |  |
| 1290 |  | ADC | \$49 |  |  |
| 1300 |  | STA | \$5F |  |  |
| 1310 |  | LDA | \$FC |  |  |
| 1320 |  | ADC | \# $0^{80}$ |  |  |
| 1330 |  | STA | \$68 |  |  |
| 1340 |  | JSR | \$A3BF | ! | ROM ROUTINE TO OPEN UP MEMORY |
| 1350 |  | CLC |  |  |  |
| 1360 |  | LDY | \# ${ }^{\text {8 }} 0$ |  |  |
| 1370 |  | LDA | \$2D | ! | SET NEW END OF PROG |
| 1380 |  | ADC: | \$BB |  |  |
| 1390 |  | STA | \$20 |  |  |
| 1400 |  | BCC | CONT3 |  |  |
| 1410 |  | INC | \$2E |  |  |
| 1420 |  | CLC |  |  |  |
| 1430 | CONT3 | LDA | \$FD |  |  |


| 1448 |  | ADC | \$BB |
| :---: | :---: | :---: | :---: |
| 1450 |  | STA | \$FD |
| 1460 |  | STA | \$57 |
| 1470 |  | STA | (\$FB), $Y$ |
| 1480 |  | LDA | \$FE |
| 1490 |  | ADC | \#\$00 |
| 1580 |  | INY |  |
| 1510 |  | STA | \$FE |
| 1520 |  | STA | \$58 |
| 1530 |  | STA | (\$FB), $Y$ |
| 1540 | NEXT3 | DEY |  |
| 1550 |  | LDA | (\$57) , Y |
| 1560 |  | STA | \$89 |
| 1570 |  | INY |  |
| 1588 |  | LDA | (\$57), Y |
| 1590 |  | STA | \$BA |
| 1608 |  | BEQ | EXIT2 |
| 1610 |  | DEY |  |
| 1628 |  | CLC |  |
| 1630 |  | LDA | \$ ${ }^{\text {9 }}$ |
| 1640 |  | ADC | \$BB |
| 1650 |  | TAX |  |
| 1668 |  | STA | (\$57) , Y |
| 1678 |  | LDA | \$BA |
| 1680 |  | ADC | \# $\mathbf{S}^{\text {0 }} 8$ |
| 1690 |  | INY |  |
| 1780 |  | STA | ( $\$ 57$ ), Y |
| 1710 |  | STA | \$58 |
| 1720 |  | TXA |  |
| 1730 |  | STA | \$57 |
| 1748 |  | JMP | NEXT3 |
| 1750 | EXIT2 | RTS |  |


| CONT | 8949 | CONT2 | 8950 |
| :--- | :--- | :--- | :--- |
| CONT3 | $898 B$ | ERROR | 8944 |
| EXIT | $892 F$ | EXIT2 | $89 C 4$ |
| NEXT1 | $890 E$ | NEXT3 | $89 A 0$ |
| NOINC | $88 C A$ | NOINCA | $88 D 4$ |
| PAGE | $88 D D$ | RECHAIN | $88 F 7$ |

## CLOSE ROUTINE

LINES 10-270: Before we can move a block of memory, we have to determine three values: the start address of the block to move, the new start address and the amount of code to move. The first thing we
work out is the number of redundant bytes. This is done, obviously, by subtracting from the original amount of data to be changed the number of bytes of the replacement code. The resultant value is stored in location \$BB. We shall need this number later for rechaining the lines. The new start of the block will be obtained by adding the line pointer, $\$ 49$, to the address of the current basic line. To this value is added the contents of $\$ B B$ which will give us the location of the first byte in the block to be moved.

To get the amount of data to be moved, the result of the last calculation is taken away from the end of program address, held in \$2D and \$2E. The answer will be held in the processor registers, the high byte in the $x$ and the low in the $y$. A page of memory is 256 bytes so the $x$ register is therefore the number of pages to be moved, increased by one for easier checking on completion. We move a complete page and then decrease $x$. $x$ will be zero when all done, checking immediately after decreasing. To summarize, we have found the amount to move, its current start and its destination.

LINES 280-470: This is the hardest part of the routine to follow, and we hope that we succeed in explaining it clearly.

We transfer the $y$ register to the accumulator. To recap briefly, this will be the number of bytes, other than complete pages, of memory to move. If the value now in the accumulator is zero, only complete pages require moving, so we skip this section completely. In closing up memory we start from the low addresses, move them, and work to the higher end addresses. We do this by setting the address of the page and moving it up, using the $y$ register as a pointer. If we have an odd number of bytes to start with, this causes a slight problem. For example, if we have $\$ 10$ bytes and the Y is set thus we would move 246 bytes by increasing Y . To compensate for this, what we do is to produce the 2's complement of the value. This is done in lines 430 to 450. The EOR \#SFF will change all the bits set to one to zero and vice versa. One is then added. So instead of $\$ 10$, we should now have $\$ \mathrm{Fb}$. This means that if we now increase $\gamma$ from $\$ F 0$ until it becomes zero it will have been incremented $\$ 10$ times.

For the same reasons we have to alter the address of the start of the block and its new start address. We add to these the original number of odd bytes, held in $\$ 58$. Finally, we decrease the high byte of the address by one. The next effect of these changes is a stalemate as the locations along with the $\gamma$ pointer value are equivalent to the original values but now allow us to increase $\gamma$ the required amount.

LINES 480-550: Having set all the values we move the data, byte by byte, until both $x$ and $y$ registers are zero. We simply load a byte from its position and store its new lower location.

LINES 560-620: The end of the BASIC program will now be shorter by
the value of location sвb. The original end address is adjusted and reset.

LINES 630-960: All that remains is to change the values of the link addresses from the current BASIC line onwards. First, we change the links in the current line and as these are also held in SFD and SFE, used by the calling routine, we change these also.

We proceed through the lines gathering the addresses, subtracting the value in $\$ B B$, and then we restore them. The end of the program is indicated when the MSB of a link address is zero. Finally, we return to the calling program, such as CODER.

## OPEN ROUTINE

LINES 970-1080: We calculate the space required by subtracting the value in $\$ 3 E$, the length of the old code, from the value in the $x$ register, the length of the new code, and store the result in $\$ B B$.

As a basic line may not exceed 255 bytes (to allow for a zero at the end making a maximum of 256), we check this by adding the line marker to the $\$ B B$ value. A set carry flag will mean the maximum has been exceeded. We then check that there will be room for the end of line zero. Failure of either of these will generate the syntax error 'STRING toolong'.

LINES 1090-1180: As we are creating space we must check that there is sufficient room available in the BASIC program area. These lines do just that by checking that we will not exceed the values in $\$ 37$ and $\$ 38$, which indicate its limit. If we do go over, we call a basic routine to generate the 'OUT OF MEMORY' error message.
LINES 1190-1340: Next on the agenda is to set the registers for the interpreter's OPEN-up memory routine at \$A3BF. On leaving this routine: $\$ 5 A$ and $\$ 5 B$ - This will hold the address of the end of the block to move. It will be the same as the end of program address before the move. $\$ 58$ and $\$ 59$ - These registers will hold the address of the end of the new block. It will also be the end of the BASIC program after the move. It is arrived at by adding the amount of move to the address in \$5A and \$5B. $\$ 5 F$ and $\$ 60$ - The start of the block to move. These hold the location of the first byte of the code to be changed. It is calculated by adding the line marker to the address of the current basic line to be processed.

LINES 1360-1420: Now that the data has been moved, we reset the end of program address to its new value.
LINES 1430-1750: A replica of the rechaining in lines 630 to 960 , except that here we increase the addresses instead of reducing them.

This concludes the new routines that we planned to introduce in this
chapter. The remainder are descriptions of some of the ROM routines we use (and hope that you will come to use).

## RECHAINING THE LINES

During our memory move routine, we did not use the ROM routine to rechain the link addresses because for our purposes it was inefficient due to the number of calls required. However, we do use the subroutine, in Delete for instance, where only one call is required. It serves another purpose in that from the addresses it exits with, one can calculate and set the end of program/start of variable registers.

## ROM LISTING

| A533 | A5 2B | LDA | \$2B |
| :---: | :---: | :---: | :---: |
| A535 | A4 2C | LDY | \$2C |
| A537 | 8522 | STA | \$22 |
| A539 | 8423 | STY | \$23 |
| A53B | 18 | CLC |  |
| A53C | A0 01 | LDY | \#\$01 |
| A53E | B1 22 | LDA | (\$22), $Y$ |
| A540 | F0 10 | BEQ | \$A55F |
| A542 | A8 04 | LDY | \#\$04 |
| A544 | C8 | INY |  |
| A545 | B1 22 | LDA | (\$22), Y |
| A547 | D0 FB | BNE | \$A544 |
| A549 | C8 | INY |  |
| A54A | 98 | TYA |  |
| A54B | $65 \quad 22$ | ADC | \$22 |
| A54D | AA | TAX |  |
| A54E | A0 00 | LDY | \# $\$ 00$ |
| A550 | 9122 | STA | (\$22), Y |
| A552 | A5 23 | LDA | \$23 |
| A554 | 6980 | ADC | \# $\$ 00$ |
| A556 | C8 | INY |  |
| A557 | 9122 | STA | (\$22) , Y |
| A559 | 3622 | STX | \$22 |
| A55B | 8523 | STA | \$23 |
| A55D | 90 DD | BCC | \$A53C |
| A55F | 60 | RTS |  |

The routine commences by getting the program start address and placing it in registers for its own use. The carry flag is cleared for addition. The first byte of a line that it picks up is the high byte of the link address and it tests for the end of the program (a zero). The $y$ register is loaded again so as to skip the addresses and line number. It
now proceeds through the line, searching for the end of line zero marker. When this is discovered, the $y$ register will contain one less than the number of bytes in the complete line. This is immediately rectified by incrementing $Y$ by one. This value is added to the line start address and placed as the link address of the line. As this is also the address of the next line, it is loaded into the locations used by the routine. The flow now branches back, (the carry flag will be clear), to process the next line. Every basic line will be processed until the end of the program.

On exiting, the program locations $\$ 22$ and $\$ 23$ will hold the address of the two end zero bytes. If this address is increased by two then the end of program address can be derived, and hence the start of variables, as they are one and the same thing vartab.

## Opening up memory

In our memory move routine we made use of a ROM routine when we required more space in a BASIC program. It will move a block up in memory even if its new start is within the original block. Six locations have to be set before entering the routine, which are, in low/high byte order:
\$5A and \$5B- End address of present block
\$5F and \$60- Start address of present block
\$58 and \$59- End address of the new block
ROM LISTING

| A3BF | 38 | SEC |  |
| :---: | :---: | :---: | :---: |
| A3C0 | A5 5A | LDA | \$5A |
| A3C2 | E5 5F | SBC | \$5F |
| A3C.4 | 8522 | STA | \$22 |
| A3C6 | A8 | TAY |  |
| A3C7 | A5 5B | LDA | \$58 |
| A3C9 | E5 60 | SBC | \$60 |
| A3CB | AA | TAX |  |
| A3CC | E8 | INX |  |
| A3CD | 98 | TYA |  |
| A3CE | F0 23 | BEQ | \$A3F3 |
| A3D0 | A5 5A | LDA | \$5A |
| A3D2 | 38 | SEC |  |
| A303 | E5 22 | SBC | \$22 |
| A3D5 | 85 5A | STA | \$5A |
| A3D7 | B0 03 | BCS | \$ 430 C |
| A3D9 | C6 5B | DEC | \$5B |
| A3DB | 38 | SEC |  |


| A3DC A5 | 58 | LDA | \$58 |
| :---: | :---: | :---: | :---: |
| A3DE E5 | 22 | SBC | \$22 |
| A3E0 85 | 58 | STA | \$58 |
| A3E2 B0 | 08 | BCS | \$A3EC |
| A3E4 C6 | 59 | DEC: | \$59 |
| A3E6 90 | 04 | BCC | \$A3EC |
| A3E8 B1 | 5A | LDA | ( $\$ 54$ ), $Y$ |
| A3EA 91 | 58 | STA | (\$58), $Y$ |
| A3EC 88 |  | DEY |  |
| A3ED D0 | F9 | BNE | \$A3E8 |
| A3EF B1 | 5A | LDA | (\$5A), $Y$ |
| A3F1 91 | 58 | STA | (\$58), Y |
| A3F3 C6 | 5B | DEC | \$5B |
| A3F5 C6 | 59 | DEC | \$59 |
| A3F7 CA |  | DEX |  |
| A3F8 D0 | F2 | BNE | \$A3EC |
| A3FA 60 |  | RTS |  |

The immediate action is to calculate the number of bytes to move. The number of low bytes is placed in the $y$ register and location $\$ 22$. The number of pages to move, the difference of the high bytes, is placed in the $x$ register and immediately increased by one. This will be the counter where the zero state is checked to determine completion. As it is decreased before being checked, increasing by one will ensure that all pages will be done. If $x$ was zero and was not incremented, then you would end up going around the circuit 256 times before a zero was discovered in the x register.

The low byte result is checked again; if there is no value, then a large chunk of instructions can be skipped. The bytes between addresses \$A3D0 and \$A3E7 deal with cases where there is an element of an incomplete page of data to move. These lines reduce the two end addresses by the number of low bytes to move. This will not effect the move as the data is loaded and stored with respect to Y and this has the number that was the reduction. The incomplete page is moved first.

Except when $Y$ is zero, all the bytes are transferred within addresses \$A3E8 and \$A3EE. The y register will start at a high value and be decremented to zero. When that is reached, the next bytes are moved separately, before the high addresses are decreased. After this has been achieved, the $x$ counter is reduced and checked, and if it is not zero, it's back to move the next page of data.

From this it can be seen that the transfer is done by taking the high addresses and moving them first. This means that the prograin will not overwrite itself.

## Find a line

This routine finds the start address of a BASIC line, given the line number. We shall use it in our renumber and delete. It uses all three processor registers and locations $\$ 5 F$ and $\$ 60$. On top of that the entry requirement is the line number in low/high byte form in locations $\$ 14$ and $\$ 15$.

ROM LISTING

| A613 | A5 2B | LDA | \$2E |
| :---: | :---: | :---: | :---: |
| A615 | Á 2C | LDX | \$2C |
| A617 | A0 81 | LDY | \#\$01 |
| A619 | 855 F | STA | \$5F |
| A618 | 8660 | STX | \$60 |
| A61D | B1 5F | LDA | (\$5F) , Y |
| A6.1F | $F 1 F$ | BEQ | \$ ${ }^{6} 640$ |
| A621 | C8 | INY |  |
| A622 | C8 | INY |  |
| A623 | A5 15 | LDA | \$15 |
| A625 | D1 5F | CMP | ( $\ddagger 5 \mathrm{~F}$ ), Y |
| A627 | 9018 | BCC | \$A641 |
| A629 | F9 03 | BEQ | \$A62E |
| A62B | 88 | DEY |  |
| A62C | D8 09 | BNE | \$A637 |
| A62E | A5 14 | LDA | 韦14 |
| A630 | 88 | DEY |  |
| A631 | D1 5F | CMP | (\$5F) , Y |
| A633 | 960 C | BCC | \$A641 |
| A635 | FB OA | BEQ | \$ 4641 |
| A637 | 88 | DEY |  |
| A638 | B1 5F | LDA | (\$5F), Y |
| A63A | AA | TAX |  |
| A63B | 88 | DEY |  |
| A63C | B1 5F | LDA | ( $\ddagger 5 \mathrm{~F}$ ) , Y |
| A63E | B0 07 | BCS | \$A617 |
| A640 | 18 | CLC |  |
| A641 | 60 | RTS |  |

Locations $\$ 5 F$ and $\$ 60$ are loaded with the start of BASIC. The high link address is again picked up first to see if the end of program has been reached. The high byte of the line number is checked first. If the value is greater than the required value, the carry will clear and the subroutine is left. If the two values are the same, we go forward to test the low byte values. Failure of either of these checks means that we have not reached the required line and have to go ahead and get the address
of the next line. When the low bytes are checked if they are equal, or the carry flag clear, the routine is terminated. On failing to find the desired line, or on finding a higher one, the link addresses are gathered in and we branch back to check the next line.

Due to the way the checks are made, the routine can be left in one of two states. In the first, the exact line number has been found, in which case the address in $\$ 5 F$ and $\$ 60$ will be what you require. The second state will be that there is no such line number and the routine returns the address of the next highest line. These conditions can be tested in the calling routine by examining the carry flag on return. If the carry is set then the actual line number was found, and if clear it was not.

## 7 Programming aid routines

## Introduction

In Chapter 5 we gave routines to help in the preparation and editing of BASIC programs. These routines were themselves in BASIC, so were slow and had to be tagged onto the end of the resident program. This chapter not only puts these routines into 6502 machine code, but also extends their capabilities. In addition the following are included: OLD, renum, Delete, merge, append, dump, trace, Coder, hex, bin, ten, two, auto, and MEM.

Our object has been to show you that with a little thought and perseverance, adding new BASIC commands is well within your grasp. Most of the routines start with an explanation of what we wish to achieve and how it is possible to do it. This is followed by the assembly listing and the label addresses used. These are provided for assemblers which do not allow the use of labels (Supermon) and with relocation in mind. Finally, a byte-by-byte explanation of the routine is given.

At the beginning of each routine, the command name and parameters are given for use in the utility.

## Renumber

## COMMAND SYNTAX

RENUM start line number or $\emptyset$, increment, new start line value
Using 0 as the first parameter will indicate that the whole program requires renumbering. If a start line is set, it will renumber from that line to the end of the program:
for example, RENUM 0,10,100
RENUM 100,10,200
Later in the chapter we will discuss an auto routine. This is of use when typing in programs where the line numbers are sequential and of a fixed step. Renumbering a program makes it easier to read and opens up space to incorporate new lines.

The system we are going to use is known as a two pass system. The first pass will renumber commands that have line numbers associated with them. This is not as straight-forward as it might at first appear as the commands THEN and RUN have optional line numbers.

There are cases where we do not need to look for a 'renumbering command'. These will be after a data or rem token is encountered, or when inside quotes. In the latter case, we just loop until the next quote or the end of the line is found, whichever is soonest. The procedure on finding the tokens is simply to go to the next statement.

On finding a line number after a command, we convert it from its stored ASCII form to a two byte number. If it is less than the 'start line number', renumbering is not required. When it is not, we calculate its new value, convert it to ASCII, and overwrite the original.

Once all the directive line numbers have been dealt with, the simple task of actually changing the line numbers themselves is carried out.

We will be using many zero page locations in the routine and so to help you to follow the routine, a list of the main ones and what they control is given below:
\$FB and \$FC Address of the current bASIC line being worked on
\$FD and \$FD- Address of the next line - the links of the current line
\$49- Stores the position in the current line, the line marker or y register
\$C9 and \$CA Line number of first line to be renumbered
$\$ 41$ and $\$ 42$ - Address of the first line to be renumbered
\$BC- Value of increment between new line numbers
$\$ B D$ and $\$ B E-$ Value of the new start line number
$\$$ B9 and \$BA- Starts with the same values as \$BD and \$BE and is changed whilst calculating the new line number for directives after keywords.
$\$ 58$ and $\$ 59$ - Starts with the same values as $\$ 41$ and $\$ 42$ and is incremented to give the actual new number of a directive command

ASSEMBLY LISTING

```
    OPEN = $8933
    CLOSE = $888B
    *=$89C5
    10 JSR $81F5 ! GET PARAMETER
    20 JSR $AEFD ! CHECK COMMA
    70 JSR $81F
    80 JSR $AEF
    90 LDA $14
100 STA $BC
110 JSR $81F5 ! GET PARAMETER
    - NEW START LINE #
```

    30
    40
    58
    60
    120
130

LDA \$14
STA \$BD
LDA $\$ 15$
STA $\$$ BE
LDA \$2B
STA \$FB
LDA \$2C
STA \$FC
LDA \$CA
BNE FINDS
LDA \$C9
BNE FINDS ! IS START LINE INPUT 0
LDY \#\$02
LDA ( $\ddagger \mathrm{FB}$ ) , Y
STA \$C9 ! GET FIRST PROG LINE \#
INY
LDA (\$FB),Y
STA $\$ C A$
LDA \$C9
STA $\$ 14$
LDA \$CA
STA $\$ 15$
JSR \$A613! FIND START LINE ADD
BCS STORE ! LINE FOUND
LDX \#\$15
JMP \#A437 ! ERROR - ILLEGAL DIRECT
LDA $\$ 5 F$ ! STORE START LINE ADD
STA $\$ 41$
LDA $\$ 60$
STA $\$ 42$
LDY \#\$00
LDA ( $\ddagger F B$ ) , $Y$
STA $\ddagger F D$ ! GET LINKS TO NEXT LINE INY
LDA ( $\$ \mathrm{FB}$ ) , $Y$
STA \$FE
BNE CONT ! NOT END OF BASIC PROG
JMP RENLM ! CHANGE LINE NUMBERS
INY ! SKIP LINE NUMBERS
INY
INY
LDA ( $\ddagger$ FB), $Y$ ! GET CHAR OF LINE
BNE CONT1! NOT END OF LINE
LDA FFD ! PUT NEXT LINE IN LINE REGISTERS
STA \$FB

| 570 |  | LDA | \$FE |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 580 |  | STA | \$FC |  |  |
| 590 |  | BNE | START | ' | ENFORCED - NEXT LINE |
| 600 | CONT 1 | CMP | \#\$22 |  | IS IT A QUOTE |
| 610 |  | BNE | COnT2 |  | NO |
| 620 | QuOTE | INY |  |  |  |
| 630 |  | LDA | ( $\ddagger \mathrm{FB}$ ) , Y | ! | LOOK FOR NEXT |
|  |  |  |  |  | QUOTE OR LINE END |
| 640 |  | BEQ | LINE | ! | END OF PROG LINE |
| 650 |  | CMP | \#\$22 | ! | QUOTE? |
| 660 |  | BNE | QUOTE | ! | NO |
| 670 |  | BEQ | NEXT |  | YES - NEXT CHAR |
| 680 | CONT 2 | CMF | \#\$8F | $!$ | REM TOKEN? |
| 690 |  | BEQ | LINE |  | YES - NEXT LINE |
| 708 |  | CMP | \#\$83 | $!$ | DATA TOKEN? |
| 710 |  | BEQ | LINE |  | YES - NEXT LINE |
| 720 |  | CMP | \#\#A7 | 1 | THEN TOKEN? |
| 730 |  | BEQ | THEN | 1 | YES |
| 740 |  | CMF | \#\$8A | 1 | RUN TOKEN? |
| 750 |  | BEQ | THEN | 1 | YES |
| 760 |  | CMP | \#\$89 | ! | GOTO TOKEN? |
| 770 |  | BEQ | CONT3 |  |  |
| 780 |  | CMP | \# ¢ CB |  | GO TOKEN? |
| 790 |  | BNE | NOGO |  | NO |
| 800 | SFACE | INY |  |  |  |
| 810 |  | LDA | (\$FB) , $Y$ |  |  |
| 820 |  | CMP | \#\$20 |  |  |
| 830 |  | BEQ | SPACE |  |  |
| 840 |  | CMP | \#\$A4 | $!$ | TO TOKEN AFTER GO? |
| 850 |  | BEQ | CONT3 |  |  |
| 868 | NOGO | CMP | \#\$8D | 1 | GOSUB TOKEN |
| 870 |  | BEQ | CONT3 | ! | YES |
| 880 |  | CMP | \#\# 66 | ! | RESET TOKEN |
| 890 |  | BEQ | CONT3 |  |  |
| 900 |  | BNE | NEXT | ! | NO RELEVENT TOKEN |
| 910 | THEN | INY |  |  |  |
| 920 |  | LDA | ( $\mathrm{F} F \mathrm{FB}$ ), Y | $!$ | GET NEXT BYTE |
| 930 |  | CMP | \#\$20 | 1 | IS IT A SPACE |
| 940 |  | BEQ | THEN | ! | YES - SKIP IT |
| 950 |  | CMP | \#\$30 | ! | LOOK FOR A NUMBER IN ASCII |
| 960 |  | BCS | NUMBER | ! | FOUND A NUMBER? |
| 970 |  | DEY |  |  |  |
| 980 |  | BNE | NEXT | ! | NOT LINE \# AFTER THEN |
| 990 | NUMBER | DEY |  |  |  |
| 1000 |  | CMP | \#\$3A | ! | IS IT A NUMBER |



| 1420 |  | LDA | \$41 | TRANS START ADD TO WORKING REGISTERS |
| :---: | :---: | :---: | :---: | :---: |
| 1430 |  | STA | \$58 |  |
| 1448 |  | LDA | \$BE |  |
| 1450 |  | STA | \$BA |  |
| 1468 |  | LDA | \$42 |  |
| 1470 |  | STA | \$59 |  |
| 1480 | FINDL | LDY | \#\$00 |  |
| 1490 |  | LDA | (\$58), Y ! | SEARCH FOR LINE NO |
| 1580 |  | STA | \$5A | SAUE LINKS |
| 1510 |  | INY |  |  |
| 1520 |  | LDA | ( $\ddagger 58$ ) , Y |  |
| 1530 |  | STA | \$5B |  |
| 1540 |  | BNE | CONT6 | NOT END OF PRQG |
| 1550 |  | LDY | \#\$02 |  |
| 1568 |  | LDA | ( $\ddagger \mathrm{FB}$ ), Y | GET LINE\# FOR ERROR MESSAGE |
| 1570 |  | STA | \$39 |  |
| 1580 |  | INY |  |  |
| 1590 |  | LDA | ( $\ddagger \mathrm{FB}$ ) , Y |  |
| 1600 |  | STA | \$3A |  |
| 1610 |  | LDX | \#\$11 |  |
| 1620 |  | JMF' | \$A437 | ERROR - UNDEF'D STATEMENT |
| 1630 | CONT 6 | INY |  |  |
| 1640 |  | LDA | ( $\ddagger 58$ ) , Y ! | GET AND STORE LINE NO |
| 1650 |  | STA | \$B7 |  |
| 1660 |  | INY |  |  |
| 1670 |  | LDA | (\$58) , Y |  |
| 1680 |  | CMF | \$C.4 | COMPARE FOR SAME LINE |
| 1690 |  | BANE | NEXTLINE! | NOT SAME |
| 1700 |  | LDA | \$B7 |  |
| 1710 |  | C.MP | \$C3 |  |
| 1720 |  | BEQ | FOUNDL |  |
| 1730 | NEXTLINE | LDA | \$ ${ }^{\text {P }}$ |  |
| 1740 |  | CLC | ! | INC REGS TO CALC NEU LINE NO |
| 1750 |  | ADC $\$$ | \$BC |  |
| 1760 |  | STA | \$89 |  |
| 1770 |  | BCC | NOINC |  |
| 1780 |  | INC | \$BA |  |
| 1798 | NOINC | LDA | \$5A |  |
| 1800 |  | STA | \$58 | PUT NEXT LINE ADD IN CURRENT REG |
| 1810 |  | LDA $\$$ | \$5B |  |
| 1820 |  | STA | \$59 |  |


| 1830 |  | BNE | FINDL |  | ENFORCED - CHECK NEXT LINE |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1840 | FOUNDL | LDX | \$89 |  |  |
| 1850 |  | LDA | \$ $\mathrm{BA}^{\text {A }}$ | ! | MSB OF NEW LINE DIRECTIUE |
| 1860 |  | JSR | \$847F | ! | CONUERT TO ASCII <br> - INPUT BUFFER |
| 1870 |  | LDA | \$ BF |  |  |
| 1880 |  | STA | \$3E |  |  |
| 1890 |  | CPX | \$ 3 E | ! | does mem have to move |
| 1908 |  | BEQ | NOMONE |  |  |
| 1910 |  | BCS | OPENUP |  |  |
| 1920 |  | JSR | CLOSE | $!$ | REQUIRES LESS SPACE |
| 1930 |  | JMP | NOMONE |  |  |
| 1940 | OPENUP | JSR | OPEN | $!$ | REQUIRES MORE SPACE |
| 1950 | NOMOUE | LDY | \$49 |  |  |
| 1960 |  | LDX | \# 000 |  |  |
| 1970 | NEXTF | LDA | \$0200, $\times$ | ! | GET NEW NO IN ASCII |
| 1980 |  | BEQ | COHMA | ! | END OF NUMBER |
| 1990 |  | STA | ( $\ddagger \mathrm{FB}$ ), Y | ! | STORE IN PROG |
| 2080 |  | INY |  |  |  |
| 2010 |  | INX |  |  |  |
| 2020 |  | BNE | NEXTF | $!$ | EJFORCED |
| 2030 | camma | LDA | (\$FB), Y | ! | COMMA MEANS ON USED |
| 2040 |  | CMF | \# ${ }^{\text {2 }}$ 2C |  |  |
| 2050 |  | BEQ | ANOTHER |  |  |
| 2060 |  | DEY |  |  |  |
| 2070 |  | JMP | NEXT | ! | GET NEXT TOKEN |
| 2080 | ANOTHER | JMP | CONT3 |  | NEXT LINE - ON COMMAND |
| 2090 | RENUM | LDY | \#\$00 |  |  |
| 2100 |  | LDA | ( $\$ 41$ ), Y | $!$ | GET AND-STARE LINKS |
| 2110 |  | STA | \$5A |  |  |
| 2120 |  | INY |  |  |  |
| 2130 |  | LDA | (\$41), Y |  |  |
| 2140 |  | STA | \$5B |  |  |
| 2150 |  | BAJE | CONT8 | $!$ | NOT END OF PROGRAM |
| 2160 |  | PLA |  | ! | REMONE RETURN ADDRESS |
| 2170 |  | PLA |  |  |  |
| 2180 |  | JMF' | \$A474 | ! | GOTO READY FOR BASIC |
| 2190 | CONT8 | INY |  |  |  |
| 2200 |  | LDA | \$BD |  | NEW LSB LINE NO |
| 2210 |  | STA | (\$41), Y | ! | CHANGE PROG |
| 2220 |  | INY |  |  |  |



| 8B5F | ANOTHER | 8AD2 | CHECK2 |
| :---: | :---: | :---: | :---: |
| 888B | CLOSE | 8B55 | COHPA |
| 842D | CONT | 843E | CONT 1 |
| 844D | COnT2 | 8A8D | C.ONT3 |
| 8AAA | CONT4 | 8AD8 | CONT 5 |
| 8805 | CONT 6 | 8B74 | CONT8 |
| 8B89 | CONT9 | 8499 | DIGITS |
| 84E8 | FINDL | 8 A 03 | FINDS |
| 8B2C | FOUNDL | 8434 | LINE |
| 8A2F | NEXT | 8B4A | NEXTF |
| $8 \mathrm{B17}$ | NEXTLINE | 8470 | NOGO |
| 8 B 22 | NOINC. | 8B46 | NOMCNE |
| 8ACA | NORE | 8488 | NLMBER |
| 8933 | OPEN | $8 \mathrm{B43}$ | OPENUP |
| 8442 | QUOTE | 8B62 | RENLM |
| 8465 | SFACE | 8A1D | START |
| 8815 | STORE | 8A7A | THEN |

LINES 10-150: The parameters are gathered here and put into their registers. Commas separating the inputs are also checked, giving 'syNtax errors' if not present.

LINES 160-190: The start of the BASIC program is now put into the current line registers, as it is also the address of the first line.

LINES 200-290: If the first parameter input was zero, indicating a full program renumber, then the first line number is found and stored in its appropriate register.

LINES 300-410: Although we do not need this at the moment - here we find the address of the start line. We use the rom routine find basic line
(see Chapter 6, page 142). If the carry flag is set, it will mean that the start line requested was not found and an 'Illegal direct' error will be printed. The address, if found, will be stored in $\$ 41$ and $\$ 42$.

LINES 420-1020: The byte-by-byte search for the appropriate keywords starts here. We start at the beginning of the BASIC program, no matter what the start line requested. As soon as the link addresses are collected and stored, the end of the program is checked for. Passing this means that only the actual line numbers require changing, so it is off to the final section of the whole routine, which is described later.

To continue finding the tokens, we skip the line numbers as they are not required here. Lines 540 to 590 set the values for the next line, after the end of line zero is discovered, and branches back to process the next line.

There is nothing of interest to us in quotes so on finding one we go into a loop to find a second quote or the end of the line. This is carried out in lines 620 to 670 .

The next two tokens checked for are data and rem. Encountering these indicates that we can proceed to the next line as there will be nothing further to renumber in these lines.

There are two keywords - RUN and THEN - that may, or may not, have line numbers. These will therefore branch to check this possibility before proceeding. The standard Commodore directive commands are next in line: goto, go to and gosub. The centre keyword is checked in two stages, first for the co. A loop is then set up to skip over spaces and then the to token is looked for. All three keywords on being found will cause the routine flow to branch further ahead than RUN or THEN, as it assumes they will have line numbers. If not, then unless you have a line number of 0 , an error will be detected later.

The last keyword to be checked for is one of our new ones: RESET.
If line 900 is reached then we have not found a relevant keyword and therefore branch back to get the next program byte.

The last part in this section is to check if the next significant byte after run or then is a number, in ASCII. Spaces are skipped over and checks are made for values between $\$ 30$ and $\$ 39$ inclusive to continue.

LINES 1030-1390: The line numbers after the keywords will be in ASCII form and the line itself in two byte form. To do our calculations, we want both in the same two byte format. ASCII numerals therefore have to be changed.

After skipping spaces, we store our y register (so we know where to write our new line number from), which is the line marker. Proceeding, we pick up bytes until a non-numeral is found, and store them in the input buffer. The $x$ register is used to count the number of digits and is stored for later use. To convert the ASCII into the form we require, we use the GET PARAMETER routine. For this to work, we perform two operations. First, we make sure that after the last line number digit
there is a non-numeric character by storing a colon there. Secondly, we set Chrget to point to the first numeral - $\$ 0200$.

The converted result is taken from registers $\$ 14$ and $\$ 15$ and stored in the two we have designated.

We only wish to renumber from the start line in the command. Lines 1290 to 1390 check this by comparing the two values. If no renumber of that particular line is required, we retrieve our line marker and increase it by the number of digits in the directive number, as we do not require to check them again. This is then transferred to the y register. It will actually point to the byte following the last digit but this is taken into account in what follows. We jump further ahead to a position noted as COMmA (described later) starting at line 2030.

LINES 1400-1830: Having got this far, we have found a number which requires a different value. To find the new value, we have to go through the program from the designated start line and find the line that it points to (remember we have not changed the actual line numbers yet). At the same time, we calculate the new value.

To start with, we take the address of the start line and store it in $\$ 58$ and $\$ 59$. We then take the new value for the start line, the third parameter in the command, and place it in \$B9 and \$BA. The line number we are checking for is held in \$C3 and \$C4.

As before, we get and store the link addresses, but here if we discover the end of the program has been reached (high link address of zero), an error is present as the line number of the directive was not found. The error produced is the same as when RuNing a program - 'undef'd statement'. To make it easier for you, we also print out the line number with the error.

As long as this is not encountered, collecting a line number and comparing it for a match comes next. If it is not the one we want, the new start line number is increased by the increment value. This will calculate the value for the following line. The value is only increased on not finding a match, which conversely means that when the line searched for is located, the value is ready and waiting. After incrementing, we transfer the links to the line registers and branch back to check further lines.

LINES 1840-1940: We now have our new line number and need only to insert it into the line after the token, overwriting the original directive.

The new value is in two byte form and so requires converting to ASCII form. We do this by using a routine earlier in the UTILITY. This requires the accumulator to be loaded with the high byte and the $x$ register with the low byte. We then call our 'convert to ASCII' at $\$ 847 \mathrm{~F}$. This will do the conversion and store the answer in the input buffer, with a zero signifying the end. Also returned is the number of characters in the $x$ register. If this value is the same as the original number,
the value stored in location \$BF, we can just overwrite with no problems and proceed to line 1950.

The number of characters is transferred to location \$3E, via the accumulator. This location is for the memory move routine described in Chapter 6. By comparing the value in the $x$ register (the new number of digits) with the accumulator figure (the old number of digits), we determine if a move is required. If the $x$ value is less, the Close routine is called; if it is greater, the OPEN routine is called; if it is the same, no move is required.

LINES 1950-2020: This leaves us one thing to do which is to write in the new number. First, we reload the $y$ register with the line marker. This points to the position of the first digit in the number. By increasing $x$, starting at zero, and $\gamma$, we take the digits from the input buffer and store them into the program line. This is repeated until we collect a zero from the buffer, when the branch in line 2020 will fail.

LINES 2030-2080: When we went through the lines checking for tokens, we did not look for the on statement. The reason for this is that the only time a comma should be used after a line directive following a gosub or goto is when the on keyword has been used.

On entering these lines the y register will point to the byte following the last one checked or stored. We load the accumulator with that byte to see if a comma is present. Finding it means that we branch back to the position just after the token search where it will commence with gathering in the line number directive for processing.

If the comma is not present, the $y$ register is decreased and we go back to start the search for the next appropriate token. The $y$ has to be decreased in case it was the end of the line, otherwise it would not have mattered.

## LINES 2090-2350: RENUMBERING THE LINE NUMBERS

All the directive line numbers have now been checked and processed where required. The only thing which remains is to renumber the program lines themselves.

This is started from the line number requested in the first parameter of the command and whose address is held in $\$ 41$ and $\$ 42$. Its new starting value is held in $\$ B D$ and $\$ B E$, and these will be incremented by the value in $\$ B C$ for each line. We progressively go through the program inserting the new numbers, in two byte form, until the end of the program is reached. After each line the number value is incremented ready for the next line.

When the renumbering is completed, we take the return address from the stack and jump to the start of BASIC to await further instructions. We do this to break out of the program in the unlikely event of RENUM being initiated within it.

## A WORD OF WARNING

We cannot stress strongly enough that a copy of the original program should be saved to tape, or disk, prior to renumbering. We all make mistakes and if the renum finds a non-existent line number after, say, a GOTO, then an error is produced. This leaves the program only partially renumbered.

## Auto

COMMAND SYNTAX
aUTO first line number, increment
To escape from auto simply press return immediately after the line number, as if deleting a single line.

This command removes the need to type in line numbers. The user just decides the start line number and the increment between consecutive lines. To achieve this we want a place to break into the normal flow of BASIC. This is made possible as every time an input line is processed it goes to a vectored jump before it is ready to receive the next. This is called the basic Warm Start Vector which is at $\$ 0302$ and $\$ 0303$. By changing the address in this vector to a routine of our own we can calculate the line number, put it into ASCII form and then insert it into the keyboard buffer just as if you typed it yourself. Having completed this, we then return to the input routine for you to make up the program line.

ASSEMBLY LISTING


| 150 |  | BEQ | EXIT |  | TURN OFF AUTO |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 160 |  | LDX | \$FB |  |  |
| 170 |  | LDA | \$FC | $!$ | NEXT LINE NLMBER |
| 180 |  | JSR | ASCII | ! | PUT LINE \# IN ASCII |
| 190 |  | STX | \$ ${ }^{\text {c } 6}$ | ! | NO OF CHARS IN KEYBOARD BUFFER |
| 280 | NEXT | L[AA | \$0208, X | $!$ | PICK UP ASCII |
| 210 |  | STA | \$0277, X | ! | PUT IN KEY BUFFER |
| 220 |  | DEX |  |  |  |
| 230 |  | BPL | NEXT |  |  |
| 240 |  | CLC: |  |  |  |
| 250 |  | LDA | \$FB |  |  |
| 260 |  | ADC | \$FD | ! | INCREMENT LINE NUMBER |
| 270 |  | STA | \$FB |  |  |
| 280 |  | BCC | NOINC: |  |  |
| 290 |  | INC | \$FC |  |  |
| 308 | NaINC: | JMP | \$A483 | ! | READY FOR BASIC |
| 310 | EXIT | LDA | \#\$83 |  |  |
| 320 |  | STA | \$0302 |  |  |
| 330 |  | LDA | \#\$A4 |  |  |
| 340 |  | STA | \$0303 |  |  |
| 350 |  | JMP | (\$0302) |  |  |
| 360 | ASCII | STX | \$63 | ! | LOW BYTE |
| 370 |  | STA | \$62 | $!$ | HIGH BYTE |
| 380 |  | LDX | \#\$90 |  |  |
| 390 |  | SEC |  |  |  |
| 400 |  | JSR | \$BC49 |  |  |
| 410 |  | JSR | \$BDDF |  |  |
| 420 |  | JSR | \$B487 |  |  |
| 430 |  | JSR | \$B6A6 |  |  |
| 440 |  | LDX |  |  |  |
| 450 | AGAIN | LDA | \$0100, X |  |  |
| 460 |  | STA | \$0208, X | ! | PUT ASCII CHARS IN INPUT BUFFER |
| 470 |  | BEQ | FINISH | ! | 2ERO CHAR WAS FOUND |
| 480 |  | INX |  |  |  |
| 490 |  | BNE | AGAIN |  |  |
| 500 | FINISH | RTS |  |  |  |
| 8494 | AGAIN |  | 847F ASCI |  |  |
| 8440 | AUT0 |  | 8472 EXIT |  |  |
| 849F | FINI EH |  | 845 B NEXT |  |  |
| 846F | NOINC |  |  |  |  |

LINES 10-130: These lines are only used when the aUto command is active. We take the parameters of the command and place the start line number in \$FB and \$FC and the increment in \$FD. As we only take the low byte of the increment, the multiples over 256 are ignored. The only syntax check is made in line 20, where we check for a comma between the two parameters. Lastly in this section we change the vector address to point to the auto numbering which starts in line 140.

LINES 140-300: The first thing is to see if the first character in the input buffer is zero. This will signify no basic coding was inserted in the line and that you want to cancel the auto routine. When a basic line is typed, the line number is put into the input buffer. During the processing stage of inserting the line into the program basic takes out the line number and moves the rest of the line back up the buffer to overwrite it. This means that if the first input after a line number was a RETURN (BASIC inserts a zero for that as an end of line marker) then the first character in the buffer will be a zero after the line number has been removed. Therefore, on finding a zero we will branch off to exit the auto mode.

Assuming that we are still auto-numbering, we take the values in \$FB and \$FC and go off to convert them to ASCII form. This we will come across shortly. On returning from that subroutine the $x$ register will have the number of ASCII characters in the line number and they will be in the input buffer. The $x$ value is stored in the register which tells the operating system how many characters will be in the keyboard buffer. Having done that, we transfer the characters from the input buffer to the keyboard buffer.

We now set the line number for next time by adding the increment in $\$ F D$ to the values in $\$ F B$ and $\$ F C$.

That is all there is to do, so we return you to the normal basic flow where the input routine will take the line number from the keyboard buffer, place it in the input buffer and print it on the screen.

LINES 310-350: These lines will be operated when you want to exit from auto. All we do is restore the basic Warm Start Vector to its initial value and then return you to BASIC to wait for your next instruction.

LINES 360-500: CONVERT TO ASCII: This subroutine will also be used by other commands when they require a one or two byte number converted into ASCII form. The subroutine is entered with the low byte of the number in the $x$ register and the high byte in the accumulator.

The conversion is carried out by four ROM routines, but before we can call them we have four items to set. First, the number to be converted is transferred to locations $\$ 63$ (the low byte) and $\$ 62$. These are part of the floating point accumulator \#1 (FAC\#1) which is the main number manipulation area for BASIC. The other two prevent certain actions in the conversion process. Setting the carry flag will bypass a
routine that will perform the complement of the number and loading $x$ with $\$ 90$ will set the Exponent byte of FAC\#1 to avoid getting an answer in exponent form.

The first ROM routine visited clears all the bytes in the fac\# (or sets them to default values) which we have not dealt with. The next routine does the actual conversion. The remaining routine puts the result into a string and places it at the bottom of the stack area. The last byte placed there will be zero to mark the end.

We cannot leave it there as BASIC often uses this area. We therefore transfer it to the input buffer where it can be taken and used by the coding calling this routine. On exit the value in the $x$ register will be the number of ASCII characters in the conversion.

## Merge and append combining BASIC programs together

COMMAND SYNTAX
MERGE "program title", device
APPEND "program title", device
The default device is tape and if there is no program title, the first program found on the tape will be used.

Merging programs together means that they are weaved together in program line order. The result is as if you typed in the lines of the merging program at the keyboard. This also means that if the programs both have lines with the same number, the ones in the merging program will overwrite the original.

Appending a program to another is simply a process of tagging it onto the end of the one in memory, irrespective of line numbers.

Merging is the more complicated of the two programs, but is not really complicated in itself. Both programs are initially loaded at the end of the current memory program, APPEND overwriting the last two bytes whilst merge comes just after them. The last two bytes of the program are the unique link address of zero, signifying the end of a BASIC program. By overwriting them on APPEND, we achieve our aim immediately and all that remains is to amend the link addresses to continue the program flow and to reset the end of program pointer.

In merging we take each new line in turn and insert it in the main program - if we had overwritten the original end links we would both merge and append which we do not want. To incorporate the new lines we make use of the normal basic input routine. After you input a program line and press Return, basic takes off the line number and then tokenizes all the basic keywords. At this point the line number is in two registers, in a two byte form rather in than the ASCII form typed in, and the line's content has been moved to the beginning of the input buffer. There is a counter of the number of bytes in the line which is four
greater to incorporate the line number and link address．BASIC therefore knows the total space required for the line．We will enter the ROM routine at this point with the appropriate data set．Unfortunately， it is not a subroutine but finishes up waiting for an input．We therefore have to change the same vectors as in the aUTO routine（the BASIC Warm Start）to point us back to continued merging until we reach the end．

ASSEMBLY LISTING
$9 *=\$ 87 \mathrm{~A} 7$

| 10 | JSR SETADD | ！SET ADDRESSES |
| :--- | :--- | :--- |
|  |  | FOR MERGE FROG |
| 20 | STX $\$ 2 B$ | ！SET MERGE PROG START |
| 30 | STY $\$ 2 C$ |  |
| 40 | LDY \＃$\# 80$ | ！PUT ZERO AT 1ST |
|  |  | LOCATION |

50
60
70
80
90
108
110
120
130
140
150
160 JOIN
170
180
190
200
210
220
230
240
250
260
270

## 280

## 300

310

TYA
STA（ 12 B ），Y
JSR LOAD
STX $\$ 2 \mathrm{D}$ ！END OF MERGE PROG
STY \＄2E
JSR $\ddagger$ A533 ！RECHAIN MERGE PROG
JSR RESET1 ！RESTORE POINTERS TO ORIG PROG
LDA \＃〈MERGE
LDX \＃） HERGE
STA $\$ 0302$ ！CHANGE WARM START UECTOR TO MERGE
STX $\$ 0303$
LDA \＃\＃ 01
STA $\$ 7 B$
LDA \＃まFF
STA $\$ 7 A \quad!$ SET UP CHRGET
LDY \＃ま00
LDA（ $\ddagger \mathrm{FB}$ ），$Y$ ！GET AND STORE LINKS
STA $\$$ FD
INY
LDA（ $\ddagger$ FB），$Y$
STA \＄FE
BEQ EXIT ！END OF MERGE PRG
INY
LDA $(\neq F B), Y$ ！GET AND STORE LINE NO
STA $\$ 14$ ！BASIC EXPECTS THEM IN THESE LOCATIONS
INY
LDA（ $\ddagger \mathrm{FB}$ ），$Y$


```
        700 RESET LDA $FB
```

710
720
730
740
750
768
770
780
790
300
810
820 RESET 1
830
840
850
860
870
880
890
900
910 SETADD
920
930
940
950
960
970
980
990
1000
1010
1020
1030
1040
1850
1060
1070
1080
1890
1180
1110
1120
1130
1140

LDA FFB
SEC
SBC \#\$02
STA $\$ 14$
LDA \$FC
SBC \#\$60
STA \$FC
LDA \#\$00
TAY
STA (\$14),Y ! RESTORE TWO ZEROS
AT END OF PROG
INY
STA ( $\$ 14$ ) , $Y$
LDA $\$ F D$
LDX FFE
STA \$2B
STX \$2C
LDA $\ddagger F B$
LDX \$FC
STA $\$ 2 \mathrm{D}$
STX $\$ 2 \mathrm{E}$
RTS
LDA $\$ 2 B$
STA \$FD
LDA $\$ 2 C$
STA \$FE
LDX $\$ 20$
LDY \$2E
STX FFB
STY \$FC
RTS
! APPEND ROUTINE
JSR SETADD
TXA
SEC
SBC \#\$02
STA \$2B
TYA
SBC \#\#00
STA \$2C
JSR LOAD
STX \$FB
STY $\$ F C$
JSR RESET 1
JSR \$A533
RTS

| 8830 ERROR | 8803 EXIT |
| :--- | :--- |
| $882 F$ EXIT1 | 87 CA JOIN |
| 8810 LOAD | $87 F 9$ MERGE |
| $87 E B$ NEXT | 8838 RESET |
| 8840 RESET1 | $885 E$ SETADD |

This time we are not going to describe the program in line number order. There are three subroutines in the body of the program used both by merge and APPEND and we will deal with these first.

LINES 910-990: SETADD: This simply takes the start and end addresses of the original program and temporarily stores them. On coming out of the subroutine, the $y$ register will contain the high byte of the end address and the $x$ register, the low value.

LINES 700-900: RESET: The first 12 lines will only be encountered when there is an error in loading the secondary program. These simply ensure that the end of program zeros are at the end of the original program. This will mean that when exiting from either command your original program is intact before starting.

The remaining lines are the reverse of SETADD, that is, they take the values in the temporary registers and place them in the program end and start registers. These last lines are called in the assembly listing as RESET1.

LINES 520-690: LOAD: The first thing this subroutine does is to call one resident in the BASIC ROM used by the standard LOAD and SAVE commands. It gathers up the parameters and sets various registers according to that information, and as it is there we also make use of it. We are going to do a relocated load and if a secondary address is present this will override our objective. To correct this we load location $\$$ b9 with zero to bring back the state for a relocated load.

The kernal load routine expects the start address of the load in the two processor registers, $x$ and $Y$, with the former holding the low byte of that address. The accumulator is the flag for either a load or a verify operation. The value for load is zero, the other being one, which was set whilst confirming the secondary address. The KERNAL LOAD routine is situated at \$FFD5. Error checking comes now in the order of operations - you may have put in the wrong tape or disk. The first indicator to a bad load is the carry flag being set; if this is so, then we branch off to deal with it. We have to check the $1 / 0$ status word if the carry is clear. This is achieved by calling another KERNAL routine, at \$FFB7. The result coming out of this call is ANDed with the value $\$$ BF and everything is fine if the zero flag is set. The error given for any other outcome is 'lOAD ERROR'.

On the first check we go to line 660 if the carry was set. The value in the accumulator will be used for the error so we temporarily store this
on the stack. The reason for this is to reset all the pointers we altered to give you back your original program. This is done by a call to RESET, described above. Once done we retrieve the accumulator value and jump to part of the BASIC loading routine for an error to be generated, based on the accumulator value.

## MERGE

LINES 10-150: First of all we call the SETADD routine. From this we can set the start address of the merging program to immediately after the memory program. BASIC expects the first byte to be zero and therefore we oblige it. Having done that, we load the program we want to merge. The address that is returned from loading is placed in the end of program address. At this point as far as BASIC knows the only program in memory is your merge program. The link addresses have to be set up so we know where the lines start. This is done by a call to the rom routine at $\$ A 533$ which will do this for us. From now on we want BASIC to respond to the original program and therefore we reset all the registers back to their original values through RESET1.

To merge the lines into the master program we use the program line input routine in the BASIC ROM. This is not a stand alone routine but ends up at the BASIC Warm Start after each line has been processed. It follows that, as in the AUTO routine, we will have to alter the vector pointing to that position to divert to this routine until all lines are merged. This is done in the lines 120 to 150 . The addresses will be in the location of line 420.

LINES 160-460: These are the instructions which actually combine your two programs. First, CHRGET is set to a position one place before the start of the input buffer. We now turn our attention to the merging program lines. The address of the first line will still be in locations \$FB and $\$ F C$ following the SETADD routine called at the beginning. Using this information, we get the link addresses and store them for later. We use this system a lot in our routines but in this case it is vital. We have left no room between the programs so that when the line is transferred the master program will be longer, overwriting the needed link addresses. The check is also made for the end of the merge program, the usual high byte zero link address.

Next we take the two byte line number and place it in registers so that BASIC will know where to find them, $\$ 14$ and $\$ 15$. The contents of the line, including the zero byte end marker, are now transferred to the input buffer starting at $\$ 0200$. The listing shows $\$ 01 \mathrm{FB}$ but there is a reason for this. The $x$ register will come out with a value of the number of bytes in the line, but to account for the four bytes holding the link addresses and the line number, it starts with a value of 4 . This is one more than required but the ROM routine will compensate. This means that with the initial value of $\$ 04$ in $x$ the first location written to will be
$\$ 0200$. The ROM routine we are using wants the number of bytes in the $Y$ register rather than the $x$, so we oblige by transferring them via the accumulator.

We are now ready to use the rom to merge our line into the main BASIC program. We join the ROM just after the coding that turns the keywords into tokens - ours are in that state already. As far as the ROM is concerned you have typed in the line and will put it in the program as such.

The basic Warm Start Vector will bring us back after inserting the line to LINE 420 in the above listing. We now put the address of the next line into the working registers and branch back to deal with it. The branch instruction will always succeed as we have checked previously for a zero value in $\$ F E$.

These lines will be repeated until all the merge program lines have been assigned to the master program.
LINES 470-510: The merge is complete so we restore the BASIC Warm Start Vector to its normal setting and return to BASIC where it will await further commands.

## APPEND

LINES 1010-1140: The first thing, as in MERGE, is to call the SETADD routine. On coming out though, it is slightly different. The new basic program start has to be reduced by two so will overwrite the end of memory program links. The appending program will load directly after the final line of the master program. On completion of the load, we store the loading end addresses (not in the end of program registers), but overwrite the original stored values, set by SETADD. This means that on resetting the values, by RESET1, we end up with the original start and the end marker corresponding to end of the appended lines. The final thing to do is to set the link addresses to follow as one program. This is carried out by the ROM routine at \$A533. The two programs have been joined together with our own form of 'superglue'.

## Delete

## COMMAND SYNTAX <br> DELETE first line number, [second line number]

The first line number to be deleted and the comma are essential and if missing will give errors. The second line number is optional in that if you want to delete to the end of the program you omit it; otherwise insert a number.

Deleting a line of BASIC program is easy, you just type in the line number followed by a return. There is no real hardship in deleting one or two lines but longer blocks become tedious and time-consuming.

DELETE will rid you of a block of lines with one command. To do this we use the same rom routine as if deleting one line. What the basic rom does is to take the address of the line to be deleted and the link address within that line. It then takes the program starting at that link address to the end of the program and moves it to the address of the line to be deleted. For example, if we have a line whose address is \$0901, its link address is $\$ 0925$ (the start of the next line), and the end of the program is \$0A45, the block to move is \$0925 to \$0A45 with its new starting address of $\$ 0901$. Hence, the line at $\$ 0901$ is overwritten or deleted. By the way, the variables and the arrays are also moved along with the block.

It therefore goes to show that if we get the address of the next line after the 'second line number' and place it in the link address of the 'first line number', a whole block of lines will be overwritten at once. Where there is no 'second line number' we take the end of program address and deduct two from it. This will give us the address of the two zero bytes at the end of a program. The first line number is placed in the input buffer, with a zero at the end of it signifying no further data, and goes to ROM as if you typed it in.

Two listings follow for this command. The reason for this is explained later.

ASSEMBLY LISTING1

```
*=$8F44
\begin{tabular}{|c|c|c|c|}
\hline BCC & DEL & 1 & PARAMETER A NUMBER \\
\hline JMP & \$AF08 & \(!\) & GENERATE SYNTAX ERROR \\
\hline JSR & \$81F5 & 1 & GET 1 ST LINE NUMBER \\
\hline JSR & \$A613 & ! & FIND LINE LOCATION \\
\hline BCS & FQUND & \(!\) & LINE NUMBER FOUND \\
\hline LDX & \#\$15 & ! & ILLEGAL DIRECT ERROR \\
\hline JMP & \$A437 & ! & ERROR ROUTINE \\
\hline LDA & \$5F & ! & PUT LINE ADDRESSES \\
\hline STA & \$FB & ! & IN STORAGE \\
\hline LDA & \$68 & & \\
\hline STA & \$FC & & \\
\hline JSR & \$0879 & ! & CHECK FOR COMMA \\
\hline CMP & \# \({ }^{\text {2 }}\) C & & \\
\hline BNE & SYNTAX & ! & NOT FOLND \\
\hline JSR & \$0073 & \(!\) & GET NEXT BYTE \\
\hline BNE & NLMERAL & ! & A SECOND LINE NUMBER \\
\hline SEC & & ! & PREFARE FOR SUBTRACT \\
\hline LDA & \$2D & & END OF PROGRAM \\
\hline SBC & \#\$02 & ! & DEDUCT BY TWO \\
\hline STA & \$5F & ! & READY FOR DELETION \\
\hline LDA & \$2E & & END OF PROGRAM \\
\hline
\end{tabular}
```

10
20
30 40
50
60 70
80 FOUND 90
100
110
120
130
140
150
160
170
180
190
200
218

| 220 |  | SBC | \#\$00 |  | in case of page CROSSING |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 230 |  | STA | \$60 | ! | READY FOR DELETION |
| 240 |  | BNE | CCHT | ! | ENFORCED $\$ 2 E$ CAN'T BE ZERO OR 1 |
| 250 | NUMERAL | BCS | SYNTAX | $!$ | NO NUMBER |
| 260 |  | JSR | \$81F5 | $!$ | GET 2ND LINE NUMBER |
| 270 |  | JNC | \$14 | $!$ | SO WE GET FOLLOWING LINE |
| 280 |  | JSR | \$A613 | $!$ | FIND LINE |
| 298 |  | LDA | \$FC | ! | CHECK IF $15 T$ NO IS SMALLER THAN 2ND |
| 300 |  | CMP | \$60 |  |  |
| 310 |  | BCC | CONT |  |  |
| 320 |  | BNE | SYNTAX |  |  |
| 330 |  | LDA | \#FB |  |  |
| 340 |  | CMP | \$5F |  |  |
| 350 |  | BCS | SYNTAX |  |  |
| 360 | CONT | LDY | \# $\ddagger 00$ |  |  |
| 370 |  | LDA | \$5F |  |  |
| 380 |  | STA | ( $\ddagger \mathrm{FB}$ ), Y | ! | STORE ADRESS |
| 390 |  | INY |  |  |  |
| 400 |  | LDA | \$60 |  |  |
| 410 |  | STA | ( $\ddagger \mathrm{FB}$ ) , Y |  |  |
| 420 |  | INY |  |  |  |
| 430 |  | LDA | ( $\ddagger \mathrm{FB}$ ), Y | ! | GET 2 BYTE LINE NO |
| 440 |  | TAX |  |  |  |
| 450 |  | INY |  |  |  |
| 460 |  | LDA | ( $\ddagger \mathrm{FB}$ ), Y |  |  |
| 470 |  | JSR | \$847F | $!$ | CONVERT TO ASCII |
| 480 |  |  |  | ! | AND FUT INTO START OF INPUT EUFFER |
| 490 |  | PLA |  | ! | REMONE RETURN ADDRESS |
| 500 |  | PLA |  |  |  |
| 510 |  | LDX | \# $\$ \mathrm{FF}$ | ! | to initailize INPUT BUFFER |
| 520 |  | LDA | \#\$01 |  |  |
| 530 |  | JMF' | \$9270 | ! | WILL DELETE LINE AND RETURN TO BASIC |
| 8F91 | CONT |  | 8 F 49 DEL |  |  |
| 8F56 | FOUND |  | 8F79 NuM | RAL |  |
| 8F46 | SYNTAX |  |  |  |  |

ASSEMBLY LISTING2


## LISTING1

LINES 10-110: These instructions deal with the 'first line number'. The routine first checks the carry flag, set or unset by CHRGET on entering, and if set a syntax error is generated as the first byte after the delete token was not a numeral. A call to our GET PARAMETERS routine is next, immediately followed by a visit to the ROM routine FIND BASIC LINE. The result from cet parameters is in the registers used to call find basic line. On returning from the latter, if the carry is not set, then the line was not found and we therefore generate a further error. The address of the first line is placed in locations \$FB and \$FC.

LINES 120-280: The remaining parameter is now dealt with. CHRGET is positioned to where the comma should be, so we call Chrgot to see if it is there. A call to CHRGET now will get the first byte of the second line number. If no line number is present, the zero flag will be set. In that
case we gather in the address of the end of the program, deduct two from it, and store the result in registers $\$ 5 F$ and $\$ 60$.

If the zero flag was not set, we make a further check as earlier to see if the byte picked up by CHRGET was a numeral. GET PARAMETER is called to get the second line number and the low byte result in $\$ 14$ is increased by one. This is done as we do not require the address of that line but rather the one following it. After the visit to FIND BASIC LINE the address in $\$ 5 F$ and $\$ 60$ will be the next line, whether its line number is one or ten greater than the 'second line number'.

LINES 290-350: These instructions check to see that the address of the 'second line number' is higher than that of the first, otherwise a syntax ERROR is given.

LINES 360-480: Here we insert the second address we found into the link address position of the first line. We then get the line number, in its two byte format, putting the low byte in the $x$ register and the high into the accumulator. We now call another routine which we coded at $\$ 847 \mathrm{~F}$ in the Utility where the number will be converted to ASCII and placed in the input buffer, starting at $\$ 0200$, with a zero at the end.

LINES 490-530: From the stack we remove the return addresses which were placed there on entering delete. The x register and the accumulator are given the address of the input buffer less one which will be the CHRGET address. The final thing is to jump to the second listing.

## LISTING2

LINES 10-230: The ROM routine that we will use is not a subroutine but ends up at the basic Warm Start Vector. We want to return here so we first store its present values and replace them to point back to these lines. We now go to ROM where it will treat the number in the input buffer as if you were deleting a single line from the keyboard, but as we have changed the link address it will delete more than one.

On returning, we restore the BASIC Warm Start Vector. We now subject the program to the rechain routine - not that it requires it, but from this routine we can calculate the end address. From the address the rechain routine ends with, we add two and set the end of program registers. A call to the CLR routine will set the remaining variable addresses. Finally, we jump to BASIC, printing 'READY' and give you back control.

The reason for two listings is due to the way in which the rom memory moving routine sets the end of program address. We came across this when testing the UTILITY. The basic normally expects lines of around 80 characters and definitely no more than 255. Mainly for the latter reason the ROM routine only decreases the end address by the maximum of a page. It does not affect the deletion, but it did not make the required
reduction in memory used. The second listing was added to overcome the times when the number of lines took more than 256 bytes. Thus in the second listing we were able to set the addresses ourselves.

Memory - Display number of bytes free
COMMAND SYNTAX
MEM
There are no parameters in this command. The command is available only in direct mode. If found in a program the routine is not carried out.

BASIC has a command that prints out the amount of space available to it. It is FRE( x$)$ where x is a dummy argument. Unfortunately it returns, when used with PRINT, an integer value which means any value over 32767 (\$7FFF) will be a negative number. For example, if the number of bytes free is 36500 , the result printed would be -29935 . If you add that, with the sign, to 65535 (SFFF), you will arrive at the true figure of 36500 . We produce here a short routine to print out the correct value straight away. Having said that, with the utility in place, the maximum space available is less than 32768 and so $\operatorname{FRE}(\mathrm{x})$ will always print out the correct value.
The first thing to do is call a rom routine to do a 'garbage collection'. It is at $\$ B 526$ and tidies up the variable and string area. It will reset the necessary registers after the compaction. The area of memory that is unused will be from the end of arrays to the beginning of the area used by strings. If we take the higher address from the lower, we will have the number of free bytes available.
The routine that we have used to print the result to the screen is a subroutine of the hex command, which is described later. Suffice to say that on calling this subroutine with the low byte in the $y$ register and the high in the accumulator, it will convert it to ASCII and print the result to the screen.

To check whether you are in direct mode, we look at location \$9D(157). This will hold $\$ 88$ (128) for direct or $\$ 00$ for program mode.

ASSEMBLY LISTING

| $\xi$ |  |  |
| :--- | :--- | :--- |
| 10 | LDA $\$ 90$ | ! DIRECT OR PROGRAM |
| 20 | BNE MEM | ! DIRECT ONLY |
| 30 | RTS | ! PROGRAM NOT EXECUTED |
| 40 MEM | JSR $\$ B 526$ | ! ROM COLLECT GARBAGE |
| 50 | SEC | ! PREPARE FOR SUBTRACT |
| 60 | LDA $\$ 33$ | ! POINTER START OF |
|  |  | STRING STORAGE |
| 70 | SBC $\$ 31$ | ! POINTER END OF ARRAYS |


| 80 | TAY | ! TEMP STORE |
| :--- | :--- | :--- |
| 90 | LDA $\$ 34$ | ! POINTER START OF |
|  |  | STRING STORAGE |
| 100 | SBC $\$ 32$ | : POINTER END OF ARRAYS |
| 110 | JMF $\$ 8540$ | I CONNERT TO ASCII |
|  |  |  |
|  |  | AND PRINT TO SCREEN |

8601 MEM

## Coder

## COMMAND SYNTAX <br> CODER

There are no parameters to this command.
How many times have you picked up a listing from a magazine and wondered what graphic symbol is in that PRINT statement? Is it a shifted N graphic or shifted L? How many have been used together, is it 2 or 3 ? You then come across a colour code and have to look it up in the manual to remember which colour to program. Owners of nonCommodore printers also have a problem as these symbols and graphics do not print.

We would like to introduce a routine that replaces these graphics with mnemonics. For example, the symbol for clear screen would be replaced by [cıs|.

Except for one, all the codes we want to change appear within quotes. That means we have to look through the program for a quote and when found look for ASCII values that we want to change until the end of the line or the second quote appears. Having found one, we also have to look to see if it has been repeated. This done, we will either calculate the new code or find one in a data table. The codes produced will be of a different length from the original, but if it repeats, may be of shorter overall length. To accommodate this, we will use the memory move routines described in Chapter 6 .

The one exception we mentioned earlier is the mathematical ' PI ' ( 3.14159 , etc). This we also found does not appear on some listings and is essential if in a mathematical equation. This is therefore coded whether in or out of quotes.

Most of the program operation is described after the assembly listing (see below), but before the listing we would like to say a word or two about the data make-up. This can be split into two sections. Graphics that are obtained by using the shift with most of the 'letter' keys can be calculated directly to the ASCII code of that particular letter. The remaining graphics and codes require the use of data tables.

We have employed two tables and stored them out of the way under the BASIC ROM. The first table, the data address table, has the three bytes
for each character we are going to encode. The first byte is the ASCII value of the character and is followed by the address within the second table where the data is stored. The second table, the data table, holds all the data for those characters. The data will be the characters printed between the [ ] brackets, and may be of differing length. Because of this the first byte is the number of bytes of data.

What are we going to produce instead of all these graphics and codes? These are listed in Appendix I and are mainly self-explanatory. However, an explanation of two of them is required. If you look at the Programmer's Reference Guide, page 74, under 'Other Special Characters', you will see five functions available. Three of these can be achieved more easily than described in the PRG. These are SWITCH TO LOWER CASE, DISABLE CASE-SWITCHING KEYS and ENABLE CASESWITCHING KEYS. They can be obtained by simply holding down the CTRL key and appropriate letter. In quotes they will print the appropriate symbol. Out of quotes the action will be carried out. The remaining two 'special characters' are implemented in the way the PRG describes. We have given them codes of [CRG>M| and ICRG>N|. These stand for CTRL revs graphic shift right and the appropriate letter.


| 180 | NEXT | INY |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 190 |  | LDA | (\$FB), $Y$ | ! | GET BYTE OF PROG LINE |
| 208 |  | BNE | COWT1 | ! | ZERO SIGNIFIES END OF LINE |
| 210 |  | LDA | \$FD | ! | GET NEXT LINE ADDRESS |
| 220 |  | STA | \$FB | ! | PUT IN CURRENT LINE REGISTERS |
| 230 |  | LDA | \$FE |  |  |
| 240 |  | STA | \$FC |  |  |
| 250 |  | LDX | \#\$00 | $!$ | RESET QUOTE COUNTER |
| 260 |  | BEQ | LINKS | ! | $\times$ SETS ZERO FLAG BRANCH ENFORCED |
| 270 | CONT1 | CMP | \#क FF | ! | IS IT pi |
| 280 |  | BNE | NOPI | $!$ | NO |
| 290 |  | STA | \$3E | ! | STORE VALUE |
| 308 |  | BEQ | CONT2 | ! | ENFORCED |
| 310 | NOPI | C.MP | \#\$22 | ! | IS BYTE A QUOTE |
| 320 |  | BNE | CHECK | $!$ | NO GO TO SEE IF IN QUOTES |
| 330 |  | INX |  | ! | IT'S A QUOTE SO INC: COUNTER |
| 340 |  | CPX | \#502 | ! | IS IT SECOND QUOTE |
| 350 |  | BNE | INQUOTES | ! | IN QUOTES CODER IN ACTION |
| 360 |  | LDX | \#\$00 | ! | RESET COUNTER |
| 370 |  | BEQ | NEXT | ! | ENFORCED |
| 380 | CHECK | CPX | \#\$81 |  |  |
| 390 |  | BNE | NEXT | ! | NOT IN QUOTES |
| 400 | INQUOTES | STA | \$3E | ! | STORE BYTE |
| 410 |  | CMP | \#\$ C 0 | ! | IS IT LESS THAN 192 |
| 420 |  | BCC | COMPARE | ! | YES |
| 430 |  | SBC | \#\$ 68 | ! | NO SUBTRAC:T 96 |
| 440 | COMPARE | CMP | \#\$60 | ! | IS IT $>0$ ( $=$ T0 96 |
| 450 |  | BCS | CONT2 |  |  |
| 460 |  | CMP | \#\$21 | ! | IS IT LESS THAN 21 |
| 470 |  | BCS | NEXT | ! | CHARS 21-95 <br> DON'T REQUIRE CODING |
| 480 | CONT 2 | STA | \$30 | ! | STGRE REUISED CHAR UALUE |
| 490 |  | STY | \$49 | ! | STARE LINE MARKER |
| 500 |  | STX | \$3C | ! | STORE QUOTE COUNTER |
| 510 |  | LDX | \# $\$^{81}$ |  |  |
| 520 | NEXT1 | INY |  |  |  |
| 530 |  | LDA | (\$FB), $Y$ | ! | GET NEXT CHAR |
| 540 |  | CMP | \$3E | ! | IS A REPEAT CHAR |
| 550 |  | BNE | NEXT 2 | ! | NO |


| 560 |  | JNX |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 570 |  | BNE | NEXT 1 | ! | ENFORCED |
| 580 | NEXT2 | STX | \$3E | ! | REG - NO OF REPEATS |
| 590 |  | CPX | \#\$ 82 |  |  |
| 600 |  | BCS | CONT3 | ! | MORE THAN ONE CHAR |
| 610 |  | DEX |  |  |  |
| 620 |  | LDA | \$3D | ! | $\begin{aligned} & \text { GET CHAR BACK } \\ & \text { AGAIN } \end{aligned}$ |
| 630 |  | CMP | \#\$20 | ! | IS IT A SPACE |
| 640 |  | BEQ | SPACE | ! | dON'T CODE SINGLE SPACE |
| 650 |  | LDA | \#\$00 |  |  |
| 660 |  | STA | \$40 | ! | RESET REG WITH ASCII FOR NO OF REPEATS |
| 670 |  | STA | \$3F |  |  |
| 680 |  | BEQ | CONT4 | ! | ENFORCED |
| 690 | SPACE | JMP | RELOAD | ! | RELOAD REGISTERS FOR GET NEXT BYTE |
| 700 | CONT3 | LDA | \# $\$ 00$ | ! | $X$ has low value NO HIGH Malue |
| 710 |  | JSR | \$847F | ! | NO OF REPEATS INTO ASCII FORM |
| 720 |  | LDX | \# $\$ 00$ |  |  |
| 730 |  | LDA | \$0200, X | ! | GET ASCII INTO REGS |
| 740 |  | STA | \$3F |  |  |
| 750 |  | INX |  |  |  |
| 760 |  | LDA | \$0200, X |  |  |
| 770 |  | STA | \$40 |  |  |
| 780 | CONT 4 | LDA | \$3D |  |  |
| 790 |  | CMP | \#\$61 |  |  |
| 800 |  | BCC | CONVERTB |  |  |
| 810 |  | CMP | \# ${ }^{\text {a }}$ 7B |  |  |
| 820 |  | BCS | CONVERTB |  |  |
| 830 | CONUERTA | SEC |  |  |  |
| 840 |  | SBC | \#\$20 | $!$ | REDUCE VALUE |
| 850 |  | STA | \$30 |  |  |
| 860 |  | LDX | \#\$07 | ! | total no of spaces REQUI RED |
| 870 |  | LDA | \$40 |  |  |
| 880 |  | BNE | CONT 5 | ! | MORE THAN 9 REPEATS |
| 890 |  | DEX |  |  |  |
| 908 |  | LDA | \$3F |  |  |
| 910 |  | BNE | CONT5 | ! | SOME REPEATS |
| 920 |  | DEX |  |  |  |
| 930 | CONT5 | CPX | \$3E | ! | FIND HOW MUCH ROOM |
| 940 |  | BEQ | NOMOVE | ! | RIGHT AMOLNT OF SPACE |


| $\begin{aligned} & 950 \\ & 960 \end{aligned}$ |  | $\begin{gathered} \text { BCS } \\ \text { JSR } \end{gathered}$ | OPENUP <br> CLOSE | ! NEEDS SPACE IN LINE |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | ! | GET RID OF |
|  |  |  |  |  | UNUNATED CHARS |
| 970 |  | JMP | NOMOVE | $!$ | CONT WITH PROG |
| 980 | OPENUP | JSR | OPEN | ! | NOT ENOUGH ROOM IN |
|  |  |  |  |  | LINE |
| 990 | NOMOVE | LDY | \$49 | ! | GET LINE POINTER |
| 1080 |  | LDA | \#\#5B | ! | ASCII FOR [ |
| 1010 |  | STA | (\$FB) , $Y$ | ! | PUT IN LINE |
| 1020 |  | LDA | \$40 |  |  |
| 1030 |  | BEQ | CONT6 | ! | NO TENS DIGIT |
| 1040 |  | INY |  |  |  |
| 1050 |  | STA | ( $\ddagger \mathrm{FB}$ ) , $Y$ |  |  |
| 1060 | CONT 6 | LDA | \$3F |  |  |
| 1070 |  | BEQ | CONT7 | ! | NO REPEATS |
| 1880 |  | INY |  |  |  |
| 1090 |  | STA | (\$FB) , $Y$ |  |  |
| 1100 | CONT7 | LDA | \# ${ }^{\text {47 }}$ | ! | ASCII FOR G |
| 1110 |  | INY |  |  |  |
| 1120 |  | STA | (まFB) , Y |  |  |
| 1130 |  | LDA | \#\$3E | ! | ASCII FOR > |
| 1140 |  | INY |  |  |  |
| 1150 |  | STA | (\$FB) , $Y$ |  |  |
| 1160 |  | LDA | \$3D | ! | CHAR |
| 1170 |  | INY |  |  |  |
| 1180 |  | STA | (\$FB) , $Y$ |  |  |
| 1190 |  | LDA | \#\$5D | $!$ | ASCII FOR ] |
| 1200 |  | INY |  |  |  |
| 1210 |  | STA | ( $\ddagger \mathrm{FB}$ ), Y |  |  |
| 1220 |  | LDX | \$3C | $!$ | RESET QUOTE COLINTER |
| 1230 |  | JMP | NEXT | ! | NEXT BYTE |
| 1240 | CONUERTB | STA | \$3D |  |  |
| 1250 |  | LDA | \#\$50 | ! | LSB OF DATA ADDRESS TABLE |
| 1260 |  | STA | \$62 |  |  |
| 1270 |  | LDA | \#\$A3 | ! | MSB OF D.A.T. |
| 1280 |  | STA | \$63 |  |  |
| 1290 |  | LDX | \# ${ }^{51}$ | ! | COUNTER MAX NO OF CHARS IN DATA TABLE |
| 1300 |  | LDY | \#\$00 |  |  |
| 1310 | NEXT3 | JSR | \$81FB | $!$ | SWITCH OF BASIC |
| 1320 |  | LDA | (\$62), Y | ! | GET ASCII CHAR NO |
| 1330 |  | PHA |  | , | TEMP STORE |
| 1340 |  | JSR | \$8202 | , | SUIITCH IN BASIC |
| 1350 |  | PLA |  |  | RETRIEVE |


| 1360 | CMP \$3D | ! IS IT THE SAME |
| :---: | :---: | :---: |
| 1370 | BEQ FOUND | $!$ YES |
| 1380 | INY | SKIP UNUANTED ADDRESS |
| 1390 | INY |  |
| 1480 | INY |  |
| 1410 | DEX | DECREASE COUNTER |
| 1420 | BPL NEXT3 | GET NEXT ASCII NO UNTIL $X<0$ |
| 1430 | LDA $\$ 30$ |  |
| 1440 | CMF \#\#1B | SEE IF IT USED WITH CTRL KEY |
| 1450 | BCC CTRL | ! YES |
| 1460 | LDX \#\#00 | ! NO |
| 1470 | JMP \$A437 | ERROR OUT OF DATA |
| 1480 CTRL | ADC: \#\$40 | ADD $\$ 40$ TO GIVE ASCII LETTER |
| 1490 | STA \$A448 | ! STGRE IN data table |
| 1500 | LDA \#\$43 | LSB OF DATA FOR <br> [CTRL?] |
| 1510 | STA $\$ 62$ |  |
| 1520 | LDA \#\#A4 | ! MSB OF DATA FOR [CRTL?] |
| 1530 | STA $\$ 63$ |  |
| 1540 | BNE CONT8 |  |
| 1550 FOUND | INY |  |
| 1560 | JSR \$81FB | SWITCH OUT BASIC. |
| 1570 | LDA (\$62), Y | GET LSB OF DATA POSITION |
| 1580 | PHA | ! TEMP STORE |
| 1598 | INY |  |
| 1600 | LDA (\$62), Y | GET MSB OF DATA POSITION |
| 1610 | STA \$63 |  |
| 1620 | JSR $\$ 8202$ | SWITCH IN BASIC |
| 1630 | PLA |  |
| 1640 | STA $\$ 62$ |  |
| 1650 CONT8 | LDY \#\$00 |  |
| 1660 | JSR \$81FB |  |
| 1670 | LDA (\$62), Y | ! GET NO OF Data chars |
| 1680 | STA $\ddagger$ C1 |  |
| 1690 | JSR $\$ 8202$ |  |
| 1700 | LDA $\$$ C1 |  |
| 1710 | CLC |  |
| 1720 | ADC \# 04 | ! FOR BRACKETS AND REPEATS |
| 1730 | TAX |  |


| 1740 |  | LDA | \$48 |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1750 |  | BNE | CONT9 |  |  |
| 1768 |  | DEX |  | ! | NO TENS IN REPEATS |
| 1770 |  | LDA | \$3F |  |  |
| 1780 |  | BNE | CONT9 |  |  |
| 1790 |  | DEX |  | ! | NO REPEATS AT ALL |
| 1888 | CONT9 | CPX | \$3E | ! | DO WE REQUIRE A |
| 1810 |  | BEQ | NOMOVE | ! | MEMORY MOVE NO |
| 1820 |  | BCS | OPEN1 | ! | MORE SPACE |
| 1830 |  | JSR | CLOSE | ! | LESS SPACE |
| 1840 |  | JMP | NOMOVE1 |  |  |
| 1850 | OPEN1 | JSR | OPEN |  |  |
| 1860 | NCHOVE1 | LDY | \$49 | $!$ | GET LINE MARKER |
| 1870 |  | LDA | \#\$5B | ! | ASCI] FOR [ |
| 1888 |  | STA | (\$FB), $Y$ |  |  |
| 1890 |  | LDA | \$3F |  |  |
| 1980 |  | BEQ | CONTA | $!$ | NO TENS DIGIT IN REPEATS |
| 1910 |  | INY |  |  |  |
| 1920 |  | STA | (FFB), Y | ! | STORE IN PROG |
| 1930 | CONTA | LDA | \$40 |  |  |
| 1940 |  | BEQ | CONTB | $!$ | NO REPEATS AT ALL |
| 1950 |  | INY |  |  |  |
| 1960 |  | STA | ( $\ddagger F B$ ) , Y | ! | Store in prog |
| 1970 | CONTB | STY | \$49 | ! | STORE LINE MARKER |
| 1980 |  | LDY | \#\$80 |  |  |
| 1990 |  | JSR | \$81FB | ! | SWITCH OUT BASIC |
| 2080 | data | INY |  |  |  |
| 2010 |  | LDA | (\$62), Y | ! | GET DATA FROM TABLE |
| 2020 |  | STY | \$C2 | ! | Store data marker |
| 2030 |  | LDY | \$49 | ! | GET LINE MARKER |
| 2040 |  | INY |  |  |  |
| 2850 |  | STA | $(\$ F B), Y$ | ! | store data in prog LINE |
| 2860 |  | STY | \$49 | ! | Store line marker |
| 2070 |  | LDY | \$C2 | 1 | GET DATA MARKER |
| 2080 |  | CPY | \$C:1 | 1 | have we got all data |
| 2090 |  | BNE | DATA | ! | NO |
| 2100 |  | JSR | \$8202 | ! | SWITCH IN BASIC |
| 2110 |  | LDY | \$49 |  |  |
| 2120 |  | LDA | \#\$50 | $!$ | ASCII FOR J |
| 2130 |  | JNY |  |  |  |
| 2140 |  | STA | (\$FB), Y |  |  |


| 2150 | LDX $\$ 3 C$ | ! GET QUOTE COUNTER |
| :--- | :--- | :--- |
| 2160 | JMF NEXT | ! NEXT BYTE TO PROCESS |
| 2170 RELOAD | LDY $\$ 49$ | GET LINE MARKER |
| 2180 | LDX $\$ 3 C$ | GET QUOTE COUNTER |
| 2190 | JMP NEXT | INEXT BYTE TO PROCESS |


| 8BD7 CHECK | 8889 CLOSE |
| :--- | :--- |
| 8BE3 COMPARE | 8BAF CONT |
| 8BC2 CONT1 | $8 B E B$ CONT2 |
| 8C15 CONT3 | $8 C 27$ CONT4 |
| 8C42 CONT5 | $8 C 5 E$ CONT6 |
| 8C65 CONT7 | $8 C C C$ CONT8 |
| 8CDF CONT9 | 8084 CONTA |
| 8D0B CONTB | 8C7E CONUERTB |
| 8CAB CTRL | 8D12 DATA |
| 8CBA FOUND | 8BDB INQUOTES |
| 8B9D LINKS | 8BB1 NEXT |
| 8BF3 NEXT1 | 8BFD NEXT2 |
| 8C8C NEXT3 | 8C51 NOMONE |
| 8CF7 NOMOUE1 | 8BCA NOPI |
| 8931 OPEN | 8CF4 OPEN1 |
| 8C4E OPENUF | 8D33 RELOAD |
| 8C12 SFACE |  |

LINES 10-150: These set up the routine and if necessary return control back to you through BASIC. There are no parameters included in the command to pick up as it codes the whole program. The address of the first line is taken from the start of BASIC program variables at $\$ 2 B$ and $\$ 2 C$. The $x$ register is initialized to zero and is used as a quote counter. We get the link address to the next line and if it is the end of the program we remove the return address from the stack, placed there on entering, and go back to BASIC with the program in memory coded for listing or saving.
LINES 160-390: We are going to start to look for our trigger codes quotes or pi. We skip the line number and start to scan. If the end of the line is encountered we transfer the links to the line register and start the next line.

We check for PI (ASCII value is \$FF). On finding it, we store it in a register for later use and branch further into the program. The check for the quote takes two forms. When one is found, we increase and check the $x$ register. A two here will indicate that it is the second quote and therefore going out of the area we are interested in. It also means we go back to look for another quote.

If $x$ is one, then the first quote has been found and we go forward to
check for codes. It will fail there the result is that we return to get the next byte.

On encountering a byte other than a quote or pi, we check to see if the $x$ register is one, indicating that we are in quotes and it will require processing.

LINES 400-470: We first store the byte. This is done as we are going to manipulate this data and possibly alter it. The original value is needed later when checking for repeat characters.

Values over $\$ C 0(192)$ are reduced by $\$ 60(96)$ and we are in a position to weed out characters that do not require any action. These will be values of $\$ 21$ to $\$ 5 F$ inclusive. This is the position that the first quote will end up in. These characters cause the flow to go back and get the next byte.

LINES 480-690: Our character value is stored again, as it may be different, in another register. We also store the line pointer (the $Y$ register) and the $x$ quote counter. The latter is stored because if pi is being changed, the $x$ register could be zero; at other times it will always be one.

The next procedure is to see if there is more than one character of the same type consecutively in the program. The x register will be used as a counter and as it is one already, it is already initialized. The following bytes are gathered in and checked against the original value. The $x$ register is increased until a byte of a different value is found.

The routine now splits up. Where there are two or more repeat characters, we jump ahead to CONT3 to put that number into ASCII.

Continuing along, the $x$ value will be one but we will not print out the number one as it implied. Registers $\$ 3 F$ and $\$ 40$ are set to zero, which as we shall see shortly will hold the ASCII value for repeats.

The action taken is to check to see if the character we are coding is a space or not. We do not want to code single spaces as it would clutter the listing unnecessarily. On finding a space the flow jumps further ahead to reload the registers and go back to get the next byte. For characters not spaces, and all single characters, the routine branches forward to skip the next section.

LINES 700-770: On finding more than one of the same character we want to convert the number into ASCII format. We already have the number in the $x$ register. To use our own conversion routine at $\$ 847 \mathrm{~F}$ we need to set the accumulator to zero, as the high byte value. The result will be in the input buffer with a zero after the last digit. As a line of BASIC program when typed into the 64 cannot be more than 80 characters, it therefore means that the number of repeats cannot be any greater. This means that the number of ASCII digits will be two at a maximum.

We therefore pick up the first two digits from the buffer and store
them in $\$ 3 F$ and $\$ 40$. If there was a single digit, that is 2 to 9 repeats, $\$ 40$ will be zero.

LINES 780-820: We said at the beginning that some characters would require the use of the data tables whilst some can be coded by calculation. These few lines divide up the flow into these two areas.

We load back the value achieved in earlier calculations (lines $410-440$ ) to the accumulator. Values of $\$ 60$ and under, or $\$ 7 B$ and over, will branch off to Conversion B, which uses data tables.

## CONVERSION A

LINES 830-1230: The first task to undertake is to subtract $\$ 20$ (32) from our value and store it. This is now the same value as the ASCII code of the letter of the key it shares. They will all be achieved using the shift with the key rather than the logo key.

The maximum number of characters we could insert is seven, two for the brackets, two for the number of repeats and three for the code. This number is placed into $x$. We check the 'repeat' digits storage for the number of numerals. A zero will indicate that there is no digit in that column. The $x$ register will be decreased accordingly. Location \$3E has the number of graphic characters to be coded and this is compared with $x$. From this we either open-up the program, close-up the program or leave it unchanged. The memory move routines are described in Chapter 6.

Now we are ready to insert the code in the order of:
i) The [bracket.
ii) The number of repeats if applicable.
iii) The letter G.
iv) The symbol $>$.
v) The letter of the key, held in location \$3D.
vi) The ] bracket.

Once completed, we load the quote counter back into $x$ and jump back to get the next character to code.

## CONVERSION B

LINES 1240-1640: This is where we have to use data tables to find the relevant code. This part is entered with the character value in the accumulator and is put into \$3D for later use. The first table we look up is the data address table. The start address of the table is placed in locations $\$ 62$ and $\$ 63$. The $x$ register has the total number of characters catered for and the Y register is used as a general pointer.

As the table is in the RAM under the BASIC ROM we have to disable that ROM, get the byte we want and then switch back the rom. The byte is placed temporarily on the stack during the enabling of the rom. The byte we have collected is compared for equality with our character
value. Succeeding forces a branch forward. Failure means we continue the search. The first thing is to increase the $y$ register three times. This will skip the address of the rejected character in the data table. The $y$ value will be in line with the next character value. If the $x$ register has been decreased to a value below zero, that is, \$FF, then all the data address table has been checked and a match not found. There is one further chance. It could be a character which uses the CTrl key along with a letter key. These will have values no greater than \$1A (26). This is checked and if it does not fall in, then an 'OUT OF DATA' error is generated and coder is exited. We think that this should never come about as we believe we have catered for all eventualities.

Supposing a CTRL value is the one found, then we add $\$ 40$ (64). This simply gives the value of the letter on the key. This is stored immediately in the data table. The start address of the start of CTRL data is placed into $\$ 62$ and $\$ 63$.

Now back to the other characters. A match has been found in the data address table and we have arrived at line 1550. The two bytes next in line in the address table are the data address in the second table. These are placed into registers $\$ 62$ and $\$ 63$.
LINES 1650-2160: We have now finished with the data address table and concentrate on the data table itself. This time we only require one byte, the first byte, which will give us the number of bytes of code. To this value we add four, the brackets and the 'repeat' digits, transfer it to $x$ and decrease it if one or both repeat digits are redundant. This final value is compared with the number of characters to be replaced to determine whether more or less space is required. This and the moves, if needed, are achieved in lines 1800 to 1850.

The insertion of data is the only thing left to do. We reload the line marker and start. The left square bracket and, if required, the repeat digits are stored first. The data insertion is slightly complicated. The line marker is stored and the $y$ register is re-initialized. The basic rom is switched off and a byte is taken from the table. The $y$ register is stored in \$C2 and the line marker is restored and incremented. The byte is now inserted in the program. Now the line marker is stored and the data marker placed back in $Y$. This is compared with the number of bytes of data in the code (\$C2). If it has not reached this number, we branch back to get further bytes to insert. Once all the data has been collected and stored, the BASIC ROM is switched back in.

Finally, the right hand square bracket is inserted and we jump to get the next character to be coded, after restoring the quote counter.

LINES 2170-2190: This simply restores the line marker and quote counter, after which the routine goes back to get the next character. Single space characters, which are not coded, are sent here.

## The data table - a program

After much thought, we have decided to supply the data tables for CODER in the form of a bASIC loader. This is mainly due to the fact that it is stored under the BASIC ROM which makes it hard to check and correct using a monitor. With the loader program we can put in a checksum which helps to see if you typed in the correct values.

A further item that the loader program does is to clear the area used by the key command (see Chapter 4) for its data. So type in the program, check it and save it.

```
10 L=41472:T=0
20 READD:IFD=-1 THEN40
30 T=T+D:POKEL,D:L=L+1:GOT020
40 IFT<>51131THENPRINT"[REV] DATA INCORR
ECT":END
50 FORL=41216T041471:POKEL,0:NEXT
60 PRINT"[REU] DATA LOADED":END
70 DATA3,87,72,84,2,67,68
80 DATA3,82,69,86,3,72,79
90 DATA77,3,82,69,68,2,67
100 DATA82,3,71,82,78,3,66
110 DATA76,85,3,83,80,67,3
120 DATA71,62,42,3,71,62,43
130 DATA3,71,60,45,3,71,62
140 DATA45,1,126,3,71,60,42
150 DATA3,79,82,71,2,70,49
160 DATA2,70,51,2,70,53,2
170 DATA70,55,2,70,50,2,70
180 DATA52,2,70,54,2,70,56
190 DATA3,66,76,75,2,67,85
200 DATA3,79,70,70,3,67,76
210 DATAB3,3,73,78,33,3,66
220 DATA82,78,5,76,32,82,69
230 DATA63,3,71,82,49,3,71
240 DATA82,50,5,76,32,71,82
250 DATA78,5,76,32,66,76,85
260 DATA3,71,82,51,3,80,85
270 DATA82,2,67,76,3,89,69
280 DATA76,3,67,89,78,5,71
290 DATA62,83,80,67,3,71,60
300 DATA75,3,71,60,73,3,71
310 DATA60,84,3,71,60,64,3
320 DATA71,60,71,3,71,60,43
330 DATA3,71,60,77,3,71,60
340 DATA92,3,71,62,92,3,71
350 DATA60,78,3,71,60,81,3
```

360
DATA71,60,68,3,71,60,90
370
380
390
400
DATA3,71,60,83,3,71,60
DATA80,3,71,60,65,3,71

410 DATA3,71,60,74,3,71,60
420 DATA $76,3,71,60,89,3,71$
430 DATA $60,85,3,71,60,79,3$
440 DATA $71,62,64,3,71,60,70$
450 DATA3,71,60,67,3,71,60
460 DATA88,3,71,60,86,3,71
470 DATA $60,66,5,67,84,82,76$
480 DATA $65,5,67,84,82,76,66$
490 DATA5,67,84,82,76,72,5
500 DATA67,84,82,76,73,5,67
510 DATA84,82,76,78,5,67,82
520 DATA71,62,78,5,67,82,71
530 DATA62,77,3,68,69,76,2
540 DATA80,73,255,0,0,0,0
550 DATA5, $0,162,17,4,162,18$
560 DATA7, 162,19,11,162,28,15
570 DATA $162,29,19,162,30,22,162$
580 DATA31,26,162,32,30,162,96
590 DATA $34,162,123,38,162,124,42$
60 DATA $162,125,46,162,126,50,162$
610 DATA127,52,162,129,56,162,133
620 DATA $60,162,134,63,162,135,66$
630 DATA162,136,69,162,137,72,162
640 DATA $138,75,162,139,78,162,140$
650 DATA81,162,144,84,162,145,88
660 DATA162,146,91,162,147,95,162
670 DATA148,99,162,149,103,162,150
680 DATA107,162,151,113,162,152,117
690 DATA162,153,121,162,154,127,162
700 DATA155,133,162,156,137,162,157
710 DATA141,162,158,144,162,159,148
720 DATA162,160,152,162,161,158,162
730 DATA162,162,162,163,166,162,164
740 DATA170,162,165,174,162,166,178
750 DATA162,167,182,162,168,186,162
760 DATA $169,190,162,170,194,162,171$
770 DATA193,162,172,202,162,173,206
780 DATA162,174,210,162,175,214,162
790 DATA176,218,162,177,222,162,178
800 DATA226,162,179,230,162,180,234
810 DATA1 $62,181,238,162,182,242,162$

```
820 DATA183,246,162,184,250,162,185
830 DATA \(254,162,186,2,163,187,6\)
840 DATA163,188,10,163,189,14,163
850 DATA190,18,163,191,22,163,1
860 DATA26,163,2,32,163,8,38
870 DATA163,9,44,163,14,50,163
880 DATA \(142,56,163,141,62,163,20\)
890 DATA \(68,163,255,72,163,5,67\)
900 DATA \(84,82,76,67,-1\)
```


## SAVING THE DATA AREA

The following listing will save the area we have used for both KEY and CODER routines. The saving of data through machine code is described in the Programmer's Reference Guide. The only extra coding is to switch the BASIC ROM out, so that we will save our data and not the BASIC interpreter. You could use this after setting up the function keys (see Chapter 4) so on reloading, the data is there and ready.

| 10 | LDX \#\$0 | DEVICE NO (TAPE=1) |
| :---: | :---: | :---: |
| 20 | LDA \#\$01 | LOGICAL FILE NO |
| 30 | LDY \#\#FF | NO SEC ADDRESS |
| 40 | $J S R$ FFFBA | SETLFS |
| 50 | LDA \#\$0C | Chars in filename |
| 60 | LDX \# <NAME | LOW AODRESS OF NAME |
| 70 | LDY \# ${ }^{\text {SNAME }}$ | HIGH ADDRESS |
| 30 | JSR FFFBD | SETNAM |
| 90 | LDA $\$ 01$ | SWITCH OFF BASIC |
| 100 | AND \#\#FE |  |
| 110 | STA ${ }^{\text {S }} 1$ |  |
| 120 | LDA \#\$00 | STURE START ADDRESS |
| 130 | STA $\ddagger$ FB |  |
| 148 | STA \#\#A1 |  |
| 150 | STA \#FC |  |
| 160 | LDX \#\$49 | Low end of save |
| 170 | LDY \#\$A4 | high end of save |
| 130 | LDA \#\#FB | LOCATION OF START ADD |
| 190 | JSR \#FFD8 | SAVE |
| 200 | LDA $\$ 01$ | SWITCH IN BASIC |
| 210 | ORA \#\$ 11 |  |
| 220 | STA $\$ 81$ |  |
| 230 | RTS |  |
| 240 NAME | TXT "UTILIT | DATA" |

## RELOCATING THE DATA TABLES

If you relocate coder may also want to relocate the data. Here is one suggested way. Using the basic loader program for the data, change the value of $L$ in line 10 to the new data start address. The data normally starts at \$A200 (41472) but the data address table starts at \$A350 (41808). From this calculate the data address table new address and put its value in lines 10 and 30 of the routine below. The end of the data address table is normally $\$ 4442$ (42050), so work out its new end, subtract one, and this is put in lines 210 and 240 . The difference between the old address and the new address should be put in lines 80 and 120 . The routine below is for a new table at a higher address; for one lower, change the addition to subtraction and set the carry instead of clearing it.

LDA \#\$50 ! START OF DATA ADDRESS TABLE
STA $\$ 14$
LDA \#\#A3
STA $\$ 15$
LDY \#\$01 ! POINTER
LDA (\$14), Y! LOW ADD IN TABLE
CLC
ADC \#\$60 ! ADD LON DIFFERENCE
STA (生14) , Y
INY
LDA ( $\$ 14$ ), Y! HIGH ADD IN TABLE
ADC \#\$25 ! ADD HIGH DIFFERENCE
STA ( 1 14), Y
CLC
LDA $\$ 14$ ! UPDATE TABLE ADDRESS
ADC \#\$03
STA $\$ 14$
LDA $\$ 15$
ADC \#\$00
STA $\$ 15$
CMP \#\$A4 ! END OF DATA ADDRESS TABLE
BNE NEXT ! NO
LDA $\$ 14$
CMP \#\$43
BNE NEXT ! NO

## Old

## COMMAND SYNTAX <br> OLD

There are no parameters with this command.
There are four ways to 'lose' a BASIC program. The first way is by switching off and then there is absolutely no way of recovering it. Another two ways are by doing a system cold start or a reset. This is as if you have just switched on but retaining data held by the rams. This can be achieved by typing SYS64738 or by a reset button, if you have fitted one. The final way to lose a program is by issuing the basic command New.
To lose a program the operating system of the 64 sets the first two bytes of the BASIC program area to zero. This would normally be $\$ 0801$ and $\$ 0802$ (2049 and 2050) and would be the link address in the first line of a BASIC program. This means that as far as BASIC is concerned no program is present as it would encounter the zeros straight away.

Now as long as no further lines of BASIC are typed in, we can reverse the process, but will lose all the variables. The way it is done is made clear in the description of the coding.

ASSEMBLY LISTING

| 10 | LDA | \# ${ }^{\text {FFF }}$ |  |  |
| :---: | :---: | :---: | :---: | :---: |
| 20 | LDY | \#\$01 |  |  |
| 30 | STA | ( $\$ 2 \mathrm{~B}$ ), Y | ! | PUT ANY LINK IN 1ST LINE |
| 40 | JSR | \$A533 | ! | RECHAIN LINES |
| 50 | LDA | \$22 |  |  |
| 60 | CLC |  |  |  |
| 70 | ADC | \#\$02 |  |  |
| 80 | STA | \$20 | ! | SET END OF PROG ADDRESS -LOW |
| 90 | LDA | \$23 |  |  |
| 100 | ADC | \# $\boldsymbol{6}^{\text {g }}$ | ! | In CASE CARRY Was SET IN 70 |
| 110 | STA | \$2E | ! | SET END OF PROG ADDRESS -HIGH |
| 120 | JMF | \$A660 | ! | PERFORM CLR |

LINES 10-40: If we change the first two bytes from zero, BASIC will no longer think it is at the end of the program. We put SFF in those, and get the address from the start of BASIC variables in $\$ 2 \mathrm{~B}$ and $\$ 2 \mathrm{C}$. Now a call to the rom routine rechain lines will achieve two things. First, it will
correctly set the link address in the first line, and secondly, we will be able to set the end of program variables.

LINES 50-120: Locations $\$ 22$ and $\$ 23$ are set to the beginning of the two zero bytes, which mark the end of the program, when the rechain routine is finished. By adding two to those, we have the end of program and can set the respective registers, \$2D and \$2E.

The final thing is a jump to the CLR routine. This will set all the variable addresses to coincide with the recovered program. The BASIC program is now restored to its original state.

## Dump

## COMMAND SYNTAX <br> DUMP

There are no parameters with this command. It will also only operate in direct mode. If used within a program it will just skip out of the command. Hitting stop will break out of DUMP and allows direct editing; typing CONT will resume at the break point in the BASIC program. Holding down any other key, apart from shift, will halt the routine until it is released.

The action of DUMP is identical to the BASIC subroutine in Chapter 5 except that as the routine is in machine code it does not add to the simple BASIC variables. The logic closely follows the BASIC routine. Output may again be directed to a printer by an OPEN and CMD sequence. The major departure is in the use of one or two ROM routines to carry out the mathematical conversions and convert the number to an ASCII string to be printed.

ASSEMBLY LISTING

```
*=$8E52
INOTE REAL ASC / ASC OR 0
STRING ASC:/ASC+128 OR 128
! INTEGER ASC+128 / ASC+128 OR 128
    FUNCTION ASC+128 / ASC OR a
!
!
!TEST FOR DIRECT MODE
!
90 LDA $9D !MSGFLG
100 CMP #$80 !DIRECT ???
110 BEQ DIRECT
120 RTS !PROGRAM MODE SO ABANDON
140 !NOTE CURRENT UARIABLE IN FILE NAME POINTER
```

130


```
550
560 !STRING VARIABLE
570!
5 8 0 ~ S T R I N G ~ A N D ~ \# \$ 7 F ~ ! M A K E ~ A S C I I ~
5 9 0 ~ J S R ~ F F F D 2
600 LDA #$24 !'$'
610 JSR $FFD2
620 JSR UPDATE
630 LDA #$22
640 JSR #FFD2
6 5 0
660 LDA ($BB),Y
670 TAX !LENGTH OF STRING IN }
6 8 0
6 9 0
7 0 0
7 1 0
7 2 0
730
7 4 0
750
760 CHAR
70
7 8 0
790
80
8 1 0 ~ Q U O T E ~
820
830
840
850 HALFSTART
860 !
870 !INTEGER AND FLNCTIONS
880
890 INTFN AND ##7F !CONUERT TO ASCII
900
910
JSR $FFD2
INY
920 LDA ($BB),Y
930
940 BCS INT !INTEGER IF SECOND
                                    CHAR>128
945!
950 !FLNCTION DEFINITION
960!
970 JSF &FFD2
980
JSR UPDATE
```



1420
1430 !PRINT $=$ AND SET POINTER IN $\$ B B / B C$ TO START OF UARIABLE
1440
1450 UPDATE LDA \#\$3D $!^{\prime}=\prime$
1460
1478
JSR \&FFD2

1480
1496
CLC
LDA $\$ B B$
ADC \# $\# 02$
1500
STA \$BB
1510
1520
ECC RETURN
INC \$BC
1530 RETURN RTS

8EBC CHAR
8EOE CONT
8E59 DIRECT
BECE HALF
8EEE INT
8EDG INTFN
8F11 NEXT
8F 2A NOT
8EC5 QUOTE
8F43 RETURN
8EG1 START
8F33 UPDATE
$8 E 97$ STRING
$8 F 21$ WAIT

The routine has been written to be easily relocatable. The only change necessary is to alter all ISR UPDATES to JSR (start address $+\$ \mathrm{EF}_{1}$ ).

The ROM routines used are as follows:

## CHROUT (\$FFD2)

A full description of this function is given in the Programmer's Reference Guide, 'The kernal B-5'. It outputs the contents of A as an ASCII character to the screen.

## STOP (\$FFE1)

See 'The kernal B-33'. Test for the stop key. udtim must be called before using this routine.

UDTIM (\$FFEA)
See the Programmer's Reference Guide, 'The kernal B-36'. This updates the system clock.

## MEMORY TO FAC\#1 (\$BBA2)

This routine takes a five byte real number and moves it to the floating point accumulator \#1. En route the sign bit of the mantissa is stripped off and the sign register FACSGN (\$66) set, the exponent put at FACEXP (\$61) and the mantissa of FACHO (\$62-65). On entry a must hold the low and $Y$ the high byte of the address of the bytes to be moved.

FAC\#1 TO STRING (\$BDDD)
Converts FAC\#1 to an ASCII string stored at the bottom of the stack (\$0100). On exit a holds \#\$00 and y holds \#\$01.

PRINT STRING FROM MEMORY (\$AB1E)
This routine prints successive characters starting at the memory location whose address is held in A (low) and $Y$ (high). The routine continues until a zero terminator is found (as will be the case at the bottom stack in this application). Note $A$ and $\gamma$ already hold the start address on exiting the previous routine and need not be changed.

EVALUATE TWO BYTE SIGNED INTEGER (\$BC44)
Evaluates a two byte signed integer held in FAC\#1 and deposits the result in floating point form back in FAC\#1. Before calling $x$ must be set to $\# \$ 90$, FACHO must hold the high and $\mathrm{FACHO}+1$ the low byte of the integer (remember integers are held in high/low format unlike addresses). Once in this form the same routines as for real may be used to convert to ASCII and print.

LINES 1450-1530: JSR UPDATE
This will be used by all types of variables. It will be used directly after the variable name has been printed. All this does is to print out the equals sign and increase the address registers by two, so that they will point directly to the next byte to be collected - the first of the actual variable.

LINES 90-120: These check for direct mode. If program mode is discovered, then the routine is exited.

LINES 160-190: The locations we are going to use to step through the variable area are initialized with the start of variable address, which also happens to be the end of BASIC program address. We are now ready to start.

LINES 220-280: Locations $\$ 2 F$ and $\$ 30$ are the address of the end of the variable block that we are going to DUMP. By checking the values in those with our registers, we can find out if we have completed all.
LINES 320-350: The routine is divided up here and will be further divided later. When the first byte is picked up, we check if it is an integer or a function by seeing if the value is $\$ 80$ (128) or over. These are dealt with further into the routine.

LINES 380-420: The first byte we have already is printed - the first letter of the variable name. The next byte is collected and this will distinguish between real and string variables.

LINES 460-540: REAL VARIABLES: Again we print what we have in A, making the whole variable name output. A call is now made to JSR UPDATE. With the address of the present position placed in the $A$ and $Y$
registers, we call three rom routines, described at the start, to print out the variable to the screen or output device. The accumulator is loaded with $\$ F F$ just so the branch following will succeed.
LINES 580-840: STRING VARIABLES: Before we print out the accumulator, we remove the negative bit, bit 7 , so that it is the pure ASCII code of the variable letter. The dollar sign is printed and UPDAtE is called. As there are no separaters between stored strings, we cannot use a similar approach to the one used in real variables. The first thing printed is the start quotes. The length of the string is the byte after the name and this is placed in the $x$ register. Now we gather in the address of where the string is stored. From this we can print out the characters directly, decreasing $x$ each time, until the counter is zero. Finally, the closing quotes are output. One of the following two branches must succeed so we can continue.
LINES 890-940: Integers and functions start here. In these lines we distinguish between them and act accordingly. By stripping off the negative bit we can print out the first character of the name, and do so. The second byte is loaded and this will tell us what type to deal with. A value of $\$ 80$ (128) or over will signify integer.

LINES 970-1030: FUNCTION VARIABLE: We cannot print any value for the function, so after the second name character is printed, UPDATE is called, and then we just print the letters ' $F$ ' and ' $N$ ' and branch off.
LINES 1070-1210: INTEGER VARIABLES: After printing the second character of the name the integer sign of ' $\%$ ' is output. Once more UPDATE is visited. The next two bytes in the variable area are the integer value and these are transferred to FAC \#1. With the x register set to normalize the result, we call a ROM routine at $\$ B C 44$. This will convert the integer value to a real number. We then convert to ASCII and print the result with the routine described at the beginning.

LINES 1250-1410: After dealing with any of the four types of variables, the flow is directed to this part of the routine. The return character is printed so that the next variable is printed on a new line. Each variable takes seven bytes of memory and as we added two bytes to our address registers in UPDATE, we only need to add five more to get to the start of the next variable in line.

The stop key is now tested for and if the negative flag is set then Dump is ended, as STOP was pressed. By examining \$CB we can see if any key is being pressed. As long as a key is held down we loop around here and then check for stop.

To continue with DUMP we set the carry flag and we use a Branch with Carry Set to line 850 where the same happens, going further back to proceed dumping variables.

## IMPROVEMENTS?

Obvious improvements are as for the BASIC subroutine. If you had trouble extending the BASIC subroutine to handle arrays, you haven't tried anything yet! Most dump routines (wisely) do not handle subscripted variables (probably because it is considered too difficult).

## Trace and Troff

## COMMAND SYNTAX TRACE and TROFF

## Speed Control

' 0 ' reset single-step
' 2 ' a relative delay of $2^{\wedge} 1$
'4' a relative delay of $2^{\wedge} 3$
' 6 ' a relative delay of $2^{\wedge} 5$
' 8 ' a relative delay of $2^{\prime} 7$
' 1 ' a relative delay of $2^{\wedge} \emptyset$
' 3 ' a relative delay of $2^{\wedge} 2$
'5' a relative delay of $2 \wedge 4$
'7' a relative delay of $2^{\wedge} 6$
' 9 ' a relative delay of $2^{\wedge} 8$

The space bar operates the single-stepping.
The delay is in addition to the normal time taken by basic to move to a new line and execute the common trace code. The delay may be changed at any time by hitting the appropriate key. It is, however, not possible to break into program execution in single step. If you wish to do this, hit a number other than $\emptyset$ first.

TRACE is a diagnostic aid which provides useful information on the path taken through a BASIC program. In this particular version the previous and the current line numbers are displayed in reverse video at the top right of the screen.

We considered it far more important to allow the user to be able to vary the speed of the trace and have single-stepping capability. When called for the first time the default will be to single stepping and thereafter at each run it will continue at the last set speed. After being disabled with troff it will, on being enabled, revert back to single stepping. In single step mode program execution halts until the space bar or a speed change key is pressed. The keys for speed change are given above.

ASSEMBLY LISTING

```
    9 *=803A
100 !TRACE ENABLE
110!
120!
130 ENABLE LDA $90 !MSGFLG CHECK FOR DIRECT
```

310 !TROFF = TRACE DISABLE

546
550
560 580

600
610
620
630
640

670

700
710 720
730
740
750
760
770
780
790
800
810
820 SPACE LDA \＃\＄20 ！CLEAR PREUIOUS NUMBERS
830
840
850
860
879
880
890
900
910
920
930
940
950
960

570 ！RESTORE ENTRY UALUES BEFORE CCNTINUING
590 BASIC LDA AREG ！REVERSE ENTRY PROCESS
LDA 19 D
BEQ RUNNING
! ！

LDY YREG ！TO ALLOW PROG to continue
LDX XREG ！UNCORRUPTED
PLF ！DON＇T FORGET FLAGS！！！
JMP（IGONE）！CONTINUE AT token dispatch

```
650 ！PROGRAM RUNNING SO CHECK IF TRACE ENABLED
660 ！FROM TRACE FLAG＝\＄FF？？？
680 RUNNING LDA TRFLAG
RUNNING LDA TRFLAG
BEQ BASIC !TRACE OFF sO RESTORE
                                    AND CONT
```

!
! trace is on so update display
!
SEC $\quad$ !READ CURSOR POSITION
JSR fFFF0 !AND SAVE BY CALLING PLOT
STX ROW ! WITH CARRY SET
STY COL
CLC !SET CURSOR POSITION
LDX \#\$00 !TO RON a COLUNN 24
LDY \#ま18
JSR $\ddagger$ FFF0
LDX \#ま 1 F
JSR $\operatorname{\text {FFFD}2}$
DEX
ENE SPACE
CLC !SET BACK TO ROW 日 COL 24
LDX \#\$00
LDY \#\$18
JSR 9 FFF
LDA \#\$12 !TURN ON REUERSE UIDEO
JSR 9 FFD2
LDA OLHIGH !LOW BYTE PREVIOUS LINE
LDX OLLOW ! HIGH
JSR $\ddagger$ BDCD ! PRINT LINE NUMBER
LDA \#\$92 !REUERSE OFF
JSR $\ddagger$ FFD2

| 970 |  | LDA \#\$20 | !BIT OF SPACE |
| :---: | :---: | :---: | :---: |
| 980 |  | JSR \#FFD2 |  |
| 990 |  | LDA \#\$12 | ! REPEAT FOR CURRENT LINE |
| 1000 |  | JSR \$FFD2 | ! GETTING ITS UALUES FROM |
| 1810 |  | LDA \$3A | ! CURLIN LOW BYTE |
| 1020 |  | STA OLHIGH | !NOW BECOMES OLD LINE |
| 1830 |  | LDX \$39 | ! CURLIN+1 |
| 1040 |  | STX OLLON |  |
| 1050 |  | JSR $\ddagger$ BDCD |  |
| 1060 |  | LDA \#\$92 |  |
| 1070 |  | JSR fFFD2 |  |
| 1080 |  | CLC | ! PREPARE TO RESET CURSOR |
| 1085 ! |  |  |  |
| 1096 | !IGNORE THI | 5 BIT AS ONL | Y TO ALLOU BRANCH TO WORK |
| 1095 |  |  |  |
| 1100 | BASIC1 | BCS BASIC | ! Halflday branch to basic: |
| 1110 |  | LDX ROW | !CONTINUE RESET CURSOR |
| 1120 |  | LDY COL |  |
| 1130 |  | JSR $\ddagger$ FFF0 | !RESTORE CURSOR POSITION |
| 1140 | $!$ |  |  |
| 1150 !CHECK FOR ANY KEYS PRESSED |  |  |  |
| 1160 |  |  |  |
| 1170 |  | JSR \$FFE4 | ! GETIN |
| 1180 | CHCHAR | BEQ SINGLE | ! NOTHING IN K/B EUFFER |
| 1190 |  | CMP \# ${ }^{\text {2 }}$ 2F | ! KEYく8??? |
| 1200 |  | BCC SINGLE | ! YES THEN OF NO INTEREST |
| 1210 |  | CMP \#\# 3A | ! KEY>G??? |
| 1220 |  | BCS SINGLE | !YES - NO INTEREST |
| 1230 |  | SBC \#\$30 | ! BETWEEN 8 AND 9 S0 - $\$ 30$ |
| 1240 |  | BNE CHDELAY | ! 1-9 |
| 1250 |  |  |  |
| 1260 | 19 PRESSED | SO RESET SIN | NGLE STEF |
| 1270 |  |  |  |
| 1280 |  | LDA \#\#FF |  |
| 1290 |  | STA SSTEP |  |
| 1300 |  | BtJE SINGLE | !NO NEED TO CALC delay |
| 1310 |  |  |  |
| 1320 | ! CALCULATE | DELAY AS PCW | ERS OF 2 |
| 1330 |  |  |  |
| 1340 | CHDELAY | TAX | ! PUT 1-9 IN X |
| 1350 |  | SEC | ! 1 IN CARRY |
| 1360 |  | LDA \#\$08 |  |
| 1370 | ROLL | ROL A | IMONE CARRY BIT $X$ TIMES |
| 1380 |  | DEX | !TO SET KEY-2 BIT |
| 1390 |  | BNE ROLL | !TO GIVE DOUBLING DELAY |
| 1400 |  | STA COLNT | ! STORE IT TO USE AS TIMER |




The ROM routines used are as follows:

## CHROUT (\$FFD2)

As Dump (see page 190).

## GETIN (\$FFE4)

See Programmer's Reference Guide, 'The kernal function B-11'. This removes one character from the current input device (usually the keyboard buffer) and returns its ASCII value in A. Zero is returned if none waiting.

## PLOT (\$FFF0)

See Programmer's Reference Guide, 'The kernal function B-19'. Reads the cursor position with the carry set and positions the cursor when the carry is clear. Misleadingly, $x$ is used for the row and $y$ for the column.

PRINT LINE NUMBER (\$BDCD)
Useful little routine, this one, and well worth noting. Not only can it be used for line numbers, but also for a two byte unsigned integer ( $\$ 0000$ to \$FFFF). Before calling it, $x$ must hold the low and $A$ the high byte. It also strips off the traditional leading and trailing spaces before printing.

## HANDLE NEW LINE (\$A734)

This routine is vectored by the page 3 vector IGONE (\$0308) and \$A7E4 is the 64's default setting. This is BASIC's token DISPATCH routine and is covered in great detail in Chapter 3. When used with the UTILITY, IGONE has been modified and hence the reason why IGONE has been first read and
stored. Doing it this way means the routine will work with or without the UTILITY. IGONE is called to tokenize each new line and is thus the ideal point at which to patch our trace.

LINE 130-290: TRACE ENABLE: These set up TRACE ready for when you RUN a BASIC program. A scan is made for direct operation only, and only if it is direct do we continue. During this initialization the interrupt will be disabled. The default speed is single step and its value is stored in the appropriate location at the end of the routine. The original value of igone, basic Character Dispatch Vector (see Chapter 6), are stored for safe-keeping and the start of TRACE replaces them. After clearing the interrupt, we return you to BASIC until the RUN command is issued.

LINES 330-290: TROFF - Trace Disabled: The reverse of the trace set up procedure.

LINES 470-690: The BASIC dispatch is used each time a BASIC command is issued whether in direct or program mode. This means that the routine can be called when not required. To avoid this, we check for program mode, after preserving the processing registers, as in the set up. If still in direct mode, we restore the registers and jump to the normal DISPATCH routine. A final check is made before operating TRACE to ensure it is enabled by looking at the trflag at the end of the routine.

LINES 730-1130: Display and updating are the purpose of these lines. We print at the top of the screen so as not to disrupt your display. We locate and save the current cursor position before setting it to the start of our print, top row and column 24, and clear the area we use by printing 15 spaces to the end of the line. After turning on reverse video, we gather in the values of the previous line number and visit the ROM routine to print it. To distinguish between the line numbers, we put a space between them, after turning off the reverse video. We now repeat the operation for the current line number. At the same time as getting the current line number, we store its value in our previous line store. As tRACE is called before every bASIC command is initiated, then when more than one command is on a line the previous and the present line numbers will be identical.

The instructions in line 1080 and 1100 are little tricks. Clearing the carry will ensure that the branch will fail. The branch is there for a later instruction when it will save a JUMP command.

Finally, we restore the original cursor position.
LINES 1170-1430: Is a speed change required? To find out, we use the KERNAL routine GETIN. If there is no character or it is not between the ASCII values of $\$ 30$ and $\$ 3 \mathrm{~A}$, we branch out of this section. On remaining here, we subtract $\$ 30$ from the character value to convert from ASCII to a real value between 0 and 9 .

The zero value signifies single-step mode is required again, so its flag is set to $\$$ FF.

The remainder have to be acted upon in order to gain our power of two figures for delay purposes. The value we have is to be used as a counter and so is transferred to the x register. To start with, we set the carry and initialize the accumulator to zero. The accumulator is rotated right $x$ times. The first time, the bit set in the carry is transferred to bit zero of the accumulator. The carry from this time on will be unset, except for the last time when nine is the value of $x$ and that will not worry us. Every consecutive rotate will shift the set bit further to the right and so increase the power of two of the value. When nine is raised to a power the accumulator will end up with a value of zero, but when the delay is explained, lines 1590-1090, you will see that this is in effect 256 , that is, $2^{\wedge} 8$. Finally, we disable the single-step mode, if set, by storing zero in its flag.

LINES 1470-1530: SINGLE-STEP MODE: To check if it is in operation, we test its flag. We not only look for the space character, which operates single-stepping, but also for others by branching back to lines 1170 to 1430 . This means to continue in single-step the space bar or a numeral key will in fact cause the program to continue, the latter ones ending single-step at the same time.

LINES 1590-1620: DELAY: This consists of a loop within a loop. The inner loop is completed the number of times calculated earlier, thereby giving variable time delays. When speed nine has been selected count is zero, the inner loop is carried out, and the COUNT is decreased before checking for zero. When zero is decreased it becomes $\$$ FF (255) so the check will fail until it is decreased to zero once more. This therefore operates the inner loop 256 times.

Finally, we set the carry and branch back, as the accumulator will be zero, to line 1100 where the Branch with Carry Set will send it further back to jump in to the normal IGONE routine to carry out the BASIC command.

## IMPROVEMENTS?

It is possible to modify the trace to list not only the previous and current line, but also to highlight the current statement being executed.

To list a line we can use the LIST routine in ROM which starts at \$A69C. There are two major problems if we try to use it. The first is that LIST uses a number of zero page locations also used during a run. The second is that on completion of LIST a warm boot of BASIC is carried out. (Try putting lIST in a program and running it.) We can overcome the first by copying zero page to elsewhere in ram before calling list. The second requires that on return from list, we must re-enter our TRACE
routine at the next instruction after performing LIST which must restore zero page. To do this we must read and store the warm start vector IMAIN (\$0302) and set it to the next instruction after LIST is called. TROFF must, of course, reset this vector to its original value.

To highlight the statement within a listed line places even greater demands on our ingenuity and would require the TRACE routine to be rewritten to use CHRGET which has purposely been avoided (because of DOS 5.1). If it were to use CHRGET, the line could be re-listed each time with a marker character printed at the current byte held in TXTPTR (\$7A/\$7B) through the use of the PRINT tokens link.

Both additions seem of little point as we can use the stop key followed by LIST line number(s). We have not even attempted to incorporate either possibility.

## Numeric conversions

In the world of the Commodore 64 we come across three main numbering systems: that of decimal, to the base of 10, hexadecimal, to the base of 16 , and binary which is to the base of 2 (octal, to the base eight is less common).

The binary number system is used because there are only two numerals used: 0 and 1. This matches the type of electronics used in the computer world, digital electronics, which has only two states, either on or off. These two positions are known as logical states. Logic 1 is on and obviously logic 0 is off. These, as you can see, fit well with the binary system.

The hexadecimal system was introduced because although binary matches the electronics, it is unwieldly and is not so easy to recognize in everyday form. Hexadecimal is easier to remember, using only two digits to the binary eight, and therefore faster to type in. Hexadecimal is nearly always entered in groups of two for example, \$FF.

Decimal is used in our everyday life and is therefore used in basic. One of its disadvantages is that numbers have varying amounts of digits. For instance, in numbers up to 255 there are one, two or three digits whereas with hex there is always two.

Some basics give you the option to enter numbers in forms other than decimal. The Commodore 64 basic does not. We are not going to rectify this but are giving you four conversion routines. These are converting decimal and binary, and decimal and hex.
TEN - Decimal to hexadecimal conversion
COMMAND SYNTAX
TEN decimal number [,decimal number,....]
The maximum decimal number that can be converted is 65535 and then only positive numbers can be converted. Multiple conversions can be
done if they are separated by commas. The result will be a four digit hex number.

## ASSEMBLY LISTING

| 9 | $*=\$ 84 \mathrm{Al}$ |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 10 | START | BEQ | SINTAX |  |  |
| 20 |  | BCS | SYNTAX |  |  |
| 30 |  | JSR | \$81F5 | ! | GET PARAMETER |
| 40 |  | LDA | \#\$20 |  |  |
| 50 |  | JSR | \$FFD2 | ! | PRINT A SPACE ON SCREEN |
| 60 |  | LDY | \#\$02 |  |  |
| 70 | NEXT | JSR | HEX | ! | CONUERT A BYTE TO HEX |
| 80 |  | DEY |  |  |  |
| 98 |  | BNE | NEXT | ! | TWO COMNERTS FOR EACH BYTE |
| 100 |  | LDY | \$14 |  |  |
| 110 |  | STY | \$15 | $!$ | place low byte FOR CONUERSION |
| 120 |  | LDY | \#\$02 |  |  |
| 130 | NEXT 1 | JSR | HEX |  |  |
| 140 |  | DEY |  |  |  |
| 150 |  | BNE | NEXT 1 |  |  |
| 160 |  | JSR | \$0079 | $!$ | GET LAST BYTE OF BLIFFER AGAIN |
| 170 |  | CMP | \#\$2C | $!$ | COMMA FOR MORE NUMBERS |
| 180 |  | ENE | EXIT | ! | NO MORE TO COMUERT |
| 190 |  | JSR | \$FFD2 | ! | PRINT COMMA AS SPACER |
| 200 |  | JSR | \$0073 |  |  |
| 210 |  | JMP | START |  |  |
| 228 | EXIT | RTS |  |  |  |
| 230 | HEX | LDX | \#\$04 | ! | SET COUNTER |
| 240 |  | LDA | \#\$00 | ! | INITIALISE ACC |
| 250 | AGAIN | ASL | \$15 |  |  |
| 268 |  | ROL | A |  |  |
| 270 |  | DEX |  |  |  |
| 280 |  | BTNE | AGAIN |  |  |
| 290 |  | CMP | \#\$0A | ! | IS IT 10 OR MORE |
| 300 |  | BCC | DIGIT+1 | ! | ASCII ADDITION FOR NLMBER |
| 310 |  | CLC |  |  |  |
| 320 |  | ADC | \#\$37 | $!$ | FOR LETTER |
| 330 | DIGIT | BIT | \$3069 |  |  |
| 340 |  | JSR | \$FFD2 | ! | PRINT RESULT |
| 350 |  | RTS |  |  |  |


| 8405 AGAIN | $84 E 2$ DIGIT |
| :--- | :--- |
| 84D0 EXIT | 8401 HEX |
| 84AE NEXT | $84 B A$ NEXT1 |
| $84 A 0$ START | $84 E 9$ SYNTAX |

LINES 30-50: Here we pick up the first decimal number to convert. The high byte will be in $\$ 15$ and the low, in $\$ 14$. We then print a space on the screen so the first digit is a character away from the border or the last PRINTed statement.

LINES 60-90: Each part of the hexadecimal number will have two characters. As we will convert our decimal in two stages, the high byte first then the low, each will require two entries to the conversion routine. We therefore set a counter to two, in this case the y register. After going to the conversion routine we decrease the counter. If it is zero then we have done it twice, if not we go back again.

LINES 100-110: We have now converted the high byte. As the conversion subroutine uses the high byte register in the transposition, we transfer the low byte of the decimal to that register.

LINES 120-150: This is the same as lines 40-70 but for the low byte.
LINES 160-220: The Get Parameters has already picked up the byte after the last decimal digit. Here we get that byte again by a call to Chrgot. We want to find out if more than one conversion is required. The syntax of the command is for a comma as a separator, so we check for that. If the check succeeds, we print the comma to the screen and go back to convert the number. On failing the check, it is back to BASIC via the RTS.

## THE CONVERSION ROUTINE

We use this routine four times for every decimal number in the command, twice for both the high and low bytes. We enter with the byte in location $\$ 15$. The hexadecimal number for a byte consists of two digits, one for the upper four bits and the other for the lower four. As we print on the screen from left to right, we print from the most significant hex digit and therefore want the high bits of the decimal number first.

Hex uses numerals 0 to 9 and letters A to F. Unfortunately, these do not follow in sequence in the ASCII table, as other characters lie between 9 and A. We therefore have to test for this when converting to ASCII for printing to the screen.

LINES 230-240: The $x$ register is initialized to 4 as a counter for taking off the required bits for each hex digit. The accumulator is used to gather in the bits so is initialized to zero before we start.

LINES 250-280: This is the main part of the conversion. We use the instruction ASL to move all the bits of the decimal byte one place to the left. The most significant bit (bit 7) is moved to the carry flag. The least significant bit (bit 0 ) is filled with zero (although that does not worry us). We need the bit we put into the carry back in the accumulator. This is achieved by the command rol. This moves the accumulator bit one place to the left, filling bit 0 with the carry value, which we have just set (or unset). Bit 7 of the accumulator goes to the carry and again it is of no use to us here.

Now the counter is decreased and checked to see if we have done the bit shifting four times. We have now taken the four high bits of $\$ 15$ (the decimal byte) and put them in the same order in the accumulator but in the low bit positions.
LINES 290-350: The answer in the accumulator is now converted to ASCII form and printed to the screen. If it is less than \$0A, it is a number so we add $\$ 30$ to it. Greater than 10 means it is a letter, so we have to add $\$ 37$, giving letters from A to $F$.

```
HEX - Converting a hexadecimal number to decimal COMMAND SYNTAX
HEX hex number [,hex number,....]
```

The hex number can be of two or four digits. More conversions can be added if separated by commas. The normal '\$' sign which preceeds hex numbers must not be used.

A four digit hex number can be split very conveniently into two parts. The two left digits are the high byte whilst the right are the low byte. Where a two digit conversion is required, we treat it as a low byte number. The two digits can be further split in that one represents the high nybble and the other the low nybble (a nibble is half a byte or four bits).

To do the conversion we collect two hex digits at a time and convert them to a one byte answer.

## ASSEMBLY LISTING

$9 *=\$ 8537$
10 START STA $\$ 63$

20
30
40
50
60
70
80
90
100

JSR $\$ 0073$ ! GET NEXT BYTE
STA $\$ 62$
JSR DECIMAL
PHA ! PUT HIGH ANS ON STACK
JSR \$0073
BEQ LOWPRINT! ONLY TWO BYTE HEX NO
CMP \#क2C ! IS IT A COMMA
BEQ COMMA ! YES \& ONLY 2 BYTE HEX STA $\$ 63$

| 110 |  | JSR | \$0073 |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 120 |  | STA | \$62 |  |  |
| 130 |  | JSR | DECIMAL |  |  |
| 140 |  | TAY |  | ! | PUT LOW ANS IN Y |
| 150 |  | PLA |  | ! | GET HIGH ANS IN ACC |
| 160 |  | JSR | PRINT | $!$ | PRINT ANSWER |
| 170 |  | JSR | \$0073 | ! | GET NEXT BYTE |
| 180 |  | CMF | \# ${ }^{\text {2 }} 2 \mathrm{C}$ | ! | IS IT A COMMA? |
| 190 |  | BEQ | COMMA1 | ! | YES |
| 200 |  | RTS |  |  |  |
| 210 | LOWFRINT | PLA |  | ! | GET HIGH ANS |
| 220 |  | TAY |  | ! | PUT IT IN LOW ANS REG |
| 230 |  | LDA | \#\$00 | ! | SET HIGH ANS TO ZERO |
| 240 |  | JSR | PRINT | ! | PRINT ANSWER |
| 250 |  | RTS |  |  |  |
| 260 | COMHA | PLA |  | ! | GET HIGH ANS |
| 270 |  | TAY |  | ! | PUT IT IN LOW ANS REG |
| 230 |  | LDA | \#\$00 | ! | SET HIGH ANS TO ZERO |
| 290 |  | JSR | PRINT | ! | PRNT ANSWER |
| 300 | commal | LDA | \#\$2C | ! | ASCII FOR COHMA |
| 310 |  | JSR | \$FFD2 | ! | PRINT IT AS SPACER |
| 320 |  | JSR | \$0973 | ! | GET NEXT BYTE |
| 330 |  | JMP | START |  |  |
| 340 | DECIMAL | LDY | \#\$01 | ! | COUNTER |
| 350 | AGAIN | LDA | \$0062, Y | ! | GET LOW CHAR |
| 360 |  | CMP | \#\$30 | ! | IS IT A NLMBER |
| 370 |  | BCC | SYNTAX | ! | NOT NLMBER OR LETTER |
| 380 |  | CMP | \#\$47 | ! | IS IT LETTER ) F? |
| 396 |  | BCS | SYNTAX |  |  |
| 400 |  | CMP | \#\$3A | ! | IS IT A NUMBER? |
| 410 |  | BCC | DIGIT | ! | YES |
| 420 |  | CMP | \#\$ 41 | ! | IS IT A LETTER? |
| 430 |  | BCC | STNTAX | ! | NO |
| 440 |  | SBC | \#\$37 | ! | CONNERT ASCII LETTER INTO REAL NLMBER |
| 450 |  | BNE | NEXT | ! | ENFORCED |
| 460 | DIGIT | SEC |  |  |  |
| 470 |  | SBC | \#\$30 | ! | CONUERT ASCII NLMBER INTO REAL NUMBER |
| 480 | NEXT | STA | \$0014, ${ }^{\text {\% }}$ |  |  |
| 496 |  | DEY |  |  |  |
| 500 |  | BPL | AGAIN | $!$ | NEXT CHARACTER |
| 510 |  | LDY | \#\$04 | ! | COUNTER |
| 520 | NEXT 1 | ASL | \$15 | ! | PUT HIGH ANS IN HIGH BITE |



LINES 10-50: The routine is entered with the first digit and is stored. Calling Chrget, we get the next byte and again store it. The decimal conversion routine is visited (this is described later), and the result comes back into the accumulator which we place on the stack.

LINES 60-90: The next byte of the command is now collected and two checks made, first to see if it is the end of the command and secondly for a comma. If the first succeeds we go off and print what we have already collected, but as a low byte answer. If it is a comma, we again print but will return to do further conversions.

LINES 100-130: This is a repeat of lines 10-50 except the result in the accumulator is put in the r register instead of the stack.

LINES 140-200: The result of the four digit conversion is printed to the screen here. The y register has the low byte and the high byte is pulled off the stack into A. The print routine described at the end is now called. The conversion is complete and we now check to see if further conversions are required by getting the next byte. If a comma is not present, the routine is ended.

LINES 210-250: The low byte answer is printed here. The byte is pulled off the stack and placed in the $y$ register and the accumulator set to zero. After printing, the HEX routine is left.

LINES 260-290: This is the same as lines 210-250 but instead of leaving,
we continue as we know there was a comma present when we arrived here.

LINES 300-330: A comma is printed on the screen to separate the answers. The first byte of the next conversion is gathered and we go back to the beginning to start converting again.
LINES 340-570: THE DECIMAL CONVERSION: The two bytes will be in locations $\$ 62$ and $\$ 63$. They will hold the ASCII values of either a numeral or a letter between A and F. Syntax errors are given if they fall outside these limits.

Taking each location in turn, we determine what it is and deduct from it $\$ 37$ for a letter of $\$ 30$ for a numeral, the value ending up between $\$ 00$ and $\$ 09$ ( 0 to 15 ). These are placed in registers $\$ 14$ and $\$ 15$. We now have to combine these into one number. Address $\$ 15$ will have the high nybble but in the wrong bit positions. To get them into the upper four bits we shift the bits left four times. To join the two together, the byte in $\$ 14$ is copied to the accumulator and is ored with location \$15. With the final result in the accumulator, we exit the subroutine.

LINES 590-640: PRINT RESULT TO SCREEN: Six subroutines are called here where the result of numeric calculations are converted to a string of ASCII characters and printed to the screen which except for one are all rom routines. The one exception is a subroutine in the DEEK routine (see Chapter 8). For convenience, we reproduce it below. We enter this PRINT routine with numeric data, the y register holding the low and the accumulator the high byte.
ROUTINE \$B391 - FIX TO FLOAT
This sets the data flag in SOD to zero signifying numeric data. The number we wish to convert is placed in FAC\#1 registers \$62, meaning that numbers over 32768 ( $\$ 80$ ) will be output as negative numbers.
ROUTINE $\$ 8401$ - CONVERT TO POSITIVE

| 10 | LDA $\$ 66$ |
| :--- | :--- |
| 20 | BPL EXIT |
| 30 | LDY \#〉DATA |
| 40 | LDA \#〈DATA |
| 50 | JSR $\$$ BA8C |
| 60 | JSR $\$ B 86 A$ |
| 70 EXIT | RTS |
| 80 DATA | BYT $\$ 91, \$ 00, \$ 00, \$ 00, \$ 00$ |

We check register \$66 of the FAC\#1 to see if it is negative. If so we load FAC\#2 with zero and set for no exponent. This is done through the ROM routine ${ }^{\$ B A B C}$, entering with the data start address in $A$ and $Y$.

Now by adding the two facs together we will end up with a result in FAC\#1 which is a real whole number; $\$ B 86 A$ will achieve this.

## ROUTINE \$AD8A - CHECK

This just checks that the data is numeric, otherwise a 'TYPE MISMATCH' error is given.

ROUTINE \$BDDD - FAC\# 1 TO STRING
This converts the contents of FAC\#1 to an ASCII string and places it at the bottom of the stack. The Y and A registers will hold this address when the routine is finished.

ROUTINE \$B487 - SET UP STRING
This sets various registers so that the PRINT routine knows where to print from and how long the string is.

ROUTINE \$AB21 - PRINT
This takes the data from the bottom of the stack and prints it to the screen. We jumped to this routine, so when it is ended, the processor will be directed back to the position calling this whole subroutine.

This routine, being a separate routine, is therefore capable of being used by other commands as in the MEM command.

```
TWO - Decimal to binary conversion COMMAND SYNTAX
TWO decimal number [,decimal number,....]
```

The maximum decimal number which can be converted is 65535 and must be positive. Multiple conversions can be done if they are separated by commas. The result will be two eight digit binary numbers separated by a space, unless the number is 255 or less, when only one binary result will be shown.

All we need to do is to test each bit and print a zero or a one according to its state.

ASSEMBLY LISTING
$9 *=\$ 84 E C$
19 START
20

| BEQ SYNTAX |  |
| :--- | :--- |
| BCS SYNTAX |  |
| JSR $\$ 81 F 5$ | GET FARAMETER |
| LDA \#\$20 |  |
| JSR $\$ F F D 2$ | ! PRINT SPACE |
| LDA $\$ 15$ |  |
| BEQ LSB |  |
| LDX \#\$08 |  |
| ASL $\$ 15$ |  |


| 100 |  | BCS SET+1 |  | TO PRINT A 1 |
| :---: | :---: | :---: | :---: | :---: |
| 110 |  | LDA \#\$30 |  | TO PRINT A O |
| 120 | SET | BIT \$31A9 | ! | SET+1 IS LDA \#\$31 |
| 130 |  | JSR \$FFD2 |  |  |
| 140 |  | DEX |  |  |
| 158 |  | BNE NEXT |  |  |
| 160 |  |  |  |  |
| 170 |  | JSR 事FFD2 | ! | PRINT A SPACE |
| 180 | LSE | LDX \#哑08 |  |  |
| 190 | NEXT 1 | ASL ${ }^{\text {a }} 14$ |  |  |
| 200 |  | BCS SETA+1 |  |  |
| 210 |  | LDA \#\#30 |  |  |
| 220 | SETA | BIT ${ }^{\text {31AF }}$ | ! | SETA+1 IS LDA \#\$31 |
| 230 |  | JSR ${ }^{\text {SFFD2 }}$ |  |  |
| 240 |  | DEX |  |  |
| 250 |  | BNE NEXT1 |  |  |
| 268 |  | JSR $\$ 6079$ | $!$ | GET LAST BYTE OFF BUFFER AGAIN |
| 270 |  | CMP \#\$2C | 1 | 15 IT A COMMA? |
| 280 |  | BNE EXIT | ! | NO COMHA |
| 290 |  | JSR \$FFD2 | $!$ |  |
| 306 |  | JSR \$0073 | ! | GET NEXT BYTE |
| 310 |  | JMP START |  |  |
| 320 | EXIT | RTS |  |  |
| 330 | SVNTAX | JMP \$AF08 |  |  |

8533 EXIT
84FE NEXT
8504 SET
84EC START

8512 LSB
8514 NEXT1
851A SETA
8534 SMTAX

LINES 10-20: On entering the first byte after the keyword is in the A register and by testing the carry and zero flags, we can check if a numeral is first.

LINES 30-150: We gather in the decimal number to convert and also print a space for presentation to move it away from the last item printed or from the border. Taking the high byte first, we check it for zero, if it is a zero we branch and just do the low byte. The x register is set to eight as there are eight bits to a byte, and we shall use it as a counter. We shift the bits of the high byte one place to the left. The leftmost bit comes off and goes to the carry flag. If the carry flag is zero, we therefore print ASCII $\$ 30$ which is zero and $\$ 31$ when the carry is set. This is repeated a further seven times for all the remaining bits.
LINES 160-250: This is a repeat of the above except for the low byte.

LINES 270-320: By calling chrgot, we test for a comma or if the end of the command has been reached. When the former is found, it is printed to the screen and we gather in the next byte before going to carry out the next conversion.
BIN - Binary to decimal conversion
COMMAND SYNTAX
BIN 8 bit binary number [, 8 bit binary number,....]
Here we will convert an eight bit binary number to decimal. We supply a value that would be a high byte and one that is the low byte. For example, if you demanded 11111111 was converted, the answer would come out as 255 . Only eight bit numbers are accepted but more conversions can be done by separating the items with commas.

This is essentially the reverse of the previous command. The 1 s and 0s you type in will be picked up in their ASCII form. These have their rightmost bits corresponding to their numeric value, so by taking those we can build up a single byte number.

ASSEMBLY LISTING

| $9 *=\$ 85 B F$ |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 10 |  | LDA | \$7A |  |  |
| 20 |  | BNE | LOW |  |  |
| 30 |  | DEC | \$ 7 B | ! | DECREASE CHARGET |
|  |  |  |  |  | POINTER BY ONE |
| 40 | LOW | DEC | \$7A |  |  |
| 50 | ANOTHER | LDX | \#\$88 | ! | COUNTER |
| 60 | NEXT | JSR | \$0073 | ! | GET BYTE |
| 70 |  | BCC | NLMBER | ! | IF NLMBER BRANCH |
| 80 | SYNTAX | JMP | \$AF08 |  |  |
| 90 | NLMMER | CMF | \#\$32 | ! | IS > ASCII FOR 2 |
| 100 |  | BCS | SYNTAX | ! | YES |
| 110 |  | ROR | A | ! | GET BIT 0 |
| 120 |  | ROL | \$14 | ! | PUT IN 144 MOUING |
|  |  |  |  |  | 1 LEFT EUERY TIME |
| 130 |  | DEX |  |  |  |
| 140 |  | BNE | NEXT | ! | DONE IT 8 TIMES? |
| 150 |  | LDY | \$14 | ! | PUT ANS IN Y REG |
| 160 |  | LDA |  | ! | NO HIGH ANS |
| 170 |  | JSR | *85AD | ! | PRINT ANS-HEX ROUTINE |
| 180 |  | LDA | \#\$ 2 F |  |  |
| 190 |  | JSR | \$FFD2 | ! | PRINT SLASH TO DIUIDE |
|  |  |  |  |  | LOUN ANS \& HIGH ANS |
| 200 |  | LDA | \$14 | ! | NOW PUT ANS IN |
|  |  |  |  |  | HIGH ANS REG |
| 210 |  | LDY | \# 000 | ! | SET LOW ANS REG TO |


| 220 | JSR \$85AD | ! PRINT ANS-HEX ROUTINE |
| :--- | :--- | :--- |
| 230 | JSR $\$ 0073$ | ! GET NEXT BYTE OF |
|  |  | INPUT BUFFER |
| 240 | CMP \#\$2C | IS IT A COMMA? |
| 250 | BNE EXIT | ! NO |
| 260 | JSR $\$ F F D 2$ | ! PRINT COMMA AS SPACER |
| 270 | JMP ANOTHER |  |
| 280 EXIT | RTS |  |


| 8507 ANOTHER | $85 F E$ EXIT |
| :--- | :--- |
| $85 C 5$ LOW | $85 C 9$ NEXT |
| 8501 NUMBER | $85 C E$ SYNTAX |

LINES 10-40: This decreases CHRGET address by one.
LINES 50-140: We want to pick up eight binary digits so the $x$ register is used as a counter. After we pick up a digit, via CHRGET, we check for a number and also if it is two or greater, in ASCII. By rotating the accumulator right, we take off bit $\emptyset$ and it ends up in the carry flag. Then if we rotate location $\$ 14$ left, we move all its bits one to the left and put the carry flag state into the lowest bit. If we do this eight times, address $\$ 14$ will have a number equivalent to the $1 s$ and 0 s you typed in.

LINES 150-220: The PRINT routine in the HEX command is used twice. For this the low byte needs to be in Y and the high in A . The value we want to print is in $\$ 14$ and by changing the register we load it into we can print out the states we want. A slash is printed as a separator by a visit to the KERNAL routine at \$FFD2.

LINES 230-280: Having done one conversion, we take a look to see if more are required. A comma is printed if so and then we go back to do it all again.

## 8 Enhancing the resident BASIC

## Introduction

In the previous chapter the commands were of a toolkit nature. In this chapter they are mainly improvements to the standard 64 commands, GOTO, GOSUb, RESTORE, PRINT, INPUT, GET, PEEK, POKE and LOAD. To that end, we are supplying CGOTO, CGOSUB, PROC, DPROC, EPROC, RESET, WRITE, ENTER, INKEY\$, DEEK, DOKE and CHAIN. There are five commands which have no 64 BASIC equivalent, but which we hope will enhance your BASIC programming. They are POP, PLOT, COLOUR, LOMEM and HIMEM. The final command given is that of QUIT and exists the UTILITY.

In comparison with the toolkit commands these are shorter, but no less useful to you. No doubt you can think of existing commands which could be enhanced and even more to add. This chapter should help you on your way.

## CGOTO and CGOSUB

## COMMAND SYNTAX

сяото variable or line number
cGosub variable or line number
A limitation of Commodore BASIC is that it does not permit the use of calculated destinations with GOto and gosub. We thought it would be nice to be able to use variables and mathematical expressions, for example A*20. To allow this, we have come up with two commands сGOTO and ccosub, the c standing for calculated or computed - whichever you prefer.
cGoto is the easiest of the two, not that either is complicated. The routine requires only two instructions. In the BASIC ROM routine of GOTO the first instruction gets the line number and is therefore the only thing we have missed out. So after getting the variable we only have to jump to that part of GOTO.
cGosub is a bit longer in that we have to copy the ROM routine for cOSUB and change the address for calling the GOTO routine as we want to use our 'computed' routine.

ASSEMBLY LISTING


8FAF CGOTO

LINES 10-20: сGoto: We use get parameters simply to find the destination line number. It will evaluate any expression, and jump to the normal GOTO routine, one instruction in.

LINES 40-60: cGOsub: These lines check if there is enough room on the stack to store the routine's information and a buffer amount for other routines. To do this the value in the accumulator is doubled and added to $\$ 3 \mathrm{E}(62 \mathrm{dec}$ ). This is then compared with the stack pointer. If the stack pointer is the lesser value, then an 'OUT OF MEMORY' error is generated. In our case, the stack pointer would have to be less than $\$ 44$ (it starts at \$FF).

LINES 70-130: There are two markers we will require when the subroutine is finished. These are the present byte's address from the CHRGET routine and the line number we must return to later. The stack is used to store this information.

LINES 140-150: Another value is put on the stack. This is used by the RETURN routine to check a cosub has been implemented.

LINES 160-170: Now we can go to our destination. To do this, we get the last byte collected by CHRGET again and go to our new computed coto routine.

LINE 180: Once the GOTO routine has been completed, in which the CHRGET has been given new values, we return to the normal flow of BASIC and the program is continued at its new address.

## Procedures

## COMMAND SYNTAX

PROC title - call a procedure
DPROC title - define a procedure
EPROC - the end of a procedure
The title is not required within quotes. If it is then the quotes will be considered as part of the name. Spaces also cannot be used as Chrcet ignores them (a space in the DPROC title will not be matched in the PROC title). On the other hand, a space in the PROC title will have no bearing on the matching. A colon is the only other character which cannot be used in a title.
You can have as many PROCs on a line as you want, but the DPROC must be on a line of its own. Everything following the DPROC to the end of the line is included as the title. EPROC follows exactly as return.

64 BASIC cannot be described as a structured language. GOTOs and cosubs do not form the basis of a structured language.

To start you on the road to 'structured programming', we are introducing procedures. We have nothing profound to offer but by giving you an introduction we hope you will be able to take it further (IF . . THEN . . ELSE WHLLE . . . WEND DO . . . UNTIL etc . .)

The form of procedures we have written are really no more than gosubs with names or variables (ccosub). In fact, they will be slower, but not that you would notice, than cosubs because of the extra code required. So what advantage will they have? Well,'they can be relocated anywhere in the program without changing any directive line numbers; adding procedures from one program to another, especially if they include procedures within them, is a simple matter. If сото was also given the same treatment, all directive line numbers could vanish. Renumbering a program would be a simple matter of changing the line numbers rather than going through the whole program and correcting destinations. A further function they perform, and one that should not be overlooked, is that they make a program easier to follow. For instance, to see PROC PERFORM-WAIT is clearer than GOSUB 2000.

Quite simply, all we do when finding a PROC is to search through for the token DPROC and then compare the named titles. On finding it, we perform a cosub. The UTILITY interpreter will action the command DPROC as a rem if it encounters one. The third command of the trio is eproc and is just a return by a different name. We actually go to the return routine. After the listing and description we suggest some improvements.

ASSEMBLY LISTING


| 360 |  |  | \＄A437 |  | UNDEF＇D STATEMENT ERROR |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 370 | CONT | LDY | \＃ O $^{\text {a }} 04$ | ！ | SKIP LINKS AND |
|  |  |  |  |  | LINE NUMBERS |
| 380 |  | LDA | （\＄FB），$Y$ |  |  |
| 390 |  | C．MP | \＃敉E1 | ！ | TOKEN OF DPROC |
| 400 |  | BEQ | PROC | $!$ | FOUND A DPROC |
| 410 | LINE | LDA | \＄FD | ！ | FUT LINKS TO |
| 420 |  | STA | 韦FB | ！ | LINE REGISTERS |
| 430 |  | LDA | कFE |  |  |
| 440 |  | STA | 缶C |  |  |
| 450 |  | BNE | NEXT | ！ | ENFORCED कFE |
|  |  |  |  |  | CHECKED EARLIER |
| 460 | PROC | LDX | \＃\＃FF |  |  |
| 470 | CHECK | INX |  |  |  |
| 480 |  | INY |  |  |  |
| 490 |  | LDA | （珄FB），Y | ！ | GET PROC TITLE |
| 500 |  | BEQ | 2ERO | $!$ | END OF LINE |
| 510 |  | CMP | \＄8200，$\times$ | ！ | COMPARE FOR MATCH |
| 520 |  | BEQ | CHECK | ！ | MATCH FOLND |
| 530 |  | BAVE | LINE | ！ | NO MATCH FIND NEXT DPROC |
| 540 | 2ERO | CMP | \＄0200 | ！ | COMPARE LAST EYTE |
| 550 |  | EANE | LINE | 1 | NO MATCH |
| 560 |  | SEC |  | ！ | PREPARE FOR SUBTRACT |
| 570 |  | LDA | \＄FB | ！ | ADDRESS OF PROCEDURE |
| 580 |  | SBC | \＃ 01 | ！ | decrease to end of LAST LINE |
| 590 |  | STA | \＄7A | ！ | PUT AS CHRGET ADD |
| 600 |  | LDA | \＄FC |  |  |
| 610 |  | SBC |  | ！ | IN CASE PAGE CROSSED |
| 620 |  | STA | \＄78 |  |  |
| 630 |  | JMP | \＄A7AE | ！ | EASIC TO CONT PROG |
| 902 C | CHECK |  | 8FDA COLL | ECT |  |
| 9018 | CONT |  | 9020 LINE |  |  |
| 8 FE 5 | NAMEEND |  | 9006 NEXT |  |  |
| 902 A | PROC |  | 9039 2ERO |  |  |

LINES 10－110：Using CHRGET，after decreasing it by one，we take the PROC title and store it at the start of the input buffer and tag a zero byte on the end for checking purposes．

LINES 120－230：Same as in CGOSUB，saving relevant details for RETURN，or in this case EPROC．

LINES 240-360: After collecting the start address of the program, we search through the program. This part gets the links and checks for the end of the program. The Undef'd statement error is given if the latter occurs without finding the procedure.

LINES 370-450: When inputting a BASIC line, any spaces between a line number and the first character are removed during tokenizing, (UST inserts a space for clarity). This means that the first token in a line is the fourth byte (starting at zero, remember), so we check only this byte for the DPROC token of \$E1. If not found, the link address is placed into the line registers and the hunt continues.

LINES 460-550: Having found a DPROC token, we have to compare each character separately and as long as they match we continue checking. When we reach the end of the DPROC program line, we check the input buffer for a matching zero. When all checks succeed, we have found the required procedure.

LINES 560-630: Knowing our destination, we take the start address of the DPROC line and reduce it by one, the end of the preceding line, and store it as the CHRGET address. Finally, we jump to BASIC to evaluate.

## IMPROVEMENTS?

One of the first questions that came to mind was: how could we speed up the search for the procedures? One solution to the problem is to form a table in RAM holding the start address where you first check to see if the PROC name is in it. This would involve setting aside an area of RAM: under ROM would be an ideal place, for such a table. Two characters would then have to be chosen: one to mark the end of an entry and the other, the end of the table. The make-up of the table could consist of the PROC name and its start address. How could the table be filled? When the interpreter finds the keyword PROC, the table would be searched for a match. If no match is found, then a program search, like our routine, could be instigated. On finding the DPROC with the correct name, it would be added to the table in case of another call.

There are, however, problems. Let us assume that a program containing PROCs has been RUN. This would mean the table has names and addresses within it. Before running it again, you add an extra line before the procedure. The line with the DPROC now has a different address from that in the table. Another action giving rise to the same problem is when you load another program. It may have a PROC with the same name as the previous program. Again the table may have another address. A further problem may arise in that more PROCS will be added making the table longer and longer.

Two solutions to this problem spring to mind; there are probably others. The first is to write a new RUN command, for example, PRUN,
where one of its actions would be to place the end of table marker back to the beginning - thereby effectively clearing the table. The other is to have a command that can be actioned to do just this and only this. It could be initiated in direct mode or from within a program.

A further improvement would be to allow parameters to be passed using variables which are local to the procedure. These variables could be used elsewhere in the program without losing their original values. We would envisage the PROC command to include, in say, brackets, the values or other variables to be used, for example, $\operatorname{PROC} \operatorname{INPUT}(2,4,6)$ or PROC INPUT $(2, A, 6)$. The DPROC, on the other hand, will define the variables to be used. For example, DPROC INPUT $(X, Y, Z)$. These variables may be used elsewhere but in the procedure they will start with values given in the PROC command.

What would have to be done is that when arriving at the procedure a search is carried out for the variables $x, Y, Z$ in the normal variable area. If they are found, their current values would have to be transferred to a keeping area, and the new values set up. If the variable is not present, then it will have to be created. The default value of a numeric variable is zero and this will also have to be stored in the keeping area. For strings, the addresses will have to be stored. The EPROC would have to reverse the situation and restore the original values.

The process would be the same if you wanted to incorporate clobal and LOCAL Commands to a BASIC extension.

Our last improvement, although we are sure you could think of more, is to allow the names of procedures to include keywords. This would be relatively simple in that all you have to do is to slightly alter the CRUNCH token routine (see Chapter 3). In that routine when it comes across a REM, for instance, it skips further crunching. All you have to do is to insert further coding to check for PROC and DPROC and follow the same path as REM.

## POP - RETURN without returning <br> COMMAND SYNTAX

POP
There are no parameters to this command. If it is activated without a GOSUB, CGOSUB or a PROC being used, a 'return WITHOUT GOSUB' error will be generated.

There are many occasions when one requires to leave a subroutine but not go back to the calling position. This is, of course, possible but leaves values on the stack; do it too often and the stack will become full and an 'OUT OF MEMORY' error will occur

POP will remove from the stack the data placed there by the last

GOSUB, or equivalent. This will mean, for example, if you called a subroutine which in turn called another and whilst in the second you called POP, then you will go back to main program when the RETURN is met, not the first subroutine. A GOTO after a POP will mean you can go anywhere from a subroutine without any worries about the stack. POP will also discharge any FOR/NEXT loops. If you happen to be in one at the time, watch out.

ASSEMBLY LISTING


## 8642 CONT

LINES 10-20: By loading \$4A with \$FF, we effectively cancel any FOR/NEXT loop.

LINES 30-80: This is the ROM routine used by RETURN to look for the gosub marker on the stack. On return the stack pointer is in the $X$ register and the accumulator has a value from the stack. If this is $\$ 8 \mathrm{D}$, the RETURN marker was present. An error will be produced if anything else is found.

LINES 90-150: To remove the GOSUB activity, we take the stack pointer, which is still in the $X$ register, and increase it by five and then use it as the new stack pointer.

## RESET－Selective data restorer

COMMAND SYNTAX
RESET［line number］
When no line number is present it behaves as the standard command restore．With the parameter it will set the data pointer to the specified line．

DATA statements are extremely useful commands，and with sprites on the 64 you no doubt use them frequently．The snag comes when you want to use the same DATA statements again．RESTORE only allows you to set the pointer back to the first DATA statement，actually the start of the program，which has the same effect．To use statements again that are not at the beginning，dummy READs have to be employed to get to the desired position．To allow you greater flexibility，RESET will allow you to specify the line the next READ will start at，whether before or ahead of the present position．The restore command takes the start of program address，subtracts one from it，and places it in the DATA pointer regis－ ters．RESET will take the line number，find its address，decrease it and set the pointers．Although the routine will give an error if the prescribed line number is not present，we do not check to see if it is a DATA line． This does not matter to bASIC as it will find the next data line when read is sanctioned．

ASSEMBLY LISTING

| 9 | ＊＝\＄8611 |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 10 |  | BNE | RESET | ！ | CHECK FOR PARAMETER |
| 20 |  | JMP | कA81D | ！ | NO－RESTORE IN BASIC ROM |
| 30 | RESET | JSR | \＄81F5 | $!$ | GET LINE NLMBER |
| 40 |  | JSR | \＄A613 | ！ | FIND BASIC LINE |
| 50 |  | BCS | CONT | ！ | FOUND LINE |
| 60 |  | LDX | 姓 15 | ！ | ILLEGAL DIRECT ERROR |
| 70 |  | JMP | 中A437 | ！ | ROM ERROR ROUTINE |
| 80 | CONT | SEC |  | ！ | PREPARE FOR SUBTRACT |
| 90 |  | LDA | \＄5F | ！ | LOW ADD OF LINE |
| 100 |  | SBC | \＃${ }^{\text {b }}$ 01 | $!$ | DECREASE BY ONE |
| 110 |  | STA | \＄41 | ！ | data reg in page o |
| 120 |  | LDA | 事 60 | ！ | HIGH ADD OF LINE |
| 130 |  | SBC | \＃${ }^{\text {b }} 0$ | ！ | IF PAGE 15 CROSSED |
| 140 |  | STA | \＄42 | ！ | DATA REGISTER |
| 150 |  | RTS |  |  |  |

LINES 10-20: If no line number, we go straight to restore in the rom.
LINES 30-70: When the line number is picked up it will be in the right location for a line search, which is immediately carried out. The carry flag set will indicate the line was found. The error given for not finding it is 'Illegal direct'.
LINES 80-150: Locations $\$ 5 F$ and $\$ 60$ will have the address of the line, and from these we subtract one and store them in the DATA pointers.

## DEEK and DOKE - BASIC Addressing

It should be clear by now that addresses are stored in two locations as a low and a high byte. In the resident basic the only way to find the address, held in locations, is to do two PEEKs, one for each location, then multiply the high byte by 256 and add in the low byte, giving the address in decimal. To set up an address, the reverse process is used, but using POKE in place of PeEk.

The UTILITY commands are therefore obvious. We wish to read or set an address, or pair of locations, with one command. These are deek and DOKE.

## DEEK - seeing double

## COMMAND SYNTAX

DEEK (low byte location)
This returns the 16 bit value held in the given address and the following one. The rules for PEEK apply in that it must be an argument to a command (that is, a function).

ASSEMBLY LISTING

| $9 *=\$ 307$ |  |  |
| :---: | :---: | :---: |
| 10 | LDA ${ }^{\text {P }} 15$ |  |
| 20 | PHA |  |
| 30 | LDA ${ }^{\text {P14 }}$ |  |
| 40 | FHA |  |
| 50 | JSR *AEFA | ! CHECK FOR < |
| 60 | JSR \$81F5 | ! GET PARAMMETER |
| 70 | ISR 亜AEF7 | ( CHECK FOR ) |
| 80 | LDY \#\$01 |  |
| 90 | LDA (\$14), Y |  |
| 100 | TAX |  |
| 110 | DEY |  |
| 120 | LDA (事14), Y |  |

130
140
150
160
170
180
190
200
210
220
230
240
250
260
270
280
290
300 EXIT
310 DATA

8401 CONUERT
$840 F$ EXIT

LINES 10-40: DEEK, being a BASIC function rather than a command, is used in conjunction with other keywords. You have no doubt gathered that keywords use a fair number of zero page locations, notably $\$ 14, \$ 15$ and the facs. We cannot take deek in isolation and also have to get its parameters. The latter means that we will use $\$ 14$ and $\$ 15$. We do not require to use these on exit, so we take the precaution of saving the current contents on the stack for the time being.
LINES 50-70: These not only get the parameters, but also check for the convention of them being in brackets. The rom routines used will give the error if they are not present.
LINES 80-180: Using the address, now in $\$ 14$ and $\$ 15$, we read the contents and store them into the registers $A$ and $Y$. We can restore the original values to $\$ 14$ and $\$ 15$, and do so.
LINES 190-230: The calling routine will expect the result in the FAC\#1 and this is all the routine at $\$ B 391$ does. Unfortunately it also stores it as a signed integer. To correct this, CONVERT is called. Having done so, we pull off the return address, but we do not require to go back to the evaluation routine, and jump back to rom. In rom it will check that it is numeric data and will return to the calling routine, say PRINT Or DOKE.

LINES 240-300: CONVERT: If the number requires converting it will have
a negative sign. With no sign we exit the routine. Failing that we load a constant into FAC\#2 which when added to the contents of FAC\#1, will change it to an unsigned number. For a more detailed explanation, see MEM in Chapter 7.

DOKE - complete addressing
COMMAND SYNTAX
DOKE low byte address, value
This turns the value to a two byte number and stores it in the given address and the following one. The value has a maximum of 65535 (\$FFFF).

## ASSEMBLY LISTING

$9 \%=\$ 8383$

| 10 | JSR \$81FS | ! GET PARAPMETER |
| :--- | :--- | :--- |
| 20 | JSR \$AEFD | ! CHECK FOR COMMA |
| 30 | JSR \$AD8A | ! GET NEXT PARAMETER |

40 LDA $\$ 66$
50 B+1! ERROR
60 CMP \#991
70
80 JSR \$BC9B
90
LDA \$65
100 LDX ${ }^{3} 64$
110 LDY \#क00
120 STA (事14),Y
130 INY
140 TXA
150 STA ( $\$ 14$ ), Y
160 RTS
170 ERROR JMP \$B248 ! ILLEGAL QUANTITY
ERROR

83D4 ERROR

LINES 10-30: The first routine called is the familiar one. The address will be in $\$ 14$ and $\$ 15$ after this call. The next routine checks for a comma. The last one collects the data for storing and puts it in the FAC\#1.

LINES 40-80: These check for the legality of the data and set up the FAC\#1 so we can take our values off.

LINES 80-160: After getting the data from FAC\#1, we store them in the addresses specified.

## OUTPUT - Setting the cursor

In the standard BASIC, the normal way to set the print position is to use the cursor control codes. Although they do the job, they are not ideal. You have to remember where the current position is, they take up bytes in the program, and TAB and SPC are not much better.

A far better way would be to specify the $X$ and $Y$ coordinates directly. To do this, three commands are included here, PLOT, WRITE and ENTER. The first will only set the cursor, the second will set the cursor and print what you want, whilst the last is INPUT with cursor positioning. The major command as far as routines go is PLOT. It really is a subroutine for the other two.

## PLOT - cursor setting

COMMAND SYNTAX

## PLOT (X,Y)

The maximum value of $X$ is 39 and of $Y$ is 24 . The top left hand corner of the screen, cursor home position, has the coordinates of $\emptyset, \varnothing$.

ASSEMBLY LISTING

|  | JSR 䦈AEFA | ! | CHECK FOR < |
| :---: | :---: | :---: | :---: |
|  | JSR \$81F5 | ! | GET PARAMETER |
|  | LDA \$14 |  |  |
|  | CMP \#\$28 | ! | IS X > 40 |
|  | ECC COMMA |  |  |
| ILLEGAL | JSR \$ $\mathrm{B}^{\text {248 }}$ | ! | ILLEGAL QUANTITY ERROR |
| COMMA | PHA |  |  |
|  | JSR ${ }^{\text {F }} \mathrm{AEFD}$ | $!$ | CHECK FOR COMMA |
|  | JSR \$81F5 | ! | GET PARAMETER |
|  | LDX \$14 |  |  |
|  | CPX \#\$19 | $!$ | IS Y > 25 |
|  | BCS ILLEGAL |  |  |
|  | PLA | $!$ | RETRIEUE $15 T$ PARAM |
|  | TAY |  |  |
|  | CLC | ! | SET NOT READ CO-OR |
|  |  |  |  |

LINES 10-70: The left hand bracket is checked for and the X coordinate of the command is picked up. It is then checked that it does not exceed the limit. On an occasion that it does, we go to a rom routine whose sole purpose is to generate the 'Illegal quantity error'. We require to use location $\$ 14$ again so the X coordinate is put on the stack for a while.

LINES 80-120: After checking for the separating comma, we get the $Y$ coordinate. This, too, is checked for legality.

LINES 130-180: the $Y$ coordinate was picked up in the $X$ register and now we retrieve the $X$ coordinate and place it in the $Y$ register. This is the opposite to what is logical but the KERNAL routine calls for them in that order. Before calling the routine, we clear the carry flag. (If we set the carry we would read the cursor position.)

After setting the cursor we get the next byte. This is for WRITE and ENTER so that they are set up for their respective rom routines.

## WRITE and ENTER

COMMAND SYNTAX
WRITE ( $\mathrm{X}, \mathrm{Y}$ ) [string or variable]
ENTER (X,Y)[string], variable
The coordinates take the syntax of PLOT. The remainder of the commands have the same syntax as their respective standard commands, WRITE as print and enter as input.

ASSEMBLY LISTING

| 9 | $*=\$ 83 A 7$ |  |  |
| ---: | ---: | :--- | :--- |
| 10 | WRITE COHMAND -FRINT |  |  |
| 20 | JSR $\$ 8381$ | ! PLOT ROUTINE |  |
| 30 | JMP \$AAAG | ! PRINT ROUTINE IN ROM |  |
| 40 | ENTER COMMAND -INFUT |  |  |
| 50 | JSR $\$ 8381$ | ! PLOT ROUTINE |  |
| 60 | JMP $\$ A B B F$ | INPUT ROUTINE IN ROM |  |

These simply call the previous plot routine and then go to their normal rom routines.

## Colour

COMMAND SYNTAX
COLOUR background[,border][,text]
The latter two parameters are optional. If they are omitted it will not affect their present values. There is no error checking on values in the command. However, only the low byte of a number is used, that is, numbers up to 255 , and of that only the lower four bits have effect ( 15 uses four bits whilst 16 uses five). The values to be used are the same as in the Programmer's Reference Cuide or if you prefer the key number less one, with the logo key the number plus seven. Variables can be used as parameters. If no parameters are used, the background only will be changed, and that will be to black.

ASSEMBLY LISTING


## 8380 EXIT

LINES 10-60: This handles the background colour. We get the first parameter of the command, and load in the low byte only. This is ANDed with $\$ 0 F$ which will set the top four bits to zero no matter what state they were in. The result is used to set the colour. Finally, we
check, by getting the last byte again, if the end of the command has been reached. If it has not, we continue.

LINES 70-130: This first checks for the comma. Then we can get the parameter and proceed as for the background, except to store the value in the border register.
LINES 140-190: This is the same as above except, of course, we set the text colour.

## CHAIN - Passing variables

COMMAND SYNTAX

## CHAIN ["filename"],[device]

The syntax for ChAIN follows that for LOAD except for the secondary address. No errors will be given for the inclusion of the secondary address, as the routine will overwrite it.

One of the problems of LOAD is that if you load a larger progrram, after running a smaller program, you overwrite any variables. Also LOAD, if initiated by a direct command, will perform a CLR so you will lose the variables anyway. Sometimes we wish to transfer as many of the variables as possible from one program to another, hence chain.

Chain differs from the normal load in two respects. First, it saves the data held in the variable and string areas before the load and restores it afterwards. Secondly, it automatically runs the program - obviously it has to be in BASIC.

CHAIN transfers the area of memory holding the variables and arrays to below the string storage. The desired program is loaded and then the data moved back down to the end of the new program. Finally all we have to do is to RUN the program.

Although a fuller, and better, explanation of the way variables are stored is given in Chapter 1, here is a reminder of areas that Chaln cannot deal with. Defined functions are held in the program, only the pointer is in the variable area, and therefore cannot be transferred. The same applies to strings unless they are concatenated or held in arrays.

There are two listings for this command. The first is entered on the command and it will call the main Chain routine. Although the Chain routine works as designed we found that, due to the memory move routines, if there were no variables to move you ended up with a page which could contain anything. The first routine will rectify that after the main routine. This also means that chain could be used as a direct command to load and run disk or tape programs.

ASSEMBLY LISTING 1
7 $=\$ 9283$

10
20

30
40
50
60
70
80 2ERO
90
100
110
120
130
140
150
160 RUN

LDA $\$ 32$ ！END OF ARRAYS
CMP $\$ 2 E$ ！CHECK WITH START OF UARIABLES
BNE ZERO 1
LDA $\$ 31$
CMP $\$ 20$
BNE ZERO＋1 ！NOT THE GAME ADDRESSEG
LDA \＃\＄80
BIT \＄0gA9

\＄80 FOR NONE
JSR $\mathbf{~} 0079$ ！GET LAST CHRGET BYTE
JMF $\$ 9080$ ！PERFOPM CHAIN
LDA \＄BC ！GET FLAG
BPL RLN ！VARIABLES
DEC $\$ 30$ ！DEC ARRAY ADDS BY PAGE
DEC $\$ 32$
JMP \＄A7AE

9203 RUN $92 C 1$ ZERO
ASSEMBLY LISTING 2
$9 *=\$ 9080$
10 JSR \＄E1D4 ！GET LOAD

30
40
50
60
70
80
90
100
110
120
130
140
150
160
170

```
PARAMETERS
20 LDA \＃\＄00 ！ENSURE RELOCATING LOAD
JSR $E1D4 ! GET LOAD
LDA #$00 ! ENSURE RELOCATING LOAD
STA $B9
ISR 韦B526 ! GARBAGE COLLECTION
LDA 事2D ! START OF BLOCK TO MOUE
STA &5F
LDA क2E ! END OF RESIDENT PROG
STA $60
SEC
LDA $31 ! END OF BLOCK
STA $5A
SBC $2F ! CALC AREA OF ARRAYS
STA $FD
LDA $32 ! ALSO END OF ARRAYS
STA 事B
SBC $30
STA $FE
```

180
190
200
210
220
230
240
250
260
270
280
290
300
310
320
330
340
356
360
370
380
390
400
410
420
430 status
440
450
460
470
480 CONT
490
500
518
520
530
540
550
560
576
580
590
600
610

LDA $\$ 33$
SEC ! NEW END OF BLOCK
SBC \#\$01
STA $\$ 58$
LDA $\$ 34$
SBC: \#\$00
STA $\$ 59$
JSR AABBF ! PERFORM MOVE
LDA $\$ 37$ ! SAVE END OF BASIC AREA
STA $\$ 41$
LDA ${ }^{\text {² }} 38$
STA $\$ 42$
lda $\$ 58$ ! save beginning of NEW BLOCK
STA कFB
STA $\$ 37$ ! SET TOF OF BASIC AREA
INC $\$ 59$ ! RECTIFY PAGE
LDA $\$ 55^{\circ}$
STA FFC
STA $\$ 38$
LDX $\$ 2 B$ ! SET LOAD ADDRESS
LDY $\$ 2 \mathrm{C}$
LDA \#\#00 ! SET FOR LOAD
JSR EFFD5 ! KEFNAL LOAD
BCC STATUS ! MAYBE GOOD LOAD
JMP SEOFG ! LOAD ERROR dEPENDING ON A
JSR \$FFB7 ! READ 1/O STATUS WORD
AND \#\$BF
BEQ CONT ! LOAD OK
LDX \#\$1D
JMP $\ddagger A 437$ ! LOAD ERROR
STX $\quad$ 2D ! SET END OF PROGRAM
STY $\$ 2 \mathrm{E}$
STX $\ddagger 5 \mathrm{~F}$ ! SET FOR VARIABLE MOVE
STY $\$ 60$
LDA AFB ! START OF BLOCK TO MOVE
STA 5 SA
LDA FFC
STA ${ }^{\$ 5 B}$
SEC
LDA $\$ 33$ ! END OF BLOCK
SBC \#\$01
TAY
LDA $\$ 34$
SBC \#\$00

| 620 |  | TAX |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 630 |  | TYA |  |  |  |
| 640 |  | SEC |  | $!$ | calc amolnt to move |
| 650 |  | SBC | \$5A |  |  |
| 660 |  | STA | \$58 |  |  |
| 670 |  | TAY |  | ! | NO OF BYTES OF INCOMPLETE PAGE |
| 680 |  | TXA |  |  |  |
| 690 |  | SBC | \$5B |  |  |
| 700 |  | TAX |  | $!$ | NO OF PAGES TO MOUE |
| 710 |  | INX |  | ! | FOR EASIER CHECKING |
| 720 |  | TYA |  |  |  |
| 730 |  | BEQ | Page | $!$ | NO SEPARATE BYTES |
| 740 |  | LDA | \$5A | $!$ | MOVE SEPARATE BYTES FIRST |
| 750 |  | CLC |  |  |  |
| 760 |  | ADC: | \$58 |  |  |
| 770 |  | STA | \$5A |  |  |
| 780 |  | BCC | NOINC. |  |  |
| 790 |  | INC | 事58 |  |  |
| 800 |  | CLC |  |  |  |
| 810 | NOINC | LDA | 害5 |  |  |
| 820 |  | ADC | \$58 |  |  |
| 830 |  | STA | \$5F |  |  |
| 840 |  | BCC | NOINCA |  |  |
| 850 |  | INC | \$60 |  |  |
| 360 | NOINCA | TYA |  |  |  |
| 870 |  | EOR | \# ${ }^{\text {b }}$ FF | ! | 1's COMPLEMENT |
| 880 |  | TAY |  |  |  |
| 890 |  | INY |  | $!$ | 2'5 COMPLEMENT |
| 900 |  | DEC | \$5B |  |  |
| 910 |  | DEC | \$60 |  |  |
| 920 | PAGE | LDA | (\$5A), |  |  |
| 930 |  | STA | (\$5F), $Y$ |  |  |
| 940 |  | INY |  |  |  |
| 950 |  | BNE | PAGE |  |  |
| 960 |  | INC. | \$5B |  |  |
| 970 |  | INC | \$60 |  |  |
| 980 |  | DEX |  | ! | POINTER FOR COMPLETION |
| 990 |  | BNE | PAGE |  |  |
| 1000 |  | SEC |  |  |  |
| 1010 |  | LDA | \$5F |  |  |
| 1020 |  | STA | \$31 | ! | NEW ARRAY END |
| 1030 |  | SBC | \$FD | $!$ | CALC ARRAY START |
| 1040 |  | STA | \$2F |  |  |
| 1050 |  | LDA | \$60 |  |  |


| 1060 | STA \$32 |  |  |
| :---: | :---: | :---: | :---: |
| 1070 | SBC \$FE |  |  |
| 1080 | STA \$30 |  |  |
| 1090 | LDA | ! | RESET END OF BASIC: |
| 1100 | STA \$37 |  |  |
| 1110 | LDA \$42 |  |  |
| 1120 | STA \$38 |  |  |
| 1130 | PLA |  |  |
| 1140 | PLA | $!$ | REMOUE RETURN ADDRESS |
| 1150 | JSR \$A533 | $!$ | RECHAIN LINES |
| 1160 | LDA \#\$00 |  |  |
| 1170 | JSR 刺F90 | ! | TURN OF KERNAL MESSAGES |
| 1180 | JSR \$FFE7 | $!$ | CLALL |
| 1190 | JSR \$A677 | ! | END OF CLR |
| 1200 | JSR \$A68E | ! | BACK UP TEXT POINTER |
| 1210 | JMP \$92CC | ! | BACK TO FIRST ROUTINE |


| $90 E 3$ CONT | 9119 NOINC |
| :--- | :--- |
| 9123 NOINCA | $912 C$ PAGE |
| 9007 STATUS |  |

As CHAIN is just moving memory and loading, it is an amalgamation of routines previously described. Where we come across lines used elsewhere, the description will direct you there. By copying lines rather than using subroutines, we make the routine more transportable.

## LISTING 1

LINES 10-110: Here we find out if there are variables to move by taking the address of the end of program away from the end of arrays address. On the result we set a flag in location $\$ 0 C$ to $\$ 80$. or $\$ 00$ denoting variables. We then jump to the main CHAIN routine, LISTING 2.

LINES 120-160: Having returned from the routine, we check our flag by loading and testing the sign flag. A positive result tells us that variables were transferred and no further adjustments are required. If the result was minus, then the addresses denoting the start and end of arrays are reduced by one page, the high byte of the address less one. The final action of CHAIN is to go to ROM, where the next bASIC line is executed. The main routine does the setting up for this just before it comes back to here.

## LISTING 2

LINES 10-40: We use the ROM routine to get and set up the loading parameters. To ensure that the load has no secondary address, we
unset that location. The garbage collection routine at $\$ 8526$ will tidy up the variable area so that it uses the least space possible.
LINES 50-250: Although the locations from which the addresses are gathered are different, these lines are discussed in Chapter 6, Memory. Moving, lines 1190-1340 (see pages 131-136). There is, however, one extra item involved. To be able to set the start of array address, after loading, we calculate its number of bytes and store it in locations \$FD and $\$ F E$.

LINES 260-360: The data has been moved and now we protect it by changing the pointer to the limit of BASIC. This value will be obtained from the move routine, locations $\$ 58$ and $\$ 59$, after increasing the latter by one. This is because in the move routine the high byte is decreased before checking for completion. Increasing rectifies this. The original end of BASIC pointer is stored for later use.

LINES 370-470: The loading sequence is covered in the mERGE and APPEND routine in Chapter 7 , lines 550-640 (see pages 158-163).

LINES 480-990: Moving the block down is virtually identical to lines $90-550$ of Memory Moving in Chapter 6.

LINES 1000-1210: With the major work done, just the clearing up remains. First we calculate the new start of arrays and set its registers. Then we restore the pointer to the end of basic. Six rom routines are visited to finish off the routine. The first two rechain the lines of the BASIC program, so that the interpreter can follow them, and turn off the KERNAL messages. The call to \$FFE7 closes all open files and sets the input/output channels to their default values. The following subroutine is made halfway into CLR. This will do a RESTORE, reset CONT locations, and amend the stack point. The last two routines will do the auto-run. The former sets the CHRGET address to one byte before the program starts. The last one returns us to the calling routine to finish off.

## INKEY\$ - A waiting GET

COMMAND SYNTAX
i INKEY\$
ii INKEY\$ "'"
iii INKEY\$ A\$ - where A\$ is predefined
iv INKEY\$ "characters"
All commands will stop and wait for a key press. The first two will wait until any key is pressed. The latter two will wait for a key press corresponding to a character within the defined string. The ASCII value of the key press will be placed in the variable st, and will remain there until an input-output is performed on cassette, serial or RS232.

In 64 BASIC there are two commands for receiving a user input from the keyboard: INPUT and GET. The last accepts a key press without a RETURN but will not wait for one. This entails checking the input and GOTOs until the key press you want is received (see Chapter 4 on checking for function keys in BASIC). INPUT waits for a key but you also have to press RETURN, and the cursor is also in operation.

INKEY\$ will sit and wait for a key press, after emptying the keyboard buffer, and, if required will check for a particular key or keys. To allow for further checks we use the reserved variable 'st' to store the input. Using ST is easy in that it has a predefined location in zero page.

ASSEMBLY LISTING

| * $=$ \$904E |  |  |
| :---: | :---: | :---: |
|  | GNE STRING ! | FARAMETERS PRESENT |
| ANYKEY | LDA \#\$00 ! | CLEAR KEY BUFFER |
|  | STA \$ CB |  |
| BYTE | JSR कFFE4 ! | GET CHARACTER |
|  | EEQ BYTE ! | NO KEY |
|  | STA $\$ 90$ | ST LOCATION |
|  | RTS |  |
| STRING | JSR कAD9E ! | GET STRING |
|  | JSR \$B6A3 | DISCARD LNWANTED STRING |
|  | CMP \#\$00 ! | NULL STRING? |
|  | BEQ ANYKEY | NULL STRING |
|  | STA 韦FB ! | NO OF CHARS IN STRING |
|  | LDA \#\$00 ! | EMPTY KEY BUFFER |
|  | STA \$C6 |  |
| BYTEI | JSR \$FFE4 | GET CHARACTER |
|  | BEQ BYTE1 ! | NO KEY |
|  | LDY कFB ! | GET NO. OF CHARS |
|  | DEY |  |
| NEXT | CMP (\$22), Y! | CHECK STRING |
|  | BEQ MATCH ! | FOUND SAME CHAR |
|  | DEY |  |
|  | BPL NEXT ! | CONTINUE SEARCH |
|  | BMI BYTE1 ! | ANOTHER KEY PRESS |
| MATCH | STA $\$ 90$ | ST LOCATION |
|  | RTS |  |

9050 ANYKEY
906 C BYTEI
9074 NEXT

9054 BYTE
907 D MATCH
905 C STRING

LINES 10-70: If there are parameters (cases ii, iii and iv of the command syntax), the zero flag will not be set and these lines are skipped over, at least for the time being. Proceeding on we set the flag for the number of characters in the keyboard buffer to zero. The KERNAL routine at \$FFE4 will return the ASCII value of key presses in the order they were placed in the buffer. If none, then the accumulator will hold zero, so we continue to call the routine until a value is returned. That value is placed in the location which the reserved variable st uses, and we return to continue the BASIC program.

LINES 80-240: The call to the ROM routine does our string work. It finds the string, especially if it is a variable, determining its length and giving syntax errors if a non-string parameter was supplied. On returning from the routine, the number of characters will be in A and the start address in locations $\$ 22$ and $\$ 23$. If there were no characters in the string, we branch back to the previous section and wait for any key.

After clearing the buffer and getting a key press value we can check it against the string. The $y$ register will be loaded with the number of characters and decreased as we check the whole string. If the complete string is checked and no match is found, then the next key press is evaluated. Once a match is found, it is stored in $5 T$ and we return to carry on with your program.

## LOMEM and HIMEM - Setting the area of work

COMMAND SYNTAX

## LOMEM address

HIMEM address
The address range that is permissible with these commands is between 1024 and 32767. 'ILLEGAL QUANTITY' errors are given outside this range. The actual start of a program will be one greater than the address given in lomem. Commands can be used in direct or program mode.

Changing the memory configuration is a useful, and indeed necessary, task. By raising the base of a program, you can store items such as sprite data, hires screens or even two normal screens and it will not be affected by a program.

At the other end you may wish to put a machine code routine and so to protect it at the top of memory from being overwritten by the variables, so you can set the limit of BASIC to below your routine.
lOMEM will set the lower and HIMEM the upper limit of basic. So that they could be used in a loader program the routine does not clear that program. Subsequent programs will be loaded to the new lomem address. The ideal place for these commands is at the beginning of a program before any variables are defined. Variables defined after these
commands will be placed in the new area．You can use Chain to load the next program if there are variables you wish to transfer．

## ASSEMBLY LISTING

        \(9 *=\$ 9169\)
    20
30 GATHER
40
50
60
70 ERROR
80
96 TOP
100
110
120
130
140
150
160
170
180
190
200
210
220
230
240
250
260
270
280
290
300
310
320
330
340
350
360
370
380
390
400
10 INPUT BNE GATHER ! PARAMETERS

BNE GATHER ！PARAMETERS
MP \＄AF08 ！SYNTAX ERROR
JSR \＄81F5 ！GET PARAMETERS
LDA $\$ 15$
CHP \＃\＄04 ：CHECK LOW LIMIT
BCS TOP ！O．K．
LDX 祘象QE ！ILLEGAL QUANTITY
JMP \＄A437 ！ERROR ROUTINE
CMP \＃\＄80 ！UPPER LIMIT
BCS ERROR ！FAILED
RTS
！START OF HIMEM
JSR INPUT
STA $\$ 38$ ！SET TOP POINTER
LDA $\$ 14$
STA $\$ 37$
JMP \＄A65E ！CLR AND RETURN
LOMEM
JSR INPUT
LDY \＃ 00
TYA
STA（\＄14），Y！CLEAR FIRST 3 BYTES
INY
STA（\＄14），Y
INY
STA（\＄14），Y
LDA \＄14
CLC
ADC \＃ 01
STA $2 B$ ！SET START OF BASIC
TAX
LDA $\$ 15$
ADC ${ }^{2}$ 束 80
STA \＄2C
TAY
TXA
ADC
STA $\$ 20$ ！SET START OF UARIABLES
TYA
ADC 制 00

| 410 | STA \$2E |
| :--- | :--- |
| 420 | JMP \$A663 ! CLR AND RETURN |
|  |  |
| 9177 ERROR |  |
| 9169 INPUT | $916 E$ GATHER |

LINES 10-110: INPUT: This subroutine is used by both commands. It deals with the gathering and checking of addresses. First we check that there is an address. No address, then no command, and a SYNTAX ERROR is given. When the address is picked up, it is first checked for the lower limit and then for the higher.

LINES 130-170: HIMEM: After visiting the input routine, we place the address in the pointers to the limit of BASIC. We then jump to the CLR routine to finish off: this will set all the remaining relevant pointers (such as the string pointer).

LINES 190-420: LOMEM: BASIC requires that the first byte of the BASIC program area is zero (normally $2048, \$ 0800$ ) and that two zeros signify the end of the program. In the new area these will be together, as there is no program, so we set those first from the address given. To set the start of the program we increase it by one, and from that we add a further two for the address to the start of the variables, or end of program if you prefer. Calling the CLR routine will set the end of variables and array pointers.

## QUIT

COMMAND SYNTAX

## QUIT

There are no arguments with this particular command.
QUIT disables the UTILITY and its commands, leaving you with the standard BASIC. It does not, however, reset the top of memory back to its original (\$A000). This will leave the UTILITY intact which can be reinitiated by SYS 32768 .

QUIT simply restores all the vectors and pointers we changed on start up to their standard values.

ASSEMBLY LISTING

```
9*=$91B7
```

50
60 70
30
90
100
116
120
130
140
150
160
170
180
190
200
210
220
230
240
250
260
270
280
290
300
310

LDA \# ${ }^{\text {事 } 1 A}$
STA 00306 ! BASIC TEXT LIST
LDA \#\#E4
STA $\$ 0308$ ! BASIC CHAR DISPATCH
LDA 脚A7
STA $\$ 0307$
STA $\$ 0309$
LDA \#\$86
STA \$030A ! BASIC TOKEN EVALUATION
LDA \#\$AE
STA $\$ 030 \mathrm{~B}$
LDA \#\#FE
STA $\$ 0317$ : BRK INTERRUPT
STA 0319 ! MMI INTERRUPT
LDA \#\$66
STA $\$ 0316$
LDA \#\$47
STA $\$ 0318$
SEI
LDA \#\$48
STA $\$ 628 \mathrm{~F}$ ! KEYBOARD TABLE SETUP
LDA \#\$EE
STA $\$ 0290$
CLI
PLA
PLA
JMP \$A474 ! READY FOR BASIC

## 9 The complete utility

## Introduction

We are going to supply the complete UTIIITY in the form of a Supermon listing．If you do not possess a monitor，you can find Supermon in the appendices．For the area $\$ 80 \mathrm{DE}$ to $\$ 81 \mathrm{F4}$ ，keywords and vectors，use the m function of the monitor．You may also find it easier to use the memory dump in Chapter 6 for that area．Save to tape or disk regularly as you go．
We had thought of also giving the UTLITY in DATA statement form．This would have come to about 690 lines，of seven items of data on each， which would have been a mammoth task of programming for anyone and very prone to error．

| 8088 | 20 |  | 80 | JSR | \＄809F | 803C |  | 38 |  | STA | \＄38 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8083 | 20 | 54 | 80 | JSR | \＄8054 | 803E | 85 | 34 |  | STA | \＄34 |
| 8066 | 20 | 41 | 80 | JSR | \＄8041 | 8048 | 60 |  |  | RTS |  |
| 8009 | 20 | 34 | 80 | JSR | \＄8034 | 8841 | A9 | $7 E$ |  | LDA |  |
| 808C | 4C | 08 | 92 | JMP | \＄9206 | 8843 | 80 | 16 | 03 | STA | ＋+316 |
| 808F | A9 | 09 |  | LDA | W09 | 8846 | A9 | 61 |  | LDA | 解61 |
| 8011 | 80 | 04 | 03 | STA | \＄0304 | 8048 | 8D | 18 | 03 | STA | ＋8318 |
| 8014 | A9 | BC |  | LDA |  | 8048 | A9 | 88 |  | LDA | 鳃880 |
| 8816 | 8 D | 06 | 03 | STA | ＋0386 | 8040 | 80 | 17 | 83 | STA | ＋ 6317 |
| 8019 | A9 | 02 |  | LDA | 㐌真2 | 8059 | 8D | 19 | 03 | STA | ¢ 0319 |
| 8018 | 80 | 08 | 03 | STA | ＋0308 | 8053 | 60 |  |  | RTS |  |
| 801E | A9 | 29 |  | LDA | 镇29 | 8054 | 78 |  |  | SE］ |  |
| 8020 | 8 D | BA | 03 | STA | ＋036A | 8855 | A9 | 22 |  | LDA | W ${ }^{\text {W }}$ 22 |
| 8023 | A9 | 82 |  | LDA | 竾82 | 8057 | 80 | 8F | 02 | STA | \＄${ }^{\text {S }} 28 \mathrm{~F}$ |
| 8025 | 80 | 05 | 03 | STA | ＋6305 | 805A | A9 | 87 |  | LDA | 能87 |
| 8028 | 80 | 07 | 03 | STA | ＋0307 | 805C | 80 | 90 | 82 | STA | \＄ 8290 |
| 8028 | A9 | 83 |  | LDA | 稱83 | 885F | 58 |  |  | CLI |  |
| 8020 | 80 | 89 | 03 | STA | ＋0309 | 8060 | 60 |  |  | RTS |  |
| 8030 | 8D | 88 | 03 | STA | ＋8308 | 8061 | 48 |  |  | PHA |  |
| 8033 | 60 |  |  | RTS |  | 8062 | 8A |  |  | TXA |  |
| 8034 | A9 | FF |  | LDA | 鳃FF | 8063 | 48 |  |  | PHA |  |
| 8036 | 85 | 37 |  | STA | ＋37 | 8864 | 98 |  |  | TYA |  |
| 8038 | 85 | 33 |  | STA | ＋33 | 8065 | 48 |  |  | PHA |  |
| 803A | A9 | 7F |  | LDA |  | 8866 | A9 | 7F |  | LDA | ＊ 7 F |


| 8068 | 80 |  | OD | STA | \$0DOD | 80 c 7 |  |  |  | ??? |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8068 | AC | 60 | DO | LDY | \$DDOD | 80 C8 | 8 C | 91 | 80 | STY | \$8091 |
| 8065 | 10 | 03 |  | BPL | \$8073 | 80 CB | 91 | 4D |  | STA | ( S $^{\text {d }}$ ) , $Y$ |
| 8076 | 4C | 72 | FE | JMP | \$FE72 | 80CD | 90 | FB |  | BCC | \$80CA |
| 8073 | 20 | BC | F6 | JSR | \$F6BC | 80CF | 85 | 6E |  | STA | 事6E |
| 8076 | 28 | E1 | FF | JSR | \$FFE1 | 8001 | 88 |  |  | DEY |  |
| 8079 | FO | 03 |  | BEC | \$807E | 8002 | 60 |  |  | RTS |  |
| 8078 | 4C | 72 | FE | JMP | \$FE72 | 8003 | 8D | FF | FF | STA | ¢FFFF |
| 307E | 20 | 15 | FD | JSR | \$FD15 | 8006 | FF |  |  | ??? |  |
| 8081 | 20 | A3 | FD | JSR | \$FDA3 | 8007 | FF |  |  | ??? |  |
| 8084 | 20 | 18 | E5 | JSR | \$E518 | 8008 | 00 |  |  | BRK |  |
| 8087 | 20 | 54 | 80 | JSR | \$8054 | 8009 | FF |  |  | ??? |  |
| 808A | 20 | 41 | 80 | JSR | \$8041 | 8004 | FF |  |  | ??? |  |
| 8080 | 6 C | 02 | A0 | JMP | ( $\ddagger 4002$ ) | 800B | FF |  |  | ??? |  |
| 8090 | 98 |  |  | TYA |  | 8000 | F6 | FF |  | INC | क FF, $^{\text {P }} \mathrm{X}$ |
| 8091 | 87 |  |  | ??? |  | 800E | 86 | F7 |  | LDX | \$F7, Y |
| 8892 | 4C | 86 | B2 | MP | \$8286 | 80 EO | 00 |  |  | BRK |  |
| 8095 | 83 |  |  | ??? |  | 80E1 | 60 |  |  | RTS |  |
| 8096 | 9F |  |  | ??? |  | 80E2 | 00 |  |  | BRK |  |
| 8097 | 84 | EB |  | STY | \$ ${ }^{\text {E }}$ | 80.3 | 00 |  |  | BRK |  |
| 8099 | 84 | 36 |  | STY | \$36 | 80E4 | D6 | 83 |  | DEC | \$83, C |
| 809 B | 85 | BE |  | STA | \$BE | 80E6 | D6 | 83 |  | DEC | \$83, X |
| 8090 | 85 | 14 |  | STA | \$14 | 80E8 | 80 |  |  | BRK |  |
| 809F | 84 | 51 |  | STY | \$51 | 80E9 | 00 |  |  | BRK |  |
| 80A1 | 83 |  |  | ??? |  | 80EA | 00 |  |  | BRK |  |
| 80A2 | A6 | 83 |  | LDX | \$83 | 80EB | 68 |  |  | PLA |  |
| 80A4 | AE | 8F | B4 | LDX | कB48F | 80EC | 00 |  |  | BRK |  |
| 80A7 | 8 F |  |  | ??? |  | 80ED | 00 |  |  | BRK |  |
| 8048 | 80 |  |  | ??? |  | 80EE | 00 |  |  | BRK |  |
| 80A9 | 83 |  |  | ??? |  | 80EF | 40 |  |  | RTI |  |
| 80AA | AC | 83 | 51 | LDY | \$5183 | 86 FO | 00 |  |  | BRK |  |
| 80AD | 8 BE | C4 | 89 | STX | \$89C4 | 80F1 | 00 |  |  | BRK |  |
| 80B0 | 43 |  |  | ??? |  | 80F2 | 80 |  |  | BRK |  |
| 80B1 | 8F |  |  | ??? |  | $80 \mathrm{F3}$ | 40 |  |  | RTI |  |
| $80 \mathrm{B2}$ | A6 | 87 |  | LDX | \$87 | 80 F 4 | 00 |  |  | BRK |  |
| $80 \mathrm{B4}$ | 92 |  |  | ??? |  | $80 \mathrm{F5}$ | 40 |  |  | RTI |  |
| 8085 | 8 B |  |  | ??? |  | 80F6 | 4F |  |  | ??? |  |
| 8086 | 20 | 84 | D1 | AND | \$0184 | 80F7 | 46 | C6 |  | LSR | \$C6 |
| 8089 | 8F |  |  | ??? |  | 80F9 | 4B |  |  | ??? |  |
| 80BA | 3A |  |  | ??? |  | 80FA | 45 | D9 |  | EOR | \$09 |
| 80BB | A9 | D1 |  | LDA | \#* 01 | 80FC | 44 |  |  | ??? |  |
| 80BD | A8 |  |  | TAY |  | 80FD | 4F |  |  | ??? |  |
| 80 BE | 30 | 86 |  | BMI | \$8046 | 80FE | 48 |  |  | ??? |  |
| 80C0 | B6 | 91 |  | LDX | \$91,Y | 80FF | C5 | 54 |  | CMP | \$54 |
| 80 C 2 | 39 | 80 | 10 | AND | \$1080, $Y$ | 8101 | 45 | CE |  | EOR | ¢ ${ }^{\text {ce }}$ |
| $80 \mathrm{C5}$ | 86 | B5 |  | STX | - ${ }^{\text {B5 }}$ | 8103 | 54 |  |  | ??? |  |


| 8104 | 57 |  | ??? |  | 8149 | 54 |  | ??? |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8105 | CF |  | ??? |  | 8144 | CF |  | ??? |  |
| 8186 | 48 |  | PHA |  | 814 B | 50 | 52 | BUC | \$819F |
| 8107 | 45 | D8 | EOR | \$D8 | 8140 | 4F |  | ??? |  |
| 8169 | 42 |  | ??? |  | 814 E | C3 |  | ??? |  |
| 810A | 49 | CE | E0R | \#\$CE | 814 F | 44 |  | ??? |  |
| 810 C | 4F |  | ??? |  | 8150 | 50 | 52 | BUC | \$81A4 |
| 810D | 4C | C4 43 | JMP | \$43C4 | 8152 | 4F |  | ??? |  |
| 8110 | 4F |  | ??? |  | 8153 | C3 |  | ??? |  |
| 8111 | 4C | 4F 55 | JMP | \$554F | 8154 | 45 | 50 | EOR | \$50 |
| 8114 | D2 |  | ??? |  | 8156 | 52 |  | ??? |  |
| 8115 | 57 |  | ??? |  | 8157 | 4F |  | ??? |  |
| 8116 | 52 |  | ??? |  | 8158 | C. 3 |  | ??? |  |
| 8117 | 49 | 54 | EOR | \# 54 | 8159 | 50 | 4F | BuC | \$81A |
| 8119 | C5 | 43 | CMP | 事43 | 815 B | D0 | 51 | BNE | \$81AE |
| 811 B | 47 |  | ??? |  | 815D | 55 | 49 | EOR | \$49, X |
| 811 C | 4F |  | ??? |  | 815 F | D4 |  | ??? |  |
| 8110 | 54 |  | ??? |  | 8160 | 54 |  | ??? |  |
| 811 E | CF |  | ??? |  | 8161 | 52 |  | ??? |  |
| 811 F | 43 |  | ??? |  | 8162 | 41 | 43 | EOR | ( $\ddagger 43, \mathrm{X}$ ) |
| 8120 | 47 |  | ??? |  | 8164 | C5 | 52 | CMP | \$ 52 |
| 8121 | 4F |  | ??? |  | 8166 | 45 | 53 | EOR | \$53 |
| 8122 | 53 |  | ??? |  | 8168 | 45 | D4 | EOR | 㭏4 |
| 8123 | 55 | C2 | EOR | \$ $\mathrm{C} 2, \mathrm{X}$ | 8164 | 43 |  | ??? |  |
| 8125 | 50 | 4C | BUC | \$8173 | 816 B | 48 |  | PHA |  |
| 8127 | 4F |  | ??? |  | 816 C | 41 | 49 | EOR | ( 4 49, X) |
| 8128 | D4 |  | ??? |  | 816 E | CE | 4C. 4F | DEC | \$4F4C |
| 8129 | 45 | $4 E$ | EOR | \$4E | 8171 | 4D | 45 CD | EOR | \$CD45 |
| 812 B | 54 |  | ??? |  | 8174 | 48 |  | PHA |  |
| 812 C | 45 | D2 | EOR | \$D2 | 8175 | 49 | 40 | EOR | \# ${ }^{\text {a }}$ |
| 812 E | 44 |  | ??? |  | 8177 | 45 | CD | EOR | \$ CD |
| 812 F | 55 | 4D | EOR | \$40, $X$ | 8179 | 49 | 4E | EOR | \#\$ $4 E$ |
| 8131 | D0 | 52 | BNE | \$8185 | 817 B | 48 |  | ??? |  |
| 8133 | 45 | 4E | EOR | \$4E | 817C | 45 | 59 | EOR | \$59 |
| 8135 | 55 | CD | EOR | \$CD, $X$ | 817 E | A4 | 40 | LDY | \$4D |
| 8137 | 44 |  | ??? |  | 8180 | 45 | CD | EOR | \$CD |
| 8138 | 45 | 4C | EOR | \$4C | 8182 | 41 | 50 | EOR | ( $550, \mathrm{~S}$ ) |
| 8134 | 45 | 54 | EOR | \$54 | 8184 | 50 | 45 | BUC | \$81CB |
| 813 C | C5 | 4D | CMP | \$4D | 8186 | 4E | C4 54 | LSR | \$54C4 |
| 813 E | 45 | 52 | EOR | \$52 | 8189 | 52 |  | ??? |  |
| 8140 | 47 |  | ??? |  | 818A | 4F |  | ??? |  |
| 8141 | C5 | 43 | CMP | \$43 | 8188 | 46 | C6 | LSR | \$ 66 |
| 8143 | 4F |  | ??? |  | 818D | 5A |  | ??? |  |
| 8144 | 44 |  | ??? |  | 818 E | 5A |  | ??? |  |
| 8145 | 45 | D2 | EOR | \$02 | 818F | 54 |  | ??? |  |
| 3147 | 41 | 55 | EOR | ( $\$ 55, \mathrm{X}$ ) | 8190 | 5 A |  | ??? |  |



| 81F1 | 00 |  | BRK |  | 8245 | BD | 80 | 82 | LDA | \$0200, X |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 81F2 | 08 |  | BRK |  | 8248 | 38 |  |  | SEC |  |
| 81F3 | 00 |  | BRK |  | 8249 | F9 | 9E | A | SBC | \$A09E, Y |
| 81F4 | 00 |  | BRK |  | 824C | F0 | F5 |  | BEQ | \$8243 |
| 81F5 | 20 | 8A AD | JSR | \$AD8A | 824E | C9 | 80 |  | CMP | \#\$80 |
| 81F8 | 4C | F7 B7 | JMP | \$B7F7 | 8250 | D0 | 30 |  | BNE | \$8282 |
| 81FB | A5 | 01 | LDA | \$01 | 8252 | 05 | 0B |  | ORA | \$0B |
| 81FD | 29 | FE | AND | \#\$FE | 8254 | A4 | 71 |  | LDY | \$71 |
| 81 FF | 85 | 01 | STA | \$01 | 8256 | E8 |  |  | INX |  |
| 8201 | 60 |  | RTS |  | 8257 | C8 |  |  | INY |  |
| 8202 | A5 | 01 | LDA | \$01 | 8258 | 99 | FB | 01 | STA | \$01FB, Y |
| 8204 | 09 | 81 | ORA | \#\$01 | 825B | B9 | FB | 01 | LDA | \$01FB, Y |
| 8206 | 85 | 81 | STA | \$01 | 825E | F0 | 59 |  | BEQ | \$82B9 |
| 8208 | 60 |  | RTS |  | 8260 | 38 |  |  | SEC |  |
| 8289 | A6 | 7A | LDX | \$7A | 8261 | E9 | 3A |  | SBC | \#\$3A |
| 820B | A0 | 84 | LDY | \#\$04 | 8263 | F0 | 84 |  | BEQ | \$8269 |
| 820D | 84 | 0F | STY | \$8F | 8265 | C9 | 49 |  | CMP | \#\$49 |
| 820F | BD | $08 \quad 02$ | LDA | \$0200, X | 8267 | D0 | 82 |  | BNE | \$826B |
| 8212 | 10 | 87 | BPL | \$821B | 8269 | 85 | 8F |  | STA | \$8F |
| 8214 | C9 | FF | CMP | \#\$FF | 826B | 38 |  |  | SEC |  |
| 8216 | F0 | 3E | BEQ | \$8256 | 826C | E9 | 55 |  | SBC | \#\$55 |
| 8218 | E8 |  | INX |  | 826E | D0 | 9F |  | BNE | \$820F |
| 8219 | D0 | F4 | BNE | \$820F | 8270 | 85 | 88 |  | STA | \$08 |
| 821 B | C9 | 20 | CMP | \#\$ 20 | 8272 | BD | 08 | 02 | LDA | \$0200, X |
| 821D | Fl | 37 | BEQ | \$8256 | 8275 | F0 | DF |  | BEQ | \$8256 |
| 821F | 85 | 08 | STA | \$08 | 8277 | C5 | 08 |  | CMP | \$08 |
| 8221 | C9 | 22 | CMP | \#\$22 | 8279 | F8 | DB |  | BEQ | \$8256 |
| 8223 | F0 | 56 | BEQ | \$827B | 827B | C8 |  |  | INY |  |
| 8225 | 24 | 8F | BIT | \$0F | 827C | 99 | FB | 01 | STA | \$01FB, Y |
| 8227 | 70 | 2D | BUS | \$8256 | 827F | E8 |  |  | INX |  |
| 8229 | C9 | $3 F$ | CMP | \#\$3F | 8288 | D0 | F0 |  | BNE | \$8272 |
| 822B | D8 | 84 | BNE | \$8231 | 8282 | A6 | 7A |  | LDX | \$7A |
| 8220 | A9 | 99 | LDA | \#\$99 | 8284 | E6 | 8B |  | INC | \$0B |
| 822F | D0 | 25 | BNE | \$8256 | 8286 | C8 |  |  | INY |  |
| 8231 | C9 | 30 | CMP | \#\$30 | 8287 | B9 | 90 | A0 | LDA | \$A09D, Y |
| 8233 | 90 | 04 | BCC | \$8239 | 828A | 10 | FA |  | BPL | \$8286 |
| 8235 | C9 | 3C | CMP | \#3C | 828C | B9 | 9E | A0 | LDA | \$A09E, Y |
| 8237 | 90 | 1D | BCC | \$8256 | 828F | D0 | B4 |  | BNE | \$8245 |
| 8239 | 84 | 71 | STY | \$71 | 8291 | A0 | FF |  | LDY | \#\#FF |
| 823B | A0 | 80 | LDY | \#\$00 | 8293 | CA |  |  | DEX |  |
| 8230 | 84 | 8B | STY | \$0B | 8294 | C8 |  |  | INY |  |
| 823F | 88 |  | DEY |  | 8295 | E8 |  |  | INX |  |
| 8240 | 86 | $7 A$ | STX | \$7A | 8296 | BD | 88 | 82 | LDA | \$0200, X |
| 8242 | CA |  | DEX |  | 8299 | 38 |  |  | SEC |  |
| 8243 | C8 |  | INY |  | 829A | F9 | F6 | 80 | SBC | \$80F6, Y |
| 8244 | E8 |  | INX |  | 8290 | F0 | F5 |  | BEQ | \$8294 |


| 829F | C9 | 80 | CMP | \# ${ }^{\text {\% }} 80$ | 182F7 | 20 | 47 | $A B$ | JSR | \$AB47 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 82A1 | D0 | 82 | BNE | \$82A5 | 82FA | D0 | F6 |  | BNE | \$82F2 |
| 82A3 | F0 | AD | BEQ | \$8252 | 82FC | 4C | F3 | A6 | JMP | \$A6F3 |
| 82A5 | A6 | 7A | LDX | \$7A | 82FF | 4C | EF | A6 | JMP | \$A6EF |
| 82A7 | E6 | 8B | INC | \$0B | 8302 | 20 | 73 | 80 | JSR | \$0073 |
| 82A9 | C8 |  | INY |  | 8305 | C9 | CC |  | CMP | \#\$CC |
| 82AA | B9 | F5 80 | LDA | \$88F5, Y | 8307 | 98 | $1 A$ |  | BCC | \$8323 |
| 82AD | 10 | FA | BPL | \$82A9 | 8309 | C9 | EE |  | CMP | \#\$EE |
| 82AF | B9 | F6 80 | LDA | \$88F6, Y | 830 B | B0 | 16 |  | BCS | \$8323 |
| 82B2 | D0 | E2 | BNE | \$8296 | 830D | 20 | 13 | 83 | JSR | \$8313 |
| 82B4 | BD | 8882 | LDA | \$0200, X | 8310 | 4C | EA | A7 | JMP | \$A7EA |
| $82 \mathrm{B7}$ | 10 | 98 | BPL | \$8254 | 8313 | 38 |  |  | SEC |  |
| 82B9 | 4C | 89 A6 | JMP | \$A609 | 8314 | E9 | CC |  | SBC | \#\$CC |
| 82BC | 10 | 3E | BPL | \$82FC | 8316 | 0A |  |  | ASL |  |
| 82BE | C9 | FF | CMP | \#\$FF | 8317 | A8 |  |  | TAY |  |
| 82C0 | F0 | 3A | BEQ | \$82FC | 8318 | B9 | 91 | 80 | LDA | \$8091, Y |
| 82C2 | 24 | 6F | BIT | \$0F | 831 B | 48 |  |  | PHA |  |
| 82C4 | 30 | 36 | BMI | \$82FC | 831C | B9 | 98 | 80 | LDA | \$8090, Y |
| 82C6 | C9 | CC | CMP | \#\$CC | 831F | 48 |  |  | PHA |  |
| 82C8 | 90 | 日E | BCC | \$82D8 | 8320 | 4C | 73 | 00 | JMP | \$0073 |
| 82CA | 38 |  | SEC |  | 8323 | 20 | 79 | 80 | JSR | \$0879 |
| 82CB | E9 | CB | SBC | \#\$ ${ }^{\text {CB }}$ | 8326 | 4C | E7 | A7 | JMP | \$A7E7 |
| 82CD | AA |  | TAX |  | 8329 | A9 | 08 |  | LDA | \#\$08 |
| 82CE | A9 | F6 | LDA | \# F $^{\text {F } 6}$ | 832B | 85 | 8D |  | STA | \$00 |
| 8200 | 85 | 22 | STA | \$22 | 832 D | 20 | 73 | 08 | JSR | \$0073 |
| 8202 | A9 | 80 | LDA | \#\$80 | 8330 | C9 | F7 |  | CMP | \#\$F7 |
| 8204 | 85 | 23 | STA | \$23 | 8332 | 90 | 18 |  | BCC | \$834C |
| 82D6 | D0 | 0C | BNE | \$82E4 | 8334 | C9 | F8 |  | CMP | \#\$F8 |
| 8208 | 38 |  | SEC |  | 8336 | B0 | 14 |  | BCS | \$834C |
| 82 D 9 | E9 | 7 F | SBC | \#\$ 7 F | 8338 | 20 | 3C | 83 | JSR | \$833C |
| 82DB | AA |  | TAX |  | 833B | 60 |  |  | RTS |  |
| 820C | A9 | 9E | LDA | \#\$9E | 833C | 38 |  |  | SEC |  |
| 820E | 85 | 22 | STA | \$22 | 833D | E9 | F6 |  | SBC | \#\$F6 |
| 82E0 | A9 | A8 | LDA | \#\$A0 | 833F | 8A |  |  | ASL |  |
| 82E2 | 85 | 23 | STA | \$23 | 8346 | A8 |  |  | TAY |  |
| 82E4 | 84 | 49 | STY | \$49 | 8341 | B9 | E5 | 80 | LDA | \$80E5, Y |
| 82E6 | A0 | FF | LDY | \#\$FF | 8344 | 48 |  |  | PHA |  |
| 82E8 | CA |  | DEX |  | 8345 | B9 | E4 | 80 | LDA | \$80E4, Y |
| 82E9 | F0 | 07 | BEQ | \$82F2 | 8348 | 48 |  |  | PHA |  |
| 82EB | C8 |  | INY |  | 8349 | 4C | 73 | 00 | JMP | \$0073 |
| 82EC | B1 | 22 | LDA | (\$22), Y | 834C | 20 | 79 | 00 | JSR | \$0079 |
| 82EE | 10 | FB | BPL | \$82EB | 834F | 4C | 8D | AE | JMP | \$AE8D |
| 82F0 | 30 | F6 | BMI | \$82E8 | 8352 | 20 | F5 | 81 | JSR | \$81F5 |
| 82F2 | C8 |  | INY |  | 8355 | A5 | 14 |  | LDA | \$14 |
| 82F3 | B1 | 22 | LDA | (\$22), $Y$ | 8357 | 29 | OF |  | AND | \#\$8F |
| 82F5 | 30 | 08 | BMI | \$82FF | 8359 | 8 D | 21 | D0 | STA | \$D021 |


| 835C | 20 | 79 | 00 | JSR | \$0079 | 83C9 | A6 | 64 |  | LDX | \$64 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 835F | F0 | $1 F$ |  | BEQ | \$8380 | 83CB | A8 | 88 |  | LDY | \#\$00 |
| 8361 | 20 | FD | AE | JSR | \$AEFD | 83CD | 91 | 14 |  | STA | (\$14), Y |
| 8364 | 20 | F5 | 81 | JSR | \$81F5 | 83CF | C8 |  |  | INY |  |
| 8367 | A5 | 14 |  | LDA | \$14 | 83D0 | 8A |  |  | TXA |  |
| 8369 | 29 | 0F |  | AND | \#\$0F | 83D1 | 91 | 14 |  | STA | (\$14), Y |
| 836B | 8D | 20 | D0 | STA | \$D020 | 8303 | 60 |  |  | RTS |  |
| 836E | 20 | 79 | 80 | JSR | \$0079 | 83D4 | 4C | 48 | B2 | JMP | \$8248 |
| 8371 | F0 | 6D |  | BEQ | \$8380 | 8307 | A5 | 15 |  | LDA | \$15 |
| 8373 | 20 | FD | AE | JSR | \$AEFD | 83D9 | 48 |  |  | PHA |  |
| 8376 | 20 | F5 | 81 | JSR | \$81F5 | 83DA | A5 | 14 |  | LDA | \$14 |
| 8379 | A5 | 14 |  | LDA | \$14 | 83DC | 48 |  |  | PHA |  |
| 837B | 29 | 8F |  | AND | \#\$0F | 830D | 20 | FA | AE | JSR | \$AEFA |
| 837D | 8D | 86 | 82 | STA | \$0286 | 83E0 | 20 | F5 | 81 | JSR | \$81F5 |
| 8380 | 60 |  |  | RTS |  | 83 E 3 | 20 | F7 | $A E$ | JSR | \$AEF7 |
| 8381 | 20 | FA | AE | JSR | \$AEFA | 83E6 | A0 | 01 |  | LDY | \#\$01 |
| 8384 | 20 | F5 | 81 | JSR | \$81F5 | $83 E 8$ | B1 | 14 |  | LDA | (\$14), $Y$ |
| 8387 | A5 | 14 |  | LDA | \$14 | 83EA | AA |  |  | TAX |  |
| 8389 | C9 | 28 |  | CMP | \#\$28 | 83EB | 88 |  |  | DEY |  |
| 838B | 90 | 03 |  | BCC | \$8390 | 83EC | B1 | 14 |  | LDA | (\$14), Y |
| 838D | 20 | 48 | B2 | JSR | \$8248 | 83EE | A8 |  |  | TAY |  |
| 8390 | 48 |  |  | PHA |  | 83EF | 68 |  |  | PLA |  |
| 8391 | 20 | FD | AE | JSR | \$AEFD | 83F0 | 85 | 14 |  | STA | \$14 |
| 8394 | 20 | F5 | 81 | JSR | \$81F5 | $83 F 2$ | 68 |  |  | PLA |  |
| 8397 | A6 | 14 |  | LDX | \$14 | $83 F 3$ | 85 | 15 |  | STA | \$15 |
| 8399 | E0 | 19 |  | CPX | \#\$19 | 83F5 | 8 A |  |  | TXA |  |
| 839B | B0 | F0 |  | BCS | \$838D | 83F6 | 20 | 91 | B3 | JSR | \$8391 |
| 839D | 68 |  |  | PLA |  | 83F9 | 20 | 81 | 84 | JSR | \$8401 |
| 839E | A8 |  |  | TAY |  | 83FC | 68 |  |  | PLA |  |
| 839F | 18 |  |  | CLC |  | 83FD | 68 |  |  | PLA |  |
| 83A8 | 20 | F0 | FF | JSR | \$FFF0 | 83FE | 4C | 8D | AD | JMP | \$AD8D |
| 83A3 | 20 | 73 | 00 | JSR | \$0873 | 8401 | A5 | 66 |  | LDA | \$66 |
| 83A6 | 60 |  |  | RTS |  | 8483 | 10 | 0A |  | BPL | \$840F |
| 83A7 | 20 | 81 | 83 | JSR | \$8381 | 8405 | A0 | 84 |  | LDY | \#\$84 |
| 83AA | 4C | A0 | AA | JMP | \$AAA0 | 8487 | A9 | 10 |  | LDA | \# 10 |
| 83AD | 20 | 81 | 83 | JSR | \$8381 | 8409 | 20 | 8C | BA | JSR | \$BABC |
| 83B6 | 4C | BF | AB | JMP | \$ABBF | 848C | 20 | 64 | B8 | JSR | \$B86A |
| 83B3 | 20 | F5 | 81 | JSR | \$81F5 | 840F | 60 |  |  | RTS |  |
| 83B6 | 20 | FD | AE | JSR | \$AEFD | 8410 | 91 | 00 |  | STA | (\$00), Y |
| 8389 | 20 | 8A | $A D$ | JSR | \$AD8A | 8412 | 08 |  |  | BRK |  |
| 83BC | A5 | 66 |  | LDA | \$66 | 8413 | 88 |  |  | BRK |  |
| 83BE | 30 | 14 |  | BMI | \$8304 | 8414 | 08 |  |  | BRK |  |
| 83C0 | C9 | 91 |  | CMP | \#\$91 | 8415 | A9 | FF |  | LDA | \#\$FF |
| 83C2 | B0 | 10 |  | BCS | \$83D4 | 8417 | A0 | 01 |  | LDY | \#\$01 |
| 83C4 | 20 | 98 | BC | JSR | \$BC9B | 8419 | 91 | 2B |  | STA | (\$2B), $Y$ |
| $83 C 7$ | A5 | 65 |  | LDA | \$65 | 841B | 20 | 33 |  | JSR | \$A533 |


| 841 E | A5 | 22 | LDA | \$22 | 8485 |  |  |  | SEC |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8420 | 18 |  | CLC |  | 8486 | 20 | 49 | BC | JSR | \$ ${ }^{\text {PC49 }}$ |
| 8421 | 69 | 02 | ADC | \# ${ }^{\text {\$ }} 02$ | 8489 | 20 | DF | BD | JSR | \$BDDF |
| 8423 | 85 | 20 | STA | \$2D | 848C | 20 | 87 | B4 | JSR | \$B487 |
| 8425 | A5 | 23 | LDA | \$23 | 848F | 20 | A6 | B6 | JSR | \$B6A6 |
| 8427 | 69 | 00 | ADC | \#\$00 | 8492 | A2 | 00 |  | LDX | \#\$00 |
| 8429 | 85 | 2 E | STA | \$2E | 8494 | BD | 80 | 01 | LDA | \$0180, x |
| 842B | 4C | 60 A6 | JMP | \$A660 | 8497 | 90 | 08 | 82 | STA | \$0200, $X$ |
| 842E | 28 | F5 81 | JSR | \$81F5 | 849A | F0 | 83 |  | BEQ | \$849F |
| 8431 | 20 | FD AE | JSR | \$AEFD | 849C | E8 |  |  | INX |  |
| 8434 | A5 | 14 | LDA | \$14 | 8490 | D0 | F5 |  | BNE | \$8494 |
| 8436 | 85 | FB | STA | \$FB | 849F | 60 |  |  | RTS |  |
| 8438 | A5 | 15 | LDA | \$15 | 84A0 | F0 | 47 |  | BEQ | \$84E9 |
| 843A | 85 | FC | STA | \$FC | 84A2 | B0 | 45 |  | BCS | \$84E9 |
| 843C | 28 | F5 81 | JSR | \$81F5 | 84A4 | 20 | F5 | 81 | JSR | \$81F5 |
| 843F | A5 | 14 | LDA | \$14 | 84 A 7 | A9 | 20 |  | LDA | \#\$20 |
| 8441 | 85 | FD | STA | \$FD | 84A9 | 20 | D2 | FF | JSR | \$FFD2 |
| 8443 | A9 | 4D | LDA | \#\$40 | 84AC | A0 | 02 |  | LDY | \#\$02 |
| 8445 | 8D | 0283 | STA | \$0382 | 84AE | 20 | D1 | 84 | JSR | \$84D1 |
| 8448 | A9 | 84 | LDA | \#\$84 | 84B1 | 88 |  |  | DEY |  |
| 844A | 8 D | 8383 | STA | \$0383 | 8482 | D0 | FA |  | BNE | \$84AE |
| 8440 | AD | 0802 | LDA | \$0200 | 84B4 | A4 | 14 |  | LDY | \$14 |
| 8450 | F0 | 28 | BEQ | \$8472 | 84B6 | 84 | 15 |  | STY | \$15 |
| 8452 | A6 | FB | LDX | \$FB | 84B8 | A0 | 02 |  | LDY | \# 002 |
| 8454 | A5 | FC | LDA | \$FC | 84BA | 20 | D1 | 84 | JSR | \$84D1 |
| 8456 | 20 | 7F 84 | JSR | \$847F | 84BD | 88 |  |  | DEY |  |
| 8459 | 86 | C6 | STX | \$C6 | 84BE | D0 | FA |  | BNE | \$84BA |
| 845B | BD | 0802 | LDA | \$0200, X | 84C0 | 20 | 79 | 00 | JSR | \$0079 |
| 845E | 9 D | 7702 | STA | \$0277, X | 84C3 | C9 | 2C |  | CMP | \#\$2C |
| 8461 | CA |  | DEX |  | 84C5 | D0 | 89 |  | BNE | \$84D0 |
| 8462 | 10 | F7 | BPL | \$845B | $84 C 7$ | 20 | D2 | FF | JSR | \$FFD2 |
| 8464 | 18 |  | CLC |  | 84CA | 20 | 73 | 00 | JSR | \$0073 |
| 8465 | A5 | FB | LDA | \$FB | 84CD | 4C | A0 | 84 | MP | \$8440 |
| 8467 | 65 | FD | ADC | \$FD | 8400 | 68 |  |  | RTS |  |
| 8469 | 85 | FB | STA | \$FB | 84D1 | A2 | 04 |  | LDX | \#\$84 |
| 846B | 90 | 82 | BCC | \$846F | 8403 | A9 | 00 |  | LDA | \#\$00 |
| 8460 | E6 | FC | INC | \$FC | 8405 | 06 | 15 |  | ASL | \$15 |
| 846F | 4C | 83 A4 | JMP | \$4483 | 8407 | 2A |  |  | ROL |  |
| 8472 | A9 | 83 | LDA | \#\$83 | 8408 | CA |  |  | DEX |  |
| 8474 | 8D | 8203 | STA | \$0302 | 8409 | D0 | FA |  | BNE | \$8405 |
| 8477 | A9 | A4 | LDA | \#\$A4 | 84DB | C9 | 8A |  | CMP | \#\$8A |
| 8479 | 8D | 0383 | STA | \$0303 | 840D | 90 | 84 |  | BCC | \$84E3 |
| 847C | 6 C | 0283 | JMP | (\$0302) | 84DF | 18 |  |  | CLC |  |
| 847F | 86 | 63 | STX | \$63 | 84 E 0 | 69 | 37 |  | ADC | \#\$37 |
| 8481 | 85 | 62 | STA | \$62 | 84E2 | 2C | 69 | 30 | BIT | \$3069 |
| 8483 | A2 | 90 | LDX | \#\$90 | $84 \mathrm{E5}$ | 20 | D2 | FF | JSR | \$FFD2 |


| $84 \mathrm{E8}$ | 60 |  |  | RTS |  | 8550 | 85 | 62 |  | STA | \$62 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 84E9 | 4C |  | AF | JMP | \$AF88 | 8552 | 20 | 7 C | 85 | JSR | \$857C |
| 84EC | F0 | 46 |  | BEQ | \$8534 | 8555 | A8 |  |  | TAY |  |
| 84EE | B0 | 44 |  | BCS | \$8534 | 8556 | 68 |  |  | PLA |  |
| 84F0 | 20 | F5 | 81 | JSR | \$81F5 | 8557 | 20 | AD | 85 | JSR | \$85AD |
| 84F3 | A9 | 20 |  | LDA | \#\$20 | 855A | 20 | 73 | 80 | JSR | \$0073 |
| 84F5 | 20 | D2 | FF | JSR | \$FFD2 | 8550 | C9 | 2 C |  | CMP | \# 2 C |
| 84F8 | A5 | 15 |  | LDA | \$15 | 855F | F0 | 10 |  | BEQ | \$8571 |
| 84FA | F0 | 16 |  | BEQ | \$8512 | 8561 | 60 |  |  | RTS |  |
| 84FC | A2 | 08 |  | LDX | \#\$08 | 8562 | 68 |  |  | PLA |  |
| 84FE | 86 | 15 |  | ASL | \$15 | 8563 | A8 |  |  | TAY |  |
| 8508 | B0 | 03 |  | BCS | \$8505 | 8564 | A9 | 80 |  | LDA | \#\$00 |
| 8502 | A9 | 30 |  | LDA | \#\$38 | 8566 | 20 | $A D$ | 85 | JSR | \$85AD |
| 8504 | 2C | A9 | 31 | BIT | \$31A9 | 8569 | 68 |  |  | RTS |  |
| 8507 | 20 | D2 | FF | JSR | \$FFD2 | 856A | 68 |  |  | PLA |  |
| 858A | CA |  |  | DEX |  | 856B | A8 |  |  | TAY |  |
| 850 B | D8 | F1 |  | BNE | \$84FE | 856C | A9 | 00 |  | LDA | \#\$00 |
| 8500 | A9 | 20 |  | LDA | \#\$20 | 856E | 20 | AD | 85 | JSR | \$35AD |
| 850 F | 20 | D2 | FF | JSR | \$FFD2 | 8571 | A9 | 2C |  | LDA | \# 2 C |
| 8512 | A2 | 08 |  | LDX | \#\$08 | 8573 | 20 | D2 | FF | JSR | \$FFD2 |
| 8514 | 86 | 14 |  | ASL | \$14 | 8576 | 20 | 73 | 00 | JSR | \$0073 |
| 8516 | B0 | 03 |  | BCS | \$851B | 8579 | 4C | 37 | 85 | JMP | \$8537 |
| 8518 | A9 | 30 |  | LDA | \#\$38 | 857C | A0 | 01 |  | LDY | \#\$01 |
| 851A | 2C | A9 | 31 | BIT | \$31A9 | 857E | B9 | 62 | 00 | LDA | \$0062, Y |
| 8510 | 20 | D2 | FF | JSR | \$FFD2 | 8581 | C.9 | 30 |  | CMP | \#\$30 |
| 8520 | CA |  |  | DEX |  | 8583 | 90 | 25 |  | BCC | \$85AA |
| 8521 | D0 | F1 |  | BNE | \$8514 | 8585 | C9 | 47 |  | CMP | \#\$47 |
| 8523 | 20 | 79 | 00 | JSR | \$0879 | 8587 | B0 | 21 |  | BCS | \$85AA |
| 8526 | C9 | 2C |  | CMP | \# $\mathbf{2}^{\text {2 }}$ | 8589 | C9 | 3A |  | CMP | \# 3 3A |
| 8528 | D0 | 09 |  | BNE | \$8533 | 858B | 90 | 88 |  | BCC | \$8595 |
| 852A | 20 | D2 | FF | JSR | \$FFD2 | 858D | C9 | 41 |  | CMP | \#\$41 |
| 3520 | 20 | 73 | 80 | JSR | \$0073 | 858F | 90 | 19 |  | BCC | \$85AA |
| 8530 | 4C | EC | 84 | JMP | \$84EC | 8591 | E9 | 37 |  | SBC | \# ${ }^{\text {37 }}$ |
| 8533 | 60 |  |  | RTS |  | 8593 | D0 | 03 |  | BNE | \$8598 |
| 8534 | 4C | 08 | AF | JMP | \$AF08 | 8595 | 38 |  |  | SEC |  |
| 8537 | 85 | 63 |  | STA | \$63 | 8596 | E9 | 30 |  | SBC | \#\$38 |
| 8539 | 20 | 73 | 00 | JSR | \$0073 | 8598 | 99 | 14 | 88 | STA | \$0014,Y |
| 853C | 85 | 62 |  | STA | \$62 | 859B | 88 |  |  | DEY |  |
| 853E | 20 | 7C | 85 | JSR | \$857C | 859C | 10 | E0 |  | BPL | \$857E |
| 8541 | 48 |  |  | PHA |  | 859E | A0 | 84 |  | LDY | \#\$04 |
| 8542 | 20 | 73 | 00 | JSR | \$0073 | 85A0 | 06 | 15 |  | ASL | \$15 |
| 8545 | Fo | 1 B |  | BEQ | \$8562 | 85A2 | 88 |  |  | DEY |  |
| 8547 | C9 | 2C |  | CMP | \#\$2C | 85A3 | D0 | FB |  | BNE | \$85A0 |
| 8549 | F0 | 1F |  | BEQ | \$856A | 85A5 | A5 | 14 |  | LDA | \$14 |
| 854B | 85 | 63 |  | STA | \$63 | 85A7 | 85 | 15 |  | ORA | \$15 |
| 854D | 20 | 73 | 08 | JSR | \$0873 | 85A9 | 68 |  |  | RTS |  |


| 85AA | 4C | 88 | AF | JMP | \$AF08 | 8611 | D0 | 83 | BNE | \$8616 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 85AD | 20 | 91 | B3 | JSR | \$8391 | 8613 | 4C | 1 D A8 | JMP | \$A810 |
| 85B0 | 20 | 81 | 84 | JSR | \$8401 | 86162 | 20 | F5 81 | JSR | \$81F5 |
| 85B3 | 20 | 80 | AD | JSR | \$AD8D | 86192 | 20 | 13 A6 | JSR | \$A613 |
| 85B6 | 20 | DD | BD | JSR | \$BDDD | 861C | B8 | 05 | BCS | \$8623 |
| 85B9 | 20 | 87 | B4 | JSR | \$B487 | 861 E | A2 | 15 | LDX | \#\$15 |
| 85BC | 4C | 21 | AB | MMP | \$AB21 | 8620 | 4C | 37 A4 | JMP | \$A437 |
| 85BF | A5 | 7A |  | LDA | \$7A | 8623 | 38 |  | SEC |  |
| 85C1 | D6 | 02 |  | BNE | \$85C5 | 8624 | A5 | 5F | LDA | \$5F |
| 85 C 3 | C6 | 7B |  | DEC | \$7B | 8626 | E9 | 81 | SBC | \#\$01 |
| 85C5 | C6 | 7A |  | DEC | \$7A | 8628 | 85 | 41 | STA | \$41 |
| $85 \mathrm{C7}$ | A2 | 88 |  | LDX | \#\$08 | 862A | A5 | 60 | LDA | \$68 |
| 85C9 | 20 | 73 | 88 | JSR | \$8073 | 862C | E9 | 80 | SBC | \#\$08 |
| 85CC | 90 | 83 |  | BCC | \$8501 | 862E 8 | 85 | 42 | STA | \$42 |
| 85C.E | 4C | 88 | AF | MMP | \$AF08 | 8638 | 60 |  | RTS |  |
| 85D1 | C9 | 32 |  | CMP | \# $\$ 32$ | 8631 A | A9 | FF | LDA | \#\$FF |
| 8503 | B0 | F9 |  | BCS | \$85CE | 86338 | 85 | 4A | STA | \$4A |
| 85D5 | 6A |  |  | ROR |  | 86352 | 20 | 8A A3 | JSR | \$A38A |
| 8506 | 26 | 14 |  | ROL | \$14 | 8638 | 9 A |  | TXS |  |
| 8508 | CA |  |  | DEX |  | 8639 | C9 | 80 | CMP | \#\$8D |
| 85D9 | D0 | EE |  | BNE | \$85C9 | 863 B | F0 | 85 | BEQ | \$8642 |
| 85DB | A4 | 14 |  | LDY | \$14 | 863D | A2 | 16 | LDX | \#\$16 |
| 85DD | A9 | 08 |  | LDA | \#\$08 | 863 F 4 | 4C | 37 A4 | JMP | \$A437 |
| 85DF | 20 | AD | 85 | JSR | \$85AD | 8642 E | E8 |  | INX |  |
| 85E2 | A9 | 2F |  | LDA | \#\$ 2 F | 8643 E | E8 |  | INX |  |
| 85E4 | 20 | D2 | FF | JSR | \$FFD2 | 8644 E | E8 |  | INX |  |
| 85 E 7 | A5 | 14 |  | LDA | \$14 | 8645 E | E8 |  | INX |  |
| 85E9 | A0 | 08 |  | LDY | \#\$00 | 8646 E | E8 |  | INX |  |
| 85EB | 20 | AD | 85 | JSR | \$85AD | 8647 | 9A |  | TXS |  |
| 85EE | 20 | 73 | 08 | JSR | \$0073 | 8648 | 68 |  | RTS |  |
| 85F1 | C9 | 2C |  | CMP | \# 2 C | 8649 | 60 |  | RTS |  |
| 85F3 | D0 | 86 |  | BNE | \$85FB | 8644 | 4C | 08 AF | JMP | \$AF08 |
| 85F5 | 20 | D2 | FF | JSR | \$FFD2 | 864D | AD 5 | 5B 80 | LDA | \$805B |
| 85F8 | 4C | C7 | 85 | JMP | \$85C7 | 8650 | C9 8 | 87 | CMP | \#\$87 |
| 85FB | 60 |  |  | RTS |  | 8652 F | F0 | 6D | BEQ | \$8661 |
| 85FC | A5 | 90 |  | LDA | \$9D | 8654 | A9 8 | 87 | LDA | \#\$87 |
| 85FE | D0 | 01 |  | BNE | \$8601 | 86568 | 8D | 5B 80 | STA | \$805B |
| 8600 | 60 |  |  | RTS |  | 8659 | A9 | 22 | LDA | \#\$22 |
| 8601 | 20 | 26 | B5 | JSR | \$8526 | 865B 8 | 8D | 5680 | STA | \$8056 |
| 8684 | 38 |  |  | SEC |  | 865E 2 | 205 | 5480 | JSR | \$8054 |
| 8685 | A5 | 33 |  | LDA | \$33 | 86612 | 20 | 7980 | JSR | \$0879 |
| 8667 | E5 | 31 |  | SBC | \$31 | 8664 F | F0 | 42 | BEQ | \$86A8 |
| 8689 | A8 |  |  | TAY |  | 86662 | 20 | F5 81 | JSR | \$81F5 |
| 860A |  | 34 |  | LDA | \$34 | 86692 | 20 | FD AE | JSR | \$AEFD |
| 860 C | E5 | 32 |  | SBC | \$32 | 866 C A | A5 | 14 | LDA | \$14 |
| 860E | 4C |  | 85 | JMP | \$85AD | 866E F | F0 1 | $1 E$ | BEQ | \$868E |


| 8670 | C9 1 | 11 | CMP | \#\$ 11 | 8607 | 90 | F4 |  | BCC | \$86BD |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8672 | B0 1 | $1 A$ | BCS | \$868E | $86 C 9$ | A9 | 31 |  | LDA | \#\$31 |
| 8674 | C6 1 | 14 | DEC | \$14 | 86CB | 85 | 22 |  | STA | \$22 |
| 8676 | A5 1 | 14 | LDA | \$14 | 86CD | A9 | 30 |  | LDA | \#\$30 |
| 8678 | 8A |  | ASL |  | 86CF | 85 | 23 |  | STA | \$23 |
| 8679 | 8A |  | ASL |  | 86D1 | 20 | DE | 86 | JSR | \$86DE |
| 867A | 8A |  | ASL |  | 8604 | E6 | 23 |  | INC | \$23 |
| 867B | 8A |  | ASL |  | 8606 | E6 | 5F |  | INC | \$5F |
| 867C | A8 |  | TAY |  | 8608 | E8 |  |  | INX |  |
| 867 D | A9 | A1 | LDA | \#\$A1 | 8609 | E0 | 11 |  | CPX | \# ${ }^{\text {P } 11}$ |
| 867F | 85 | 15 | STA | \$15 | 860B | 90 | F4 |  | BCC | \$8601 |
| 8681 | A9 | 80 | LDA | \#\$00 | 860D | 60 |  |  | RTS |  |
| 8683 | 85 | 14 | STA | \$14 | 86DE | A0 | 05 |  | LDY | \#\$85 |
| 8685 | A6 | 0A | LDX | \$0A | 86E0 | B9 | 1C | 87 | LDA | \$871C,Y |
| 8687 | 20 | 7980 | JSR | \$8879 | 86 E3 | 20 | D2 | FF | JSR | \$FFD2 |
| 868A | C9 | 22 | CMP | \#\$22 | 86E6 | 88 |  |  | DEY |  |
| 868C | F0 | 83 | BEQ | \$8691 | 86 E 7 | D0 | F7 |  | BNE | \$86E0 |
| 868 E | 4C | 08 AF | JMP | \$AF08 | $86 E 9$ | A5 | 22 |  | LDA | \$22 |
| 8691 | 20 | 7300 | JSR | \$0873 | 86EB | 20 | 02 | FF | JSR | \$FFD2 |
| 8694 | F0 | 8A | BEQ | \$86A8 | 86EE | A5 | 23 |  | LDA | \$23 |
| 8696 | C9 | 22 | CMP | \#\$22 | 86F0 | 20 | D2 | FF | JSR | \$FFD2 |
| 8698 | F0 | 86 | BEQ | \$86A | 86F3 | A9 | 2C |  | LDA | 甡2C |
| 869A | 91 | 14 | STA | (\$14), Y | 86F5 | 20 | D2 | FF | JSR | \$FFD2 |
| 869 C | C8 |  | INY |  | 86F8 | A9 | 22 |  | LDA | \#\$22 |
| 8690 | CA |  | DEX |  | 86FA | 20 | D2 | FF | JSR | \$FFD2 |
| 869E | D0 | F1 | BNE | \$8691 | 86FD | A5 | 5F |  | LDA | \$5F |
| 86A0 | A9 | 80 | LDA | \#\$00 | 86FF | 6A |  |  | ASL |  |
| 86A2 | 91 | 14 | STA | (\$14), Y | 8700 | 0A |  |  | ASL |  |
| 86A4 | 20 | 7380 | JSR | \$0073 | 8781 | 8A |  |  | ASL |  |
| 86A7 | 60 |  | RTS |  | 8702 | 8A |  |  | ASL |  |
| 86A8 | A2 | 00 | LDX | \#\$08 | 8703 | A8 |  |  | TAY |  |
| 86AA | 86 | 5F | STX | \$5F | 8764 | 20 | FB | 81 | JSR | \$81FB |
| 86AC | E8 |  | INX |  | 8707 | B1 | 14 |  | LDA | (\$14), Y |
| 86AD | A9 | 20 | LDA | \#\$20 | 8709 | 48 |  |  | PHA |  |
| 86AF | 85 | 22 | STA | \$22 | 870A | 20 | 82 | 82 | JSR | \$8202 |
| 86B1 | A9 | 31 | LDA | \#\$31 | 870D | 68 |  |  | PLA |  |
| 8683 | 85 | 23 | STA | \$23 | 870E | F8 | 86 |  | BEQ | \$8716 |
| 86B5 | A9 | 00 | LDA | \# $\$ 00$ | 8710 | 20 | D2 | FF | JSR | \$FFD2 |
| 86B7 | 85 | 14 | STA | \$14 | 8713 | C8 |  |  | INY |  |
| 8689 | A9 | A1 | LDA | \#\$A1 | 8714 | D8 | EE |  | BNE | \$8704 |
| 86BB | 85 | 15 | STA | \$15 | 8716 | A9 | 22 |  | LDA | \# \$ $^{\text {2 }}$ |
| 86BD | 20 | DE 86 | JSR | \$86DE | 8718 | 20 | D2 | FF | JSR | \$FFD2 |
| 86C0 | E6 | 23 | INC | \$23 | 871B | 60 |  |  | RTS |  |
| 86 C 2 | E6 | 5F | INC | \$5F | 871 C | 20 | 59 | 45 | JSR | \$4559 |
| $86 C 4$ | E8 |  | INX |  | 871F | 4B |  |  | ??? |  |
| 8655 | E0 | 8A | CPX | \# ${ }^{\text {a }}$ | 8720 | 20 | 8D | A4 | JSR | \$A40D |


| 8723 | CB |  | ??? |  | 8781 | 68 |  | PLA |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8724 | C0 | 03 | CPY | \#\$03 | 8782 | F0 | 8D | BEQ | \$8791 |
| 8726 | 90 | 04 | BCC | \$872C | 8784 | C9 | 5 F | CMP | \#\$5F |
| 8728 | C0 | 87 | CPY | \#\$07 | 8786 | D0 | 02 | BNE | \$878A |
| 872A | 90 | 83 | BCC | \$872F | 8788 | A9 | 6D | LDA | \#\$0D |
| 872C | 4C | 48 EB | JMP | \$EB48 | 878A | 90 | 7782 | STA | \$0277, X |
| 872F | AD | 8D 02 | LDA | \$028D | 878D | E8 |  | INX |  |
| 8732 | C4 | C5 | CPY | \$C5 | 878E | C8 |  | INY |  |
| 8734 | D0 | 85 | BNE | \$873B | 878F | D0 | E7 | BNE | \$8778 |
| 8736 | CD | 8 E 02 | CMP | \$828E | 8791 | 86 | C6 | STX | \$C6 |
| 8739 | F0 | F1 | BEQ | \$872C | 8793 | A9 | 7F | LDA | \#\$ 7F |
| 873B | 84 | C5 | STY | \$C5 | 8795 | 8D | 00 DC | STA | \$DC00 |
| 873D | 8D | 8 E 82 | STA | \$028E | 8798 | 60 |  | RTS |  |
| 8740 | C0 | 04 | CPY | \#\$04 | 8799 | A9 | 48 | LDA | \#\$48 |
| 8742 | F0 | 8B | BEQ | \$874F | 879B | 8 D | 5680 | STA | \$8056 |
| 8744 | C0 | 05 | CPY | \#\$05 | 879E | A9 | EB | LDA | \#\$EB |
| 8746 | F8 | 8A | BEQ | \$8752 | 87A0 | 8D | 5B 80 | STA | \$805B |
| 8748 | CO | 86 | CPY | \#\$06 | 87A3 | 20 | 5480 | JSR | \$8054 |
| 874A | F0 | 09 | BEQ | \$8755 | 87A6 | 68 |  | RTS |  |
| 874C | A0 | 87 | LDY | \#\$07 | 87A7 | 20 | 5E 88 | JSR | \$885E |
| 874E | 2C | A0 81 | BIT | \$01A0 | 87AA | 86 | 2B | STX | \$2B |
| 8751 | 2C | A8 03 | BIT | \$03A8 | 87AC | 84 | 2C | STY | \$2C |
| 8754 | 2C | A8 85 | BIT | \$05A0 | 87AE | A0 | 00 | LDY | \#\$08 |
| 8757 | C9 | 02 | CMP | \#\$02 | 87B6 | 98 |  | TYA |  |
| 8759 | 90 | 87 | BCC | \$8762 | 87B1 | 91 | 2B | STA | (\$2B), Y |
| 875B | F0 | 03 | BEQ | \$8760 | 87B3 | 20 | 1088 | JSR | \$8810 |
| 8750 | A9 | 89 | LDA | \#\$89 | 87B6 | 86 | 2 D | STX | \$2D |
| 875F | 2C | A9 08 | BIT | \$08A9 | 87B8 | 84 | 2E | STY | \$2E |
| 8762 | 84 | BB | STY | \$BB | 87BA | 20 | 33 A5 | JSR | \$A533 |
| 8764 | C6 | BB | DEC | \$BB | 87BD | 20 | 4D 88 | JSR | \$884D |
| 8766 | 18 |  | CLC |  | 87C0 | A9 | F9 | LDA | \#\$F9 |
| 8767 | 65 | BB | ADC | \$BB | 87 C 2 | A2 | 87 | LDX | \#\$87 |
| 8769 | QA |  | ASL |  | 87C4 | 8D | 8283 | STA | \$0362 |
| 8764 | 8A |  | ASL |  | $87 \mathrm{C7}$ | 8E | 0303 | STX | \$0303 |
| 876B | BA |  | ASL |  | 87CA | A9 | 01 | LDA | \#\$81 |
| 876C | QA |  | ASL |  | 87CC | 85 | 7B | STA | \$7B |
| 876D | A0 | A1 | LDY | \#\$A1 | 87CE | A9 | FF | LDA | \#\$FF |
| 876F | 84 | 15 | STY | \$15 | 8700 | 85 | 7A | STA | \$7A |
| 8771 | A0 | 80 | LDY | \# $\$ 80$ | 8702 | A0 | 80 | LDY | \#\$08 |
| 8773 | 84 | 14 | STY | \$14 | 8704 | B1 | FB | LDA | (\$FB), Y |
| 8775 | A8 |  | TAY |  | 8706 | 85 | FD | STA | FFD |
| 8776 | A2 | 00 | LDX | \#\$00 | 8708 | C8 |  | INY |  |
| 8778 | 20 | FB 81 | JSR | \$81FB | 8709 | B1 | FB | LDA | (\$FB), $Y$ |
| 877B | B1 | 14 | LDA | (\$14), Y | 870B | 85 | FE | STA | \$FE |
| 877D | 48 |  | PHA |  | 87DD | F0 | 24 | BEQ | \$8803 |
| 877E | 20 | 0282 | JSR | \$8202 | 87DF | C8 |  | INY |  |


| 87E0 | B1 | FB | LDA | (\$FB), Y | 883F | A5 | FC | LDA | \$FC |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 87E2 | 85 | 14 | STA | \$14 | 8841 | E9 | 00 | SBC | \#\$00 |
| 87E4 | C8 |  | INY |  | 8843 | 85 | FC | STA | \$FC |
| 87E5 | B1 | FB | LDA | (\$FB), Y | 8845 | A9 | 00 | LDA | \#\$00 |
| 87E7 | 85 | 15 | STA | \$15 | 8847 | A8 |  | TAY |  |
| 87E9 | A2 | 04 | LDX | \#584 | 8848 | 91 | 14 | STA | (\$14), Y |
| 87EB | E8 |  | INX |  | 884A | C8 |  | INY |  |
| 87EC | C8 |  | INY |  | 884B | 91 | 14 | STA | (\$14), Y |
| 87ED | B1 | FB | LDA | (\$FB), $Y$ | 884D | A5 | FD | LDA | \$FD |
| 87EF | 9 D | FB 01 | STA | \$01FB, $X$ | 884F | A6 | FE | LDX | \$FE |
| 87F2 | D0 | F7 | BNE | \$87EB | 8851 | 85 | 2B | STA | \$2B |
| 87F4 | 8A |  | TXA |  | 8853 | 86 | 2C | STX | \$2C |
| 87F5 | A8 |  | TAY |  | 8855 | A5 | FB | LDA | \$FB |
| 87F6 | 20 | A2 A4 | JSR | \$ ${ }^{\text {4 }} 4 \mathrm{~A} 2$ | 8857 | A6 | FC | LDX | \$FC |
| 87F9 | A5 | FD | LDA | \$FD | 8859 | 85 | 20 | STA | \$20 |
| 87FB | A6 | FE | LDX | \$FE | 885B | 86 | 2E | STX | \$2E |
| 87FD | 85 | FB | STA | कFB | 885D | 60 |  | RTS |  |
| 87FF | 86 | FC | STX | \$FC | 885E | A5 | 2B | LDA | \$2B |
| 8801 | D0 | C7 | BNE | \$87CA | 8868 | 85 | FD | STA | \$FD |
| 8803 | A9 | 83 | LDA | \#\$83 | 8862 | A5 | 2C | LDA | \$2C |
| 8805 | A2 | A4 | LDX | \#\$A4 | 8864 | 85 | FE | STA | \$FE |
| 8807 | 8D | 0283 | STA | \$8302 | 8866 | A6 | 2D | LDX | \$2D |
| 880A | 8 E | 8383 | STX | \$8303 | 8868 | A4 | 2E | LDY | \$2E |
| 880D | 20 | 74 A4 | JSR | \$A474 | 886A | 86 | FB | STX | \$FB |
| 8810 | 20 | D4 E1 | JSR | \$E1D4 | 886C | 84 | FC | STY | कFC |
| 8813 | A9 | 08 | LDA | \#\$00 | 886E | 60 |  | RTS |  |
| 8815 | 85 | B9 | STA | \$89 | 886F | 20 | 5E 88 | JSR | \$885E |
| 8817 | A6 | 2B | LDX | \$2B | 8872 | 8A |  | TXA |  |
| 8819 | A4 | 2C | LDY | \$2C | 8873 | 38 |  | SEC |  |
| 881 B | 20 | D5 FF | JSR | \$FFD5 | 8874 | E9 | 02 | SBC | \#\$02 |
| 881 E | B0 | 10 | BCS | \$8830 | 8876 | 85 | 2B | STA | \$2B |
| 8820 | 20 | B7 FF | JSR | \$FFB7 | 8878 | 98 |  | TYA |  |
| 8823 | 29 | BF | AND | \#\$BF | 8879 | E9 | 80 | SBC | \#\$08 |
| 8825 | F0 | 08 | BEQ | \$882F | 887B | 85 | 2C | STA | \$2C |
| 8827 | 20 | 3888 | JSR | \$8838 | 887D | 20 | 1088 | JSR | \$8810 |
| 882A | A2 | 1 D | LDX | \#\$1D | 8880 | 86 | FB | STX | \$FB |
| 882C | 4C | 37 A4 | JMP | \$A437 | 8882 | 84 | FC | STY | \$FC |
| 882F | 60 |  | RTS |  | 8884 | 20 | 4D 88 | JSR | \$884D |
| 8830 | 48 |  | PHA |  | 8887 | 20 | 33 A5 | JSR | \$A533 |
| 8831 | 20 | 3888 | JSR | \$8838 | 888A | 60 |  | RTS |  |
| 8834 | 68 |  | PLA |  | 888B | 86 | C2 | STX | \$C2 |
| 8835 | 4C | F9 E0 | JMP | \$E0F9 | 888D | A5 | 3E | LDA | \$3E |
| 8838 | A5 | FB | LDA | \$FB | 888F | 38 |  | SEC |  |
| 883A | 38 |  | SEC |  | 8890 | E5 | C2 | SBC | \$C2 |
| 883B | E9 | 82 | SBC | \#\$82 | 8892 | 85 | BB | STA | \$BB |
| 8830 | 85 | 14 | STA | \$14 | 8894 | 18 |  | CLC |  |


| 8895 | A5 | FB | LDA | \$FB | 88E6 | E6 | 60 | INC | \$60 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8897 | 65 | 49 | ADC | \$49 | 88 E 8 | CA |  | DEX |  |
| 8899 | 85 | 5 F | STA | \$5F | 88E9 | D0 | F2 | BNE | \$88DD |
| 889B | A5 | FC | LDA | \$FC | 88EB | 38 |  | SEC |  |
| 889D | 69 | 80 | ADC | \#\$88 | 88EC | A5 | 2 D | LDA | \$2D |
| 889 F | 85 | 60 | STA | \$60 | 88EE | E5 | BB | SBC | \$BB |
| 88A1 | A5 | 5 F | LDA | \$5F | 88F0 | 85 | 2 D | STA | \$2D |
| 88A3 | 65 | BB | ADC | \$BB | 88F2 | B0 | 83 | BCS | \$88F7 |
| 88A5 | 85 | 5A | STA | \$5A | 88F4 | C. 6 | 2E | DEC | \$2E |
| 88A7 | A5 | 68 | LDA | \$60 | 88F6 | 38 |  | SEC |  |
| 88A9 | 69 | 80 | ADC | \#\$80 | 88F7 | A0 | 80 | LDY | \#\$00 |
| 88AB | 85 | 5B | STA | \$5B | 88F9 | A5 | FD | LDA | \$FD |
| 88AD | A5 | 2D | LDA | \$2D | 88FB | E5 | BB | SBC | \$BB |
| 88AF | 38 |  | SEC |  | 88FD | 85 | FD | STA | \$FD |
| 88B6 | E5 | 5A | SBC | \$5A | 88FF | 91 | FB | STA | (\$FB), Y |
| 88B2 | 85 | 58 | STA | \$58 | 8901 | 85 | 57 | STA | \$57 |
| 88B4 | A8 |  | TAY |  | 8903 | A5 | FE | LDA | \$FE |
| $88 \mathrm{B5}$ | A5 | 2 E | LDA | \$2E | 8905 | E9 | 08 | SBC | \#\$00 |
| 88B7 | E5 | 5B | SBC | \$5B | 8987 | C8 |  | INY |  |
| 8889 | AA |  | TAX |  | 8908 | 85 | FE | STA | \$FE |
| 88BA | E8 |  | INX |  | 890A | 85 | 58 | STA | \$58 |
| 88BB | 98 |  | TYA |  | 890C | 91 | FB | STA | (\$FB), Y |
| 88BC | F0 | $1 F$ | BEQ | \$88DD | 890E | 88 |  | DEY |  |
| 88BE | A5 | 5A | LDA | \$5A | 890F | B1 | 57 | LDA | (\$57) , Y |
| 88C0 | 18 |  | CLC |  | 8911 | 85 | B9 | STA | \$89 |
| $88 \mathrm{C1}$ | 65 | 58 | ADC | \$58 | 8913 | C8 |  | INY |  |
| 88C3 | 85 | 5A | STA | \$5A | 8914 | B1 | 57 | LDA | (\$57), Y |
| 88C5 | 90 | 83 | BCC | \$88CA | 8916 | 85 | BA | STA | \$BA |
| $88 \mathrm{C7}$ | E6 | 5B | INC | \$5B | 8918 | F0 | 18 | BEQ | \$8932 |
| 88C9 | 18 |  | CLC |  | 891A | 88 |  | DEY |  |
| 88CA | A5 | 5 F | LDA | \$5F | 891 B | 38 |  | SEC |  |
| 88CC | 65 | 58 | ADC | \$58 | 891C | A5 | B9 | LDA | \$89 |
| 88CE | 85 | 5 F | STA | \$5F | 891 E | E5 | BB | SBC | \$8B |
| 88D0 | 98 | 82 | BCC | \$88D4 | 8920 | AA |  | TAX |  |
| 88 D 2 | E6 | 60 | INC | \$60 | 8921 | 91 | 57 | STA | (\$57) , Y |
| 88 D 4 | 98 |  | TYA |  | 8923 | A5 | BA | LDA | \$BA |
| 8805 | 49 F | FF | EOR | \#\$FF | 8925 | E9 | 80 | SBC | \#\$80 |
| 88 D 7 | A8 |  | TAY |  | 8927 | C8 |  | INY |  |
| 88D8 | C8 |  | INY |  | 8928 | 91 | 57 | STA | (\$57) , Y |
| 8809 | C6 5 | 5 B | DEC | \$5B | 892A | 85 | 58 | STA | \$58 |
| 880B | C6 | 60 | DEC | \$60 | 892C | 8A |  | TXA |  |
| 880D | B1 | 5 A | LDA | (\$5A), Y | 8920 | 85 | 57 | STA | \$57 |
| 88DF | 915 | 5F | STA | (\$5F), Y | 892F | 4C | 8E 89 | JMP | \$890E |
| 88E1 | C8 |  | INY |  | 8932 | 60 |  | RTS |  |
| 88 E 2 | D8 F | F9 | BNE | \$88DD | 8933 | 8A |  | TXA |  |
| 88 E 4 | E6 5 | 5B | INC | \$5B | 8934 | 38 |  | SEC |  |


| 8935 E5 3E | SBC \$3E | 898F 85 FD | STA \$FD |
| :---: | :---: | :---: | :---: |
| 893785 BB | STA \$BB | 89918557 | STA \$57 |
| 893918 | CLC | 899391 FB | STA (\$FB), $Y$ |
| 893A A5 49 | LDA \$49 | 8995 A5 FE | LDA \$FE |
| 893C 65 BB | ADC \$ ${ }^{\text {B }}$ | 89976960 | ADC \#\$08 |
| 893E B0 84 | BCS $\$ 8944$ | 8999 C8 | INY |
| 8940 C9 FE | CMP \#\$FE | 899A 85 FE | STA \$FE |
| 89429085 | BCC $\$ 8949$ | 899C 8558 | STA $\$ 58$ |
| 8944 A2 17 | LDX \#\$17 | 899E 91 FB | STA (\$FB), $Y$ |
| 8946 4C 37 A4 | JMP \$A437 | 89A0 88 | DEY |
| 8949 A5 2D | LDA \$ 2 D | 89A1 B1 57 | LDA (\$57),Y |
| 894 B 65 BB | ADC \$ BB | 89A3 85 B9 | STA \$B9 |
| 894D A | TAX | 89A5 C8 | INY |
| 894 A A 2E | LDA \$2E | 89A6 B1 57 | LDA (\$57), Y |
| 89566980 | ADC \#\$08 | 89A8 85 BA | STA \$BA |
| 8952 C5 38 | CMP \$38 | 89AA F0 18 | BEQ \$89C4 |
| 8954 D0 07 | BNE \$895D | 89AC 88 | DEY |
| 8956 E4 37 | CPX $\$ 37$ | 89AD 18 | CLC |
| 89589803 | BCC \$895D | 89AE A5 B9 | LDA \$B9 |
| 895A 4C 35 A4 | JMP \$A435 | 89B6 65 BB | ADC \$BB |
| 895D 18 | CLC | 89B2 AA | TAX |
| 895E A5 2D | LDA \$2D | 89B3 9157 | STA (\$57), Y |
| 8968 85 5A | STA \$ 5A | 89B5 A5 BA | LDA \$BA |
| 896265 BB | ADC \$ ${ }^{\text {B }}$ | 89B7 6900 | ADC \#\$00 |
| 89648558 | STA $\$ 58$ | 89B9 C8 | INY |
| 8966 A5 2E | LDA \$2E | 89BA 9157 | STA (\$57),Y |
| 896885 5B | STA \$5B | 89BC 8558 | STA \$58 |
| 896A 6900 | ADC \#\$00 | 89BE 8A | TXA |
| 896C 8559 | STA $\$ 59$ | 89BF 8557 | STA \$57 |
| 896 A A FB | LDA \$FB | 89 Cl 4 C A0 89 | JMP \$89A0 |
| 89706549 | ADC \$49 | 89C4 68 | RTS |
| 897285 5F | STA \$5F | 89 C5 20 F5 81 | JSR \$81F5 |
| 8974 A5 FC | LDA \$FC | $89 \mathrm{C8} 20$ FD AE | JSR \$AEFD |
| 89766980 | ADC \#\$00 | 89CB A5 14 | LDA \$14 |
| 89788560 | STA \$60 | 89CD 85 C9 | STA \$C9 |
| 897A 20 BF A3 | JSR \$A3BF | 89CF A5 15 | LDA \$15 |
| 897018 | CLC | 89D1 85 CA | STA \$CA |
| 897 E A0 86 | LDY \#\$00 | 890320 F5 81 | JSR \$81F5 |
| 8980 A5 2D | LDA \$2D | 890620 FD AE | JSR \$AEFD |
| 898265 BB | ADC \$ $\mathrm{BB}^{\text {d }}$ | 8909 A5 14 | LDA \$14 |
| 898485 2D | STA \$ 2 D | 890B 85 BC | STA \$BC |
| 89869083 | BCC \$898B | 89DD 20 F5 81 | JSR \$81F5 |
| 8988 E6 2E | INC \$2E | 89E0 A5 14 | LDA \$14 |
| 898A 18 | CLC | 89E2 85 BD | STA \$BD |
| 898B A5 FD | LDA FFD | 89E4 A5 15 | LDA \$15 |
| 898D 65 BB | ADC \$ $\mathrm{BB}^{\text {d }}$ | 89E6 85 BE | STA \$ BE |


| 89E8 | A5 | 2B | LDA | \$2B | 8A42 | C8 |  | INY |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 89EA | 85 | FB | STA | \$FB | 8 843 | B1 | FB | LDA | (\$FB), Y |
| 89EC | A5 | 2C | LDA | \$2C | 8A45 | F0 | ED | BEQ | \$8A34 |
| 89EE | 85 | FC | STA | \$FC | 8 A 47 | C9 | 22 | CMP | \# ${ }^{\text {b }} 22$ |
| 89F0 | A5 | CA | LDA | \$CA | 8A49 | D0 | F7 | BNE | \$8A42 |
| 89F2 | D0 | 日F | BNE | \$8A03 | 8 A 4 B | F0 | E2 | BEQ | \$8A2F |
| 89F4 | A5 | C9 | LDA | \$C9 | 8A4D | C9 | 8F | CMP | \#\$8F |
| 89F6 | D0 | 0B | BNE | \$8A03 | 8A4F | F8 | E3 | BEQ | \$8A34 |
| 89F8 | A 0 | 02 | LDY | \#\$02 | 8A51 | C9 | 83 | CMP | \#\$83 |
| 89FA | B1 | FB | LDA | (\$FB), Y | 8453 | F0 | DF | BEQ | \$8434 |
| 89FC | 85 | C9 | STA | \$ C 9 | 8 A 55 | C. 9 | A7 | CMP | \#\$A7 |
| 89FE | C8 |  | INY |  | 8A57 | F0 | 21 | BEQ | \$8A7A |
| 89FF | B1 | FB | LDA | ( $\ddagger \mathrm{FB}$ ), Y | 8459 | C9 | 8A | CMP | \# $\$ 8 \mathrm{~A}$ |
| 8 A 01 | 85 | CA | STA | \$CA | 8A5B | F0 | 1 D | BEQ | \$8A7A |
| 8 A 03 | A5 | C9 | LDA | \$ ${ }^{\text {c }}$ | 8A5D | C9 | 89 | CMP | \#\$89 |
| 8A05 | 85 | 14 | STA | \$14 | 8A5F | F0 | 2C | BEQ | \$8A8D |
| 8407 | A5 | CA | LDA | \$ $C A$ | 8A61 | C9 | CB | CMP | \#\$ CB |
| 8A09 | 85 | 15 | STA | \$15 | 8A63 | D0 | 8B | BNE | \$8A70 |
| $8 \mathrm{A0B}$ | 20 | 13 A6 | JSR | \$4613 | 8A65 | C8 |  | INY |  |
| 8 AOE | B0 | 85 | BCS | \$8A15 | 8466 | B1 | FB | LDA | (\$FB) , Y |
| $8 \mathrm{Al0}$ | A2 | 15 | LDX | \# ${ }^{\text {1 }} 5$ | 8A68 | C9 | 20 | CMP | \#\$20 |
| $8 A^{12}$ | 4C | 37 A4 | JMP | \$A437 | 8A6A | F0 | F9 | BEQ | \$8465 |
| 8A15 | A5 | 5 F | LDA | \$5F | 8A6C | C9 | A4 | CMP | \# $\ddagger$ A4 |
| 8 A17 | 85 | 41 | STA | \$41 | 3A6E | F0 | 1 D | BEQ | \$8A8D |
| 8 A19 | A5 | 60 | LDA | \$60 | 8A70 | C9 | 8D | CMP | \#\$8D |
| 8 AlB | 85 | 42 | STA | \$42 | 8 A72 | F0 | 19 | BEQ | \$8A8D |
| 8A1D | A0 | 80 | LDY | \#\$00 | 8A74 | C9 | E6 | CMP | \#\$E6 |
| 8A1F | B1 | FB | LDA | ( $\ddagger \mathrm{FB}$ ), Y | 8476 | F0 | 15 | BEQ | \$8A8D |
| 8A21 | 85 | FD | STA | \$FD | 8A78 | D0 | B5 | BNE | \$8A2F |
| 8 A23 | C8 |  | INY |  | 8A7A | C8 |  | INY |  |
| 8A24 | B1 | FB | LDA | (\$FB), Y | 8A7B | B1 | FB | LDA | (\$FB), Y |
| 8 A26 | 85 | FE | STA | \$FE | 8A7D | C9 | 20 | CMP | \# ${ }^{\text {2 }} \mathbf{2 0}$ |
| 8A28 | D0 | 83 | BNE | \$8A2D | 8A7F | F0 | F9 | BEQ | \$8A7A |
| 8A2A | 4C | 62 8B | JMP | \$8B62 | 8A81 | C9 | 30 | CMP | \#\$38 |
| 8A2D | C8 |  | INY |  | 8A83 | B0 | 83 | BCS | \$8A88 |
| 8A2E | C8 |  | INY |  | 8485 | 88 |  | DEY |  |
| 8A2F | C8 |  | INY |  | 8A86 | D0 | A7 | BNE | \$8A2F |
| 8 A30 | B1 | FB | LDA | (\$FB), Y | 8488 | 88 |  | DEY |  |
| 8 A32 | D0 | 8A | BNE | \$8A3E | 8489 | C9 | 3 A | CMP | \# ${ }^{\text {S }}$ 3A |
| 8 A34 | A5 | FD | LDA | \$FD | 8A8B | B0 | A2 | BCS | \$8A2F |
| 8 836 | 85 | FB | STA | \$FB | 8A8D | C8 |  | INY |  |
| 8 838 | A5 | FE | LDA | \$FE | 8A8E | B1 | FB | LDA | (\$FB), Y |
| 8А3A | 85 | FC | STA | \$FC | 8498 | C9 | 20 | CMP | \#\$20 |
| 8АЗС | D0 | DF | BNE | \$8A1D | 8492 | F0 | F9 | BEQ | \$8480 |
| 8A3E | C9 | 22 | CMP | \#\$22 | 8494 | 84 | 49 | STY | \$49 |
| 8 A40 | D0 | 0B | BNE | \$8A4D | 8 A 96 | 88 |  | DEY |  |


| 8A97 | A2 | 88 | LDX | \#\$88 | 8AF3 | D0 | 10 | BNE | \$8B65 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8499 | C8 |  | INY |  | 8AF5 | A0 | 02 | LDY | \#\$02 |
| 8A9A | B1 | FB | LDA | (\$FB), Y | 8AF7 | B1 | FB | LDA | (\$FB) , $Y$ |
| 8A9C | C9 | 30 | CMP | \#\$30 | 8AF9 8 | 85 | 39 | STA | \$39 |
| 8A9E | 90 | 8A | BCC | \$8AAA | 8AFB | C8 |  | INY |  |
| 8AAD | C9 | 3A | CMP | \#\$3A | 8AFC B | B1 | FB | LDA | (\$FB), Y |
| 8AA2 | B0 | 06 | BCS | \$8AAA | 8AFE | 85 | 3A | STA | \$3A |
| 8AA4 | 9 D | $00 \quad 02$ | STA | \$0200, X | $8 \mathrm{B00}$ A | A2 | 11 | LDX | \#\$11 |
| 8AA7 | E8 |  | INX |  | $8 \mathrm{B0} 2$ | 4C | 37 A4 | JMP | \$A437 |
| 8AA8 | D8 | EF | BNE | \$8A99 | 8 B 05 | C8 |  | INY |  |
| 8AAA | A9 | 3A | LDA | \#\$3A | 8B66 | B1 | 58 | LDA | (\$58), $Y$ |
| 8AAC | 90 | 0802 | STA | \$0200, X | $8 \mathrm{B08}$ | 85 | B7 | STA | \$B7 |
| 8AAF | 86 | BF | STX | \$BF | 8B0A | C8 |  | INY |  |
| 8AB1 | A9 | 82 | LDA | \#\$02 | 8B0B B | B1 | 58 | LDA | (\$58) , Y |
| 8AB3 | 85 | 7B | STA | \$7B | 8B0D | C5 | C. 4 | CMP | \$C4 |
| 8 AB5 | A9 | 88 | LDA | \# 0 $^{\text {0 }}$ | 8B6F | D0 | 86 | BNE | \$8B17 |
| 8AB7 | 85 | 7A | STA | \$7A | 8B11 | A5 | B7 | LDA | \$B7 |
| 8 8B9 | 20 | F5 81 | JSR | \$81F5 | 8 B 13 | C5 | C3 | CMP | \$C3 |
| 8ABC | A5 | 14 | LDA | \$14 | $8 \mathrm{B15}$ F | F0 | 15 | BEQ | \$8B2C |
| 8ABE | 85 | C3 | STA | \$ C 3 | 8 B 17 A | A5 | B9 | LDA | \$89 |
| 8AC0 | A5 | 15 | LDA | \$15 | 8B19 18 | 18 |  | CLC |  |
| 8AC2 | 85 | C4 | STA | \$C4 | 8B1A 6 | 65 | BC | ADC | \$BC |
| 8AC4 | C5 | CA | CMP | \$ $C$ A | 8B1C 8 | 85 | B9 | STA | \$89 |
| 8AC6 | F0 | BA | BEQ | \$8AD2 | 8B1E 9 | 90 | 02 | BCC | \$8822 |
| 8AC8 | B6 | 9E | BCS | \$8AD8 | $8 \mathrm{B20}$ | E6 | BA | INC | \$BA |
| 8ACA | A5 | 49 | LDA | \$49 | 8 B 22 A | A5 | 5A | LDA | \$5A |
| 8ACC | 65 | BF | ADC | \$ BF | 8 B 248 | 85 | 58 | STA | \$58 |
| 8ACE | A8 |  | TAY |  | $8 \mathrm{B26}$ | A5 | 5B | LDA | \$5B |
| 8ACF | 4C | 55 8B | JMP | \$8B55 | $8 \mathrm{B28} 8$ | 85 | 59 | STA | \$59 |
| 8AD2 | A5 | C3 | LDA | \$C3 | 8B2A | D0 | BC | BNE | \$8AE8 |
| 8AD4 | C5 | C9 | CMP | \$ ${ }^{\text {c9 }}$ | 8B2C | A6 | B9 | LDX | \$B9 |
| 8AD6 | 90 | F2 | BCC | \$8ACA | 8B2E | A5 | BA | LDA | \$BA |
| 8AD8 | A5 | BD | LDA | \$BD | 8836 | 20 | 7F 84 | JSR | \$847F |
| 8ADA | 85 | B9 | STA | \$89 | 8 B 33 A | A5 | BF | LDA | \$BF |
| 8ADC | A5 | 41 | LDA | \$41 | 88358 | 85 | 3E | STA | \$3E |
| 8ADE | 85 | 58 | STA | \$58 | 8 B 37 E | E4 | 3E | CPX | \$3E |
| 8AE0 | A5 | BE | LDA | \$BE | 8839 F | F0 | 8B | BEQ | \$8B46 |
| 8AE2 | 85 | BA | STA | \$BA | 8B3B | B0 | 06 | BCS | \$8843 |
| 8AE4 | A5 | 42 | LDA | \$42 | 8B3D 20 | 20 | 8 B 88 | JSR | \$888B |
| 8AE6 | 85 | 59 | STA | \$59 | $8 \mathrm{B40} 4$ | 4C | 46 8B | JMP | \$8B46 |
| 8AE8 | A0 | 00 | LDY | \#\$80 | 8B43 2 | 20 | 3389 | JSR | \$8933 |
| 8AEA | B1 | 58 | LDA | (\$58) , Y | 8B46 A | A4 | 49 | LDY | \$49 |
| 8AEC | 85 | 5A | STA | \$5A | $8 B 48$ | A2 | 80 | LDX | \#\$08 |
| 8AEE | C8 |  | INY |  | 8B4A | BD | 0802 | LDA | \$0200, X |
| 8AEF | B1 | 58 | LDA | (\$58), Y | 8B4D F | F0 | 86 | BEQ | \$8855 |
| 8AF1 | 85 | 5B | STA | \$5B | 8B4F 9 |  | FB | STA | (\$FB) , $Y$ |


| $8 \mathrm{B51}$ | C8 |  | INY |  | 8BA6 8 | 85 F | FE | STA | \$FE |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8852 | E8 |  | INX |  | 8BA8 D | D0 8 | 05 | BNE | \$8BAF |
| $8 \mathrm{B5} 3$ | D0 | F5 | ENE | \$8B4A | 8BAA 6 | 68 |  | PLA |  |
| $8 \mathrm{B55}$ | B1 | FB | LDA | (\$FB) , Y | 8BAB 6 | 68 |  | PLA |  |
| 8857 | C9 | 2C | CMP | \# ${ }^{\text {2 }}$ 2C | 8BAC 4 | 4C7 | 74 A4 | JMP | \$A474 |
| $8 \mathrm{B59}$ | F0 | 04 | BEQ | \$8B5F | 8BAF C | c8 |  | INY |  |
| 8B5B | 88 |  | DEY |  | 8BB6 | c8 |  | INY |  |
| 8B5C | 4C | 2F 8A | JMP | \$8A2F | $8 \mathrm{BB1}$ C | C8 |  | INY |  |
| 8B5F | 4C | 8D 8A | JMP | \$8A8D | 8BB2 B | B1 F | FB | LDA | (\$FB), $Y$ |
| $8 \mathrm{B62}$ | A0 | 80 | LDY | \#\$00 | 8BB4 D | D0 0 | 0C | BNE | \$8BC2 |
| $8 \mathrm{B64}$ | B1 | 41 | LDA | (\$41), Y | 8BB6 ${ }^{\text {A }}$ | A5 F | FD | LDA | \$FD |
| 8 B 66 | 85 | 5A | STA | \$5A | 8BB8 8 | 85 F | FB | STA | \$FB |
| $8 \mathrm{B68}$ | C8 |  | INY |  | 8BBA A | A5 F | FE | LDA | कFE |
| $8 \mathrm{B69}$ | B1 | 41 | LDA | (\$41), Y | 8BBC 8 | 85 F | FC | STA | कFC |
| 8B6B | 85 | 5B | STA | \$5B | 8BBE A | A2 | 80 | LDX | \#\$00 |
| 8B6D | D0 | 85 | BNE | \$8B74 | 8BC0 | F0 D | DB | BEQ | \$8B9D |
| 8B6F | 68 |  | PLA |  | 8 BC 2 C | C9 F | FF | CMP | \#\$FF |
| $8 \mathrm{B70}$ | 68 |  | PLA |  | 8BC4 D | D0 8 | 04 | BNE | \$8BCA |
| 8B71 | 4C | 74 A4 | JMP | \$A474 | 8BC6 8 | 853 | 3E | STA | \$3E |
| $8 \mathrm{B74}$ | C8 |  | INY |  | $8 \mathrm{BC8}$ F | F0 2 | 21 | BEQ | \$8BEB |
| 8B75 | A5 | BD | LDA | \$BD | 8BCA | C9 2 | 22 | CMP | \#\$22 |
| 8877 | 91 | 41 | STA | (\$41), Y | 8BCC D | D0 8 | 89 | BNE | \$8BD7 |
| 8B79 | C8 |  | INY |  | 8BCE E | E8 |  | INX |  |
| 8B7A | A5 | BE | LDA | \$BE | 8BCF E | E0 0 | 02 | CPX | \#\$02 |
| 8B7C | 91 | 41 | STA | (\$41), Y | $8 \mathrm{BD1} 1$ | D0 0 | 08 | BNE | \$8BDB |
| 8B7E | 18 |  | CLC |  | 8BD3 A | A2 0 | 08 | LDX | \#\$00 |
| 8B7F | A5 | BD | LDA | \$BD | 8BD5 F | F0 D | DA | BEQ | \$8BB1 |
| 8B81 | 65 | BC | ADC | \$BC | 8BD7 | E0 0 | 01 | CPX | \#\$01 |
| 8883 | 85 | BD | STA | \$BD | 8BD9 D | D0 D | D6 | BNE | \$8BB1 |
| $8 \mathrm{B85}$ | 90 | 02 | BCC | \$8B89 | 8BDB 8 | 85 | 3E | STA | \$3E |
| $8 \mathrm{B87}$ | E6 | BE | INC | \$ BE | 8BDD | C9 | C8 | CMP | \#SC0 |
| 8B89 | A5 | 5A | LDA | \$5A | 8BDF 9 | 90 | 02 | BCC | \$8BE3 |
| 8B8B | 85 | 41 | STA | \$41 | 8BE1 E | E9 | 60 | SBC | \#\$60 |
| 8B8D | A5 | 5B | LDA | \$5B | 8BE3 | C9 | 60 | CMP | \#\$60 |
| 8B8F | 85 | 42 | STA | \$42 | 8BE5 | B0 | 84 | BCS | \$8BEB |
| 8891 | D0 | CF | BNE | \$8B62 | 8BE7 | C9 | 21 | CMP | \# ${ }^{\text {2 }} 21$ |
| 8B93 | A2 | 00 | LDX | \#\$88 | 8BE9 | B8 | C6 | BCS | \$8BB1 |
| 8B95 | A5 | 2B | LDA | \$2B | 8BEB 8 | 85 | 30 | STA | \$3D |
| $8 \mathrm{B97}$ | 85 | FB | STA | \$FB | 8BED 84 | 84 | 49 | STY | \$49 |
| 8B99 | A5 | 2C | LDA | \$2C | 8BEF 8 | 86 | 3C | STX | \$3C |
| 8B9B | 85 | FC | STA | \$FC | 8 BF 1 | A2 | 81 | LDX | \#\$01 |
| 8B9D | A0 | 80 | LDY | \#\$00 | 8BF3 | C8 |  | INY |  |
| 8B9F | B1 | FB | LDA | (\$FB), $Y$ | 8BF4 | B1 | FB | LDA | (\$FB), $Y$ |
| 8BA1 | 85 | FD | STA | \$FD | 8BF6 | C5 | 3E | CMP | \$3E |
| 8BA3 | C8 |  | INY |  | 8BF8 | D0 | 83 | BNE | \$8BFD |
| 8BA4 | B1 | FB | LDA | (\$FB), $Y$ | 8BFA | E8 |  | INX |  |


| 8BFB | De | F6 | ENE | \$8BF3 | 8 C 59 | F8 | 83 | BEQ | \$8C5E |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8BFD | 86 | 3E | STX | \$3E | 8C5B | C8 |  | INY |  |
| 8BFF | E0 | 02 | CPX | \#\$62 | 8C5C | 91 | FB | STA | (\$FB), Y |
| 8 CO 1 | B0 | 12 | BCS | \$8C15 | 8C5E | A5 | 3F | LDA | \$3F |
| 8 C 03 | CA |  | DEX |  | 8C60 | F0 | 03 | BEQ | \$8C65 |
| 8 C 04 | A5 | 3D | LDA | \$3D | 8C62 | C8 |  | INY |  |
| 8 CO 6 | C9 | 20 | CMP | \#\$20 | 8C63 | 91 | FB | STA | (\$FB), Y |
| 8 C 08 | F8 | 88 | BEQ | \$8C12 | 8C65 | A9 | 47 | LDA | \#\$47 |
| 8C0A | A9 | 80 | LDA | \#\$00 | 8C67 | C8 |  | INY |  |
| 8 COC | 85 | 40 | STA | \$48 | 8C68 | 91 | FB | STA | (\$FB), Y |
| 8 COE | 85 | $3 F$ | STA | \$3F | 8C6A | A9 | 3E | LDA | \#\$3E |
| 8 Cl 10 | F0 | 15 | BEQ | \$8C27 | 8C6C | C8 |  | INY |  |
| 8 C 12 | 4C | 3380 | JMP | \$8033 | 8C6D | 91 | FB | STA | (\$FB), Y |
| 8 C 15 | A9 | 08 | LDA | \#\$00 | 8C6F | A5 | 3D | LDA | \$3D |
| 8 C 17 | 28 | 7F 84 | JSR | \$847F | 8 C 71 | C8 |  | INY |  |
| 8C1A | A2 | 00 | LDX | \#\$00 | $8 C 72$ | 91 | FB | STA | (\$FB), Y |
| 8C1C | BD | 0882 | LDA | \$0200, x | $8 C 74$ | A9 | 5D | LDA | \#\$5D |
| 8C1F | 85 | 3 F | STA | \$3F | 8 C 76 | C8 |  | INY |  |
| 8C21 | E8 |  | INX |  | 8 C 77 | 91 | FB | STA | (\$FB), Y |
| 8C22 | BD | 0082 | LDA | \$0200, X | 8 C 79 | A6 | 3C | LDX | \$3C |
| 8C25 | 85 | 40 | STA | \$48 | 8С7B | 4C | B1 8B | JMP | \$8BB1 |
| 8C27 | A5 | 3D | LDA | \$3D | 8C7E | 85 | 30 | STA | \$3D |
| 8C29 | C9 | 61 | CMP | \#\$61 | $8 \mathrm{C80}$ | A9 | 50 | LDA | \#\$50 |
| 8C2B | 90 | 51 | BCC | \$8C7E | 8 C 82 | 85 | 62 | STA | \$62 |
| 8C2D | C9 | 7B | CMP | \#\$7B | 8C84 | A9 | A3 | LDA | \#きA3 |
| 8C2F | B0 | 4D | BCS | \$8C7E | 8 C 86 | 85 | 63 | STA | \$63 |
| 8C31 | 38 |  | SEC |  | $8 \mathrm{C88}$ | A2 | 51 | LDX | \#\$51 |
| 8C32 | E9 | 20 | SBC | \#\$20 | 8C8A | A0 | 80 | LDY | \#\$00 |
| 8C34 | 85 | 3 D | STA | \$3D | 8C8C | 20 | FB 81 | JSR | \$81FB |
| 8 C 36 | A2 | 07 | LDX | \#\$07 | 8C8F | B1 | 62 | LDA | (\$62), Y |
| 8 C 38 | A5 | 40 | LDA | \$40 | 8C91 | 48 |  | PHA |  |
| 8С3A | D0 | 06 | BNE | \$8C42 | 8 C 92 | 28 | 0282 | JSR | \$8202 |
| 8С3C | CA |  | DEX |  | 8C95 | 68 |  | PLA |  |
| 8C3D | A5 | 3 F | LDA | \$3F | 8C96 | C5 | 3D | CMP | \$3D |
| 8C3F | D0 | 81 | ENE | \$8C42 | 8 C 98 | F8 | 20 | BEO | \$8CBA |
| 8C41 | CA |  | DEX |  | 8C9A | C8 |  | INY |  |
| $8 \mathrm{C42}$ | E4 | 3 E | CPX | \$3E | 8С9B | C8 |  | INY |  |
| 8C44 | F0 | 0B | BEQ | \$8C51 | 8С9C | C8 |  | INY |  |
| 8C46 | B8 | 06 | BCS | \$8C4E | 8C90 | CA |  | DEX |  |
| 8C48 | 20 | 8888 | JSR | \$888B | 8C9E | 10 | EC | BPL | \$8C8C |
| 8C4B | 4C | 518 C | JMP | \$8C51 | 8CAO | A5 | 30 | LDA | \$3D |
| 8C4E | 20 | 3389 | JSR | \$8933 | 8CA2 | c9 | 1B | CMP | \#\$1B |
| 8C51 | A4 | 49 | LDY | \$49 | 8CA4 | 98 | 85 | BCC | \$8CAB |
| 8 C 53 | A9 | 5B | LDA | \#\$5B | 8CA6 | A2 | 8D | LDX | \#\$0D |
| 8C55 | 91 | FB | STA | (\$FB), Y | 8CA8 | 4C | 37 A4 | JMP | \$A437 |
| 8 C 57 | A5 | 48 | LDA | \$40 | 8CAB | 69 | 48 | ADC | \#\$48 |



| 8068 | 8 D | 46 | 8E | STA | \$8E46 | 80D4 | 20 | D2 | FF | JSR | \$FFD2 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8D6B | AD | 50 | 8E | LDA | \$8E50 | 8DD7 | A9 | 12 |  | LDA | \# ${ }^{\text {1 }} 12$ |
| 806E | 8 D | 88 | 03 | STA | \$0388 | 8DD9 | 20 | D2 | FF | ISR | \$FFD2 |
| 8071 | AD | 51 | 8E | LDA | \$8E51 | 8DDC | A5 | 3A |  | LDA | \$3A |
| 8074 | 8 D | 09 | 03 | STA | \$8309 | 8DDE | 8D | 4F | 8E | STA | \$8E4F |
| 8077 | 58 |  |  | CLI |  | 8DE1 | A6 | 39 |  | LDX | \$39 |
| 8078 | 68 |  |  | RTS |  | 80E3 | 8 E | 4E | 8E | STX | \$8E4E |
| 8079 | 80 | 49 | 8E | STA | \$8E49 | 8DE6 | 20 | CD | BD | JSR | \$BDCD |
| 807C | 88 |  |  | PHP |  | 8DE9 | A9 | 92 |  | LDA | \#\$92 |
| 8D7D | 8E | 4A | 8E | STX | \$8E4A | 8DEB | 20 | D2 | FF | JSR | \$FFD2 |
| 8080 | 8C | 4B | 8E | STY | \$8E4B | 8DEE | 18 |  |  | CLC |  |
| 8 D 83 | A5 | 9 D |  | LDA | \$90 | 8DEF | B0 | 96 |  | BCS | \$8D87 |
| 8085 | F0 | 0D |  | BEQ | \$8094 | 8DF1 | AE | 4D | 8E | LDX | \$8E4D |
| 8D87 | AD | 49 | 8E | LDA | \$8E49 | 8DF4 | AC | 4C | 8E | LDY | \$8E4C |
| 808A | $A C$ | 4B | 8E | LDY | \$8E4B | 8DF7 | 20 | F8 | FF | JSR | कFFF0 |
| 8D8D | AE | 4A | 8 E | LDX | \$8E4A | 8DFA | 20 | E4 | FF | JSR | \$FFE4 |
| $8 \mathrm{D9} 0$ | 28 |  |  | PLP |  | 8DFD | F0 | 25 |  | BEQ | \$8E24 |
| 8091 | 6C | 50 | 8E | JMP | (\$8E50) | 8DFF | C9 | 2 F |  | CMP | \# ${ }^{\text {2 }} 2 \mathrm{~F}$ |
| 8094 | A5 | 39 |  | LDA | \$39 | 8E01 | 90 | 21 |  | BCC | \$8E24 |
| 8D96 | AD | 46 | 8E | LDA | \$8E46 | 8 E 03 | C9 | 3A |  | CMP | \#\$3A |
| $8 \mathrm{D99}$ | F0 | EC |  | BEQ | \$8D87 | 8 E 05 | B0 | 1D |  | BCS | \$8E24 |
| 809B | 38 |  |  | SEC |  | 8 E 07 | E9 | 30 |  | SBC | \#\$30 |
| 809C | 20 | F0 | FF | JSR | \$FFF0 | 8 E 99 | D0 | 07 |  | BNE | \$8E12 |
| 8D9F | 8 E | 4 D | 8E | STX | \$8E4D | 8E0B | A9 | FF |  | LDA | \#\$FF |
| 8DA2 | 8C | 4 C | 8E | STY | \$8E4C | 8E0D | 8 D | 47 | 8 E | STA | \$8E47 |
| 8DA5 | 18 |  |  | CLC |  | 8 E 10 | D0 | 12 |  | BNE | \$8E24 |
| 8DA6 | A2 | 00 |  | LDX | \#\$80 | 8 E 12 | AA |  |  | TAX |  |
| 8DA8 | A8 | 18 |  | LDY | \# \$ $18^{\text {d }}$ | 8 E 13 | 38 |  |  | SEC |  |
| 8DAA | 20 | F0 | FF | ISR | \$FFF0 | 8E14 | A9 | 80 |  | LDA | \#\$00 |
| 8DAD | A2 | 8 F |  | LDX | \#\$0F | 8 E 16 | 2 A |  |  | ROL |  |
| 8DAF | A9 | 20 |  | LDA | \# 20 | 8 E 17 | CA |  |  | DEX |  |
| 8DB1 | 20 | D2 | FF | JSR | \$FFD2 | 8E18 | D0 | FC |  | BNE | \$8E16 |
| 8DB4 | CA |  |  | DEX |  | 8E1A | 8D | 48 | 8 E | STA | \$8E48 |
| 8DB5 | D0 | F8 |  | BNE | \$8DAF | 8 E 1 D | A9 | 00 |  | LDA | \# 00 |
| 80B7 | 18 |  |  | CLC |  | 8E1F | 8D | 47 | 8 E | STA | \$8E47 |
| 80B8 | A2 | 08 |  | LDX | \#\$00 | 8 E 22 | F0 | 14 |  | BEQ | \$8E38 |
| 8DBA | A0 | 18 |  | LDY | \#\$18 | 8E24 | AE | 47 | 8 E | LDX | \$8E47 |
| 8DBC | 20 | F0 | FF | JSR | \$FFF0 | 8 E 27 | F0 | 0F |  | BEQ | \$8E38 |
| 8DBF | A9 | 12 |  | LDA | \#\$12 | 8E29 | C9 | 20 |  | CMP | \# ${ }^{\text {2 }} 20$ |
| 8DC1 | 20 | D2 | FF | JSR | \$FFD2 | 8E2B | F0 | 16 |  | BEQ | \$8E43 |
| 80C.4 | AD | 4F | 8 E | LDA | \$8E4F | 8E2D | 20 | E4 | FF | JSR | \$FFE4 |
| 8DC7 | AE | 4E | 8 E | LDX | \$8E4E | 8 E 30 | F0 | FB |  | BEQ | \$8E2D |
| 8DCA | 20 | CD | BD | JSR | \$BDCD | 8 E 32 | C9 | 20 |  | CMP | \#\$ 20 |
| 8DCD | A9 | 92 |  | LDA | \#\$92 | 8 E 34 | F0 | 8D |  | BEQ | \$8E43 |
| 8DCF | 20 | D2 | FF | JSR | \$FFD2 | 8 E 36 | D0 | C5 |  | BNE | \$8DFD |
| 8DD2 | A9 | 20 |  | LDA | \# $\$ 20$ | 8 E 38 | AE | 48 | 8 E | LDX | \$8E48 |



| 8EEE | 29 | 7F |  | AND | \#\$ 7 F | 8F51 | A2 | 15 |  | LDX | \#\$15 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8EF0 | 20 | D2 | FF | JSR | \$FFD2 | 8 F 53 | 4C | 37 | A4 | JMP | \$A437 |
| 8EF3 | A9 | 25 |  | LDA | \#\$25 | 8 F 56 | A5 | 5 F |  | LDA | \$5F |
| 8EF5 | 20 | D2 | FF | JSR | \$FFD2 | 8F58 | 85 | FB |  | STA | \$FB |
| 8EF8 | 20 | 33 | 8F | JSR | \$8F33 | 8F5A | A5 | 68 |  | LDA | \$68 |
| 8EFB | A0 | 00 |  | LDY | \#\$00 | 8F5C | 85 | FC |  | STA | \$FC |
| 8EFD | B1 | BB |  | LDA | (\$BB), Y | 8F5E | 20 | 79 | 00 | JSR | \$0079 |
| 8EFF | 85 | 62 |  | STA | \$62 | 8F61 | C9 | 2C |  | CMP | \# ${ }^{\text {2C }}$ |
| 8 F 01 | C8 |  |  | INY |  | 8 F 63 | D0 | E1 |  | BNE | \$8F46 |
| 8 F 02 | B1 | BB |  | LDA | (\$BB) , Y | 8F65 | 20 | 73 | 08 | JSR | \$0073 |
| 8 F 04 | 85 | 63 |  | STA | \$63 | 8F68 | D0 | 8 F |  | BNE | \$8F79 |
| 8 F 06 | A2 | 90 |  | LDX | \#\$90 | 8F6A | 38 |  |  | SEC |  |
| 8 F 98 | 20 | 44 | BC | JSR | \$BC44 | 8F6B | A5 | 2 D |  | LDA | \$2D |
| 8F6B | 20 | DD | BD | JSR | \$BDDD | 8F60 | E9 | 82 |  | SBC | \#\$02 |
| 8FgE | 20 | 1 E | $A B$ | JSR | \$ABIE | 8F6F | 85 | 5F |  | STA | 韦5 |
| 8F11 | A9 | 8D |  | LDA | \#\$0D | 8F71 | A5 | 2 E |  | LDA | \$2E |
| 8F13 | 20 | D2 | FF | JSR | \$FFD2 | 8F73 | E9 | 80 |  | SBC | \#\$08 |
| 8F16 | 18 |  |  | CLC |  | 8F75 | 85 | 60 |  | STA | \$60 |
| 8F17 | A5 | BB |  | LDA | \$8B | 8F77 | D0 | 18 |  | BNE | \$8F91 |
| 8F19 | 69 | 05 |  | ADC | \#\$05 | 8 F 79 | B0 | CB |  | BCS | \$8F46 |
| 8F1B | 85 | BB |  | STA | \$BB | 8F7B | 20 | F5 | 81 | JSR | \$81F5 |
| 8F1D | 90 | 02 |  | BCC | \$8F21 | 8F7E | E6 | 14 |  | INC | \$14 |
| 8F1F | E6 | BC |  | INC | \$BC | 8F80 | 28 | 13 | A6 | JSR | \$A613 |
| 8 F 21 | 20 | E4 | FF | JSR | \$FFE4 | 8F83 | A5 | FC |  | LDA | \$FC |
| 8F24 | 20 | E1 | FF | JSR | \$FFE1 | 8F85 | C5 | 60 |  | CMP | \$60 |
| 8 F 27 | D0 | 01 |  | BNE | \$8F2A | 8F87 | 90 | 88 |  | BCC | \$8F91 |
| 8F29 | 60 |  |  | RTS |  | 8F89 | DB | BB |  | BNE | \$8F46 |
| 8F2A | A5 | CB |  | LDA | \$CB | 8F8B | A5 | FB |  | LDA | \$FB |
| 8F2C | C9 | 40 |  | CMP | \#\$40 | 8F8D | C5 | 5 F |  | CMP | \$5F |
| 8F2E | D6 | F1 |  | BNE | \$8F21 | 8F8F | B0 | B5 |  | BCS | \$8F46 |
| 8F30 | 38 |  |  | SEC |  | 8F91 | A0 | 08 |  | LDY | \#\$08 |
| 8F31 | B0 | 9 B |  | BCS | \$8ECE | $8 \mathrm{F93}$ | A5 | 5 F |  | LDA | \$5F |
| 8F33 | A9 | 3D |  | LDA | \#\$3D | 8F95 | 91 | FB |  | STA | (\$FB), $Y$ |
| 8F35 | 20 | D2 | FF | JSR | \$FFD2 | $8 \mathrm{F97}$ | C8 |  |  | INY |  |
| 8F38 | 18 |  |  | CLC |  | 8F98 | A5 | 68 |  | LDA | \$60 |
| 8F39 | A5 | BB |  | LDA | \$BB | 8F9A | 91 | FB |  | STA | (\$FB) , Y |
| 8F3B | 69 | 02 |  | ADC | \#\$82 | 8F9C | C8 |  |  | INY |  |
| 8F3D | 85 | BB |  | STA | \$BB | 8F90 | B1 | FB |  | LDA | (\$FB), Y |
| 8F3F | 98 | 02 |  | BCC | \$8F43 | 8F9F | AA |  |  | TAX |  |
| 8F41 | E6 | BC |  | INC | \$BC | 8FA0 | C8 |  |  | INY |  |
| 8F43 | 60 |  |  | RTS |  | 8FA1 | B1 | FB |  | LDA | (\$FB), Y |
| 8F44 | 90 | 83 |  | BCC | \$8F49 | 8FA3 | 320 | 7F | 84 | JSR | \$847F |
| 8F46 | 4C | 88 | AF | JMP | \$AF08 | 8FA6 | 68 |  |  | PLA |  |
| 8F49 | 20 | F5 | 81 | JSR | \$81F5 | 8FA7 | 68 |  |  | PLA |  |
| 8F4C | 20 | 13 | A6 | JSR | \$A613 | 8FA8 | A2 | FF |  | LDX | \#\$FF |
| 8F4F | B0 | 05 |  | BCS | \$8F56 | 8FAA | A9 | 01 |  | LDA | \# ${ }^{\text {d }} 1$ |


| 8FAC | 4C |  | 92 | JMP | \$927D | 9008 | B1 | FB |  | LDA | (\$FB), Y |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8FAF | 20 | F5 | 81 | JSR | \$81F5 | 900A | 85 | FD |  | STA | \$FD |
| 8FB2 | 4C | A3 | A8 | JMP | \$A8A3 | 900C | C8 |  |  | INY |  |
| 8FB5 | A9 | 03 |  | LDA | \#\$03 | 9000 | B1 | FB |  | LDA | (\$FB) , Y |
| 8FB7 | 20 | FB | A3 | JSR | \$A3FB | 980 F | 85 | FE |  | STA | \$FE |
| 8FBA | A5 | 7B |  | LDA | \$7B | 9011 | D0 | 05 |  | BNE | \$9018 |
| 8FBC | 48 |  |  | PHA |  | 9813 | A2 | 11 |  | LDX | \#\$11 |
| 8FBD | A5 | 7A |  | LDA | \$7A | 9015 | 4C | 37 | A4 | JMP | \$A437 |
| 8FBF | 48 |  |  | PHA |  | 9818 | A 0 | 84 |  | LDY | \#\$84 |
| 8FC0 | A5 | 3A |  | LDA | \$3A | 901 A | B1 | FB |  | LDA | (\$FB) , Y |
| 8FC2 | 48 |  |  | PHA |  | 901 C | C9 | E1 |  | CMP | \# ${ }^{\text {E }}$ E1 |
| 8FC3 | A5 | 3B |  | LDA | \$3B | 901 E | F0 | 日A |  | BEQ | \$902A |
| 8FC5 | 48 |  |  | PHA |  | 9020 | A5 | FD |  | LDA | \$FD |
| 8FC6 | A9 | 8D |  | LDA | \#\$8D | 9022 | 85 | FB |  | STA | \$FB |
| 8FC8 | 48 |  |  | PHA |  | 9824 | A5 | FE |  | LDA | \$FE |
| 8FC9 | 28 | 79 | 00 | JSR | \$0079 | 9026 | 85 | FC |  | STA | \$FC |
| 8FCC | 20 | AF | 8F | JSR | \$8FAF | 9828 | D0 | DC |  | BNE | \$9806 |
| 8FCF | 4C | AE | A7 | JMP | \$A7AE | 902A | A2 | FF |  | LDX | \#\$FF |
| 8FD2 | A2 | 80 |  | LDX | \#\$00 | 902C | E8 |  |  | INX |  |
| 8FD4 | C6 | 7A |  | DEC | \$7A | 9020 | C8 |  |  | INY |  |
| 8FD6 | B0 | 02 |  | BCS | \$8FDA | 902 E | B1 | FB |  | LDA | (\$FB), Y |
| 8FD8 | C6 | 7B |  | DEC | \$7B | 9030 | F0 | 07 |  | BEQ | \$9039 |
| 8FDA | 20 | 73 | 88 | JSR | \$0873 | 9032 | DD | 80 | 02 | CMP | \$8200, X |
| 8FDD | F0 | 86 |  | BEQ | \$8FE5 | 9035 | F0 | F5 |  | BEQ | \$902C |
| 8FDF | 90 | 80 | 02 | STA | \$0200, X | 9837 | D0 | E7 |  | BNE | \$9820 |
| 8FE2 | E8 |  |  | INX |  | 9839 | DD | 80 | 02 | CMP | \$0200, X |
| 8FE3 | D0 | F5 |  | BNE | \$8FDA | $983 C$ | D8 | E2 |  | BNE | \$9820 |
| 8FE5 | A9 | 08 |  | LDA | \#\$00 | 903 E | 38 |  |  | SEC |  |
| 8FE7 | 9 D | 00 | 02 | STA | \$8280, X | $903 F$ | A5 | FB |  | LDA | \$FB |
| 8FEA | A9 | 83 |  | LDA | \#\$03 | 9041 | E9 | 01 |  | SBC | \#\$01 |
| 8FEC | 20 | FB | A3 | JSR | \$A3FB | 9843 | 85 | 7A |  | STA | \$7A |
| 8FEF | A5 | 7B |  | LDA | \$7B | 9045 | A5 | FC |  | LDA | \$FC |
| 8FF1 | 48 |  |  | PHA |  | 9847 | E9 | 80 |  | SBC | \#\$80 |
| 8FF2 | A5 | 7A |  | LDA | \$7A | 9049 | 85 | 7B |  | STA | \$7B |
| 8FF4 | 48 |  |  | PHA |  | 904 B | 4C | AE | A7 | MMP | \$A7AE |
| 8FF5 | A5 | 3A |  | LDA | \$3A | 904 E | D0 | 0C |  | BNE | \$905C |
| 8FF7 | 48 |  |  | PHA |  | 9850 | A9 | 00 |  | LDA | \#\$00 |
| 8FF8 | A5 | 39 |  | LDA | \$39 | 9052 | 85 | CB |  | STA | \$CB |
| 8FFA | 48 |  |  | PHA |  | 9854 | 20 | E4 | FF | JSR | \$FFE4 |
| 8FFB | A9 | 8D |  | LDA | \# $\$^{\text {8 }}$ | 9057 | F0 | FB |  | BEQ | \$9054 |
| 8FFD | 48 |  |  | PHA |  | 9859 | 85 | 98 |  | STA | \$98 |
| 8FFE | A5 | 2B |  | LDA | \$2B | 905B | 60 |  |  | RTS |  |
| 9000 | 85 | FB |  | STA | \$FB | 905C | 20 | 9 E | AD | JSR | \$AD9E |
| 9002 | A5 | 2C |  | LDA | \$2C | 905 F | 20 | A3 | B6 | JSR | \$B6A3 |
| 9084 | 85 | FC |  | STA | \$FC | 9862 | C9 | 88 |  | CMP | \#\$80 |
| 9006 | A0 | 00 |  | LDY | \#\$00 | 9064 | F0 | EA |  | BEQ | \$9050 |


| 9066 | 85 | FB | STA | \$FB | 900.1 | E6 | 59 | INC | \$59 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 9068 | A9 | 00 | LDA | \#\$00 | 90.3 | A5 5 | 59 | LDA | \$59 |
| 9064 | 85 | C6 | STA | \$ 66 | 9065 | 85 F | FC | STA | \$FC |
| 906 C | 20 | E4 FF | JSR | \$FFE4 | 90.78 | 85 | 38 | STA | \$38 |
| 906 F | F0 | FB | BEQ | \$906C | 90.9 | Aó | 2B | LDX | \$2B |
| 9071 | A4 | FB | LDY | \$FB | 90 CB A | A4 2 | 2C | LDY | \$2C |
| 9073 | 88 |  | DEY |  | 90 CD | A9 | 00 | LDA | \#\$08 |
| 9074 | D1 | 22 | CMP | (\$22), Y | 90 CF 2 | 20 | D5 FF | JSR | \$FFD5 |
| 9076 | Fb | 85 | BEQ | \$907D | 9002 | 90 | 83 | BC.C | \$9007 |
| 9078 | 88 |  | DEY |  | 9004 | 4C | F9 E0 | JMP | \$E0F9 |
| 9079 | 10 | F9 | BPL | \$9074 | 90072 | 20 | B7 FF | JSR | 韦FFB7 |
| 907 B | 30 | EF | BMI | \$906C | 90082 | 29 | BF | AND | \#\$BF |
| 907 D | 85 | 90 | STA | \$98 | 900 C | F0 | 85 | BEQ | \$90E3 |
| 907 F | 68 |  | RTS |  | 900 A | A2 1 | 1 D | LDX | \#\$1D |
| 9080 | 20 | D4 E1 | JSR | \$E1D4 | 98 ED | 4C | 37 A4 | JMP | \$4437 |
| 9083 | A9 | 88 | LDA | \#\$00 | 90.83 | 862 | 2 D | STX | \$2D |
| 9885 | 85 | B9 | STA | \$B9 | 90 ES 8 | 84 | 2 E | STY | \$2E |
| 9087 | 20 | 26 B5 | JSR | \$8526 | 90 E 78 | 86 | 5F | STX | \$5F |
| 9084 | A5 | 2D | LDA | \$2D | $90 \mathrm{E9} 8$ | 84 | 60 | STY | \$60 |
| 908 C | 85 | 5F | STA | \$5F | 90 EB A | A5 | FB | LDA | \$FB |
| 988 E | A5 | 2 E | LDA | \$2E | $90 E D 8$ | 85 | 5 A | STA | \$5A |
| 9090 | 85 | 68 | STA | \$68 | 90 EF | A5 | FC | LDA | \$FC |
| 9892 | 38 |  | SEC |  | 90 F 18 | 85 | 5 B | STA | \$5B |
| 9093 | A5 | 31 | LDA | \$31 | 9073 | 38 |  | SEC |  |
| 9095 | 85 | 5A | STA | \$5A | 90 F 4 | A5 3 | 33 | LDA | \$33 |
| 9097 | E5 | 2F | SBC | \$2F | 90F6 E | E9 | 01 | SBC | \#\$01 |
| 9899 | 85 | FD | STA | \$FD | 9078 | A8 |  | TAY |  |
| 909 B | A5 | 32 | LDA | \$32 | 9079 | A5 3 | 34 | LDA | \$34 |
| 999 D | 85 | 5B | STA | \$5B | 90 FB | E9 | 00 | SBC | \#\$00 |
| 909 F | E5 | 30 | SBC | \$30 | 90FD A | A |  | TAX |  |
| 90A1 | 85 | FE | STA | कFE | 90 FE 9 | 98 |  | TYA |  |
| 90A3 | A5 | 33 | LDA | \$33 | 90 FF 3 | 38 |  | SEC |  |
| 90A5 | 38 |  | SEC |  | 9100 | E5 5 | 5A | SBC | \$5A |
| 90A6 | E9 | 01 | SBC | \#\$01 | 91028 | 85 | 58 | STA | \$58 |
| 9048 | 85 | 58 | STA | \$58 | 9104 | A8 |  | TAY |  |
| 90AA | A5 | 34 | LDA | \$34 | 91058 | 8A |  | TXA |  |
| 90 AC | E9 | 08 | SBC | \#\$08 | 9106 | E5 | 5B | SBC | \$5B |
| 90AE | 85 | 59 | STA | \$59 | 9108 A | AA |  | TAX | - |
| 90 BO | 28 | BF A3 | JSR | \$A3BF | 9109 E | E8 |  | INX |  |
| 9083 | A5 | 37 | LDA | \$37 | 910A 9 | 98 |  | TYA |  |
| 9085 | 85 | 41 | STA | \$41 | 918 B F | F0 1 | 1 F | BEQ | \$912C |
| 9087 | A5 | 38 | LDA | \$38 | 910 D A | A5 | 5A | LDA | \$5A |
| $98 \mathrm{B9}$ | 85 | 42 | STA | \$42 | 918 F 1 | 18 |  | CLC |  |
| 90 BB | A5 | 58 | LDA | \$58 | 91186 | 655 | 58 | ADC | \$58 |
| 90 BD | 85 | FB | STA | \$FB | 91128 | 85 | 5 A | STA | \$5A |
| 90 BF | 85 | 37 | STA | \$37 | 91149 | 90 | 03 | BCC | \$9119 |


| 9116 E6 | 58 | INC \$58 | 9171 A5 15 | LDA \$15 |
| :---: | :---: | :---: | :---: | :---: |
| 911818 |  | CLC | 9173 C9 04 | CMP \#\$04 |
| 9119 A5 | 5 F | LDA $\$ 5 \mathrm{~F}$ | 9175 B0 65 | BCS \$917C |
| 911 B 65 | 58 | ADC $\$ 58$ | 9177 A2 DE | LDX \#\$0E |
| 911085 | 5 F | STA \$5F | 9179 4C 37 A4 | JMP \$A437 |
| 911 F 90 | 02 | BCC $\$ 9123$ | 917C C9 80 | CMP \#\$80 |
| 9121 E6 | 68 | INC $\$ 60$ | 917 EBCF | BCS $\$ 9177$ |
| 912398 |  | TYA | 918060 | RTS |
| 912449 | FF | EOR \#\$FF | 9181206991 | JSR \$9169 |
| 9126 A8 |  | TAY | 91848538 | STA $\$ 38$ |
| 9127 C 8 |  | INY | 9186 A5 14 | LDA \$14 |
| 9128 C6 | 5B | DEC \$5B | 91888537 | STA $\$ 37$ |
| 912 C C6 | 60 | DEC \$60 | 918A 4C 63 A6 | MMP \$A663 |
| 912 C B1 | 5A | LDA ( $\$ 5.5$ ), Y | 9180206991 | JSR \$9169 |
| 912 E 91 | 5F | STA (\$5F), Y | 9190 A0 80 | LDY \#\$00 |
| 9130 C8 |  | INY | 919298 | TYA |
| 9131 D0 | F9 | BNE $\$ 912 \mathrm{C}$ | 91939114 | STA (\$14), Y |
| 9133 E6 | 5B | INC \$5B | $9195 \mathrm{C8}$ | INY |
| 9135 E6 | 60 | INC \$ 60 | 91969114 | STA (\$14), Y |
| 9137 CA |  | DEX | $9198 \mathrm{C8}$ | INY |
| 9138 D0 | F2 | BNE \$912C | 91999114 | STA (\$14), Y |
| 913A 38 |  | SEC | 919 B A5 14 | LDA \$14 |
| 913 A A5 | 5F | LDA $\$ 5 F$ | 919D 18 | CLC |
| 913 D 8 | 31 | STA \$31 | 919 E 691 | ADC \#\$01 |
| 913 F E5 | FD | SBC \$FD | 91 A 0852 B | STA \$ 2B |
| 914185 | 2F | STA \$2F | 91 A 2 AA | TAX |
| 9143 A5 | 60 | LDA \$ 60 | 91A3 A5 15 | LDA $\$ 15$ |
| 914585 | 32 | STA $\$ 32$ | 91A5 6900 | ADC \#\$00 |
| 9147 E5 | FE | SBC \$FE | 91 A 782 C | STA \$ 2 C |
| 914985 | 30 | STA \$30 | 91A9 A8 | TAY |
| 914 B A5 | 41 | LDA \$41 | 91AA 8A | TXA |
| 914 D 85 | 37 | STA $\$ 37$ | 91 AB 6902 | ADC \# ${ }^{\text {O }} 2$ |
| 914F A5 | 42 | LDA \$42 | 91 AD 852 D | STA \$ 2 D |
| 915185 | 38 | STA \$38 | 91 AF 98 | TYA |
| 915368 |  | PLA | $91 \mathrm{B0} 6980$ | ADC \#\$08 |
| 915468 |  | PLA | $91 \mathrm{B2} 852 \mathrm{E}$ | STA \$2E |
| 915520 | 33 A5 | JSR \$A533 | $91 \mathrm{B4} 4 \mathrm{C} 63 \mathrm{~A} 6$ | JMP \$A663 |
| 9158 A9 | 00 | LDA \#\$00 | $91 \mathrm{B7}$ A9 7C | LDA \#\$7C |
| 915A 20 | 90 FF | JSR \$FF90 | 91B9 8D 0483 | STA \$0304 |
| 915020 | E7 FF | JSR \$FFE7 | 91 BC A9 A5 | LDA \#\$A5 |
| 916020 | 77 A6 | JSR \$A677 | 91 BE 8 D 0583 | STA \$0305 |
| 916320 | 8E A6 | JSR \$A68E | $91 \mathrm{C1}$ A9 1A | LDA \#\#1A |
| 9166 4C | CF 92 | JMP \$92CF | 91 C 38 D 868 | STA \$0306 |
| 9169 D0 | 83 | BNE \$916E | $91 \mathrm{C6}$ A9 E4 | LDA \#\$E4 |
| 916 BC | 88 AF | JMP \$AF08 | $91 \mathrm{C8} 8 \mathrm{D} 8803$ | STA $\$ 0308$ |
| 916 E 20 | F5 81 | JSR \$81F5 | $91 \mathrm{CB} \mathrm{A9} \mathrm{A7}$ | LDA \#\$A7 |


| $91 . C D$ | 8D | 07 | 03 | STA | \$0307 | 9239 | 8D | 11 | 86 | STA | \$061 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 9100 | 8D | 09 | 03 | STA | \$8309 | 923 C | 8D | D6 | 05 | STA | \$05D6 |
| 9103 | A9 | 86 |  | LDA | \#\$86 | 923F | 8D | 26 | 86 | STA | \$0626 |
| 9105 | 8D | 0A | 03 | STA | \$030A | 9242 | A9 | 05 |  | LDA | \#\$05 |
| 9108 | A9 | AE |  | LDA | \#\$AE | 9244 | 8 D | C1 | D9 | STA | \$D9C1 |
| 910A | 8 D | 8B | 83 | STA | \$030B | 9247 | 80 | 11 | DA | STA | \$DA11 |
| 9100 | A9 | FE |  | LDA | \#S FE | 924 A | 8D | D6 | D9 | STA | \$D906 |
| 91 DF | 8D | 17 | 83 | STA | \$0317 | 9240 | 8D | 26 | DA | STA | \$DA26 |
| 91E2 | 8D | 19 | 83 | STA | \$0319 | 9258 | A9 | 8D |  | LDA | \#\$0D |
| 9155 | A9 | 66 |  | LDA | \#\$66 | 9252 | 20 | D2 | FF | JSR | \$FFD2 |
| 9157 | 8 D | 16 | 03 | STA | \$0316 | 9255 | 4C | 74 | A4 | JMP | \$ ${ }^{\text {4 } 474}$ |
| 91EA | A9 | 47 |  | LDA | \#\$47 | 9258 | 18 |  |  | CLC |  |
| 91EC | 8 D | 18 | 83 | STA | \$0318 | 9259 | 20 | F0 | FF | JSR | \$FFF0 |
| 91 EF | 78 |  |  | SEI |  | 925C | A9 | 2A |  | LDA | \#\$2A |
| 91 Fb | A9 | 48 |  | LDA | \#\$48 | 925E | A2 | 16 |  | LDX | \#\$16 |
| 91F2 | 8 D | 8F | 02 | STA | \$028F | 9268 | 20 | D2 | FF | JSR | \$FFD2 |
| 91F5 | A9 | EB |  | LDA | \#\$EB | 9263 | CA |  |  | DEX |  |
| $91 F 7$ | 8 D | 90 | 02 | STA | \$0290 | 9264 | D0 | FA |  | BNE | \$9260 |
| 91FA | 58 |  |  | CLI |  | 9266 | 60 |  |  | RTS |  |
| 91FB | 68 |  |  | PLA |  | 9267 | 2A |  |  | ROL |  |
| 91FC | 68 |  |  | PLA |  | 9268 | 28 | 59 | 54 | JSR | \$5459 |
| 91FD | 4C | 74 | A4 | JMP | \$A474 | 926 B | 49 | 4C |  | EOR | \#\$4C |
| 9208 | 20 | 63 | A6 | JSR | \$ 4663 | 9260 | 49 | 54 |  | EOR | \#\$54 |
| 9203 | A9 | 93 |  | LDA | \#\$93 | 926F | 55 | 20 |  | EOR | \$20, X |
| 9205 | 20 | D2 | FF | JSR | \$FFD2 | 9271 | 43 |  |  | ??? |  |
| 9208 | A9 | 80 |  | LDA | \#\$00 | 9272 | 49 | 53 |  | EOR | \#\$53 |
| 920A | 8D | 20 | D0 | STA | \$0020 | 9274 | 41 | 42 |  | EOR | ( $\$ 42, X$ ) |
| 920 D | 8 D | 21 | D0 | STA | \$D021 | 9276 | 20 | 20 | 4E | JSR | \$4E20 |
| 9210 | A9 | 85 |  | LDA | \#\$85 | 9279 | 43 |  |  | ??? |  |
| 9212 | 8D | 86 | 82 | STA | \$0286 | 927A | 50 | 20 |  | BUC | \$929C |
| 9215 | A2 | 8A |  | LDX | \#\$8A | 927 C | 2A |  |  | ROL |  |
| 9217 | A0 | 89 |  | LDY | \#\$09 | 9270 | AD | 02 | 83 | LDA | \$0302 |
| 9219 | 20 | 58 | 92 | JSR | \$9258 | 9280 | 8 D | 97 | 92 | STA | \$9297 |
| 921 C | A2 | 0C |  | LDX | \#\$0C | 9283 | AD | 03 | 83 | LDA | \$0303 |
| 921 E | A0 | 09 |  | LDY | \#\$89 | 9286 | 8D | 9C | 92 | STA | \$929C |
| 9220 | 18 |  |  | CLC |  | 9289 | A9 | 96 |  | LDA | \#\$96 |
| 9221 | 20 | F0 | FF | JSR | \$FFF0 | 928 B | 8 D | 02 | 83 | STA | \$0302 |
| 9224 | A2 | 15 |  | LDX | \# ${ }^{\text {d }} 15$ | 928 E | A9 | 92 |  | LDA | \#\$92 |
| 9226 | BD | 67 | 92 | LDA | \$9267, X | 9290 | 8D | 03 | 83 | STA | \$8383 |
| 9229 | 20 | D2 | FF | JSR | \$FFD2 | 9293 | 4C | 86 | A4 | JMP | \$A486 |
| 922C | CA |  |  | DEX |  | 9296 | A9 | 83 |  | LDA | \# ${ }^{\text {\$ }} 83$ |
| 922 D | 10 | F7 |  | BPL | \$9226 | 9298 | 80 | 82 | 03 | STA | \$0302 |
| 922 F | A2 | 8 E |  | LDX | \# ${ }^{\text {¢ }}$ E | 929 B | A9 | A4 |  | LDA | \#\$A4 |
| 9231 | A0 | 09 |  | LDY | \#\$09 | 9290 | 8D | 83 | 03 | STA | \$0303 |
| 9233 | 20 | 58 | 92 | JSR | \$9258 | 92 AB | 20 | 33 | A5 | JSR | \$A533 |
| 9236 | 80 | Cl | 05 | STA | \$85C1 | 92 A 3 | 18 |  |  | CLC |  |


| 9244 | A5 | 22 | LDA | \$22 | 928E | C5 | 2 D |  | CMP | \$2D |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 92A6 | 69 | 02 | ADC | \#\$02 | 92C0 | D0 | 83 |  | BNE | \$92C5 |
| 92 AB | 85 | 20 | STA | \$20 | 92C2 | A9 | 80 |  | LDA | \#\$80 |
| 92AA | A5 | 23 | LDA | \$23 | $92 C 4$ | 2C | A9 | 80 | BIT | \$00A9 |
| 92AC | 69 | 88 | ADC | \#\$08 | $92 C 7$ | 85 | 8C |  | STA | \$0C |
| 92AE | 85 | 2E | STA | \$2E | $92 C 9$ | 20 | 79 | 00 | JSR | \$0079 |
| $92 \mathrm{B0}$ | 20 | 68 A6 | JSR | \$A660 | 92CC | 4C | 88 | 90 | JMP | \$9880 |
| 92B3 | 4C | 74 A4 | JMP | \$A474 | 92CF | A5 | 0C |  | LDA | \$0C |
| $92 \mathrm{B6}$ | A5 | 32 | LDA | \$32 | 9201 | 10 | 84 |  | BPL | \$9207 |
| $92 \mathrm{B8}$ | C5 | 2E | CMP | \$2E | 9203 | C6 | 30 |  | DEC | \$30 |
| 92BA | D0 | 09 | BNE | \$92C5 | 9205 | C6 | 32 |  | DEC | \$32 |
| 92BC | A5 | 31 | LDA | \$31 | 9207 | 4C | AE | A7 | JMP | \$A7AE |

Loading the utility
Once the utility and the data for Coder have been set up, a loader program something like the following should be used.

```
10 A=A+1:IF A=1 THEN LOAD"UTILITY DATA",8,1
20 IF A=2 THEN LOAD"UTILITY'',8,1
30 SYS32768
```


## 10 Bits 'n' pieces

## General

This chapter is a collection of snippets of information we have found out since acquiring our 64 s about 18 months ago. No detailed code here, just the bare facts and a few ideas.

## AUTO-REPEATS and INTERRUPTS

We have seen two articles on the subject of providing a repeat on all keys. Both articles were based upon the same idea used on the pre-8000 series PETs. In essence the normal IRQ service routine is patched to include additional code by changing the vector at cinv from its default of \$EA31. The additional routine simple scans SFDX - if a key is being pressed then it is reset to no key (\$40) - and ends with a Jump to \$EA31 to process the normal interrupt. This will then detect a key as being pressed and enter the appropriate character in the keyboard buffer. Alternatively, for a repeat on all keys, simply POKE 650,128 . To disable the repeat, POKE 650,1. (For a full description of the IRQ service routine, see Chapter 4.)

The second method is obviously far easier, but the first does allow a selective auto-repeat to be implemented.

Whilst on the subject of the hardware interrupt (see Chapter 4 for its implementation in the KEY commands), here is a short example to demonstrate what can be done. The following program patches IRQ to scan for function keys 1 and 3 . These keys are used to increment the border and background colours respectively. The routine only takes the appropriate action once every 60 interrupts (about a second). If you remove the interrupt counter from \$C01F to \$C02D, the effects produced are quite unusual, but the routine becomes of little practical use as it is too fast to exercise selective control.

To enable: SYS 49152
To disable: SYS 49170

| C000 | 78 | SEI |  | ENABLE ENTRY |
| :--- | :--- | :--- | :--- | :--- |
| C001 | A91F | LDA $\# \$ 1 F$ | SET CINN TO POINT TO $\$ C 01 F$ |  |
| C003 | 801403 | STA | $\$ 0314$ |  |


| C.008 | 801503 | STA | \$8315 |  |
| :---: | :---: | :---: | :---: | :---: |
| C00B | A900 | LDA | \# $\ddagger 00$ | SET IRQ COUNTEF TO ZERO |
| C000 | 8000c1 | STA | +C100 |  |
| C010 | 58 | CLI |  |  |
| C011 | 60 | RTS |  |  |
| C012 | 78 | SEI |  | DISABLE ENTRY |
| C013 | A931 | LDA | \#\$31 | RESTORE CINM TO \$EA31 |
| C015 | 3D1403 | STA | \$0314 |  |
| Ca18 | A9EA | LDA | \#\# EA |  |
| C01A | 801503 | STA | \$0315 |  |
| C01D | 58 | CLI |  |  |
| CalE | 60 | RTS |  |  |
| ColF | EE00C1 | INC: | \$C100 | NEW IRQ ENRTY |
| $\mathrm{C0} 22$ | ADOOC1 | LDA | \$C180 |  |
| Ca25 | C93C | CMP | \#\$3C | 60 INTERRUPTS ??? |
| C027 | D02A | BNE | \$C053 | NO - SKIP KEY SCAN |
| C029 | A900 | LDA | \#\$00 | YES - SO RESET COLNTER |
| C02B | 3000C1 | STA | \$C100 |  |
| C02E | A5CB | LDA | \$CB | SFDX - CURRENT KEY PRESS |
| C030 | 18 | CLC. |  |  |
| C031 | C904 | CMP | \#\$04 | F1 ??? |
| C033 | D00F | BNE | \$C044 |  |
| C035 | 18 | CLC. |  |  |
| c036 | AD20D0 | LDA | \$0020 | BDR COLOUR |
| C039 | 290 F | AND | \# $\$ 0 \mathrm{~F}$ | BITS 0-3 ONLY (0-15 DEC) |
| C03B | 6901 | ADC | \#\$01 | INCREMENT IT |
| C03D | 802000 | STA | \$0020 |  |
| C040 | A900 | LDA | \#\$00 | ENSURE SKIP TAKEN |
| C042 | F90F | BEQ | \$ 6053 | SKIP BKD COLOUR |
| C044 | C.905 | CMP | \#\#05 | BKD COLOUR |
| C646 | D00B | BNE | \$C053 |  |
| C048 | 18 | CLC |  |  |
| C849 | AD2100 | LDA | \$D021 |  |
| C04C | 2907 | AND | \#\$0F |  |
| C04E | 6901 | ADC | \#\$01 |  |
| C050 | 3D2100 | STA | \$0021 |  |
| C053 | $4[31 \mathrm{EA}$ | JMP | \$ EA31 | CONTINUE NOPIAL IRQ |

The BAsic loader:
1 DATA $120,169,31,141,20,3,169,19$
2, 141, 21, 3
2 DATA $169,0,141,0,193,88,96,120$, 169, 49
3 DATA 141, 20, 3, 169, 234, 141, 21, 3, 88,96

```
4 DATA 238, 0, 193, 173, 0, 193, 201, 60
    208, }4
5 DATA 169, 0, 141, 0, 193, 165, 203, 24
    201, 4
6 DATA 208, 15, 24, 173, 32, 208, 41, 15
    105, 1
7 DATA 141, 32, 208, 169, 0, 240, 15, 20
1, 5, 208
8 DATA 11, 24, 173, 33, 208, 41, 15, 105
, 1, 141
9 DATA 33, 208, 76, 49, 234, 237, 61, 3,
    170, 173
10 FOR I=49152 TO 49238:READ A:POKE I,A:
NEXT I
```

The IRQ vector can be used to great advantage. One common use is to provide interrupt driven music (see The Companion to the Commodore 64 pub. by Pan/PCN) and as in the Utility to make the function keys programmable (Chapter 4).

## Simple program protection

Some basics include commands to 'unlist' or generate protected files; the 64 's, however, does not. In order to protect our software we have to resort to programming tricks.
There are many ways to afford a program some degree of protection from unauthorized change. Most of these are well known and do little to prevent the experienced user from gaining access. Chapter 1 showed how the link addresses could be modified to make program lines invisible and list out of sequence. Another way to hide areas of code is to end lines with a rem"idelidel . . . . ." sequence. On listing, the 'deletes' will erase characters to their left. Most other techniques require a program to be RUN.
Once a program has been run we can destroy some of the vectors from 50300 to $\$ 0333$. These include the PRINT TOKENS LINK, IQPLOP, which could be directed to, say, print 'syntax error' at \$afar, the save vector at ISAVE to prevent saving and also disable the RUN/STOP at istop. We could also put a specified value somewhere in memory which, if not found, will cause the program to crash, erase itself or even perform a cold boot of the system. Unfortunately (or perhaps fortunately) any BASIC program can be loaded without being run. To produce programs which auto-run on loading requires knowledge of both machine code and the operating system of the 64 (see Chain in Chapter 8). Nearly all commercial software uses a number of levels of protection, one of which is usually auto-running. We have covered a number of ways to accomplish this for your own software in Chapter 2, but purposely
leave out many of the techniques used by commercial software houses. (Remember it is illegal to reproduce commercial software.)

Commodore Computing International Volume 2 No.II has an article on program protection. It contains the usual:

```
DISABLE RUN/STOP POKE 808,251
DISABLE LIST POKE 774,131:POKE 775,164
DISABLE SAVE
POKE 818,131:POKE 819,164
```

The first simply bypasses the test by jumping to the end of the routine (RTS). The latter two jump to 'ready for Basic'. Similar changes may be made to run/restore at nminv (see Chapter 6). The article does give a program to generate auto-run programs from your own code. If you are interested then, as the program is copyrighted, we suggest you get hold of a copy of the magazine. A further tape protection idea was given in Commodore User Volume 1 No. 10.

## Specified input

One of the most difficult and time-consuming tasks in producing software for use by others is in making it 'crash-proof'. basic does not allow the programmer to specify which keys are valid during input. The results of incorrect entries in type, size or number can spoil a well-thought-out, pleasing display or even crash the program. The way round this problem is to write your own input routine.

Commodore Horizons magazine of February 1984 published a very good machine code 'Keyscan' input program written by Adrian Warman which does just about everything you could ask for. We see little point in re-inventing the wheel, so we suggest that you read that article. However, we have approached the problem from a different angle and produced a simple routine entirely in basic. This is intended to be called when input is required. The type of input expected is set using the variable ' $F$ ' which is set to 0 for a real number, 1 for an integer and 2 for a string. Strings may contain commas and quotes. Editing an input may only be carried by using the delete key. The returned value may, if required, be converted to a number by a simple val(). If the routine is to be used more than once, as must be emptied by: A $\$=\times \prime$ before each use.

60000- Generate a flashing cursor.
60030 If it is a delete check chars are there to be removed.
60040 'Return' marks end of input and the resulting string As is passed back to the main program.
60050 Real numbers.
60060 Integer only.
60070 String.
60090 Update the display and wait for next char.

60110－Real numbers may begin with + or－and may contain only numerals and a decimal point．
60170－As for real but may not have a decimal point and must lie in the range given．
60230－Allows any of the standard alphanumerics．To provide for lower case mode where uppercase characters have their high bit set some graphics are permitted（ $128+32$ to $128+64$ ）．

```
60000 POKE 204,0:POKE 207,0
60010 GET Yक:IF Y$="" GOTO 60000
60020 A=LEN(A$)
60030 IF Y}\ddagger=CHR$(20) AND A>0 THEN A क=LE
T$(A$,A-1):GOTO 60090
60040 IF Y
60050 IF F=0 THEN GOSUB 60110:GOT0 60090
60060 IF F=1 THEN GOSUB 60170:GOTO 60090
60070 IF F=2 THEN GOSUB 60230:GOTO 60090
60080 GOTO 60000
60090 PRINT Y$;:GOTO 60000
60100 RETURN
60110 REM REAL
60120 IF Y方="+" OR Y$="-" AND A=0 THEN
A$=Y$:GOTO 60160
60130 IF Y&>"/" AND Y&<":" THEN A A=A &+Y官
:GOTO 60160
60140 IF Y&="." THEN A A=A = +Y*:GOTO 60160
60150 Y$=""
60160 RETURN
60170 REM INTEGER
60180 IF Yक="+" OR Y$="-" AND A=0 THEN
A$=Y各:GOTO 60220
```



```
:GOTO 60220
60200 IF UAL(Aま)>32767 OR UAL(A$)<-32768
    THEN A$=LEFT$(A$,A):GOTO 60220
60210 Y$=""
60220 RETURN
6 0 2 3 0 ~ R E M ~ S T R I N G ~
60240 IF (ASC(Y婁) AND 127)<32 OR (ASC(Y击
) AND 127)>95 THEN Y $=" ":GOTO 60260
60250 A $=A $+Y利
6 0 2 6 0 ~ R E T U R N
```

The above is intended only as a starting point．Obvious improvements would be to allow the use of the cursor keys by manipulating the string with LEFTS，MIDS and RIGHTS．The maximum length of the string field could also be set to prevent overwriting an existing display．

## Invisible characters

Most readers will no doubt be aware that character data may be directly poked to the screen. They will also have discovered that on occasions no effect is apparent. A screen character (normal mode) takes on the colour set in the corresponding location of the colour map (\$D800 on). If no character has been printed at this location nor a colour set since the last clear screen the adopted colour is that of the background. To see if the character is there, simply move the cursor to the location.

We can use this to good effect by making displays change quite quickly from BASIC with just pokes to the colour map. It is important to remember that on an INPUT even though the character cannot be seen it is still there and active.

## PET-64/64-PET

Commodore has maintained almost complete compatibility in the storage of programs on tape and to a lesser degree on disk. A tape prepared on any machine may be read by another. The 1541 disk drive uses an identical format to the 4040 unit. It can also read 2040 and 3040 formatted disks, but will corrupt these disks if it writes to them. You may also find that you get write problems on 4040 formatted disks. This compatibility does not mean to say that a program written for one machine will work directly on another.

A BASIC program saved on a PET can be loaded and run directly on the 64, whereas the reverse is not true. A word of warning to cassette users that the secondary address of 3 available on the 64 is not recognised by the PET. This is due to the different start address of bASIC on the two machines ( $\$ 0401$ on PET). Loading a program with a secondary address of zero will not allow it to load below the current start of BASIC. This means that on the 64 the load will be forced to $\$ 0801$. A 64 program will normally have a start greater than $\$ 0401$ and will go in above the start of BASIC and is not directly usable.

There are two ways to overcome the problem. The first is raise the start of BASIC on the PET to $\$ 0801$ by (BASIC 2 \& 4)

## POKE 41,08:POKE 2048, 0 :NEW

before loading. All pokes to the screen will have to be adjusted for the PET screen which starts at $\$ 8000$ (32768) and all pokes to the colour map removed. Defining the start of the screen with a variable and using offsets from this simplifies the conversion. The easiest way is to avoid using anything other than PRINT for output.

The second technique involves configuring the 64 to look like a PET BASIC at $\$ 0401$ and the screen in bank 2 at $\$ 8000$. The Programmer's Reference Guide (Chapter 3, 'Screen Memory') tells us how to relocate
the screen by setting bits 7 to 4 of \$D018 (53272), remembering to tell the screen editor where it has gone by setting HIBASE ( $\$ 0288 / 648$ ). The start of BASIC is lowered by setting TXTTAB and the top of memory set using MEMTOP and string storage with FRETOP. The following program if run will carry out the necessary changes. Syntax errors may result until the screen is cleared due to invisible characters.

10 POKE 51,000:POKE 55,000
20 POKE 52,128:POKE 56,128
30 POKE 56578,PEEK(56578) OR 3:POKE56576 , (PEEK (56576)AND252) OR 1
40 POKE 53272, (PEEK (53272)AND15) OR0:POKE 648,128:POKE 1024,0:POKE 44,4:NEW

There are some very sophisticated PET emulators on the market and even cross-assemblers for machine code and cross-compilers for BASIC. If a lot of your work is in an area where portability is important, it might well be worthwhile pursuing the matter.

## Load and run

Pressing SHIFT and RUN/STOP will load and run the first program on tape providing it is in BASIC. This may also be performed by:

POKE 631,131: POKE 198,1
The advantage of the second is that it can be used from within a program to avoid the problems associated with chaining if variables are not to be retained. Less well known is its use with disk. The format is:

LOAD"PROG" 8 : [Press SHIFT \& RUN/STOP]

## Disk bugs

When using sequential data files problems may be encountered if the same logical file number is used for both read and write operations. Typically, error 63 FILE EXISTS is reported. The only way to be sure is to use different numbers for input and output.

A less annoying feature is that null strings written to a data file are ignored on reading back. One way to overcome the problem is to always default null strings to a set value which is recognized on reading back. Alternatively, GET\# may be used to pick up returns and commas (a bit laborious).

## Appendices

## APPENDIX A: Storage of bASIC text

Standard CBM 64 tokens

| hex | dec |  | hex dec |  |  | hex dec |  |  | hex dec |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| \$20 | 32 | sp | \$40 | 64 | @ | \$80 | 128 | END | \$A6 | 166 | SPC( |
| \$21 | 33 | + | \$41 | 65 | A | \$81 | 129 | FOR | \$A7 | 167 | THEN |
| \$22 | 34 | " | \$42 | 66 | B | \$82 | 130 | NEXT | \$A8 | 168 | NOT |
| \$23 | 35 | \# | \$43 | 67 | C | \$83 | 131 | DATA | \$A9 | 169 | STEP |
| \$24 | 36 | \$ | \$44 | 68 | D | \$84 | 132 | INPUT\# | \$AA | 170 | + add |
| \$25 | 37 | \% | \$45 | 69 | E | \$85 | 133 | INPUT | \$AB | 171 | - minus |
| \$26 | 38 | \& | \$46 | 70 | F | \$86 | 134 | DIM | \$AC | 172 | * multi |
| \$27 | 39 |  | \$47 | 71 | G | \$87 | 135 | READ | \$AD | 173 | / div |
| \$28 | 40 | ( | \$48 | 72 | H | \$88 | 136 | LET | \$AE | 174 | $\uparrow$ power |
| \$29 | 41 | ) | \$49 | 73 | 1 | \$89 | 137 | GOTO | \$AF | 175 | AND |
| \$2A | 42 | * | \$4A | 74 | J | \$8A | 138 | RUN | \$B0 | 176 | OR |
| \$2B | 43 | + | \$4B | 75 | K | \$8B | 139 | IF | \$B1 | 177 | $>\mathrm{gt}$ |
| \$2C | 44 | , | \$4C | 76 | L | \$8C | 140 | RESTORE | \$B2 | 178 | $=\mathrm{eq}$ |
| \$2D | 45 | - | \$4D | 77 | M | \$8D | 141 | GOSUB | \$B3 | 179 | $<\mathrm{lt}$ |
| \$2E | 46 |  | \$4E | 78 | N | \$8E | 142 | RETURN | \$B4 | 180 | SGN |
| \$2F | 47 | 1 | \$4F | 79 | O | \$8F | 143 | REM | \$B5 | 181 | INT |
| \$30 | 48 | 0 | \$50 | 80 | P | \$90 | 144 | STOP | \$B6 | 182 | ABS |
| \$31 | 49 | 1 | \$51 | 81 | Q | \$91 | 145 | ON | \$B7 | 183 | USR |
| \$32 | 50 | 2 | \$52 | 82 | R | \$92 | 146 | WAIT | \$B8 | 184 | FRE |
| \$33 | 51 | 3 | \$53 | 83 | S | \$93 | 147 | LOAD | \$B9 | 185 | POS |
| \$34 | 52 | 4 | \$54 | 84 | T | \$94 | 148 | SAVE | \$BA | 186 | SQR |
| \$35 | 53 | 5 | \$55 | 85 | U | \$95 | 149 | VERIFY | \$BB | 187 | RND |
| \$36 | 54 | 6 | \$56 | 86 | V | \$96 | 150 | DEF | \$BC | 188 | LOG |
| \$37 | 55 | 7 | \$57 | 87 | W | \$97 | 151 | POKE | \$BD | 189 | EXP |
| \$38 | 56 | 8 | \$58 | 88 | X | \$98 | 152 | PRINT\# | \$BE | 190 | COS |
| \$39 | 57 | 9 | \$59 | 89 | Y | \$99 | 153 | PRINT | \$BF | 191 | SIN |
| \$3A | 58 | : | \$5A | 90 | Z | \$9A | 154 | CONT | \$C0 | 192 | TAN |
| \$3B | 59 | ; | \$5B | 91 | [ | \$9B | 155 | LIST | \$C1 | 193 | ATN |
| \$3C | 60 | < | \$5C | 92 | £ | \$9C | 156 | CLR | \$C2 | 194 | PEEK |
| \$3D | 61 | $=$ | \$5D | 93 | ] | \$9D | 157 | CMD | \$C3 | 195 | LEN |
| \$3E | 62 | > | \$5E | 94 | $\uparrow$ | \$9E | 158 | SYS | \$C4 | 196 | STR\$ |
| \$3F | 63 | ? | \$5F | 95 | $\leftarrow$ | \$9F | 159 | OPEN | \$C5 | 197 | VAL |
|  |  |  |  |  |  | \$A0 | 160 | CLOSE | \$C6 | 198 | ASC |
|  |  |  |  |  |  | \$A1 | 161 | GET | \$C7 | 199 | CHR\$ |
|  |  |  |  |  |  | \$A2 | 162 | NEW | \$C8 | 200 | LEFT\$ |
|  |  |  |  |  |  | \$A3 | 163 | TAB | \$C9 | 201 | RIGHT\$ |
|  |  |  |  |  |  | \$A4 | 164 | TO | \$CA | 202 | MID\$ |
|  |  |  |  |  |  | \$A5 | 165 | FN | \$CB | 203 | GO |

## Extended basic tokens

| hex | dec |  | hex | dec |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| \$CC | 204 | OFF | \$E6 | 230 | RESET |
| \$CD | 205 | KEY | \$E7 | 231 | CHAIN |
| \$CE | 206 | DOKE | \$E8 | 232 | LOMEM |
| \$CF | 207 | TEN | \$E9 | 233 | HIMEM |
| \$D0 | 208 | TWO | \$EA | 234 | INKEY\$ |
| \$D1 | 209 | HEX | \$EB | 235 | MEM |
| \$D2 | 210 | BIN | \$EC | 236 | APPEND |
| \$D3 | 211 | OLD | \$ED | 237 | TROFF |
| \$D4 | 212 | COLOUR | \$EE | 238 | unused |
| \$D5 | 213 | WRITE | \$EF | 239 | unused |
| \$D6 | 214 | CGOTO | \$F0 | 240 | unused |
| \$D7 | 215 | CGOSUB | \$F1 | 241 | unused |
| \$D8 | 216 | PLOT | \$F2 | 242 | unused |
| \$D9 | 217 | ENTER | \$F3 | 243 | unused |
| \$DA | 218 | DUMP | \$F4 | 244 | unused |
| \$DB | 219 | RENUM | \$F5 | 245 | unused |
| \$DC | 220 | DELETE | \$F6 | 246 | unused |
| \$DD | 221 | MERGE | \$F7 | 247 | DEEK |
| \$DE | 222 | CODER | \$F8 | 248 | unused |
| \$DF | 223 | AUTO | \$F9 | 249 | unused |
| \$E0 | 224 | PROC | \$FA | 250 | unused |
| \$E1 | 225 | DPROC | \$FB | 251 | unused |
| \$E2 | 256 | EPROC | \$FC | 252 | unused |
| \$E3 | 227 | POP | \$FD | 253 | unused |
| \$E4 | 228 | QUIT | \$FE | 254 | unused |
| \$E5 | 229 | TRACE |  |  |  |

## APPENDIX B: Hex to decimal and decimal to hex converter

| hex | decimal |  | hex | decimal |  | hex | decimal |  | hex | decim |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | low | high |  | low | high |  |  | high |  |  | , |
|  | 0 | 0 | \$29 | 41 | 10496 | \$52 | 82 | 20992 | B | 23 | 1488 |
|  |  | 256 | 2 A | 42 | 10752 | \$53 | 83 | 21248 | \$7C | 124 | 31744 |
| \$02 | 2 | 512 | 2 B | 43 | 11008 | \$54 | 84 | 21504 | \$7D | 125 | 32000 |
| \$03 | 3 | 768 | \$2C | 44 | 11264 | \$55 | 85 | 21760 | \$7E | 126 | 32256 |
| \$04 | 4 | 1024 | \$2D | 45 | 11520 | \$56 | 86 | 22016 | \$7F | 127 | 32512 |
| \$05 | 5 | 1280 | \$2E | 46 | 11776 | \$57 | 87 | 22272 | \$80 | 128 | 32768 |
| \$06 | 6 | 1536 | \$2F | 47 | 12032 | \$58 | 88 | 22528 | \$81 | 129 | 33024 |
| \$07 | 7 | 1792 | \$30 | 48 | 12288 | \$59 | 89 | 22784 | \$82 | 130 | 33280 |
| \$08 | 8 | 2048 | \$31 | 49 | 12544 | \$5A | 90 | 23040 | \$83 | 131 | 33536 |
| \$09 | 9 | 2304 | \$32 | 50 | 12800 | \$5B | 91 | 23296 | \$84 | 132 | 33792 |
| A | 10 | 2560 | \$33 | 51 | 13056 | \$5C | 92 | 23552 | \$85 | 133 | 34048 |
| B | 11 | 2816 | \$34 | 52 | 13312 | \$5D | 93 | 23808 | \$86 | 134 | 34304 |
| \$0C | 12 | 3072 | \$35 | 53 | 13568 | \$5E | 94 | 24064 | \$87 | 135 | 34560 |
| \$0D | 13 | 3328 | \$36 | 54 | 13824 | \$5F | 95 | 24320 | \$88 | 136 | 34816 |
| \$0E | 14 | 3584 | \$37 | 55 | 14080 | \$60 | 96 | 24576 | \$89 | 137 | 35072 |
| \$0F | 15 | 3840 | \$38 | 56 | 14336 | \$61 | 97 | 24832 | \$8A | 138 | 35328 |
| \$10 | 16 | 4096 | \$39 | 57 | 14592 | \$62 | 98 | 25088 | \$8B | 139 | 35584 |
| 1 | 17 | 4352 | \$3A | 58 | 14848 | \$63 | 99 | 25344 | \$8C | 140 | 35840 |
| \$12 | 18 | 4608 | \$3B | 59 | 15104 | \$64 | 100 | 25600 | \$8D | 141 | 36096 |
| 13 | 19 | 4864 | \$3C | 60 | 15360 | \$65 | 101 | 25856 | \$8E | 142 | 36352 |
| 14 | 20 | 5120 | \$3D | 61 | 15616 | \$66 | 102 | 26112 | \$8F | 143 | 36608 |
| 5 | 21 | 5376 | \$3E | 62 | 15872 | \$67 | 103 | 26368 | \$90 | 144 | 36864 |
| \$16 | 22 | 5632 | \$3F | 63 | 16128 | \$68 | 104 | 26624 | \$91 | 145 | 37120 |
| \$17 | 23 | 5888 | \$40 | 64 | 16384 | \$69 | 105 | 26880 | \$92 | 146 | 7376 |
| \$18 | 24 | 6144 | \$41 | 65 | 16640 | \$6A | 106 | 27136 | \$93 | 147 | 37632 |
| \$19 | 25 | 6400 | \$42 | 66 | 16896 | \$6B | 107 | 27392 | \$94 | 148 | 37888 |
| \$1A | 26 | 6656 | \$43 | 67 | 17152 | \$6C | 108 | 27648 | \$95 | 149 | 38144 |
| \$1B | 27 | 6912 | \$44 | 68 | 17408 | \$6D | 109 | 27904 | \$96 | 150 | 38400 |
| \$1C | 28 | 7168 | \$45 | 69 | 17664 | \$6E | 110 | 28160 | \$97 | 151 | 38656 |
| \$1D | 29 | 7424 | \$46 | 70 | 17920 | \$6F | 111 | 28416 | \$98 | 152 | 38912 |
| \$1E | 30 | 7680 | \$47 | 71 | 18176 | \$70 | 112 | 28672 | \$99 | 153 | 39168 |
| \$1F | 31 | 7936 | \$48 | 72 | 18432 | \$71 | 113 | 28928 | \$9A | 154 | 39424 |
| \$20 | 32 | 8192 | \$49 | 73 | 18688 | \$72 | 114 | 29184 | \$9B | 155 | 39680 |
| \$21 | 33 | 8448 | \$4A | 74 | 18944 | \$73 | 115 | 29440 | \$9C | 156 | 39936 |
| 22 | 34 | 8704 | \$4B | 75 | 19200 | \$74 | 116 | 29696 | \$9D | 157 | 40192 |
| \$23 | 35 | 8960 | \$4C | 76 | 19456 | \$75 | 117 | 29952 | \$9E | 158 | 40448 |
| \$24 | 36 | 9216 | \$4D | 77 | 19712 | \$76 | 118 | 30208 | \$9F | 159 | 40704 |
| \$25 | 37 | 9472 | \$4E | 78 | 19968 | \$77 | 119 | 30464 | \$A0 | 160 | 40960 |
| \$26 | 38 | 9728 | \$4F | 79 | 20224 | \$78 | 120 | 30720 | \$A1 | 161 | 41216 |
| \$27 | 39 | 9984 | \$50 | 80 | 20480 | \$79 | 121 | 30976 | \$A2 | 162 | 41472 |
| \$28 | 40 | 10240 | \$51 | 81 | 20736 | \$7A | 122 | 31232 | \$A3 | 163 | 41728 |


| hex | decimal |  | hex | decima |  | hex | cima |  | hex | decimal |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | low | high |  | low | high |  | low | high |  | low | high |
| \$A4 | 164 | 41984 | \$BB | 187 | 47872 | D2 | 210 | 53760 | \$E9 | 233 | 59648 |
| \$A5 | 165 | 42240 | \$BC | 188 | 48128 | \$D3 | 211 | 54016 | \$EA | 234 | 59904 |
| \$A6 | 166 | 42496 | \$BD | 189 | 4838 | \$D4 | 212 | 54272 | \$EB | 235 | 60160 |
| \$A7 | 167 | 42752 | \$BE | 190 | 48640 | \$D5 | 213 | 54528 | \$EC | 236 | 60416 |
| \$A8 | 168 | 43008 | \$BF | 191 | 48896 | \$D6 | 214 | 54784 | \$ED | 237 | 60672 |
| \$A9 | 169 | 43264 | \$C0 | 192 | 49152 | \$D7 | 215 | 55040 | \$EE | 238 | 60928 |
| \$AA | 170 | 43520 | \$C1 | 193 | 49408 | \$D8 | 216 | 55296 | \$EF | 239 | 61184 |
| \$AB | 171 | 43776 | \$C2 | 194 | 49664 | \$D9 | 217 | 55552 | \$F0 | 240 | 61440 |
| \$AC | 172 | 44032 | \$C3 | 195 | 49920 | \$DA | 218 | 55808 | \$F1 | 241 | 61696 |
| \$AD | 173 | 44288 | \$C4 | 196 | 50176 | \$DB | 219 | 56064 | \$F2 | 242 | 61952 |
| \$AE | 174 | 44544 | \$C5 | 197 | 50432 | \$DC | 220 | 56320 | \$F3 | 243 | 62208 |
| \$AF | 175 | 44800 | \$C6 | 198 | 50688 | \$DD | 221 | 56576 | \$FA | 244 | 62464 |
| \$B0 | 176 | 45056 | \$C7 | 199 | 50944 | \$DE | 222 | 56832 | \$F5 | 245 | 62720 |
| \$B1 | 177 | 45312 | \$C8 | 200 | 51200 | \$DF | 223 | 57088 | \$F6 | 246 | 62976 |
| \$B2 | 178 | 45568 | \$C9 | 201 | 51456 | \$ED | 224 | 57344 | \$F7 | 247 | 63232 |
| \$B3 | 179 | 45824 | \$CA | 202 | 51712 | \$E1 | 225 | 57600 | \$F8 | 248 | 63488 |
| \$B4 | 180 | 46080 | \$CB | 203 | 5196 | \$E2 | 226 | 57856 | \$F9 | 249 | 63744 |
| \$B5 | 181 | 46336 | \$CC | 204 | 5222 | \$E3 | 227 | 58112 | \$FA | 250 | 64000 |
| \$B6 | 182 | 46592 | \$CD | 205 | 52480 | \$E4 | 228 | 58368 | \$FB | 251 | 64256 |
| \$B7 | 183 | 46848 | \$CE | 206 | 52736 | \$E5 | 229 | 58624 | \$FC | 252 | 64512 |
| \$B8 | 184 | 47104 | \$CF | 207 | 52992 | \$E6 | 230 | 58880 | \$FD | 253 | 64768 |
| \$B9 | 185 | 47360 | \$D0 | 208 | 53248 | \$E7 | 231 | 59136 | \$FE | 254 | 65024 |
| \$BA | 186 | 47616 | \$D1 | 209 | 53504 | \$E8 | 232 | 59392 | \$FF | 255 | 65280 |

## APPENDIX C: Machine code mnemonics and hex values

## 6510 OP-CODES

The tables below are a quick reference guide only and for more detailed information a 6502 assembler book should be consulted.

The tables should be read row then column. If in doubt, remember LDA immediate mode is \$A9. The following abbreviations have been used:
\# - immediate mode
\$-absolute address
Z-zero page
I - indirect address
$A=$ accumulator
$X=X$ inde $x$ register
$\mathrm{Y}=\mathrm{Y}$ index register

| 0 | 1 | 2 | 4 | 5 | 6 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 BRK | ORA ( $1, X$ ) |  |  | ORA Z | ASL Z |
| 1 BPL | ORA (I), Y |  |  | ORA Z, X | ASL Z, X |
| 2 JSR | AND (I, X) |  | BIT Z | AND Z | ROL Z |
| 3 BMI | AND (1), Y |  |  | AND $\mathrm{Z}, \mathrm{X}$ | ROL Z, X |
| 4 RTI | EOR (I, X) |  |  | EOR Z | LSR Z |
| 5 BVC | EOR (I), Y |  |  | EOR Z, X | LSR Z, X |
| 6 RTS | ADC ( $1, \mathrm{X}$ ) |  |  | ADC Z | ROR Z |
| 7 BVS | ADC (I), Y |  |  | ADC Z, X | ROR Z, X |
| 8 | STA (I,X) |  | STY Z | STA Z | STX Z |
| 9 BCC | STA ( 1 , Y |  | STY Z, X | STA Z, X | STX Z, Y |
| A LDY\# | LDA ( $1, \mathrm{X}$ ) | LDX \# | LDY Z | LDA Z | LDXZ |
| B BCS | LDA (I), Y |  | LDY Z, X | LDA Z, X | LDX Z, Y |
| C CPY\# | CMP (I,X) |  | CPY Z | CMP Z | DEC Z |
| D BNE | CMP (I), Y |  |  | CMP Z, X | DEC Z, X |
| E CPX \# | SBC (1,X) |  | CPX Z | SBC Z | INC Z |
| F BEQ | SBC (1), Y |  |  | SBC Z, X | INC Z, X |


| 8 | 9 | A | C | D | E |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 PHP | ORA \# | ASL A |  | ORA \$ | ASL \$ |
| 1 CLC | ORA \$, Y |  |  | ORA \$ , X | ASL \$, X |
| 2 PLP | AND \# | ROLA | BIT \$ | AND \$ | ROL \$ |
| 3 SEC | AND \$, Y |  |  | AND \$, X | ROL \$, X |
| 4 PHA | EOR \# | LSR A | JMP \$ | EOR \$ | LSR \$ |
| 5 CLI | EOR \$, Y |  |  | EOR \$ , X | LSR \$, X |
| 6 PLA | ADC \# | ROR A | JMP I | ADC \$ | ROR \$ |
| 7 SEI | ADC \$, Y |  |  | ADC \$, X | ROR \$, X |
| 8 DEY |  | TXA | STY \$ | STA \$ | STX \$ |
| 9 TYA | STA \$, Y | TXS |  | STA \$, X |  |
| A TAY | LDA \# | TAX | LDY \$ | LDA \$ | LDX \$ |
| B CLV | LDA \$, Y | TSX | LDY \$, X | LDA \$, X | LDX \$, Y |
| CINY | CMP \# | DEX | CPY \$ | CMP \$ | DEC \$ |
| D CLD | CMP \$, Y |  |  | CMP \$, X | DEC \$, X |
| E INX | SBC \# | NOP | CPX \$ | SBC \$ | INC \$ |
| F SED | SBC \$, Y |  |  | SBC \$, X | INC \$, X |

## APPENDIX D: BASIC loader for Supermon

This is the BASIC program to produce Jim Butterfield's Supermon monitor. Type it in and SAVE it before running. There are lots of numbers so it is easy to make a mistake.

We have put in some checksums to help isolate errors. Once loaded correctly you will be able to save the machine code version from the monitor itself. Instructions for using Supermon are given in Appendix E.

Supermon normally loads to the top of BASIC memory. We have modified it to sit at $\$ 0000$ on to allow you to enter the code for the utility. Once you have it up and running, you can use Supermon to modify itself to sit anywhere. There are addresses which require changing so we have included a relocater program after the loader.

The Loader

```
10 A=49152:C=0
20 READB:IFB=-1 THEN40
30 POKEA, B:A=A+1:C=C+B:GOT020
4 0 ~ I F C = 2 7 9 1 4 T H E N 6 0 ~
50 PRINT"DATA ERROR IN 1000 - 1300":END
6 0 ~ C = 0
70 READB: IFB=-1 THEN90
80 POKEA, B:A=A+1:C=C+B:GOTO78
90 IFC=26078THEN110
100 PRINT"DATA ERROR IN 1310 - 1600":END
110 C=0
120 READB:IFB=-1 THEN140
130 POKEA, B:A=A+1:C=C+B:GOTO1 20
140 IFC=26897THEN160
150 PRINT"DATA ERROR IN 1610 - 1900":END
160 C=0
170 READB:IFB}=-1\mathrm{ THEN190
180 POKEA, B:A=A+1:C=C+B:GOTO170
190 I FC=28055THEN210
200 PRINT"DATA ERROR IN 1910 - 2200":END
210 C=0
220 READB:IFB=-1 THEN240
230 POKEA, B:A=A+1:C=C+B:GOT0220
240 IFC=25343THEN260
250 PRINT"DATA ERROR IN 2210 - 2500":END
260 C=0
270 READB:IFB=-1THEN290
280 POKEA,B:A=A+1:C=C+B:GOT0270
290 IFC=25432THEN310
300 PRINT"DATA ERROR IN 2510 - 2800":END
310 C=0
320 READB: I FB=-1 THEN346
330 POKEA,B:A=A+1:C=C+B:GOT0320
340 IFC=27324THEN360
350 PRINT"DATA ERROR IN 2810 - 3100":END
```

360
370
380 POKEA, $B: A=A+1: C=C+B: G O T 0370$
390 IFC=25335THEN410
400 PRINT"DATA ERROR IN 3110 - 3400": END
$410 \mathrm{C}=0$
420 READB: IFB $=-1$ THEN440
430 POKEA, $B: A=A+1: C=C+B: G O T 0420$
440 IFC=28057THEN460
450 PRINT"DATA ERROR IN 3410 - 3700": END
460 C=0
470 READB: I $F B=-1$ THEN490
480 POKEA, $B: A=A+1: C=C+B: G O T 0470$
490 IFC=20514THEN510
500 PRINT"DATA ERROR IN 3710-4008": END
$510 \mathrm{C}=0$
520 READB: IFB $=-1$ THEN540
536 POKEA, $B: A=A+1: C=C+B: G O T 0520$
540 IFC=22061 THEN560
550 PRINT"DATA ERROR IN $4010-4290^{\circ}:$ END
560 PRINT"DATA CORRECT":PRINT
570 PRINT"SYS49152 TO USE"
580 END
1800 DATA76, 233, 192, 255, 0, 0, 255
1010 DATA255,0,0,255,255,0,0
1026 DATA255, 255, 0, 0, 255, 255,0
1030 DATA日, 255, 255, 0, 0, 255, 255
1040 DATA0, $0,255,255,0,0,255$
1050 DATA255, $0,0,255,255,0,0$
1060 DATA $255,255,0,0,255,255,0$
1070 DATA0, 255, 255, 0, 0, 255, 255
1080 DATA0, $0,255,255,0,0,255$
1090 DATA255, $0,0,255,255,0,0$
1100 DATA $255,255,0,0,255,255,0$
1110 DATA0, 255, 255,0,0,255,255
1120 DATA0, $0,255,255,0,0,255$
1130 DATA255, $0,0,255,255,0,0$
1140 DATA $255,255,0,0,255,255,0$
1150 DATA0, 255, 255, 0, 0, 255, 255
1160 DATA $0,0,255,255,0,0,255$
1170 DATA $255,0,0,255,255,0,0$
1186 DATA255, 255, $128,0,255,255,8$
1190 DATA0, 255, 255, 0, $0,255,255$
1200 DATA $0,0,255,255,0,0,255$
1210 DATA255, $0,0,255,255,0,0$
1220 DATA $255,255,0,0,255,255,0$

1230
1240 DATAQ,0,255,255,0,0,255
1250 DATA255,0,0,255,255,0,0
1260 DATA255,255,0,0,255,255,0
1270 DATA0,255,255,0,0,255,255
1280 DATA0,0,255,255,0,0,255
1290 DATA255,0,0,255,255,0,0
1300 DATA255,255,0,0,255,255,0,-1
1310 DATA0,255,255,0,0,255,255
1320 DATA0,0,255,255,0,0,255
1330 DATA255,0,169,160,133,56,173
1340 DATA230,200,141,22,3,173,231
1350 DATA200,141,23,3,169,128,32
1360 DATA $144,255,0,216,104,141,62$
1370 DATA2,104,141,61,2,104,141
1380 DATA60,2,104,141,59,2,104
1390 DATA $170,164,168,56,138,233,2$
1400 DATA $141,58,2,152,233,0,141$
1410 DATA57,2,186,142,63,2,32
1426 DATA87,198,162,66,169,42,32
1430 DATA87,195,169,82,208,52,230
1440 DATA $193,208,6,230,194,268,2$
1450 DATA $230,38,96,32,207,255,201$
1460 DATA $13,208,248,104,104,169,154$
1470 DATA32,210,255,169,0,133,38
1480 DATA $162,13,169,46,32,87,195$
1490 DATA $169,159,32,210,255,32,62$
1500 DATA193,201,46,240,249,201,32
1510 DATA240,245,162,14,221,183,200
1520 DATA208,12,138,10,170,189,199
1530 DATA200,72,189,198,200,72,96
1540 DATA202,16,236,76,237,195,165
1550 DATA193,141,58,2,165,194,141
1560 DATA57,2,96,169,8,133,29
1570 DATA $160,8,32,84,198,177,193$
1580 DATA $32,72,195,32,51,193,198$
1590 DATA29,208,241,96,32,136,195
1600 DATA $144,11,162,0,129,193,193,-1$
1610 DATA193,240,3,76,237,195,32
1620 DATA51,193,198,29,96,169,59
1630 DATA133,193,169,2,133,194,169
1640 DATA5,96,152,72,32,87,198
1650 DATA104,162,46,76,87,195,169
1660 DATA154,32,216,255,162,0,189
1670 DATA $234,200,32,210,255,232,224$
1680 DATA22,208,245,160,59,32,194

| 1690 | DATA193,173,57,2,32,72,195 |
| :---: | :---: |
| 1700 | DATA173,58,2,32,72,195,32 |
| 1710 | DATA183,193,32,141,193,240,92 |
| 1720 | DATA32,62,193,32,121,195,144 |
| 1730 | DATA51, 32,105,195,32,62,193 |
| 1740 | DATA32,121,195,144,40,32,105 |
| 1750 | DATA195,169,154,32,210,255,32 |
| 1760 | DATA $225,255,240,60,166,38,208$ |
| 1770 | DATA56,165,195,197,193,165,196 |
| 1780 | DATA $229,194,144,46,160,58,32$ |
| 1790 | DATA194,193,32,65,195,32,139 |
| 1800 | DATA193,240,224,76,237,195,32 |
| 1810 | DATA1 $21,195,144,3,32,128,193$ |
| 1820 | DATA32,183,193,208,7,32,121 |
| 1830 | DATA195,144,235,169,8,133,29 |
| 1840 | DATA32,62,193,32,161,193,208 |
| 1850 | DATA248,76,71,193,32,207,255 |
| 1860 | DATA201, 13,240,12,201,32,208 |
| 1870 | DATA209, 32, 121,195,144,3,32 |
| 1880 | DATA128,193,169,154,32,210,255 |
| 1890 | DATA $174,63,2,154,120,173,57$ |
| 1900 | DATA2, $72,173,58,2,72,173,-1$ |
| 1910 | DATA59,2,72,173,60,2,174 |
| 1920 | DATA $61,2,172,62,2,64,169$ |
| 1930 | DATA $154,32,210,255,174,63,2$ |
| 1940 | DATA $154,108,2,160,160,1,132$ |
| 1950 | DATA186,132,185,136,132,183,132 |
| 1960 | DATA144, 132,147,169,64,133,187 |
| 1970 | DATA169,2,133,188,32,207,255 |
| 1980 | DATA201,32,240,249,201, 13,240 |
| 1990 | DATA56, 201, 34, 208, $20,32,207$ |
| 2000 | DATA255, $201,34,240,16,201,13$ |
| 2010 | DATA $240,41,145,187,230,183,200$ |
| 2020 | DATA192,16,208,236,76,237,195 |
| 2030 | DATA32, 207, 255, 201, 13, 240, 22 |
| 2040 | DATA $201,44,208,220,32,136,195$ |
| 2050 | DATA41, 15,240,233,201,3,240 |
| 2060 | DATA229,133,186,32,207,255,201 |
| 2070 | DATA $13,96,108,48,3,108,50$ |
| 2080 | DATA3, 32,150,194,208,212,169 |
| 2090 | DATA154,32,210,255,169,0,32 |
| 2100 | DATA239,194,165,144,41,16,208 |
| 2110 | DATA196,76,71,193,32,150,194 |
| 2120 | DATA201,44,208,186,32,121,195 |
| 2130 | DATA32, 105,195,32,207, 255,201 |
| 2140 | DATA $44,208,173,32,121,195,165$ |


2610 DATA51,193,208,235,32,40,196
2620 DATA 24,165,30,101,195,133,195
2630 DATA152,101,196,133,196,32,12
2640 DATA196,166,38,208,61,161,193
2650 DATA1 29,195,32,40,196,176,52
2660 DATA $32,184,195,32,187,195,76$
2670 DATA125,196,32,212,195,32,105
2680 DATA195,32,229,195,32,165,195
2690 DATA32,62,193,32,136,195,144
2790 DATA20,133,29,166,38,208,17
2710 DATA32,47,196,144,12,165,29
2720 DATA129,193,32,51,193,208,238
2730 DATA76,237,195,76,71,193,32
2740 DATA212,195,32,105,195,32,229
2750 DATA195,32,105,195,32,62,193
2760 DATA $162,0,32,62,193,201,39$
2770 DATA208,20,32,62,193,157,16
2780 DATA2,232,32,267,255,201,13
2790 DATA240,34,224,32,208,241,240
2800 DATA28,142,0,1,32,143,195,-1
2810 DATA144,198,157,16,2,232,32
2820 DATA207,255,201,13,240,9,32
2830 DATA136,195,144,182,224,32,208
2840 DATA236,134,28,169,154,32,210
2850 DATA $255,32,87,198,162,0,160$
2866 DATA0,177,193,221,16,2,208
2870 DATA12,200,232,228,28,208,243
2880 DATA32,65,195,32,84,198,32
2890 DATA51,193,166,38,208,141,32
2900 DATA47,196,176,221,76,71,193
2910 DATA $32,212,195,133,32,165,194$
2920 DATA133,33,162,0,134,40,169
2930 DATA147,32,210,255,q69,154,32
2940 DATA210,255,169,22,133,29,32
2950 DATA106,197,32,202,197,133,193
2960 DATA1 32,194,198,29,208,242,169
2970 DATA145,32,210,255,76,71,193
2980 DATA $160,44,32,194,193,32,84$
2990 DATA198,32,65,195,32,84,198
3000 DATA162,0,161,193,32,217,197
3010 DATA $72,32,31,198,104,32,53$
3020 DATA198,162,6,224,3,208,18
3036 DATA $64,31,240,14,165,42,201$
3040 DATA232,177,193,176,28,32,194
3050 DATA197,136,208,242,6,42,144
3060 DATA14,189,42,200,32,165,198

| 3070 | DATA189, $48,200,240,3,32,165$ |
| :---: | :---: |
| 3080 | DATA198,202,208,213,96,32,205 |
| 3090 | DATA197, 170, 232, 208, 1, 200, 152 |
| 3180 | DATA $32,194,197,138,134,28,32,-1$ |
| 3110 | DATA $2,195,166,28,96,165,31$ |
| 3120 | DATA56,164,194,170,16,1,136 |
| 3130 | DATA $101,193,144,1,200,96,168$ |
| 3140 | DATA $74,144,11,74,176,23,201$ |
| 3150 | DATA34, 240,19,41,7,9,128 |
| 3160 | DATA $74,178,189,217,199,176,4$ |
| 3170 | DATA $74,74,74,74,41,15,208$ |
| 3180 | DATA4,160,128,169,0,170,189 |
| 3190 | DATA29,200, 133,42,41,3,133 |
| 3200 | DATA31, 152,41,143,170,152,160 |
| 3210 | DATA3, 224,138,240,11,74,144 |
| 3220 | DATA8, 74,74,9,32,136,208 |
| 3230 | DATA $250,200,136,208,242,96,177$ |
| 3240 | DATA $193,32,194,197,162,1,32$ |
| 3250 | DATA254,195,196,31,200,144,241 |
| 3260 | DATA $162,3,192,4,144,242,96$ |
| 3270 | DATA $168,185,55,200,133,40,185$ |
| 3280 | DATA $119,200,133,41,169,0,160$ |
| 3290 | DATA5, 6, 41, 38,40,42,136 |
| 3300 | DATA $208,248,105,63,32,210,255$ |
| 3310 | DATA $202,208,236,169,32,44,169$ |
| 3320 | DATA13,76,210,255,32,212,195 |
| 3330 | DATA32,105,195,32,229,195,32 |
| 3340 | DATA105,195,162,0,134,40,169 |
| 3350 | DATA154,32,210,255,32,87,198 |
| 3360 | DATA32,114,197,32,202,197,133 |
| 3370 | DATA $193,132,194,32,225,255,240$ |
| 3380 | DATA5, 32,47,196,176,233,76 |
| 3390 | DATA $1,193,32,212,195,169,3$ |
| 3400 | DATA $133,29,32,62,193,32,161,-1$ |
| 3410 | DATA193,208,248,165,32,133,193 |
| 3420 | DATA $165,33,133,194,76,70,197$ |
| 3430 | DATA197,40,240,3,32,210,255 |
| 3440 | DATA96,32,212,195,32,105,195 |
| 3450 | DATA142,17,2,162,3,32,204 |
| 3460 | DATA195,72,202,208,249,162,3 |
| 3470 | DATA $184,56,233,63,160,5,74$ |
| 3480 | DATA $110,17,2,110,16,2,136$ |
| 3490 | DATA $208,246,202,268,237,162,2$ |
| 3500 | DATA32, $207,255,201,13,240,30$ |
| 3510 | DATA201,32,240,245,32,208,199 |
| 3520 | DATA $176,15,32,156,195,164,193$ |

3530 DATA $132,194,133,193,169,48,157$
3540 DATA16,2,232,157,16,2,232
3550 DATA208,219,134,40,162,0,134
3560 DATA $38,240,4,230,38,240,117$
3570 DATA $162,0,134,29,165,38,32$
3580 DATA217,197,166,42,134,41,170
3590 DATA188,55,200,189,119,200,32
3600 DATA185,199,208,227,162,6,224
3610 DATA3,208,25,164,31,240,21
3620 DATA $165,42,201,232,169,48,176$
3630 DATA33,32,191,199,208,204,32
3640 DATA193,199,208,199,136,208,235
3650 DATA6,42,144,11,188,48,200
3660 DATA189,42,200,32,185,199,208
3670 DATA181,202,208,209,240,10,32
3680 DATA $184,199,268,171,32,184,199$
3690 DATA208,166,165,40,197,29,208
3700 DATA160,32,105,195,164,31,240,-1
3710 DATA40,165,41,201,157,208,26
3720 DATA32,28,196,144,10,152,208
3730 DATA4,165,30,16,10,76,237
3740 DATA195,200,208,250,165,30,16
3750 DATA246,164,31,208,3,185,194
3760 DATA0,145,193,136,208,248,165
3770 DATA38,145,193,32,202,197,133
3780 DATA193,132,194,169,154,32,210
3790 DATA255,160,65,32,194,193,32
3800 DATA84,198,32,65,195,32,84
3810 DATA198,169,159,32,210,255,76
3820 DATA176,198,168,32,191,199,208
3830 DATA17,152,240,14,134,28,166
3840 DATA29,221,16,2,8,232,134
3850 DATA29,166,28,40,96,201,48
3850 DATA144,3,201,71,96,56,96
3870 DATA $64,2,69,3,208,8,64$
3880 DATA9,48,34,69,51,208,8
3890 DATA64,9,64,2,69,51,208
3900 DATA8,64,9,64,2,69,179
3910 DATA208,8,64,9,0,34,68
3920 DATA51,208,140,68,0,17,34
3930 DATA $68,51,208,140,68,154,16$
3940 DATA34,68,51,208,8,64,9
3950 DATA $16,34,68,51,208,8,64$
3960 DATA9,98,19,120,169,0,33
3970 DATA129,130,0,0,89,77,145
3980 DATA $146,134,74,133,157,44,41$

```
3990 DATA44,35,40,36,89,0,88
4000 DATA36, 36,0,28,138, 28,35,-1
4010 DATA93,139,27,161,157,138,29
4 0 2 0 ~ D A T A 3 5 , 1 5 7 , 1 3 9 , 2 9 , 1 6 1 , 0 , 4 1
4 0 3 0 ~ D A T A 2 5 , 1 7 4 , 1 0 5 , 1 6 8 , 2 5 , 3 5 , 3 6 ~
4 0 4 0 \text { DATA83,27,35,36,83,25,161}
4050 DATA0, 26,91,91,165,105,36
4060 DATA36,174,174,168,173,41,0
4 0 7 0 \text { DATA1 24,0,21,156,109,156,165}
4 0 8 0 ~ D A T A 1 0 5 , 4 1 , 8 3 , 1 3 2 , 1 9 , 5 2 , 1 7
4 0 9 0 ~ D A T A 1 6 5 , 1 0 5 , 3 5 , 1 6 0 , 2 1 6 , 9 8 , 9 0 ~
4100 DATA72, 38,98,148,136,84,68
4110 DATA200,84,104,68,232,148,0
4 1 2 0 \text { DATA1 80,8,132,116,180,40,110}
4130 DATA116,244,204,74,114,242,164
4140 DATA1 38,0,170,162,162,116,116
4150 DATA116,114,68,104,178,50,178
4168 DATA0,34,0,26,26,38,38
4170 DATA114,114,136,200,196,202,38
4 1 8 0 ~ D A T A 7 2 , 6 8 , 6 8 , 1 6 2 , 2 0 0 , 5 8 , 5 9
4190 DATA82,77,71,88,76,83,84
4 2 0 8 \text { DATA78,72,68,80,44,65,66}
4210 DATA194,53,194,204,193,247,193
4 2 2 0 \text { DATA86,194,137,194,244,194,12}
4230 DATA195,62,196,146,196,192,196
4 2 4 0 \text { DATA56,197,91,198,138,198,172}
4 2 5 0 ~ D A T A 1 9 8 , 7 0 , 1 9 3 , 2 5 5 , 1 9 2 , 2 3 7 , 1 9 2 ,
4 2 6 0 \text { DATA1 3,32,32,32, 80,67,32}
4270 DATA32,83,82,32,65,67,32
4 2 8 0 \text { DATA88,82,32,89,82,32,83}
4290 DATA80,108,239,255,254,221, 222,-1
```


## Relocation

If you need to relocate Supermon to another area of memory, the following program will assist. First load in SUPERMON and use the transFER option to copy it to the desired location. Now run the program below and it changes the appropriate locations so that it will run at its new address.

```
10 PRINT"[CLS][REU]SUPERMON RELOCATION":
PRINT:PRINT
20 INPUT"NEW START ADDRESS - DECIMAL";AD
30 IFAD>40960THENPOKEAD+234,160:GOT090
40 PRINT"DO YOU WANT IT PROTECTED FROM BASIC"
50 PRINT"Y OR N"
```

60 GETA $:$ :IFA $\$\rangle$ "Y"ANDA $\langle<>" N " T H E N 60$
70 IF A $\ddagger=$ "N"THEN POKE AD,160:GOT090
80 POKEAD +234 , INT ( (AD $/ 256+.5)$ )
90 READ OF:IF OF=-1 THEN130
100 READ DA: DA=DA+AD:D2=INT (DA/256):D1=D A-(D2*256)
110 POKE OF+AD,D1:POKE OF+AD+1,D2
120 GOTO99
130 END
$140 \mathrm{CO}=0$
150 READ A:IF A=-1 THEN170
160 CO=CO+A: GOTO150
170 IF CO=451387THENPRINT "[REU]DATA OK" : END
180 PRINT"[REU]DATA INCORRECT"
190 REM $* * * * *$ START OF DATA $* * * * *$
200 DATA $, 233,238,2278,244,2279,294,1623$ ,301,855,341,855,349,318,370,2247
210 DATA $374,2246,382,1005,402,1620,407,8$
$40,410,307,418,904,431,1005,434,307$
228 DATA $453,1623,459,855,469,2282,482,45$ $0,488,840,494,840,497,439,500,397$
230 DATA505,318,508,889,513,873,516,318, $519,889,524,873,553,450,556,833,559,395$
240 DATA564,1065,567,889,572,384,575,439 ,580,889,589,318,592,417,597,327,611,889 256 DATA $16,384,719,1005,733,904,758,662$ ,770,751,779,327,782,662,789,889
268 DATA $792,873,862,889,813,873,828,754$, 831,327,836,840,846,864,853,864
270 DATA890,904,897,904,909,318,925,943, 934,318,937,943,973,318,986,972
280 DATA989,911,992,892,998,318,1001,889 , 1020,327,1023,1620,1069,1075
290 DATA1088,980,1091,873,1094,997,1097,
1036,1100,997,1103,1071,1106,873
306 DATA1115,1064,1124,1029,1127,367,113
2,1064,1147,1036,1158,1064
310 DATA1163,952,1166,955,1169,1149,1172 ,980,1175,873,1178,997,1181,873
320 DATA1 184,318,1187,904,1198,1071, 1207 ,307,1212,1005,1215,327,1218,980
330 DATA1 221, 873, 1224,997, 1227,873, 1230, 318,1235,318,1242,318,1265,911
340 DATA $1281,904,1297,1623,1317,833,1320$ ,1620,1323,307,1335,327,1338,980

```
350 DATA1365,1386,1368,1482,1384,327,138
9,450,1392,1620,1395,833,1398,1620
360 DATA1405,1497,1409,1567,1413,1589,14
34,1474,1444,2090,1447,1701
370 DATA1450,2096,1455,1701,1462,1485,14
71,1474,1477,840,1515,2009,1533,2077
380 DATA1570,1474,1575,1022,1591,2103,15
96,2167,1629,980,1632,873,1635,997
390 DATA1638,873,1650,1623,1653,1394,165
6,1482,1668,1071,1673,327,1676,980
400 DATA1683,318,1686,417,1699,1350,1710
,980,1713,873,1721,972,1762,2000
416 DATA1767,924,1866,985,1814,2103,1817
,2167,1820,1977,1843,1983,1848,1985
420 DATA1860,2096,1863,2090,1866,1977,18
76,1976,1881,1976,1892,873,1905,1052
430 DATA1917,1005,1943,1482,1957,450,196
0,1620,1963,833,1966,1620
446 DATA1974,1712,1978,1983,1330,1071
450 DATA362,2231,916,318,1806,1497
460 DATA2246,578,2248,565,2250,460
470 DATA2252,503,2254,598,2256,649
480 DATA2258,756,2260,780,2262,1086
490 DATA2264,1170,2266,1216,2268,1336
500 DATA2270,1627,2272,1674,2274,1708
510 DATA2276,326,2278,255,2280,237,-1
```

APPENDIX E: Instructions for the use of Supermon

TO USE SUPERMON (relocated version)

LOAD "SUPERMON", DEVICE, 1
NEW
SYS49152

## GENERAL NOTE

On entering Supermon it will save the stack which is restored on exit. It further changes the BREAK vector so when a BRK is met in a program Supermon is called.

All values are entered in hex. Only in ASSEMBLER mode do they have to be prefixed with ' $\$$ ' and then only for the operand.

Once Supermon has been loaded, it is resident until the 64 is either turned off or a program loaded which uses memory from \$cu00.

## INSTRUCTIONS

A - ASSEMBLER - Allows simple assembly of machine code.
A \$START OPCODE OPERAND.
For example, A 8000 LDA \#\$0A
Supermon will prompt with the next address.
Entering <RETURN> after address will exit assembler mode.
Branches are written with the destination address and not its displacement, that is, BEQ \$C456.

D - DISASSEmbler - Disassembles 22 instructions from the address specified.

D \$Start
for example, D 8000
Hex values may be changed by overtyping and on <RETURN> the same locations will again be disassembled.
Typing D on the bottom line will disassemble the next 22 instructions.
Typing <SPACE> <RETURN> will exit disassembler mode.
F - FILL MEMORY - Fill an area of memory with a specified value.
F \$FROM \$TO BYTE
for example, F5000 6000 FF
Useful to set up defaults prior to assembly, in particular to fill with NOPs (\$EA).

G-GORUN - execute machine code.
G - Starts execution at address currently in the Program Counter Register (PC).

G \$Start - Starts execution at specified address.
H - HUNT - search memory for specified bytes.
H \$start \$to data
for example, H50006000 'READ - Hunts for ASCII string "READ"
H50006000 A90A - Hunts for LDA \#\$0A
A maximum of 32 bytes may be set.
L - LOAD - Loads a program at its secondary address, leaving BASIC pointers unchanged.

L "filename" ,DEVICE - Device in hex.
08 disk 01 tape
M-MEMORY DISPLAY - Displays hex values.
M \$FROM $\$$ TO
for example, м 0801, 0891
11 Bytes may be overtyped to change.
$\mathbf{P}$ - PRINT DISASSEMBLY - Output hard copy of disassembled listing.
If in Supermon then exit (see below) and set up printer as for normal listing. Re-enter Supermon with subsequent output being directed to the specified device.

For example OPEN4,4:CMD4:SYS49152
P\$FROM \$TO
When complete, exit Supermon and close printer channel.
R - REGISTER DISPLAY - Displays current register values. This displays the PC, IRQ, Status Register (SR), A, X, Y and Stack Pointer (SP).

R
Values can be overtyped to change. This is of particular use in debugging operations where any of the registers may be altered and program execution continued with a co command.

S - SAVE - Saves an area of memory to tape or disk.
S "filename", DEVICE,\$START,\$END
Saves from the start up to, but not including the end address.
For example, S "NAME", $08,5000,6001$ - Saves from $\$ 5000$ to $\$ 6000$, but not the byte at $\$ 6001$.

T - TRANSFER - Transfers an area of memory to another leaving the original intact.

T \$FROM $\$$ TO \$START
for example, T 500060001000
You can also use Memory display this way.
This option may be used in conjunction with the relocator to generate versions of Supermon for use at other locations.

X-EXIT SUPERMON
X
The stack saved when Supermon was entered will be restored. A CLR from BASIC should tidy up any stack problems.

## COPYING SUPERMON

Use the save command as normal with the following addresses:
SUPERMON \$C000 \$C900

## APPENDIX F: Extended basic memory map

The following gives the main entry points for the UTILITY:

| Address | Description |
| :--- | :--- |
| $\$ 8000$ | Initialize Extended BASIC |
| $\$ 800 \mathrm{~F}$ | Set up Keyword Vectors |
| $\$ 8034$ | Set Top of Memory |
| $\$ 8041$ | Set NMI and BRK Vectors |
| $\$ 8054$ | Set Keyboard Table Set-up Vector |
| $\$ 8061$ | NMI Routine |
| $\$ 807 \mathrm{E}$ | BRK Routine |


| Address | Description |
| :---: | :---: |
| \$8090 | Keyword Vector Table |
| \$80F6 | Keyword Table - Command Keywords |
| \$81C7 | Keyword Table - Function Keywords |
| \$81F5 | Routine-GET PARAMETER |
| \$81FB | Switch off basic |
| \$8202 | Switch on basic |
| \$8209 | CRUNCH Tokens |
| \$82BC | PRINT Tokens |
| \$8302 | Token DISPATCH - Command Keywords |
| \$8329 | Token DISPATCH - Function Keywords |
| \$8352 | Perform COLOUR |
| \$8381 | Perform PLOT |
| \$83A7 | Perform WRITE |
| \$83AD | Perform ENTER |
| \$83B3 | Perform DOKE |
| \$83D7 | Perform DEEK |
| \$8401 | Routine - Convert to Positive |
| \$8415 | Perform OLD |
| \$842E | Perform AUTO |
| \$847F | Routine - Convert to ASCII |
| \$84A0 | Perform TEN |
| \$84EC | Perform TWO |
| \$8537 | Perform HEX |
| \$85BF | Perform BIN |
| \$85FC | Perform MEM |
| \$8611 | Perform RESET |
| \$8631 | Perform POP |
| \$864D | Perform KEY |
| \$8722 | KEY Interrupt Routine |
| \$8799 | Perform OFF |
| \$87A7 | Perform MERGE |
| \$886F | Perform APPEND |
| \$888B | Routine to close up memory and rechain |
| \$8933 | Routine to open up memory and rechain |
| \$89C5 | Perform RENUM |
| \$8B93 | Perform CODER |
| \$8D3A | Perform TRACE |
| \$8D61 | Perform TROFF |
| \$8E52 | Perform DUMP |
| \$8F44 | Perform DELETE |
| \$8FAF | Perform CGOTO |
| \$8FB5 | Perform CGOSUB |
| \$8FD2 | Perform PROC |
| \$904E | Perform INKEY\$ |


| Address | Description |
| :--- | :--- |
| $\$ 9080$ | CHAIN routine |
| \$9169 | HIMEM/LOMEM routine |
| \$9181 | Perform HIMEM |
| \$918D | Perform LOMEM |
| \$91B7 | Perform QUIT |
| \$9200 | Start up message |
| \$927D | Completion of DELETE |
| \$92B6 | Perform CHAIN |
| \$92DA- |  |
| \$9FFF | unused (expansion for sound/graphic/disk) |

## Appendix G: Reading an assembler listing

The machine code routines in this book have been presented in two formats. The first was generated using Supermon which is given in Appendix D and instructions for its use in Appendix E. The second was produced using Supersoft's mikro assembler cartridge. This appendix deals with listings generated using MIKRO as we feel they require some explanation.

## PSEUDO-OPS

These are instructions to the assembler and are not directly executable op-codes.

$$
\star=\$ C 000
$$

This tells the assembler to start its assembly at address \$c000.

> WOR, BYT, and TXT

These instructions reserve bytes in memory. Both WOR and BYT may be followed by any number of arguments separated by commas up to the limit of two screen lines. WOR reserves two bytes and is used to store absolute addresses in low/high format.

WOR \$C000,\$0100
Puts the four bytes $\$ 00, \$ 00, \$ 00, \$ 01$ in four consecutive addresses. BYT reserves single bytes.

## BYT \$A9,\$FF

TXT is followed by a quoted string and places the hex values of the ASCII codes sequentially in memory.
TXT "ABCD"

Puts $\$ 41, \$ 42, \$ 43$ and $\$ 44$ into memory.

Each of these directives allocates bytes from the address at which it appears.

## LABELS

Labels are used to identify an absolute address in memory. They are normally used as the destination for branches and jumps. They may also be used as operands.
LDA \#\$00 BYTE = \#\$FF
BEQ ZERO or LDX BYTE

## ZERO RTS

The absolute value of an address may be divided into low/high format by the use of '\#<' and '\#>' operators.

$$
\star=\$ C 000
$$

START LDA \# < START
LDX \# > START
This loads $A$ with $\$ 00$ and $X$ with $\$ C 0$.
Simple numerical calculations may be performed.
STORE BYT \$00,\$FF
LDA STORE
LDX STORE+1
This loads A with the value held in STORE which has been set to $\$ 00$ and $X$ with the byte from the next location, that is, $\$ F F$.

## ADDRESS TABLES

Where labels have been used their values, starting at an arbitrary address, have been given. This is useful to determine the hex values for all branches and jumps.

## LINE NUMBERS

The assembler code is entered exactly as one would type in a BASIC program. The same editing rules apply to a MIKRO program as to a BASIC program. Generally, we have retained these line numbers in the listings given for clarity and to aid description.

## COMMENTS

An exclamation mark is used in the same way as a REM from BASIC and
prefixs comment statements. It tells mikRO to ignore anything which follows it.

This is by no means the definitive mikro manual. We have limited ourselves to using only a few of the options available to allow easier conversion to other assemblers.

## Appendix H: Mnemonics generated by CODER

| [BLK] | - BLACK | [GR1] | - GRAY1 | [DEL] | - DELETE |
| :---: | :---: | :---: | :---: | :---: | :---: |
| [WHT] | - WHITE | [GR2] | - GRAYS2 | [INS] | - INSERT |
| [RED] | - RED | [LT GRN] | - LIGHT GREEN | [REV] | - REVERSE ON |
| [CYN] | - CYAN | [LT BLU] | - LIGHT BLUE | [OFF] | - REVERSE OFF |
| [PUR] | - PURPLE | [GR3] | - GRAY3 | [SPC] | - SPACE |
| [GRN] | - GREEN | [CLS] | - CLEAR SCREEN | $[\mathrm{G}>\mathrm{SPC}]$ | $\begin{aligned} & \text { - SHIFTED } \\ & \text { SPACE } \end{aligned}$ |
| [BLU] | - BLUE | [HOM] | - HOME CURSOR | $[G>$ ?] | $\begin{aligned} & \text { - GRAPHIC } \\ & \text { WITH } \\ & \text { SHIFT } \end{aligned}$ |
| [YEL] | - YELLOW | [CU] | - CURSOR UP | $[\mathrm{G}<$ ?] | $\begin{aligned} & \text { - GRAPHIC } \\ & \text { WITH LOGO } \end{aligned}$ |
| [ORG] | - ORANGE | [CD] | - CURSOR DOWN | [CTRL?] | $\begin{aligned} & \text { - CONTROL } \\ & \text { WITH } \\ & \text { LETTER } \end{aligned}$ |
| [BRN] | - BROWN | [CR] | - CURSOR RIGHT | [F?] | $\begin{aligned} & \text { - FUNCTION } \\ & \text { KEYS } \end{aligned}$ |
| [LT RED] | - LIGHT RED | [CL] | - CURSOR LEFT | [PI] | $\begin{aligned} & -\mathrm{PI} \\ & 3.1416 \end{aligned}$ |

MULTIPLE CHARACTERS are coded as [10CD]
SPACES
Single, unshifted spaces are not coded. We thought it unnecessary as it detracted from the legibility of the listing.

## SPECIAL CODES

The following is an extract from the Programmer's Reference Guide, page 74:

There are some other characters that can be PRINTed for special functions, although they are not easily available from the keyboard. In order to get these into quotes, you must leave empty spaces for them in the line, hit $<$ RETURN $>$ or $<$ SHIFT $><$ RETURN $>$, and go back to the spaces with the cursor controls. Now you must hit $<$ CTRL $><$ RVS/ON $>$, to start typing reversed characters, and type the keys shown below:


Functions 1 and 3 of the above are achieved as stated. CODER replaces them with:

```
[CRG>M] - SHIFT RETURN
[CRG>N] - SWITCH TO UPPER CASE
```

The other three can be achieved far more easily. Whilst PRINTing in quotes mode, press <CTRL> and the appropriate letter.

## Appendix I: Key codes

The following are the values assigned to keys in locations SFDX and LSTX (\$CB/203 \& \$C5/197) :

| dec | key | dec key | dec key |  |
| :--- | :--- | :--- | :--- | :--- |
|  | INST/DEL | 22 | T | 44 |
| 1 | RETURN | 23 | X | 45 |
| 2 | CRSR R/L | 24 | 7 | 46 |
| 3 | F7 | 25 | Y | 47 |
| 4 | F1 | 26 | G | 48 |
| 5 | F3 | 27 | 8 | 49 |
| 6 | F5 | 28 | B | 50 |
| 7 | CRSR U/D | 29 | H | 51 |
| 8 | 3 | 30 | CLR/HOME |  |
| 9 | W | 31 | V | 52 |
| 10 | A | 32 | 9 | 53 |
| 11 | 4 | 33 | I | 54 |
| 12 | 34 | J | 55 | $/$ |
| 12 | Z | 35 | 0 | 56 |
| 13 | S | 36 | M | 57 |
| 14 | E | 37 | K | 58 |
| 15 | None | 38 | O | 59 |
| 16 | 5 | 39 | N | 60 |
| 17 | R | 40 | + | 61 |
| 18 | D | 41 | P | 62 |
| 19 | 6 | 42 | L | 63 |
| 20 | 43 | - | 64 | RUN |
| 20 | C |  |  |  |
| 21 | F |  |  |  |

The following are the values of the shift registers SHFLAG and LSTSHF (\$028D/653 and \$028E/654):

| dec key pattern |  |
| :--- | :--- |
| 0 | NO SHIFTS |
| 1 | SHIFT |
| 2 | LOGO |
| 3 | SHIFT AND LOGO |
| 4 | CTRL |
| 5 | CTRL AND SHIFT |
| 6 | CTRL AND LOGO |
| 7 | CTRL, SHIFT AND LOGO |

## Appendix J: Summary of UTILITY commands

APPEND "program name", device
As for merge except that the appended program is tagged on the end of the memory program. Line numbers are not altered. Peculiar listings can be the result. Use renum after an append.

AUTO first line number, increment
Automatic line numbering when entering code.
BIN 8 bit binary number[,....
Prints out decimal conversion of binary number in two forms. The first as a low byte conversion and then, separated by an oblique, the high byte conversion (low $\star 256$ ). The binary number must be of eight bits.

CGOTO variable, calculation or line number
Line numbers can be mathematical equations.
CGOSUB variable, calculation or line number
Line numbers can be mathematical equations.
CHAIN ['filename"][,device]
Will load and run a BASIC program. It also transfers most variables from one program to another.

## CODER

Will replace non-standard ASCII and graphic codes with mnemonics. See Appendix H for full list.

COLOUR screen[,border][,text]
Values over 15 can be input, but only the lower four bits will be considered. Border and text parameters are optional.

DEEK(address)
Two byte PEEK. Returns memory location held in address and address+1.

DELETE first line to be deleted, [last line to be deleted]
Deletes lines in the range specified. No last line parameter, it will delete to the end of program.

DPROC name
Start of procedure called 'name'.
DOKE address, value
Two byte poke. Stores value ( $0-65535$ ) in address and address +1 .
DUMP
Displays the values of all simple variables currently in use.
ENTER ( $\mathrm{x}, \mathrm{y}$ )
Same as INPUT, but first sets cursor position as in PLOT.
EPROC
End of a procedure.
HEX hexadecimal number[,hex number][,...
Prints decimal conversion of hex input. The hex input can be of either two or four characters, but does not require a prefix of ' $\$$ '.

HIMEM address
Will set the top of memory to the given address, within the range of 1024 to 32767.

INKEY\$ [string or string variable]
Will wait for a key press. With no parameter, it will wait for any key. With a parameter, it will wait for a key to correspond to any character in the string. The ASCII value of the key press is placed in reserved variable 'ST'.

KEY 1 to 16 ,"data"
Loads function keys with data. Maximum of ten characters per key is permissible. Inputs over ten characters will generate a SYNTAXERROR, but the first ten characters of data will be assigned to the particular key. To generate a return in the data, use " $\leftarrow$ ".
for example, KEY $7, " L I S T \leftarrow "$
KEY
Will display the data assigned to all 16 keys in the format they were first entered. This will allow you to overtype the displayed data to amend the key assignations.

To obtain keys: KEY 1-8 as marked on keys
KEYS 9,11,13,15 key with logo
KEYS 10,12,14,16, key with shift and logo

Note: any key command will enable the keys if they have previously been disabled.

## LOMEM address

Will set the start of BASIC to the given address, within the range of 1024 to 32767 .

MEM
Display amount of memory free as an unsigned number.
MERGE ["program name"],[device]
Merges a stored program with that currently in memory according to their line numbers. Lines numbers of the merging program take precedence. If no program name and/or device then the command will default to tape. With no name then first program on tape will be merged.
OFF
Disable the function keys.
OLD
Restores a BASIC program after a NEW or system reset have been actioned. This will not work if an edit has been carried out before old is actioned.

## PLOT ( $x, y$ )

Sets the cursor column and row position. $x-\eta$ to 39 and $y-\emptyset$ to $24.0, \varnothing$ is the top left hand corner of the screen (cursor home).

POP
Rectifies stack on leaving a subroutine before a RETURN has been called.
PROC name
Calls a procedure called 'name'.
QUIT
Disables the utility and its commands, but protects the area that it uses. The UTILITY can be initialized again by SYS 32768 .

RENUM first line number to be changed or 0 , increment, new start line number
If first parameter is 0 , the whole program will be renumbered, otherwise, from designated line to the end of program. Renumbers the following tokens: GOTO,GO TO, GOSUB, IF THEN, RUN, ON GOTO, ON GOSUB and RESET. It will not renumber CGOTO or CGOSUB.

RESET [line number]
Restore dATA pointer to specific line or start of program.
TEN decimal number[,.....
Prints hex conversion of a decimal number.

TRACE
A diagnostic to follow the execution of a BASIC program as it runs.

## TROFF

Disables trace function.
TWO decimal number[,....
Prints binary conversion of decimal number.
WRITE ( $\mathrm{x}, \mathrm{y}$ )
Same as PRINT, but sets cursor position first as in PLOT.
Note: All commands performing number conversions will do more than one conversion if the values are separated by commas.

## ERRORS

These are particular to the Utility.

## CODER

STRING TOO LONG - more than 254 bytes have been generated by the mnemonics for one program line.
OUT OF DATA - found a character not handled by CORDER.

## RENUM

IILEGAL DIRECT - line at which to start renumbering does not exist. UNDEF'D STATEMENT - no destination found for a GOTO or GOSUB directive.

## APPENDIX K: 64 low memory map

The following is the first few pages of the memory map in the Programmer's Reference Guide (PRG), Chapter 5. It is included to avoid continual reference to the PRG to look up label addresses. Some of the descriptions have been changed through personal experience or preference to those in J. Butterfield's map.

| LABEL | hex | decimal | Description |
| :--- | :--- | :--- | :--- |
| D6510 | 0000 | 0 | 6510 direction register |
| R6510 | 0001 | 1 | 6510 I/O, memory and tape |
| ADRAY1 | $0003-0004$ | $3-4$ | Float to fixed vector |
| ADRAY2 | $0005-0006$ | $5-6$ | Fixed to float vector |
| CHARAC | 0007 | 7 | Search character |
| ENDCHR | 0008 | 8 | End of quote flag |
| TRMPOS | 0009 | 9 | Save screen last TAB |
| VERCK | $000 A$ | 10 | Flag: LOAD=0 VERIFY=1 |
| COUNT | $000 B$ | 11 | Ptrinput buffer/\#subscripts |
| DIMFLG | $000 C$ | 12 | Default DIM to 10 flag |
| VALTYP | $000 D$ | 13 | DATA type:string=255 numeric=0 |
| INTFLG | $000 E$ | 14 | $\quad$ :integer=128 float $=0$ |


| LABEL | hex | decimal | Description |
| :---: | :---: | :---: | :---: |
| GARBFL | 0005 | 15 | DATA scan/LIST quote/garbage collection flag |
| SUBFLG | 0010 | 16 | subscript/user fn call |
| INPFLG | 0011 | 17 | $\$ 00=$ INPUT $\$ 40=$ GET $\$ 80=$ READ |
| TANFLG | 0012 | 18 | TAN sign/comparison |
|  | 0013 | 19 | current l/O prompt |
| LINNUM | 0014-0015 | 20-21 | integer value |
| TEMPPT | 0016 | 22 | pointer:temp string stack |
| LASTPT | 0017-0018 | 23-24 | last temp string address |
| TEMPST | 0019-0021 | 25-33 | stack for temp strings |
| INDEX | 0022-0025 | 34-37 | utility pointer area |
| RESHO | 0026-002A | 38-42 | product area for multiply |
| TXTTAB | $002 \mathrm{~B}-002 \mathrm{C}$ | 43-44 | pointer start of BASIC |
| VARTAB | 002D-002E | 45-46 | pointer start of variables |
| ARYTAB | 002F-0030 | 47-48 | pointer start of arrays |
| STREND | 0031-0032 | 49-50 | pointer end of arrays |
| FRETOP | 0033-0034 | 51-52 | pointer bottom of strings |
| FRESPC | 0035-0036 | 53-54 | utility string pointer |
| MEMSIZ | 0037-0038 | 55-56 | pointer highest address used by BASIC |
| CURLIN | 0039-003A | 57-58 | current BASIC line number |
| OLDLIN | 003B-003C | 59-60 | previous BASIC line number |
| OLDTXT | 003D-003E | 61-62 | BASIC statement for CONT |
| DATLIN | 003F-0040 | 63-64 | current DATA line |
| DATPTR | 0041-0042 | 65-66 | current DATA address |
| INPPTR | 0043-0044 | 67-68 | INPUT vector |
| VARNAM | 0045-0046 | 69-70 | pointer current variable name |
| VARPNT | 0047-0048 | 71-72 | pointer current variable data |
| FORPNT | 0049-004A | 73-74 | pointer variable for FOR/NEXT |
|  | $004 \mathrm{~B}-004 \mathrm{C}$ | 75-76 | Y-save/op-save/bASIC pointer save |
|  | 904D | 77 | comparison symbol accumulator |
|  | 004E-0050 | 78-83 | misc work area |
|  | 0054-0056 | 84-86 | jump vectors for functions |
|  | 0057-0060 | 87-96 | misc numeric work area |
| FACEXP | 0061 | 97 | FPACC\#1: exponent |
| FACHO | 0062-0065 | 98-101 | FPACC\#1:mantissa |
| FACSGN | 0066 | 102 | FPACC\#1:sign |
| SGNFLG | 0067 | 103 | pointer series evaluation constant |
| BITS | 0068 | 104 | FPACC\#1: overflow digit |
| ARGEXP | 0069 | 105 | FPACC\#2: exponent |
| ARGHO | 006A-006D | 106-109 | FPACC\#2:mantissa |
| ARGSGN | 006 E | 110 | FPACC\#2:sign |
| ARISGN | $006 F$ | 111 | sign comparison result |
| FACOV | 0070 | 112 | FPACC\#1:low order rounding |
| FBUFPT | 0072-0072 | 113-114 | pointer cassette buffer |


| LABEL | hex | decimal | Description |
| :---: | :---: | :---: | :---: |
| CHRGET | 0073-008A | 115-138 | subroutine: get next byte of BASIC |
| CHRGOT | 0079 | 121 | entry point to get same byte |
| TXTPTR | 007A-007B | 122-123 | pointer current byte of BASIC |
| RNDX | 008B-008F | 139-143 | RND seed value |
| STATUS | 0090 | 144 | KERNAL I/O status ST |
| STKEY | 0091 | 145 | switch: STOP and RVS keys |
| SVXT | 0092 | 146 | timing constant for tape |
| VERCK | 0093 | 147 | LOAD $=\emptyset$ VERIFY $=1$ |
| C3PO | 0094 | 148 | serial output: deferred char flag |
| BSOUR | 0095 | 149 | serial output deferred char |
| SYNO | 0096 | 150 | tape EOT received |
|  | 0097 | 151 | register save |
| LDTND | 0098 | 152 | how many open files\# |
| DFLTN | 0099 | 153 | input device (default=0) |
| DFLTO | 009A | 154 | output device ( default=3) |
| PRTY | 009B | 155 | tape char parity |
| DPSW | 009C | 156 | tape byte received flag |
| MSGFLG | 009D | 157 | BASIC mode flag $\$ 00=$ program $\$ 80=$ direct |
| PTR1 | 009 E | 158 | tape pass 1 error log |
| PTR2 | 009 F | 159 | pass 2 error log |
| TIME | 00Ab-00A2 | 160-162 | real-time jiffy clock |
|  | 00A3 | 163 | serial bit count/EOI flag |
|  | 00A4 | 164 | cycle count |
| CNTDN | 00A5 | 165 | tape sync countdown/bit count |
| BUFPNT | 00A6 | 166 | pointer tape I/O buffer |
| INBIT | 0047 | 167 | RS232 input bits tape wrt ldr/rd count |
| BITCI | 00A8 | 168 | RS232 input bit count tape wrt new byte/rd error |
| RINONE | 00A9 | 169 | RS232 start bit flag |
| RIDATA | 00AA | 170 | RS232 input byte buffer tape scan/counter/ldr |
| RIPRTY | 00AB | 171 | RS232 input parity tape wrt Idr length'rd checksum |
| SAL | 00AC-00AD | 172-173 | pointer tape buffer/scrn scroll |
| EAL | 00AE-00AF | 174-175 | tape end address/end program |
| CMPO | 00B0-00B1 | 176-177 | tape timing constants |
| TAPE1 | 00B2-00B3 | 178-179 | pointer start of tape buffer |
| BITTS | $00 \mathrm{B4}$ | 180 | RS232 out bit count/tape enabled=1 |
| NXTBIT | $00 \mathrm{B5}$ | 181 | RS232 next bit to send/tape EOT |
| RODATA | 00B6 | 182 | RS232 out byte buffer/rd char error |
| FNLEN | $00 \mathrm{B7}$ | 183 | Length current file name |
| LA | $00 \mathrm{B8}$ | 184 | Current logical file number |


| LABEL | hex | decimal | Description |
| :---: | :---: | :---: | :---: |
| SA | 00B9 | 185 | Current secondary address |
| FA | 00BA | 186 | Current device number |
| FNADR | 00BB-00BC | 186-187 | Ptr current file name address |
| ROPRTY | 00BD | 189 | RS232 out parity/tape rd input char |
| FSBLK | 00BE | 190 | tape \#blocks left to wrt/rd |
| MYCH | OUBF | 191 | Serial word buffer |
| CAS1 | 00C0 | 192 | Tape motor control |
| STAL | 00C1-00C2 | 193-194 | 1/O start address |
| MEMUSS | 00C3-00C4 | 195-196 | KERNAL setup ptr/tape temp address |
| LSTX | 00 C 5 | 197 | Last key pressed |
| NDX | 00 C 6 | 198 | \#characters in k/b queue |
| RVS | 00 C 7 | 199 | RVS char print flag $1=$ yes $\boldsymbol{\theta}=$ no |
| INDX | 00 C 8 | 200 | Ptr end of line for INPUT |
| LXSP | 00C9-00CA | 201-202 | Cursor row, col at start of INPUT |
| SFDX | 00 CB | 203 | Current key pressed $64=$ no key |
| BLNSW | 00CC | 204 | $0=$ blink cursor |
| BLNCT | 00 CD | 205 | Cursor countdown timer |
| GDBLN | 00CE | 206 | Character at cursor pos |
| BLNON | 00CF | 207 | Cursor blink flag on/off |
| CRSW | 00D0 | 208 | Flag: INPUT from screen or GET from keyboard |
| PNT | 00D1-00D2 | 209-210 | Ptr current start of screen line add |
| PNTR | 00 D 3 | 211 | Cursor col on above line |
| QTSW | 00 D 4 | 212 | Flag: $\boldsymbol{\theta}=$ cursor in edit mode else in quote mode |
| LNMX | 00 D 5 | 213 | Physical screen line length |
| TBLX | 00D6 | 214 | Current row where cursor lives |
|  | 00D7 | 215 | Last inkey/checksum/buffer temp data |
| INSRT | $00 \mathrm{D8}$ | 216 | \#inserts outstanding |
| LDTB1 | 00D9-00F2 | 217-242 | Screen line link table |
| USER | 00F3-00F4 | 243-244 | Ptr screen colour |
| KEYTAB | 00F5-00F6 | 245-246 | $\mathrm{K} / \mathrm{b}$ decode table vector |
| RIBUF | 00F7-00F8 | 247-248 | RS232 input buffer ptr |
| ROBUF | 00F9-00FA | 249-250 | RS232 output buffer ptr |
| FREKZP | 00FB-00Fe | 251-254 | Free zero page area |
| BASZPT | 00 FF | 255 | BASIC temp data area |
|  | 0100-010A | 256-266 | Float to ASCII work area |
|  | 0100-013E | 256-318 | Tape error log |
|  | 0100-01FF | 256-511 | Processor stack |
| BUF | 0200-0258 | 512-600 | System input buffer |
| LAT | 0259-0262 | 601-610 | Logical file table |
| FAT | 0263-026C | 611-620 | Device number table |
| SAT | 026D-0276 | 621-630 | Secondary address table |
| KEYD | 0277-0280 | 631-640 | Keyboard buffer |


| LABEL | hex | decimal | Description |
| :---: | :---: | :---: | :---: |
| MEMSTR | 0281-0282 | 641-642 | Start of BASIC memory |
| MEMSIZ | 0282-0283 | 643-644 | Top of basic memory |
| TIMOUT | 0285 | 645 | Serial bus time out flag |
| COLOR | 0286 | 646 | Current character colour |
| GDCOL | 0287 | 647 | Background colour under cursor |
| HIBASE | 0288 | 648 | Start of screen memory:page number |
| XMAX | 0289 | 649 | Size of k/b buffer |
| RPTFLG | 028A | 650 | Flag: 128=repeat all keys |
| KOUNT | 028B | 651 | Repeat speed counter |
| DELAY | 028D | 653 | Flag: shift/ctrl/logo key |
| LSTSHF | 028 E | 654 | Last shift pattern |
| KEYLOG | 028F-0290 | 655-656 | $\mathrm{K} / \mathrm{b}$ table setup ptr |
| MODE | 0291 | 657 | Flag: $0=$ disable shift keys $128=$ enable |
| AUTODN | 0292 | 658 | $\theta=$ scroll down enable |
| M51CTR | 0293 | 659 | RS232 control register |
| M51CDR | 0294 | 660 | RS232 command register |
| M51AJB | 0295-0296 | 661-662 | RS232 non-standard baud rate |
| RSSTAT | 0297 | 663 | RS232 status register |
| BITNUM | 0298 | 664 | RS232 bits left to send |
| BAUDOF | 0299-029A | 665-666 | RS232 Baud rate |
| RIDBE | 029B | 667 | RS232 index to end of input buffer |
| RIDBS | 029C | 668 | RS232 page number of start of input buffer |
| RODBS | 029D | 669 | RS232 page number of start of output buffer |
| RODBE | 029 E | 670 | RS232 index to end of output buffer |
| IRQTMP | 029F-02A0 | 671-672 | IRQ save during tape I/O |
| ENABL | 02 A 1 | 673 | RS232 enable/CIA 2 (NMI) interrupt control |
|  | 02A2 | 674 | CIA 1 timer A control log during tape I/O |
|  | 02 A 3 | 675 | CIA 1 interrupt log tape read |
|  | 02A4 | 676 | CIA 1 Timer A enable log tape read |
|  | 02A5 | 677 | Screen line marker |
|  | 02 A 6 | 678 | PAL/NTSC flag $0=$ NTSC $1=$ PAL |
|  | 02A7-02FF | 679-767 | Unused |
|  | 02Cb-02FE | 704-766 | Block 11 for sprites |
| IERROR | 0300-0301 | 768-769 | Vector: basic error message (\$E3B8) |
| IMAIN | 0302-0303 | 770-771 | Vector: basic warm start(\$A483) |
| ICRNCH | 0304-0305 | 772-773 | Vector:Crunch basic tokens(\$A57C) |
| IQPLOP | 0306-0307 | 774-775 | Vector:Print basic tokens(\$A71A) |
| IGONE | 0308-0309 | 776-777 | Vector: Start new basic line(\$A7E4) |
| IEVAL | 030A-030B | 778-779 | Vector: bASIC token evaluate(\$AE86) |
| SAREG | 030C | 780 | Save A register |


| LABEL | hex | decimal | Description |
| :--- | :--- | :--- | :--- |
| SXREG | 030 D | 781 | Save X register |
| SYREG | 030 E | 782 | Save Y register |
| SPREG | 030 F | 783 | Save status register |
| USRPOK | 0310 | 784 | USR function jump instrn (\$4C) |
| USRADD | $0311-0312$ | $785-786$ | USR address low/high form(\$B248) |
|  | 0313 | 787 | Unused |
| CINV | $0314-0315$ | $788-789$ | Vector:Hardware IRQ(\$EA31) |
| CBINV | $0316-0317$ | $790-791$ | Vector:BRK interrupt(\$FE66) |
| NMINV | $0318-0319$ | $792-793$ | Vector:NMI(\$FE47) |
| IOPEN | $031 A-031 B$ | $794-795$ | Vector:KERNAL OPEN(\$F34A) |
| ICLOSE | $031 C-031 D$ | $796-797$ | Vector:KERNAL CLOSE(\$F291) |
| ICHKIN | $031 \mathrm{E}-031 \mathrm{~F}$ | $798-799$ | Vector:KERNAL CHKIN(\$F20E) |
| ICKOUT | $0320-0321$ | $800-801$ | Vector:KERNAL CHKOUT(\$F250) |
| ICLRCH | $0322-0323$ | $802-803$ | Vector:KERNAL CLRCHN(\$F333) |
| IBASIN | $0324-0325$ | $804-805$ | Vector:KERNAL CHRIN(\$F157) |
| IBSOUT | $0326-0327$ | $806-807$ | Vector:KERNAL CHROUT(\$F1CA) |
| ISTOP | $0328-0329$ | $808-809$ | Vector:KERNAL STOP(\$F6ED) |
| IGETIN | $032 A-032 B$ | $810-811$ | Vector:KERNAL GETIN\$F13E) |
| ICLALL | $032 C-032 D$ | $812-813$ | Vector:KERNAL CLALL(\$F32F) |
| USRCMD | $032 E-032 F$ | $814-815$ | Vector:Warm start(\$FE66) |
| ILOAD | $0330-0331$ | $816-817$ | Vector:KERNAL LOAD(\$F4A5) |
| ISAVE | $0332-0333$ | $818-819$ | Vector:KERNAL SAVE(\$F5ED) |
|  | $0334-033 B$ | $820-827$ | Unused |
| TBUFFR | $033 C-03 F B$ | $828-1019$ | Tape I/O buffer |
|  | $03 F C-03 F F$ | $1020-1023$ | Unused |
|  | $0340-037 E$ | $832-894$ | Block 13 sprite data |
|  | $0380-03 B E$ | $896-958$ | Block 14 sprite data |
|  | $03 C 0-03 F E$ | $960-1022$ | Block 15 sprite data |
| VICSGN | $0400-07 F F$ | $1024-2047$ | Screen memory |
|  | $0400-07 E 7$ | $1024-2023$ | Visible memory |
|  | $07 F 8-07 F F$ | $2040-2047$ | Sprite block data pointers $0-7$ |
|  | 0800 | 2048 | Start of BASIC (TXTTAB-1) |

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